

~~if  $b$  in  $\text{Im } f$   $\cdot b \cdot f(a)$~~

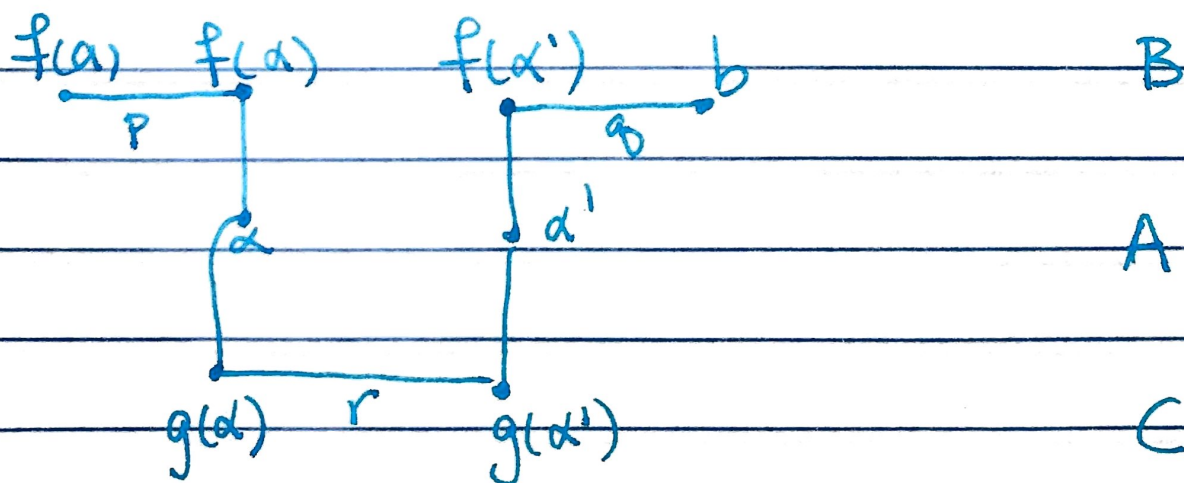
HOTT Feb 23

$\text{decode}(\text{code}(f(a), b)) \xleftarrow{\text{show these are equal}} \text{decode}(\text{code}(g(a), b))$   
 $\text{decode} \parallel \quad \text{decode} \parallel$   
 $\left( \sum_{\alpha': A} (b = f(\alpha')) \times (g(a) = g(\alpha')) \right) \quad \text{decode} \left( \sum_{\alpha': A} (g(a) = g(\alpha')) \times (b = f(a)) \right)$

Ask Michael Schulman: How can we formalize our argument that the definition of  $\text{code}(f(a), b)$  is the bottom left of  $\text{code}(b, b')$  definition.

~~send Amelia research statement~~ ✓

$\text{decode}(p, q, r) = \underbrace{\text{ap}_{\text{inl}}(q) \cdot \text{glue}_{\alpha'} \cdot \text{ap}_{\text{inr}}(r) \cdot \text{glue}_{\alpha} \cdot \text{ap}_{\text{inl}}(p)}_{\text{has to be refl.}}$



Next: show that this path is equiv. to  $\text{decode}(\text{code}(g(a), b))$