

LARGE-SCALE SENSOR/ACTUATOR NETWORKS AS DATABASES OF  
DYNAMIC HETEROGENEOUS ACTIONS

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## ABSTRACT

# LARGE-SCALE SENSOR/ACTUATOR NETWORKS AS DATABASES OF DYNAMIC HETEROGENEOUS ACTIONS

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Wireless sensor networks (WSN) provide a robust and versatile solution to distributed low-power computing over a range of diverse network topographies. One-way data aggregation from network to data center is traditionally the fundamental data acquisition feature for many wireless sensor networks. With the increasing prevalence of IoT devices and the large number of WSN application possibilities, network and nodal reconfiguration to suit changing environmental conditions and sensing needs via outbound commands is not only an attractive feature, but in many cases a growing necessity. The work of this thesis comprises the implementation of a web-based application allowing end users to identify and reconfigure sets or subsets of WSN nodes based on real-time network state information. This end-to-end network reconfiguration procedure relies upon an existing two-way cyber-infrastructure and a relational database to store and track network state information. Additionally, the streaming middleware solution, integral to this cyber-infrastructure, provides a tracking mechanism for network reconfiguration commands. This tracking ability along with the reconfiguration mechanisms

will not only allow end-users to enact behavioral network changes, but also provide the foundation for virtualized users to autonomously monitor and enact behavioral changes to the network in response to changing network or environmental conditions. The specific aim of this work is the design and implementation of the mechanisms upon which these features will be built.

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*To my family*

# Chapter 1

## Introduction

### 1.1 Wireless Sensor Networks

Transducers can be used to sense radiation and measure soil moisture, temperature, and other physical phenomena. Transducers are often built into devices referred to as wireless sensors, with wireless networking capabilities. Wireless sensor networks (WSNs) use wireless networking technologies to enable groups of wireless sensors to communicate with each other and connect to the Internet. Low power distributed computing, intelligent wireless networking, and sophisticated data management form the core of WSN technology. The integration of these technologies are enabling the development of modern applications like the Internet of Things (IoT).

Some collections of wirelessly networked sensing devices form a cyber-physical system. A cyber-physical system is one where computation and networking processes operate in conjunction with a physical system, such as a smart garden or an advanced industrial process. While there are a broad range of implementations and platforms that fall under the umbrella of WSN, they all have a particular set of design restrictions and offer fundamental baseline features. WSN devices generally lack sophisticated processing power, run on a finite power supply (typically batteries), and use a diverse range of network topologies. Due to the power constraints of many wireless sensor nodes, transmission ranges are often quite short, and networks require a base station with greater networking capabilities, enabling access to the rest of the Internet; this is typically accomplished via a cellular or satellite connection.

WSNs allow researchers, scientists, and analysts to embed transducers into the environment which manage and perform sampling duties on a diverse range of sensors. Additionally, actuators allow for changes to be imposed upon the environment at given scheduled times or in reaction to certain events. In a WSN, the data collected from the various devices is aggregated and routed through the network to a base station where it is accessible to applications or archived at a data center. Depending on the diversity of the data types, the sophistication of the sensors, or the high-level data management schema, there is a diverse range of operations such as data conversions, checksum error detection, and stream creation that must be performed. These operations, data conversion for instance, might be

performed at the sensor node, at the base station, or even at the data center.

### **1.1.1 Difficulties of Network Reconfiguration**

The continued persistent functionality of a WSN requires that it have a certain level of autonomy. Basic WSN operation requires data aggregation from the nodes at the data center by whatever infrastructure is in place. There is also, however, a need for direct network interaction allowing a user to observe and enact behavioral changes within subsets of a network or the network as a whole. Researchers rely upon WSN to provide the means for extended monitoring or experimentation that span lengthy periods of time; changing circumstances may require for users to change the configuration of WSN devices based upon changing experimental needs or environmental conditions. For example, if a researcher studying the effects of soil moisture at a particular geographic location predicts a rain event, he or she might want to reconfigure the behavior of the WSN devices to gather soil moisture readings at an increased sampling rate to achieve greater data resolution.

There are a number of different design implementations for heterogeneous wireless sensor network control and reorganization, such as TinyOS and MagnetOS. These employ systems that allow modification of the sensors' behavior. It can be difficult to perform reconfigurations at the device level when WSNs are in remote or difficult to access locations such as rural areas or dangerous industrial facilities. Providing a user interface for researchers which accommodates making

such modifications would expand the functionality of the network in potentially new and unique ways.

## **1.2 Statement of the Problem**

When physical access to a device is needed to perform device or network modification, this can place a significant burden on researchers by necessitating a large number of hours of travel time to remote locations or the traversal of dangerous environments to access the devices, such as in an industrial setting. The wireless nature of WSN systems offers a natural mechanism for users to communicate with the devices remotely.

Enabling users to remotely change the behavior of WSN nodes would allow researchers to rapidly tailor the behavior of the network to changing needs or environmental conditions. The creation of a web-based application which provides users with real-time network state information and the ability to enact changes would significantly reduce the amount of travel time, increase the amount of control a user has to rapidly reconfigure WSN nodes, and allow the mechanisms necessary to automate the reconfiguration process.

An agent is a software entity that can utilize the same network device reconfiguration capabilities that a person would have, though its actions can be defined by software that can respond to changes in network conditions. An example is an

agent which monitors the power consumption of WSN devices in the network and reconfigures the devices to conserve energy. For example, if a specific WSN node is low on battery voltage, an agent might send the node commands that lower a transducer sampling rate.

The specific objectives of this work are as follows:

- Implementation of an application which allows users to remotely identify and select sets of WSN nodes based on static and dynamic search criteria;
- Implementation of an interface allowing users to enact behavioral changes in individual or groups of nodes; and
- Generalization of the reconfiguration software such that agents can monitor and reconfigure network nodes.

**This thesis presents a system for both users and automated agents to reconfigure arbitrary subsets of nodes in large, complex collections of WSNs using (1) a database paradigm for selection of node subsets based on both static and dynamic data, and (2) an interface that allows both users and automated agents to identify subsets and issue valid reconfiguration commands.**

## **1.3 Summary of Contents**

Chapter 2 summarizes the academic literature discussing reconfiguration of sensor networks and the various platforms where reconfiguration has been implemented.



This chapter also discusses the challenges associated with reconfiguring wireless sensor networks and the design features that researchers have utilized to overcome these challenges.

Chapter 3 discusses the WSN platform and cyber-infrastructure named WiS-ARDNet.

Chapter 4 explains the end to end network reconfiguration software and execution via a relational database to track network state information.

In chapter 5, node reconfiguration use cases, software validation, and user access control are discussed in detail.

In chapter 6, node reconfiguration procedures and automated users are discussed in detail.

The final chapter summarizes the results of the research and software implementation. Additionally, opportunities for future work are discussed, followed by conclusions.

# Chapter 2

## Literature Review

### 2.1 Introduction

WSN technology has matured such that it is a core component of the developing Internet of Things (IoT). The IoT describes the way in which networking capabilities formerly restricted to more traditional computers are being extended to physically-embedded devices, allowing for the acquisition and transfer of data on a much larger scale. WSN are geographically clustered networked devices equipped with one or more transducers which can range from a simple temperature probe to an implanted medical device. In general, these devices have limited computational resources and operate on supplies with limited power and energy. Due to these limitations, developers need to consider the design constraints of WSN applications. Many WSN devices are constrained such that their computational resources may not support a complete operating system (OS), which have been

fundamental to traditional computing systems. The lack of a traditional OS may complicate the development or limit the functionality of an application.

The range of devices being added to the IoT which generate data grows daily. Even though extending the operational lifespan and overcoming computational limitations remains at the forefront of the challenges facing WSN technology, numerous developments are enabling greater WSN functionality. Some of the features being developed within the domain of WSN are intelligent routing and communication protocols, fault tolerance and network health monitoring, service oriented design, and application execution frameworks. In recent years, developers and researchers have streamlined many of these features and have created new abstractions for WSNs. These features allow developers to overcome many of the hurdles facing modern WSNs.

## **2.2 WSN Execution**

The functionality necessary for WSN devices includes two critical features: distributed sensing and wireless data aggregation. WSN nodes have the ability to sample data from transducers and the ability to enact changes through the use of actuators. Ideally, researchers would be able to automate their experiments through an application program that interacts with individual or groups of WSN nodes. Clever use of abstraction to manage the complexities of the WSN will dictate how an application is executed; an application might run as an executable binary directly on a sensor node's hardware or perhaps within an abstraction of

the network at a higher layer in the form of scripted events or command sequences. WSN hardware is often extremely restrictive in its computational resources, and therefore application deployment and execution must be carefully designed.

Ivester et al. in [1] describe ISEE, an interactive execution framework for monitoring network services and specifying operation of applications executing on a WSN node using a design paradigm known as service oriented architecture (SOA). Network monitoring functionality is an important set of features, especially as experiments increase in complexity. The generalization of these features into modular services allows them to be used with greater versatility and independent of specific application code. Researchers who have developed other modern WSN platforms have also adopted the SOA paradigm of implementing WSN functionality. Hammoudeh et al. in [2] describe a SOA approach to WSN fault tolerance and inter-node cooperation.

Another category of WSN execution paradigms includes those that use a virtual machine (VM). A virtual machine emulates a particular hardware architecture by translating instructions to another hardware architecture, thus allowing software designed for a particular platform to execute atop a different platform. Two of the more prominent WSN reconfiguration approaches that utilize virtual machines are Maté [3] and MagnetOS [4]. Maté is a virtual machine that is created to overcome many of the challenges of implementing heterogeneous network execution on

resource-constrained wireless nodes. The fundamental feature of Mat   is a byte-code interpreter. Programs and applications, even those with minimal complexity, can be hundreds of kilobytes in size. The virtual machine takes high-level operational instructions and organizes them into capsules that can be smaller than 100 bytes in size. The Mat   byte code interpreter running on individual nodes allows for the execution of capsules that are disseminated throughout the network. This allows for a scalable approach for specifying network execution on a distributed system of resource constrained nodes. Similarly, the creators of MagnetOS use a virtual machine to specify network-level execution. These researchers use a version of the Java Virtual Machine (JVM) designed to appear as though it is operating on a single computer; however, the computer is actually a WSN. This abstraction of the WSN as a platform which can run the JVM allows for Java application code to be written by developers. The Java objects produced by the JVM can be sent throughout the network to different physical hosts, enabling the distributed execution of a single Java application. This is performed by a byte-code level translation into instructions that the nodes can execute. Mat   and MagnetOS can compress large programs into small device operations, due to the abstractions which they were able to make in their design. Though a complete virtual machine abstraction of a WSN is not utilized in the work of this thesis, the core principles of energy and resource aware network execution through abstraction is one of the primary design goals.

Flikkema et al. [5] describe the design and implementation of a cyber-physical system which utilizes an ultra low power sensing platform, streaming data middleware, and a control-oriented approach to performing complex ecological experiments. Several core principles from the previous works such as energy-aware network execution and abstracted application development through high-level programming languages have been designed into this platform. An important distinction between previously described works and the platform used in this thesis is that the WSN cyber-infrastructure is a network of networks. An individual cluster of nodes at a geographic location form a WSN, but behavior needs to be tracked and controlled across multiple WSN locations. This platform forms the foundation of the work of this thesis and will be described in greater detail in Chapter 3.

As the complexity of WSNs and diversity of applications increase, so increases researchers' need to reconfigure specific hardware components or WSN properties. A simple sensing network might only support unidirectional communication where the nodes pass their data to a base station. More sophisticated sensing networks support bi-directional communication where nodes both send and receive information. Bidirectional communication is a critical feature for any WSN system to support remote device reconfiguration and reprogramming. Depending on the size of a WSN supporting a particular application, a change could affect one node or even an entire cluster of nodes, depending on the type or magnitude of the change. For example, a researcher monitoring temperature fluctuations over a large geographic area with a cluster of WSN nodes might want to increase the sampling rate

to observe fine-scale environmental responses to a weather event. The difficulty of this task depends on how device reconfiguration was designed into the WSN platform and accompanying cyber-infrastructure.

Reprogramming and reconfiguration are often used interchangeably, as they might have the same effect on a researcher’s experiment, but there is an important distinction between them. According to some interpretations, like e.g. [6], if changes are made exclusively to the software components, then the device has been reprogrammed. Alternatively, if there is some change in the device hardware, the device is said to have been reconfigured. For the work of this thesis, we will define reprogramming as involving transfer of new source code or application software to one or more devices, whereas reconfiguration implies that software parameters or hardware functionality can be altered without the addition of new program code. Ivester et al. [1] describe the role of control information in a WSN as a means for reconfiguration. Control information describes commands or network state information that are injected into a WSN, causing change in one or more nodes. Software called Ismanage used in [1] accompanies ISEE by providing several features, the most important of which is the ability to disseminate control information into the WSN for the purpose of reprogramming and reconfiguration. This approach is similar to the reconfiguration methodology used in WiSARDNet, the platform used for this thesis.

Eronu et al. [6] summarize many of the issues and challenges surrounding WSN reconfiguration. Given that WSN hardware is often limited in power availability, the energy overhead involved in WSN reconfiguration is a major concern, especially when the number of nodes to be reconfigured increases [6]. The authors of [7] describe a wireless data collection protocol named ORiNoCo, designed specifically to utilize the existing bidirectional messaging capabilities of a WSN for the purpose of energy efficient network reconfiguration. Efficient network reconfiguration is also a core feature of the Maté and MagnetOS platforms discussed in the previous section [3], [4]. Nodes that run the Maté byte-code converter execute capsules which contain small control sequences. In this way, new configurations can be encoded and transmitted to a WSN node. The target node, upon receiving the encoded configuration, can decode the capsule with an intermediate software layer that translates the capsule's sequence into specific values to be stored in the device's memory. The reconfiguration of nodes in the MagnetOS platform allows for software changes to be instigated within an image that runs on the Java Virtual Machine. The virtual machine then interprets the image and sends the appropriate byte-code values to each hardware device in the network. Reconfiguration of these nodes then becomes no more difficult than executing a new application. We can see that when the sensing hardware is made available as a service to higher layers as in [1], and bidirectional communication protocols are leveraged effectively, network-wide reconfiguration of many nodes can be achieved with low energy overhead.



## 2.3 Relational Databases for WSN Applications

Abstractions and unique design paradigms can guide the development of complex systems in ways that are simple to understand. Storing information in a database enables access to specific data via query statements. A database abstraction can be used to simplify the complexities of data collection from a WSN [8], [9], [10]. [8] uses a WSN design paradigm where the data from each node can be aggregated at a data sink, and the sink acquires the data from each node through SQL-like queries. The nodes act like individual databases, and therefore the complexity of distributed data collection can be managed through simple queries. In this way, data is acquired on demand rather than continuously streaming towards the sink, effectively reducing the amount of data transmitted to only what is desired as opposed to all available samples. The authors in [10] utilize this approach for data acquisition as well.

Database abstractions provide data archival, organization, and access features that support a variety of applications. These features can greatly improve the efficiency and usefulness of WSN systems by providing scalable storage solutions across a variety of platforms. The division of data into multiple tables which can be managed separately and linked via foreign key constraints allows complex yet versatile access by experimenters and analysts. These ideas also pair well with modern software design paradigms such as behavioral parameterization [11], where

methods can be written in such a way that they define different behaviors to be decided at runtime. Behavioral parameterization is a powerful design paradigm that will allow WSN data to be easily accessed and utilized in the ever-changing landscape of IoT applications.

The aim of this thesis is the use of a database abstraction to simplify the complexity of network reconfiguration by storing network state information and other parameters that specify WSN execution, and utilizing this information to generate control messages for dissemination into the network for reconfiguration. Modular software design principles such as behavioral parameterization are used extensively in this thesis to accommodate rapid network reconfiguration as well as changing design requirements of future applications. This approach is described in further detail in Chapter 4.

# Chapter 3

## WiSARDNet Cyberinfrastructure

### 3.1 Overview

This chapter describes the platform on which the work of this thesis is built. WiSARDNet is a WSN platform which consists of hardware and software that connect sensor/actuator nodes to users and their ecological experiments. Cyberinfrastructure describes computational systems with advanced data acquisition, processing, and management capabilities, according to one of the more widely used definitions of the term [12]. According to this definition, WiSARDNet can be thought of as cyberinfrastructure. Each network component is described in this chapter, as well as the data management and archival processes. The cyberinfrastructure hardware and software components can be divided into four categories: WSN, base station, real-time data center (RTDC), and end-user. Figure 3.1 shows how the hardware and software components for each of these categories fit together. The

core component which links the WSN, the base station, and the real-time data center is a TCP/IP based message broker that uses the MQTT protocol. MQTT will be described in further detail in this chapter. Additionally, database archival components will be described in terms of their role in the cyberinfrastructure in this chapter. The technical specifics of the database will be described in greater detail in Chapter 4.

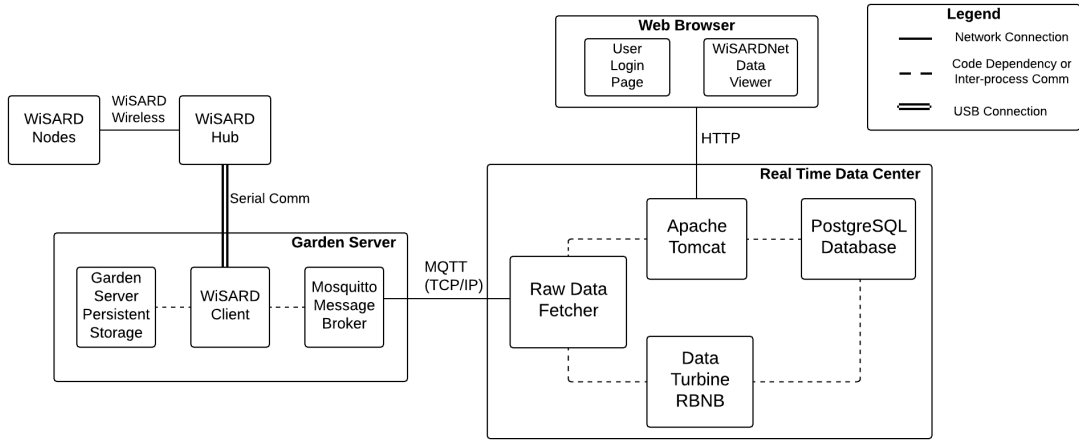


FIGURE 3.1: Cyberinfrastructure hardware and software component hierarchy

## 3.2 WiSARD Nodes

A Wireless Sensor/Actuator and Relay Device (WiSARD) is modular, adapting easily to the different sensing and actuation needs of users. Flexibility and energy awareness are the two driving design motivations in the WiSARD development. At the heart of each WiSARD device is a central processor (CP) which governs the operation of the device and its peripherals. The CP dictates when and how tasks will be scheduled and executed, establishes and manages wireless links to other

WiSARDs, facilitates the storage of sensor readings, and offloads the sensing and actuation tasks to daughterboards called satellite processors (SP). A WiSARD can accommodate up to four different SPs. This enables each WiSARD to perform a variety of tasks specifically tailored to meet the needs of researchers and experimenters. The modular design allows the CP to offload sensing operations to its SPs which can execute tasks in parallel with CP operation.

Both CP and SP boards use a Texas Instruments MSP430 16-bit ultra-low-power microcontroller. This microcontroller has a variety of features and can be placed in various low-power modes. The low-power modes of operation, accompanied by energy-conscious software design, allow for a WiSARD to operate on a single pack of three AAA batteries for many weeks or even months, depending on the rate at which the devices dispatch their sampling operations. The CP also connects to a radio board which uses an Analog Devices ADF7020 RF transceiver module that operates in the 902-928 MHz license-free ISM band. When WiSARDs are powered up, they autonomously form a self-organizing and self-healing multi-hop wireless ad-hoc network [5]. A special WiSARD which is assigned the software role of Hub acts as the base station, and is connected to an embedded Linux computer referred to as a garden server.

### **3.3 Middleware**

Middleware is a term which describes a software abstraction of the transfer of information from point A to point B. By hiding the complexities of data transport

behind a simple interface that connects two pieces of software, the development of powerful data processing and management tools can be accelerated. WiSARDNet relies heavily upon the use of middleware to make the WSN data available to the rest of the world. The central architectural component that connects the garden servers to the RTDC is MQTT. MQTT is a data transfer protocol created by IBM that utilizes the publish/subscribe middleware paradigm [13].

In order for MQTT to be explained in the context of this work, a few more terms need to be defined. This work is concerned with the reconfiguration of WiSARD nodes. Each WiSARD is comprised of multiple devices. A device is a singular piece of hardware; a transducer, a CP, a SP, and a radio are each an example of a device. A deployment is the placement of a device in a specific location. This set of operational parameters that control a deployed device is called the configuration of the device. Reconfiguration is the process of changing the operational parameters of one or more deployed devices.

MQTT facilitates the sharing of data between applications through the use of brokers. A broker is an application that accepts data from other applications. The data accepted is called a message. When a program sends a message to a broker, it labels the message with a topic name. The program sending the message is called a publishing client. Any program that is interested in data from a publishing client can subscribe to a topic through the broker; this is called a subscription client.

All messages from publishing clients that are labeled with a matching topic will be sent by the broker to all subscription clients that have subscribed to that topic.

WiSARDNet uses Mosquitto [14], an open source implementation of a message broker that uses the MQTT protocol. A message broker installed on each garden server keeps all of the data from that site in non-volatile memory. The data is organized into archive files at regular time intervals. To transfer the data from the WiSARDs to the broker, a Java application referred to as a WiSARD client uses MQTT messages to publish data. The RTDC can then retrieve the data using a subscription client that connects to the garden server message brokers over cellular or satellite connections.

An MQTT subscription client for each garden site runs at the RTDC and is responsible for subscribing to all the data published to each broker. These subscription clients are what is referred to in WiSARDNet as a raw data fetcher (RDF). The purpose of the RDFs is to retrieve all of the data from all of the gardens so that the data is accessible for processing, storage, or viewing. These processes are discussed in detail in the following section.

### 3.4 Real-time Data Center

The RTDC is a collection of compute and data servers that host a Tomcat Apache web server, a Mosquitto MQTT subscription client, an instance of an open-source data streaming middleware called Data Turbine, a PostgreSQL relational database, and all of the data management processes and applications which interact with the WSN. The RTDC has several functions, the first of which is to retrieve all of the WSN data from each of the gardens. The RTDC is a centralized destination for all of the network's data streams. From this location, applications can access, store, and modify data for their various needs, as opposed to having to interact with multiple networks individually. For instance, a process which reads in raw data streams from temperature sensors might need to calculate human-readable values from the raw sensor readings. A single data converting processor which moves data from raw transducer values to human-readable values, can easily acquire all of the raw data streams for temperature sensors at all network locations with a single fetch, rather than requiring that the process fetch the raw data from each specific garden individually. Additionally, aggregating all data at the RTDC greatly simplifies archival and backup procedures.

Another function that the RTDC performs is the execution of applications and processes which interact with the WSNs and their data. One example of an application running on a server at the RTDC might be an experiment which



attempts to maintain equal soil moisture readings at two different geographic locations. The RTDC servers possess high-performance computing hardware that many applications and processes can use to interact with the networks in real-time.

### **3.4.1 Relational Database**

Relational databases provide a flexible way to store and access large amounts of data. A relational database is composed of one or more tables; tables are 2 dimensional matrices of rows and columns. According to Rockoff [15], rows and columns are referred to as records and fields, respectively. An entry in a database table occupies a single row and each field corresponding to a different attribute of that entry constitutes a new column. For example, a table which stores people might have columns for an identification number, a name, an address, and a telephone number. An entry of a new person into this table would occupy a single row, with the data matching each of the attributes would occupy the rows' intersections with each column. Formally, these columns are referred to as primary keys, as they uniquely identify database records. The database at the RTDC has numerous tables which store information about every device, garden location, and sensor reading.

In WiSARDNet, data streams are archived in a PostgreSQL table. Other tables in this database store all of the meta-data and logistical information regarding

all WSN devices, as well as every data point sampled from all transducers. Additionally, there are many data streams that need to be tracked and archived such as diagnostic and control data, error reporting, and other useful information. An example of diagnostic data might be the logging of a device restart event; such events can be crucial in detecting and analyzing network performance issues.

The PostgreSQL database where all of this data is archived is physically located on its own server hardware at the RTDC, separate from the other data management and processing clients. By placing the database on its own server hardware, the database has exclusive access to dedicated processing resources to allow for the quickest possible query and insert times.

### **3.4.2 Data Schema**

WiSARD hardware components are subject to a design hierarchy that specifies how they interact. Figure 3.2 shows the hierarchy of the WiSARD hardware components.

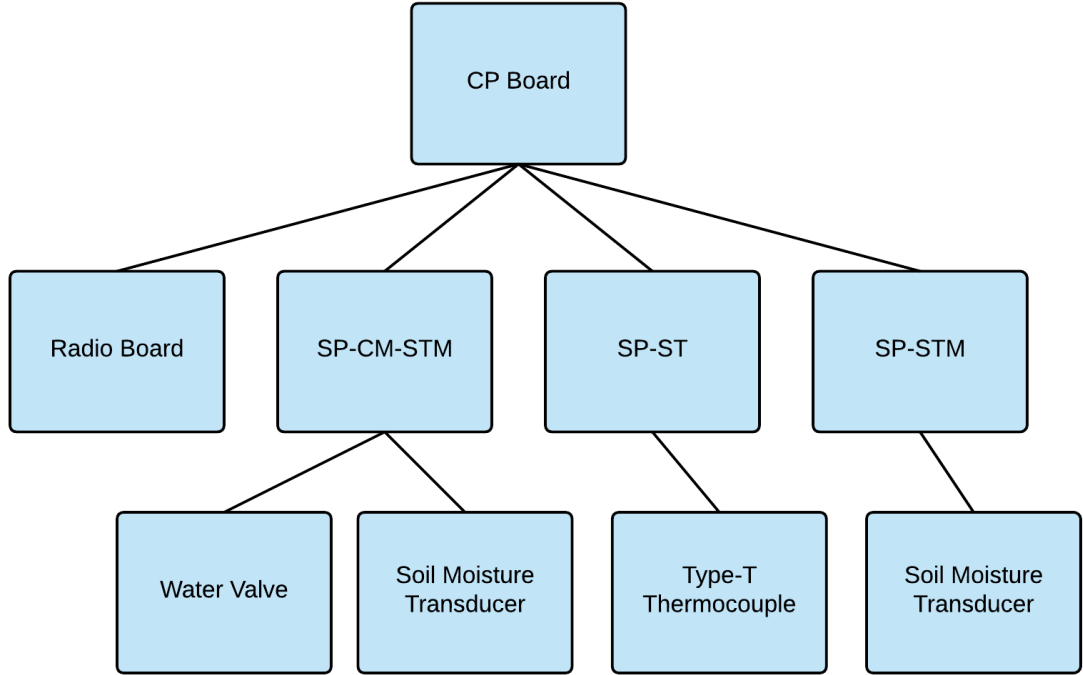


FIGURE 3.2: Device Relationship Structure

In a similar manner, the data gathered in WiSARDNet follows a design hierarchy that specifies how data is organized. The organizational structure of the the data tables and their relations, or the data schema as it is commonly referred to, was specifically designed to best accommodate changes to the configuration status of WiSARDNet devices. Every piece of hardware is referred to in the data schema as a device, and therefore the primary table in the data schema is the Device Table. A record in the Device Table will describe any device of any type. For instance, a radio, a soil-moisture transducer, a satellite processor, or a garden server are all devices that have a record in the Device Table. The configuration of a device might change at some point in time; it might be moved to another location, it might be upgraded or replaced, or it might receive a new firmware

version. In this schema, the device’s software parameters, hardware configuration, and state of operation are all encapsulated in a Deployment. Using the hierarchical structure of the WiSARD hardware, we use a tree structure in grouping devices together through their deployment records. Figure 3.3 shows how devices and deployments are structured in the database. Each deployment record has a field for a parent device deployment.

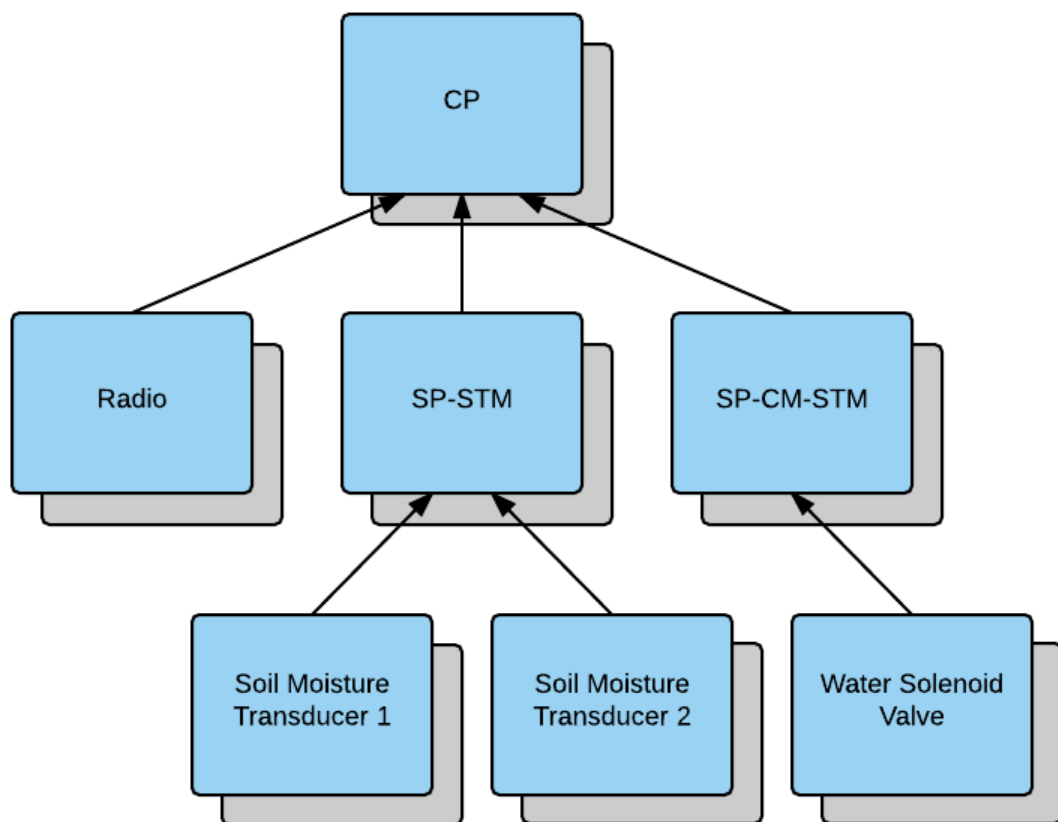


FIGURE 3.3: A visual showing each device (background) and its active deployment (foreground) and how the deployments are connected

When a transducer is attached to a satellite processor, then its active deployment record references the satellite processor’s deployment record in the parent deployment field. These are recorded and tracked through the Deployment Table

where each record is a particular deployment that references a device record in the device table. When a change is made to any of the parameters tracked in the deployment table, the device is considered as having a new deployment, and therefore a new record is inserted into the deployment table, and the device record is referenced by the new deployment. The key field in the deployment table is a binary field which stores a true/false value of whether or not the deployment is active. Once a deployment is changed, the old deployment becomes inactive and the new deployment becomes active. Over time, many deployment records may accumulate for a particular device as changes are made; the entire history of the device is stored so that for any given data point, the deployment configuration of the device associated with that data sample is accessible.

Having the ability to store and track each device's configuration and related meta-data in the database is of great value in the management of the networks and the various experiments. When a WiSARD is ready for deployment at a garden site, a deployment record is created with all of its configuration data, the deployment record is set to active, and the start date and time is set. As long as the device remains in this configuration, the deployment object pointing to that device will remain active. If a device is changed in any substantive way, the following actions are taken:

1. The deployment record referencing the changed device is set to inactive.
2. The current time is inserted into the stop-time field of the deployment record

3. A new deployment record is inserted into the deployment table
4. The parent field of the new deployment record references the deployment record of the parent device
5. The active status field of the new deployment record is set to active
6. The current time is inserted into the start-time field of the new deployment record

By following this procedure, device configurations are easily accessible from the database. A previous deployment exists as a single record in the Deployment Table which is trivial with regards to computational complexity and storage resources. This approach is simple, intuitive, and scales well as the number of devices deployed in the field increase. With this paradigm, a data sample is not merely a sample from a device, it is a sample from a specific deployment of a device. When a device's configuration is altered, a new deployment is created for that device and the data generated by that device references the new deployment record.

When new data from WiSARDs arrive at the RTDC from the MQTT broker, they need to be accessible for users and other services to request. Data Turbine's network accessible ring buffer data structure works well for this purpose. At the RTDC, the data streams arriving via MQTT are placed into Data Turbine's ring buffer in accordance with the stream name which identifies the device from which the sample was taken. Since the complexities of WiSARD configurations

are handled via storing their configuration data in database tables, archival of sampled data in the database becomes a simple procedure. All samples from all streams are placed as individual records in a table named Data. In addition to the sampled value and the time that it was sampled, the value can be related to the specific device and deployment records via its Data Turbine stream name. In this way, data archival of acquired samples is extremely simple and all information regarding a device and its samples is easily accessible and intuitively obtained through the thoughtful design of the data schema.

### **3.5 Summary**

WiSARDNet is a WSN CI which was designed from the ground up to be modular, scalable, and accessible to users. The Mosquitto MQTT broker, the data streaming middleware Data Turbine, and the RTDC enable data sampled by sensors to be retrieved, managed, processed, and archived into a PostgreSQL database. The way in which WiSARD meta-data and configuration information is stored and accessed is critical for the development of a network configuration software which comprises the work of this thesis. How these features are utilized in the network management software is described in further detail in Chapter 4.

# Chapter 4

## Network Reconfiguration

## Software

### 4.1 Overview

This chapter discusses the software developed in this thesis. First, this chapter discusses the procedures involved in network reconfiguration and the features required to achieve that goal. Next, the features are discussed in terms of the software functionality that needs to be implemented. Finally, the chapter discusses how the different functionality is grouped into software modules. The specifics of how each module works is described in Chapter 5.



## 4.2 Approach

To achieve network reconfiguration functionality within WiSARDNet, the following features are required:

- The user needs an up to date understanding of a WSN's configuration.
- The user needs to be able to specify the desired configurations or features they want to make to a set or subset of WSN nodes.
- Users need to be able to communicate configuration changes to the WSN.
- The WSN needs to execute the new behaviors that the user specifies.
- The system needs protections that ensure continued and correct WSN operation.

As was shown in Chapter 2, there are many ways that other WSN platforms have approached network reconfiguration. For the work in this thesis, a database abstraction is used as the general approach for structuring WSN reconfiguration. The specific abstraction used that the WSNs in WiSARDNet can be viewed as a database of dynamic heterogeneous behaviors. By viewing the WSNs from a database perspective, queries and inserts become the language that users can use to interact with the WSN. The database paradigm simplifies the complexity of WiSARDNet and guides the design of WSN reconfiguration.

In addition to the database abstraction, another abstraction is used to approach WSN reconfiguration. As described in Chapter 3, the components that comprise a WiSARD are represented in the WiSARDNet PostgreSQL database as individual devices with deployments that describe their current state. Since a WiSARD is a collection of these physical devices, it is useful to represent the device deployments as WiSARD objects. This is especially true given the number of database relationships that are managed, such as those to parent deployments, device types, sites. Figure 4.1 shows how a WiSARD can encapsulate multiple related deployments.

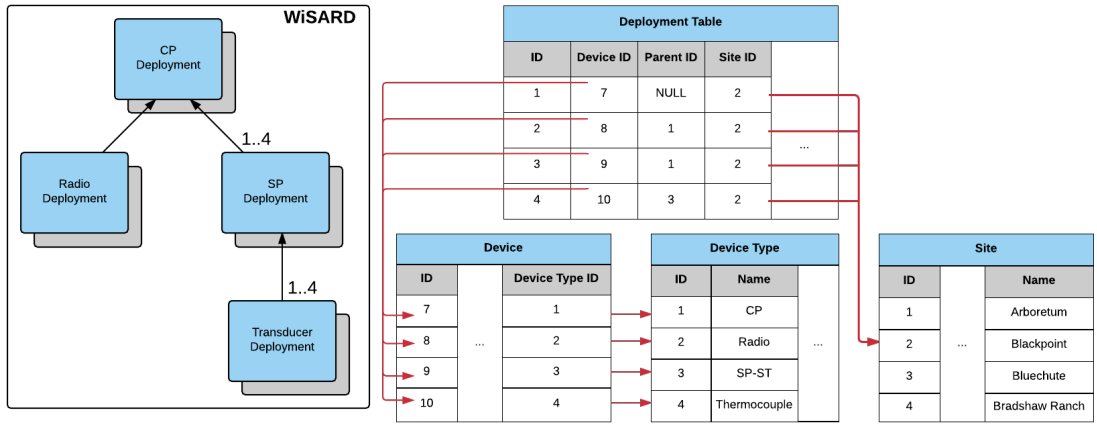


FIGURE 4.1: WiSARD abstraction and corresponding database representation

Using WiSARD objects as an encapsulation of different device deployments means that a WSN can be reconfigured by updating sets of WiSARDs. The database abstraction of the WSNs and the WiSARD abstraction of the WSN node hardware provide a conceptual foundation for the reconfiguration features to be implemented in software. The WiSARD abstraction is implemented in code as

a Java class. Each instance of the WiSARD class encapsulates all of the elements of a WiSARD object as shown in Figure 4.1.

## **4.3 Solution Description**

Implementing WSN reconfiguration functionality into an existing software system requires that new code modules be developed that account for all of the known use cases and as many future use cases as possible. A practical example is a scenario where a researcher wants all soil moisture transducers attached to WSN nodes in a specific region to be sampled at twice their current sampling rate. Performing this procedure involves the user specifying the region and sensor types needed to produce a matching set of WSN nodes, and that he/she wants the sampling rate of those transducers to be changed. The software needs to validate that there are no conflicts between that user's requested changes and the existing operational configurations of the node. Finally, the nodes will each be signaled via command packets, execute the specified change in configuration, and will then alert the user of the changes made.

### **4.3.1 Implementing Functionality**

Each of the features described in the previous section needs to be implemented as software functionality to be used in the WiSARDNet system. This section describes the software tasks for each feature.

#### **4.3.1.1 Identification of WSN Nodes**

The first stage in reconfiguring a WSN is the specification of the node or nodes whose configurations will be modified. This operation can be completed by querying the network state information stored in PostgreSQL database tables. Transducer samples, sensor operational data, software profile information, and hardware meta-data are all stored within the database and are necessary to determine which nodes will need to be reconfigured.

#### **4.3.1.2 Describing the New Configuration**

A WiSARD's task execution is governed by a set of configuration parameters stored within the device's non-volatile memory; this area of memory is called the task control block. Changing a device's behavior is achieved by overwriting specific values within the task control block with values that represent different behaviors. The user must describe the changes to the system so that the correct fields in the task control block can be overwritten with appropriate values corresponding to the new behaviors. This is achieved by encoding different behavior or behavior profiles into control sequences which the WiSARDs are able to decipher.

#### **4.3.1.3 Validating the New Configuration**

As the number of configuration changes and the number of affected WSN nodes increases, the possibility of errors or resource conflicts between experiments also increases. To prevent conflicts, there is a need to validate that the intended changes

will not interfere with other experiments each node is associated with. This is achieved with user access control and command validation.

User access control is the process of restricting the access of users to specific pieces of hardware or software based on a set of rules and permissions that govern the extent to which they can interact with the WSN. For example, a researcher should not be allowed to reconfigure another researcher's hardware or reconfigure the WSN in a way that will adversely affect other experiments.

Command validation is the process of analyzing the intended changes and the state of the network and determining whether or not the the command is feasible. User access control is the process of making sure the user or process issuing the command has the appropriate permissions or membership access to the affected hardware. Both procedures are essential in validating a command before execution.

#### **4.3.1.4 Command Synthesis**

Once the commands are created, the specified nodes within the WSN need to receive the desired changes. The command packets are added to MQTT messages and flushed from the MQTT broker at the data center to the destination garden server or servers. Each garden server has a WiSARD client which retrieves the command packet from the local MQTT instance and sends it into the WSN via the hub node.

#### **4.3.1.5 Executing the Reconfiguration**

When a WiSARD receives a command message from a parent node, the WiSARD parses the command message and schedules a task to execute the command. When a WiSARD executes a command, a report message is created and sent back through the network to the RTDC. The report message will confirm that the task was successfully executed or report that it failed if the task did not execute successfully.

#### **4.3.1.6 Verifying the New Configuration**

Report messages that arrive at the RTDC are treated as data streams, the same way sensor readings or diagnostic data are. Reports are archived and made accessible.

After making a change a WiSARD will report its status to the RTDC. That change is then reflected in the user's view of the network. The changes can then be compared against the user's selected command to verify that the reconfiguration was successfully performed.

### **4.3.2 Software Modules**

The software implementation of the functionality described in the previous section is grouped into three distinct modules:

- WiSARD browser module
- Command generator module

- Validator module

The WiSARD browser module contains the functionality that can identify sets of WiSARDs that match a user's search criteria. Additionally, this module is responsible for providing an up to date view of a WSN. The functionality which allows a user to describe a new WSN configuration and signal those changes to the WSN is grouped into the command generator module. The functionality which validates the safety of a new configurations, checks for conflicts, and implements user access control into the system is grouped into the validation module. Figure 4.2 illustrates where these new software modules fit into the existing WiSARDNet CI.

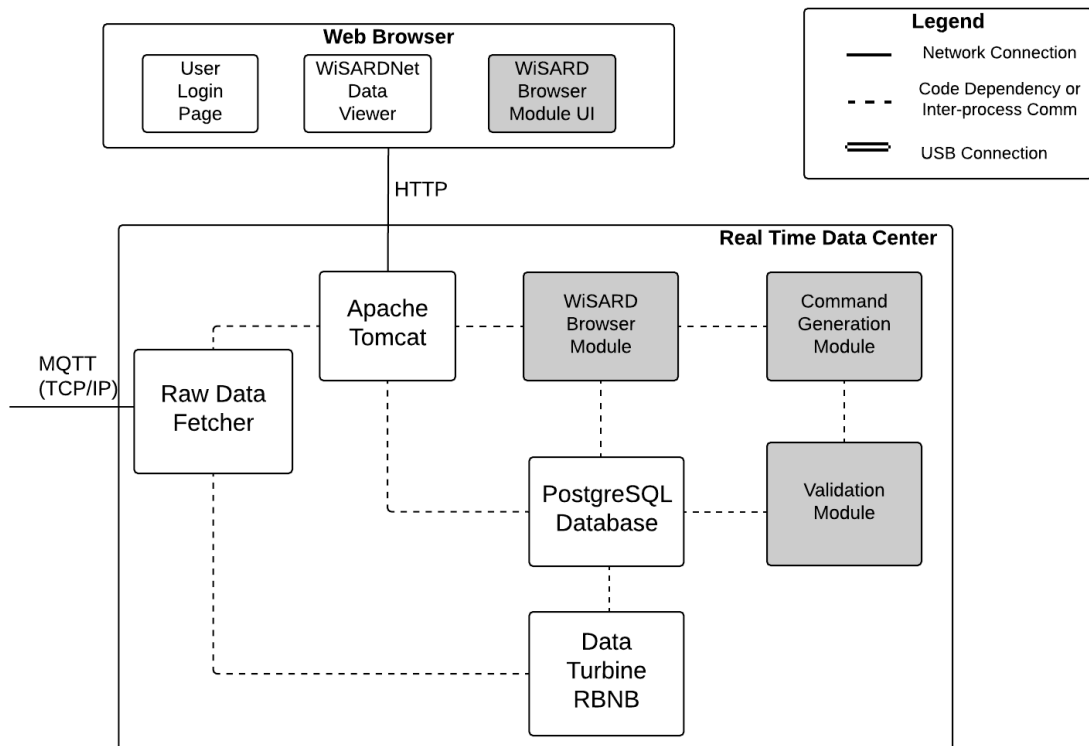


FIGURE 4.2: The modules being added to the WiSARDNet cyberinfrastructure

Each software module contains separate functionality, but they interact and share information with each other as well as the other components of WiSARDNet. The WiSARD browser module interacts with the web portal user interface, it gathers information about the WSN nodes from the PostgreSQL database, provides sets of WiSARDs to the other modules. The command generator module interacts with the WiSARDNet web portal to generate commands based on a user's input. Additionally, it accesses the list of WiSARDs that the WiSARD browser module produces. This module accesses the PostgreSQL database to enable stored behaviors to be turned into commands. Lastly, this module provides information to the validator module so that the commands can be verified. The validator module gets session variables from the web portal to authenticate users. It also takes in a set of commands from the command module, that can validate and report its results. To verify that the user is permitted to run a specified command, the user and command need to be compared against permission records stored in PostgreSQL database tables. These tables are discussed in detail in Chapter 5.

In summary, the functionality is grouped into three distinct software modules: WiSARD browser, command generator, and validator. These modules need to be able to interact with the different components of the WiSARDNet CI as well as with each other. The software implementations of the modules are discussed in Chapter 5.



# Chapter 5

## Software Module Implementation

### 5.1 Overview

This chapter explains how each of the modules discussed in the previous chapter are implemented in software. First, the WiSARD Browser is discussed. The command generator module will be discussed next, followed by the validation module. The validation module is discussed last, as it has the broadest impact on the other modules and the system as a whole.

### 5.2 WiSARD Browser

The WiSARD Browser module is a collection of software components which address four major objectives.

- Allow a user to specify WiSARD search parameters through a user interface

- Search the WiSARDNet PostgreSQL database for WSN state information
- Produce a set of WiSARD data objects matching the search criteria specified by the user
- The module should integrate with the existing WiSARDNet software

Figure 5.1 shows the flow of execution for this module. Each procedure shown is described in this section.

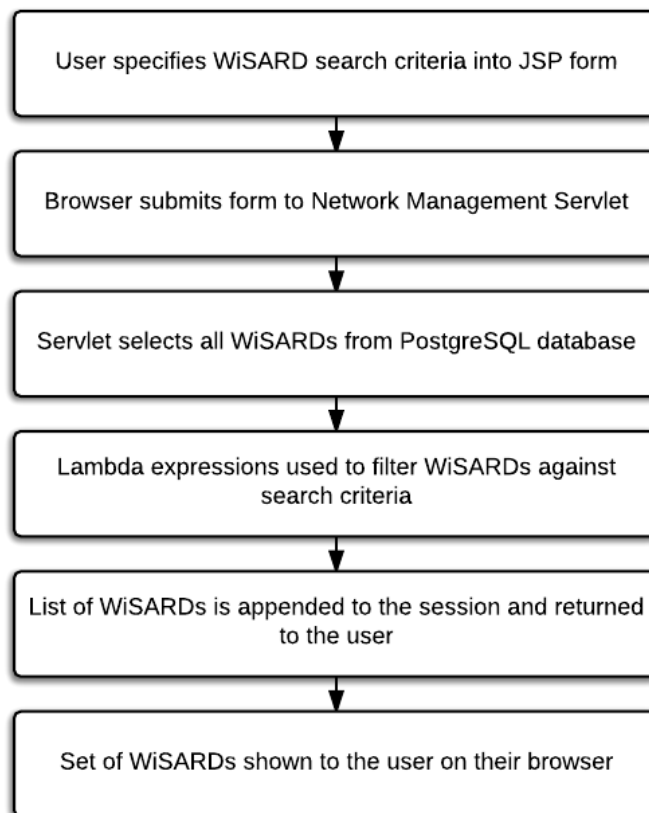


FIGURE 5.1: A flow diagram showing each action the WiSARD Browser Module makes

### 5.2.1 Obtaining the User's Search Criteria

A user of this system needs to have an interface with which to send and receive information. To create this interface, JSP (JavaServer Pages) and Java servlets are used. A JSP page can present information to a user through a web browser as well as gather user input necessary for back-end operations. The JSP page can send information back and forth with a servlet which is a java class that implements methods to respond to Get and Post requests from a web browser. Figure 5.2 shows how a servlet is used to deliver information from the users to the appropriate software module, and vice versa.

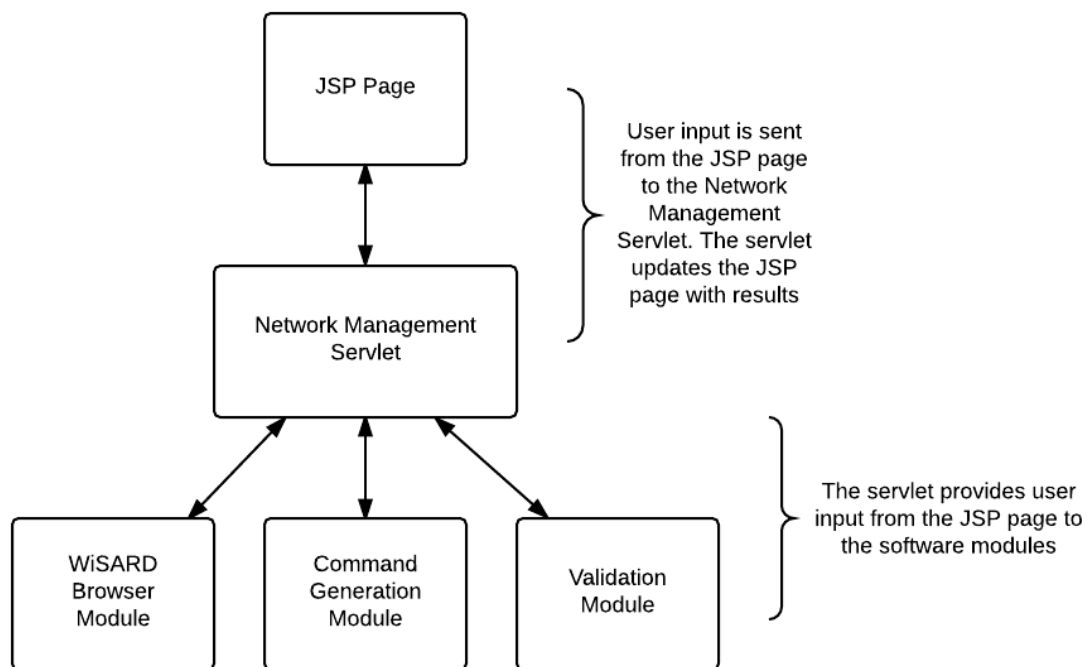


FIGURE 5.2: A depiction of the way in which the servlet relays information between the user interface and the core software modules

A JSP web page provides a simple interface for the user to populate an HTML form with attributes that can be used to select a set of WiSARDs:

- WiSARD Network ID
- Device Serial ID
- Garden Site
- Associated Experiment
- SP type
- Transducer type
- WiSARD role
- Deployment Type

These attributes are static or dynamic pieces of information that together represent the current state of the WSN nodes. Traditionally, this information can be obtained from a database query that can be written into a function or a method. With the functional needs of the WiSARDs changing, and the types of experimentation that WiSARDNet will perform still in development, it is important for this software to be flexible and support changes to the system. For example, if new WiSARD features are added, resulting in new fields in a PostgreSQL table, the new features should not prompt extensive modification of the existing code.

PostgreSQL supports a wide variety of query features that enable sophisticated data processing and filtration of queried data sets. If achieving the fastest possible query speed was of great importance for this module, then all of the searching

and filtration of WiSARDs from the user's search criteria should be handled by complex PostgreSQL queries. While query time is important, the benefits that can be gained at the sacrifice of a small amount of processing speed make other solutions desirable. Chapter 2 briefly explained some of the ways in which modern software development techniques such as behavioral parameterization can add value to software. For the WiSARD browser module, a functional data processing approach would not only allow for this module to meet the needs of the system, but would allow for development flexibility at the expense of some very minor performance overhead. This does not impact the performance of the system in a significant way is because the amount of WiSARD state information compared to the amount of data stored in the database is trivial.

WiSARD data objects translate well from design concept to a practical software entity. A Java class with member variables that correspond to the different attributes that define a WiSARD's state, and methods to translate higher-level queries into database queries and results provides a foundation for objects to be created and passed from module to module. To achieve the desired flexibility, the module performs a single fetch to the database and retrieves all of the state information for all of the WiSARDs. It then formats the information from the query into WiSARD objects before adding them to a list data structure. All of this functionality is built into the WiSARD class.

Once the servlet has obtained a list of all WiSARD objects, it can then filter that list against the attributes which the user specified in his/her form submission. Utilizing functional data processing techniques supported by lambda expressions in Java 8, obtaining a list that matches a user's search criteria is simple and flexible. Below is an example of the way in which lambda expressions are used in this software to produce powerful results with a very small amount of code.

```
return wisards.where((Wisard w) -> 'Arboretum'.equals(w.getSite()));
```

This statement returns an `ArrayList` named *wisards* that contains WiSARD objects. The *where()* method iterates over every entry in the `ArrayList` and only returns those who are not filtered out against the logic in the lambda expression. The lambda expression in this example compares the WiSARD site against a string containing the site name *Arboretum*. Any WiSARD in the `ArrayList` that does not have *Arboretum* as its site name will not be added to the `ArrayList` that is returned.

With one line of code, a data structure of WiSARD objects can be filtered based on its member variables that would otherwise have required a specific database query. Additionally, values can be added to the same expression to create more complex data filtration without requiring that the underlying functionality be altered. This functionality provides a quick and versatile way to support changing project requirements for managing networks of WiSARDs without being restricted by specific database queries.

When a list of WiSARDs has been filtered to match a user's search, the servlet appends it to the browser session where it is accessible to the JSP page. The JSP page displays the WiSARDs and their state information for the user. From this point a user can decide to continue with reconfiguring the WiSARDs that this search has returned, or return to the form and specify a different set of WiSARDs. Should the user choose to continue on with the set of WiSARDs they have selected, the data structure is accessible within the browser session, ready to send to the command module.

### **5.3 Command Generation**

The command generation module is similar to the WiSARD browser module, in that it relies on user input being sent from the JSP page to the servlet. The user will utilize the same JSP page's user interface as the WiSARD browser to select the change he/she would like to apply to the selected WiSARDs. Figure 5.3 shows the flow of execution for this module. Each procedure shown in the diagram is described in this section.

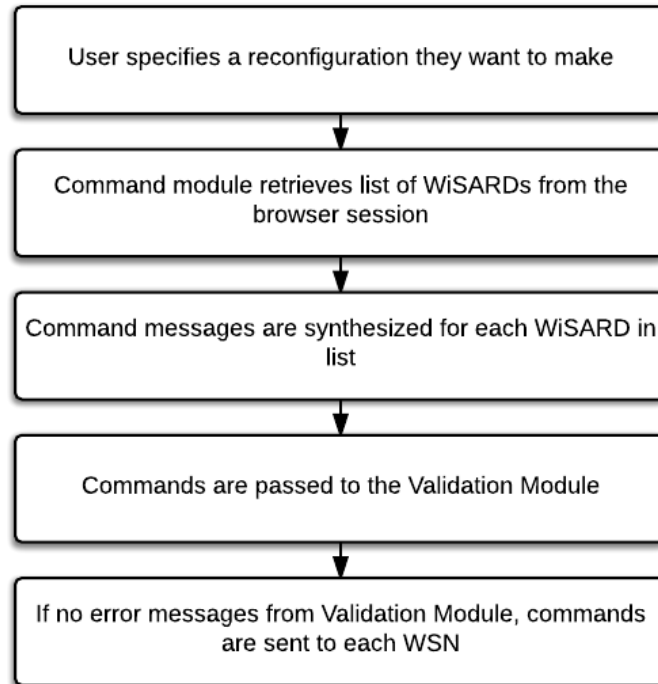


FIGURE 5.3: A flow diagram showing each action the Command Generation Module makes

To perform a reconfiguration in WiSARDNet, a specific command is generated and sent to each WiSARD that the change will affect. A command in the WiSARDNet communication protocol consists of multiple components. A hub address and a destination address describe which WSN the command should go to, and which WiSARD in that tree that the command is intended for.

These changes include a soft reset on the device, updating a sampling rate, changing a software role or ID, or modifying a task parameter for a transducer. The list of commands that are possible with a WiSARD continues to grow with each new experiment and feature that WiSARDNet supports. It is for this reason that commands need to be implemented in a flexible way such that the current



design will not present an obstacle to future developments. In WiSARDNet, a command consists of a payload, an expiration, and a checksum. The payload is the byte stream which contains the command and its parameters the WiSARD will interpret and execute. The expiration is a value which the sender can set to prevent out of date commands from being executed. For example, if a command for some reason takes an hour to traverse the network, it may no longer be relevant and should not be executed. The checksum allows the destination WiSARD to verify that it received the command correctly. A command class encapsulates the elements of a command and provides a foundation to build on for future work.

To perform a reconfiguration, a command object is generated for each WiSARD that the change will affect. The change that the user specifies will be enacted for all of the WiSARDs in the list that was added to the session by the servlet. The module loops through each WiSARD and chooses the correct command parameters to include in the payload. Once the command objects have been successfully generated they are sent to the Validation module in a data structure, similar to the way that the WiSARD browser module provided the WiSARD data structure to the Command module.

## 5.4 Validation

Before commands can be sent out to reconfigure WiSARDs, there are two sets of checks which the commands must pass. First, the Validator module confirms that

the user that submitted the commands has appropriate permissions to reconfigure the selected WiSARD nodes. Second, the module assess the commands for their impact on the system, and ensures that none of the changes will cause a WiSARD to misbehave. It is likely that certain WiSARDs will have hardware that is shared between multiple experiments, and a researcher reconfiguring that WiSARD could have adverse on one or more experiments. The procedures that this module performs are shown in Figure 5.4. Each of these procedures are described in this section.

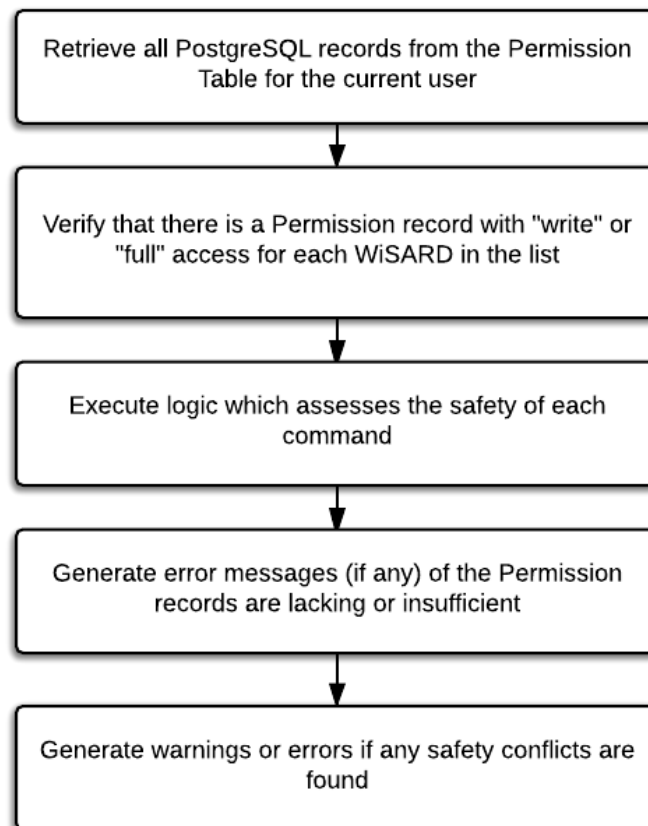


FIGURE 5.4: A flow diagram showing each action the Validation Module makes

### 5.4.1 Permissions and User Access Control

The PostgreSQL database schema prior to this thesis had certain tables which could have been used for user access control, but the schema was not robust. To implement user access control in a flexible way, a better approach was needed. Two core abstractions of the existing database schema are used to design the permissions system for WiSARDNet. The first abstraction is that anyone to whom a permission might be granted is considered a permission entity. The second abstraction is that anything which would need its access restricted is to be considered a permission resource. After applying these two abstractions to the database, the Person and Organization tables are defined as permission entities. Likewise, the Deployment, Site, and Experiment tables are defined as deployment resources. With these two abstractions, granting permission simply becomes a mapping between a permission entity and a permission resource with a value designating the level of access the entity has to the resource.

Three access levels are used to denote permissions to resources: read, write, and full. Table 5.1 describes what each permission level allows. An important note is that an entity with full access can override changes made by other entities with full access, but a warning message is generated to inform the user of the scenario. In this case, the user should proceed at their own discretion. In the greater context of the user access control paradigm, full permissions also grant a user delete privileges. Even though that use case doesn't directly correlate to the

reconfiguration module, the paradigm still supports that feature for other applications. These three permission levels are sufficient to regulate how users interact with the WiSARDs. This approach to user access control also provides a flexible way to manage other hardware or software resources which might be incorporated into WiSARDNet in the future.

Permission Level	Description
read	An entity can view a resource and any data associated it is associated with
write	An entity has access to modify and reconfigure a resource
full	An entity is given full authority to override the configurations of resources set by other entities

TABLE 5.1: A description of each access level an entity can be granted to a resource

To implement user access control with this approach, three tables were added to the database. First, a Permission Entity table was created to reference the Person and Organization tables. Next, a Permission Resource table was created to reference the Deployment, Experiment, and Site tables. Lastly, a Permission table was created to map records from the Permission Entity table to the Permission Resource table. The references between these tables are governed by foreign key constraints, or referential integrity constraints as they are sometimes referred. Figure 5.5 shows these tables and their foreign key constraints. Each foreign key

constraint references the primary key of the table that the arrow points. The unnamed table on the right shows how the schema can be used to manage new resources in the future.

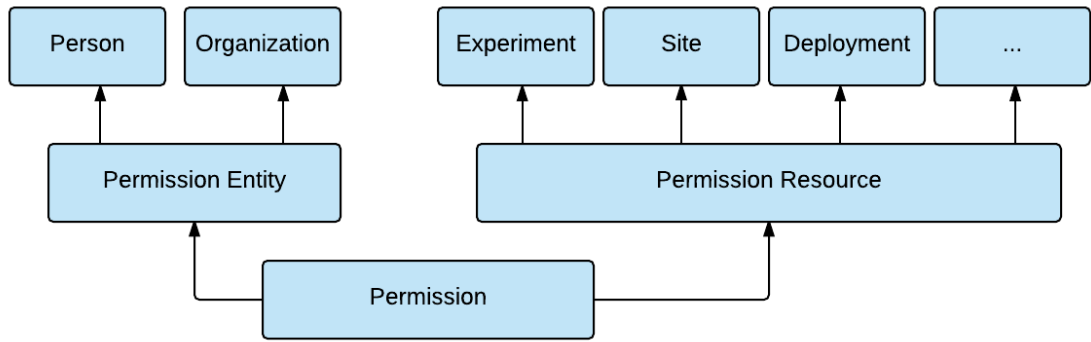


FIGURE 5.5: A diagram that shows how the three new Permission tables are used to create a user access control schema

When the validation module receives a data structure of command objects from the command generation module, the software in the validation module first checks that there are no user access violations in any of the commands. This is done by iterating over each command object and performing a lookup to the permission table. The objective is to gather all of the Permission records whose permission entity matches the user that created the commands, and verify that a permission resource with at least write access is present, for the WiSARD in question. If the user is found to have insufficient permissions to execute a command, an error message is generated and added to a queue of messages that is displayed once the validation module has iterated over every command.

### 5.4.2 Safety Validation

If all of the commands pass the user access permission checks, then the module performs a configuration safety test on each command. The difficulty of the safety check comes from the different types of commands that can be executed, and the diversity of experiments and their needs. It is for this reason that the command class was designed such that each command type will implement its own safety logic. This way, when a new command is developed, the designer can build in a method which will perform a safety check within the context of that method. Figure 5.6 shows the structure of these command classes and how they inherit and override behaviors.

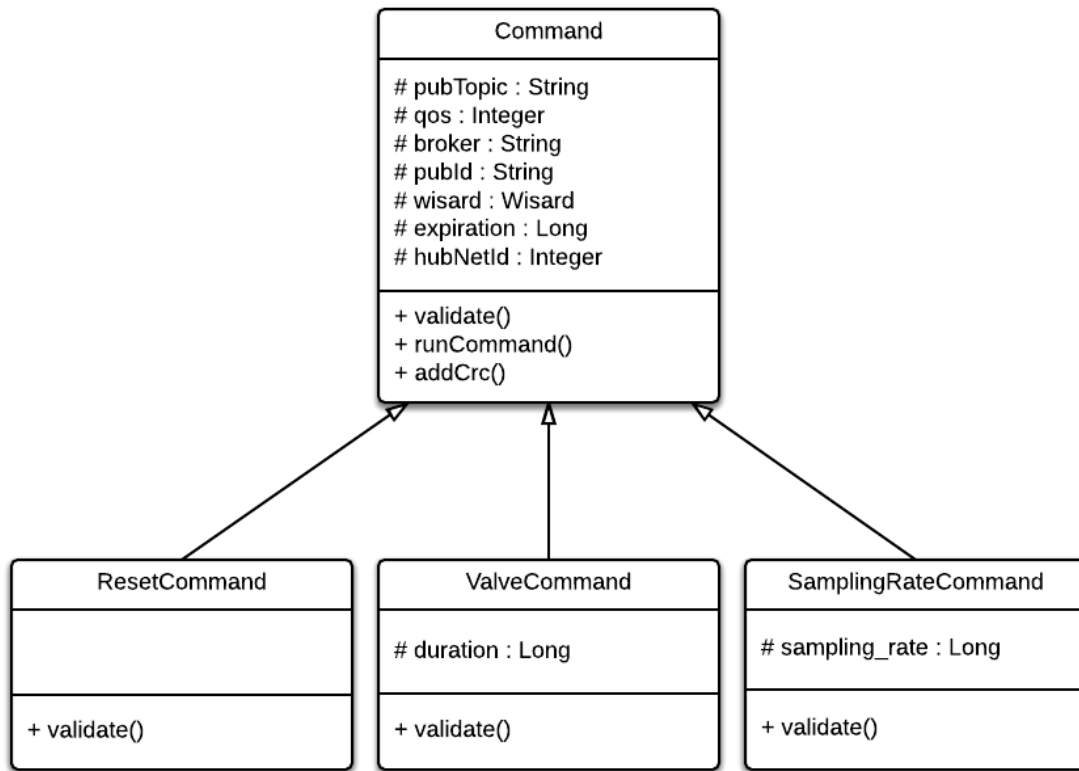


FIGURE 5.6: A class diagram showing how new commands inherit from their parent class, yet implement their own validate method

A practical example of a safety validation method would be a case where a particular command is designed to increase the sampling rate of a particular transducer type on a WiSARD, then an appropriate safety check would be to ensure that the new sampling rate will be a reasonable value. If the user enters a non-numeric value, a negative value, or a value corresponding to a sampling rate which the WiSARDs cannot physically perform, those should result in a failed safety check. Another example is a command to reconfigure a WiSARD that one or more experiments are using. If a configuration change could impact an experiment, a warning message is generated and added to a queue that is shown

to the user. In this scenario, the command is not rejected but the user will be warned of the conflict.

## 5.5 Summary

Each of the three software modules developed for this thesis serves a distinct role in enabling network reconfiguration. The WiSARD Browser is a flexible module that enables users to search for and select a group of WiSARDs based on a wide range of network state information. The command generator module allows for a user to select a change he/she would like to make to the set of selected WiSARDs, and automatically generate the commands necessary to make them happen. The validation module performs user access control and network safety checks on commands that users create, to ensure that all reconfigurations are made in a way that minimizes the potential for conflicts. The design and implementation of these three modules provides great flexibility to developers who will work on the WiSARDNet platform in the future.



# Chapter 6

## Automated Reconfiguration

### 6.1 Overview

This chapter discusses automated network reconfiguration and how it is supported by the work of this thesis. First, the concept of automated reconfiguration and the motivation for its use is described. Then the design and implementation of the automated reconfiguration software in the work of this thesis is described. Finally, an example is shown that demonstrates the practical uses of this software as a proof of concept.

### 6.2 Design

The work in this thesis allows network reconfiguration through user interaction. User interaction can sometimes be inadequate because of slow response time, a lack of understanding of system behavior, or human error. Creating a system that

can automatically respond to incoming sensor data (or calculations derived from incoming sensor data) by reconfiguring the network allows overcoming many of the limitations encountered by requiring user input.

The software applications designed in this work to autonomously monitor and reconfigure WiSARDs are called automated agents. Agents have three main requirements that they must fulfill. First, the agents must adhere to the user access and safety validations that govern the rest of the system. More specifically, the developers of automated agents should not be allowed to intentionally or unintentionally bypass, improperly implement, or override the security provisions. Second, they must be able to receive and handle data messages, including being able to handle software exceptions that may occur. Finally, the agents should be able to interact with the cyberinfrastructure using the WiSARD abstraction. For instance, the agent should be able to interact with the PostgreSQL database using the methods that the WiSARD class provides.

The automated agents are defined by a Java class designed to address these three requirements. The first requirement of preventing agents from circumventing system security is solved by separating the validation logic from the agent class. If the validation logic is external to the agent class, then agents will not be able to dictate their own security privileges. The second requirement of receiving and handling data streams and handling exceptions specific to each agent is solved by

the design decision that each agent must be provided message handling logic and exception handling logic specific to each agent. The final requirement of preventing an agent from directly accessing the PostgreSQL database is solved by using the WiSARD class discussed in Chapter 5 as a means of interacting with the database.

To accommodate these solutions, an agent controller class is used that encapsulates much of the functionality described in the previous chapters, primarily validation and safety. The controller class will log in an agent, validate that it has sufficient privileges to access the data streams it uses, and then create an agent object. The agent class implements the Java runnable interface, that allows it to be put into a thread and executed. The design and implementation of the agents is generalized, so that there is common code from agent to agent. Every agent uses the same logic to listen for and handle incoming messages. Message processing and error logic that distinguishes one agent from another are passed to the agent as parameters in the form of lambda expressions. A lambda expression is an anonymous function.

Using inheritance in object-oriented software design is typically a good way to reuse code. Inheritance has been used extensively in the code described in the previous chapters. However, in the case of designing the automation behavior, using composition over inheritance resulted in a cleaner design. Composition is the object-oriented design concept where the functionality of an object is defined

through the creation of member objects rather than through inheriting methods. This approach is accomplished with the use of two interfaces referred to as the Message Processor and the Error Handler. These interfaces are what allow the operational logic for each agent to be passed in as lambda expressions. Figure 6.1 is a class diagram that shows how these classes are implemented using a composition approach.

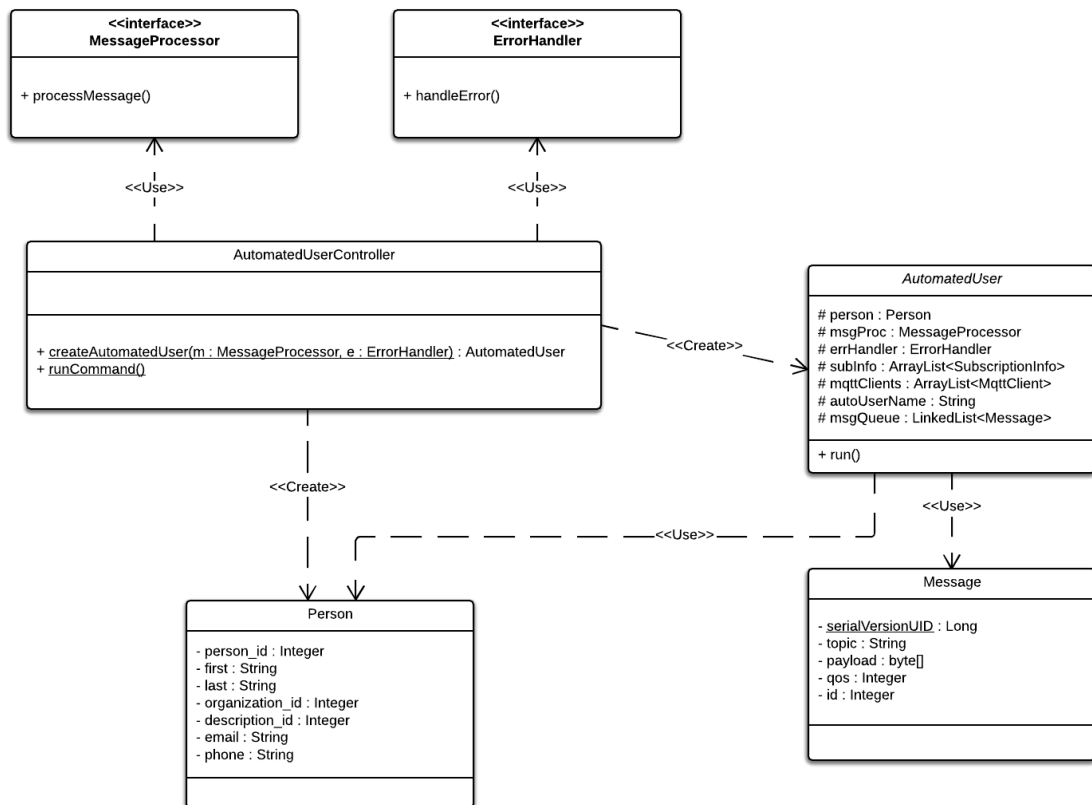


FIGURE 6.1: A class diagram showing the composition of the agents and their relevant classes

The text on each of the arrows denote a class using the item it is pointing to, or whether it is creating an instance of the item it is pointing to. The Automated

User Controller class uses the Message Processor and Error Handler interfaces to create Automated User and Person objects. Each Automated User object can then use its Person object, the functionality in the Message class, and whatever functions that were passed in for the Message Processor and Error Handler to complete its tasks. A detailed example of the software classes showing how this software actually works is described in the next section.

## 6.3 Example Agent

Using the design described in the previous section, there are many different ways that an automated agent can be used to manage WiSARDs. The following code snippets and class definitions demonstrate how this functionality is used in practice.

Below is the class definition for the agent controller that handles authentication and creation of the agents.

```
/*
 * class definition for automated user controllers which create, authenticate,
 * and return automated user objects
 */
public class AutomatedUserController{
    public static AutomatedUser createAutomatedUser(String username, String password, ArrayList<
        ↳ SubscriptionInfo> subscriptions, MessageProcessor msgProc, ErrorHandler errHandler){
        // logs in user, validates permissions to subscriptions, creates person and
        ↳ automated user
    }

    public static runCommand(NetManagementCommand nmc, AutomatedUser user){
        // validates user's permission to run command and performs safety validation
    }
}
```

This is the class definition for the agents. An instance of this class is created by the agent controller and returned. In the constructor, all of the necessary information to establish connections to MQTT message brokers as well as define the agent's behavior is passed in as a set of parameters. The run() method shows how an agent establishes its subscriptions with the MQTT brokers it will be listening to, and then waits for messages to arrive. Once messages arrive, they are given to the processMessage() method whose functionality was passed into the constructor as a lambda expression.

```

/*
 * class definition for automated users which encapsulates a person and message
 *   ↪ processor
 * objects, as well as an array list of subscription objects. This will allow
 *   ↪ automated
 * agents to monitor data streams and enact changes upon the WiSARDs they have
 *   ↪ sufficient
 * permissions for
 */
public abstract class AutomatedUser implements Runnable, MqttCallback{
    protected Person person;
    protected MessageProcessor msgProc;
    protected ErrorHandler errHandler;
    protected ArrayList<SubscriptionInfo> subInfo;
    protected ArrayList<MqttClient> mqttClients;
    protected String autoUserName;
    protected LinkedList<Message> msgQueue;

    public AutomatedUser(Person p, ArrayList<SubscriptionInfo> s, String autoUserName,
        ↪ MessageProcessor m, ErrorHandler e){
        this.person = p;
        this.msgProc = m;
        this.errHandler = e;
        this.subInfo = s;
        this.autoUserName = autoUserName;
        this.mqttClients = new ArrayList<MqttClient>();
        this.msgQueue = new LinkedList<Message>();
    }

    public void run(){
        try{
            // create connections for all subscription clients and add to an ArrayList
            for(SubscriptionInfo subscription : subInfo){
                MqttConnectOptions connOpts = new MqttConnectOptions();
                connOpts.setCleanSession(false);
                connOpts.setConnectionTimeout(180);
                MqttClient subClient = new MqttClient(subscription.broker, subscription.gsCommonName
                ↪ + "/" + autoUserName);
                subClient.subscribe(subscription.subTopic);
                mqttClients.add(subClient);
            }
        }
    }
}

```

```

    }

    // while thread is running, listen for messages
    while(!Thread.currentThread().isInterrupted()){
        synchronized(msgQueue){
            Message m = msgQueue.peekFirst();
            if(m != null){
                msgProc.processMessage(m);
                msgQueue.removeFirst();
            }
        }
    }
} catch(Exception e){
    errHandler.handleError(e);
}
}

@Override
public void connectionLost(Throwable arg0) {
    // empty
}

@Override
public void deliveryComplete(IMqttDeliveryToken arg0) {
    // empty
}

@Override
public void messageArrived(String arg0, MqttMessage arg1){
    synchronized(msgQueue){
        msgQueue.add(new Message(arg0,arg1));
    }
}
}

```

These are the definitions for the functional interfaces that allow lambda expressions to be passed to the agents to define their behavior.

```

/*
 * Classes who implement this interface must define a method to handle
 * messages
 */
public interface MessageProcessor{
    public void processMessage(Message msg);
}

/*
 * Classes who implement this interface must define a method to handle
 * exceptions
 */
public interface ErrorHandler{
    public void handleError(Exception e);
}

```

This code snippet shows an example of how to utilize the classes defined above in a runnable program that uses automated agents. The comments demonstrate

what code a user would need to add to customize their agent.

```
// An example of a class with a main method – how you would use these classes in a
↪ program
public class MyAutomatedAgent{
    public static void main(String[] args){
        // create and fill list of SubscriptionInfo objects
        ArrayList<SubscriptionInfo> subscriptions = new ArrayList<SubscriptionInfo>();

        // an example of how to use the AutomatedUserController to create an
        ↪ AutomatedUser
        AutomatedUser u = AutomatedUserController.createAutomatedUser("username", "password",
        ↪ subscriptions, (Message msg) -> {
            // write logic here to process message
        }, (Exception e) -> {
            // handle any exceptions here
            // use instanceof to check for specific exception types
            // thread will exit after this returns
        });
        new Thread(u).start();
    }
}
```



# Chapter 7

## Conclusion and Future Work

### 7.1 Overview

This work introduces new functionality into the WiSARDNet platform that enhances the usability of the system for researchers. Its design also provides a foundation for more features to be added in the future. This chapter summarizes the accomplishments of this thesis and assesses the impact each component has on the WiSARDNet platform and its users. Finally, this chapter discusses opportunities for future work.

### 7.2 Conclusions

This work presents a design and implementation for adding network reconfiguration capabilities to the WiSARDNet platform according to the requirements

established in Chapter 1. Each design requirement is listed below along with a description of how it was accomplished in this work.

- **Implementation of an application that allows users to remotely identify and select sets of WSN nodes based on static and dynamic search criteria.** The WiSARD browser module fulfills this objective by allowing users to specify information that corresponds to a variety of WiSARD attributes that are both static and dynamic in nature. This module produces a set of WiSARDs that match the search criteria entered by the user.
- **Implementation of an interface allowing users to enact behavioral changes in individual or groups of nodes.** This objective is satisfied by the command generation module that generates commands to reconfigure WiSARDs based on changes by the user. The extent to which a user can reconfigure the WiSARDs is governed by a user access control and network safety check that is performed for every command that is generated.
- **Generalization of the reconfiguration software such that agents can monitor and reconfigure network nodes.** This is satisfied with the implementation of automated agents that are able to utilize each of the three modules to reconfigure WiSARDs in reaction to different events that can be specified for each agent. Additionally, automated agents are subject to the same user access control and validation rules as human users.

## 7.3 Future Work

WiSARDNet is a platform that supports complex distributed ecological experiments. In this work, the addition of user access control, remote reconfiguration, and the ability to automate reconfiguration adds valuable functionality to the WiSARDNet platform. It also provides a foundation supporting additional capabilities.

### 7.3.1 Stored Behaviors and Profiles

The central paradigm of this thesis is that a WSN can be viewed as a database of heterogeneous behaviors, and that network reconfiguration can be achieved using a database approach. The process of network reconfiguration utilizes a command generator to build messages that invoke WSN configuration changes. A natural extension of this functionality would be for a user to have access to reconfiguration profiles, or groups of reconfiguration commands that produce a more complex defined behavior for a WiSARD. These profiles could be stored in a PostgreSQL data table, and could be selected and pushed out to a set of WiSARDs.

One example is a low power profile that reduces the sampling rates of all transducers in a WiSARD. This profile could also change the WiSARD's role from a node that relays messages to other nodes into a leaf node that does not need to expend energy communicating with other nodes. Once operational, a user would have access to this profile that could then be executed via the commands it

encapsulates. Additionally, profiles could be governed by the permission system, due to the generalized nature of the user access control schema. An extension of the permission resource table to include profiles would allow a permission entity to be granted access to specific profiles.

### **7.3.2 OTA Reprogramming**

Currently, new firmware versions must be physically uploaded to the WiSARDs because there currently is no over-the-air (OTA) reprogramming capability for the WiSARDs. The software implemented for this thesis could be expanded to support OTA reprogramming. The compiled source code for a new firmware update could be broken down into a series of messages, where each message payload contains a numbered piece of the new firmware. These messages could be sent sequentially to the intended WiSARD; once the WiSARD has received all of the messages, it could begin to overwrite the existing firmware in flash memory with the new firmware.

Implementing these features would not be a trivial task. OTA reprogramming would require expanding the command generation with features that would support breaking down a file containing a new firmware into many messages, numbering these messages, and sending these messages to the intended WiSARD. Additionally the firmware on WiSARDs would need to be upgraded to support these behaviors as well. The WiSARD would need to be able to recognize that each

message is a piece of a new firmware, and would need to store that piece of the message. Additionally, the WiSARD would need to verify that it has received all of the messages correctly, assemble the firmware pieces from each message into the entire file, and then overwrite the existing firmware with the new version. The process of overwriting the old firmware with the new one would also need to have a way to revert back to the old firmware if the new firmware was unable to be correctly downloaded or stored.

OTA reprogramming would require all of these new features, but the work of this thesis provides a foundation for these features to be built upon.

### **7.3.3 Creation of Automated Agents**

The addition of automated agents to the WiSARDNet system provides a way for the WSNs to be monitored and reconfigured in real time based on different events that a developer or researcher must specify for each agent. The use case described in Chapter 6 is a fairly simple example which manages the sampling rate of a WiSARD based on the time of day, but the system supports complex agents that monitor and reconfigure large numbers of WiSARDs in response to many different events. The agents are designed to have access to all of the monitoring and reconfiguration features that a human user does, so any improvements or additions to the modules such as the ones described in the previous sections would also benefit the agents in expanding their capabilities. Everything is in place for a user to

build a complex agent by following the example code in Chapter 6.

### **7.3.4 Improved User Interface**

As is the case with many software systems, there are plenty of opportunities for improvements to be made to the user interface. The user interface for the WiSARD Browser and Command Generation Module presents users with a simple way to enter WiSARD search criteria and specify changes. This can be improved with the inclusion of a geographic map that allows a user to view all of the WiSARDs that he/she has permission to reconfigure. Users with experiments spanning large geographic regions would have an easier time visualizing the effects of a reconfiguration if they had access to a user interface that displayed WiSARDs in this way.

Another opportunity for improvement is a user interface that lets a user build an automated agent visually, rather than requiring that he/she develop an agent by programming it directly. This would make automating network reconfiguration much more accessible to researchers without sufficient Java programming experience. This would require generalizing the agent code and allowing boiler-plate features to be specified and populated from either a form or some other type of user input. Additionally, allowing users to edit and modify agents from this user interface would also be beneficial.

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