## MEETING NOTES

## 1. Dec 1 2017

## 1.1. Chandrika's notes.

1.1.1. Rushed Introduction and summary. Consider the following diagram, where A, B, C are sets, and D is the homotopy pushout of B and C along A.



We would like to show that D is a homotopy set as well using the encode-decode method. Therefore, we defined code :  $D \to D \to \text{Type}$  so that code(x, y) is a proposition.

Amelia spoke of a picture to think of for D. This is my understanding of the picture. D is the following topological space ("the double mapping cylinder").  $(B \cup A \times [0,1] \cup C)/\sim$ , where  $\sim$  is defined by:

- $(a,1) \sim b \iff f(a) = b, \forall a \in A, b \in B$
- $(a,0) \sim c \iff g(a) = c, \forall a \in A, c \in C$

We have attempted to define code(x,y) to be the homotopy structure of the paths from x to y in this double mapping cylinder.

In Utah, we successfully defined encode and began to define decode : code  $\to$  Id<sub>D</sub>. We defined decode(code(x, y)) for all  $x, y \in (B \cup C) \subset D$ .

1.1.2. What's next? It remains to define  $ap_{decode}$ . I think that these are the types of the various parts of  $ap_{decode}$ .

$$\begin{split} \operatorname{ap_{\operatorname{decode}}}\Big(\operatorname{ap_{\operatorname{code}}}(\operatorname{inl}b,\operatorname{glue}(a))\Big) : \left(\operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}b,\operatorname{inl}f(a)\big)\Big) = \operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}b,\operatorname{inr}g(a)\big)\Big) \right) \\ \operatorname{ap_{\operatorname{decode}}}\Big(\operatorname{ap_{\operatorname{code}}}(\operatorname{glue}(a),\operatorname{inl}b)\Big) : \left(\operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}f(a),\operatorname{inl}b\big)\Big) = \operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}g(a),\operatorname{inl}b\big)\Big) \right) \\ \operatorname{ap_{\operatorname{decode}}}\Big(\operatorname{ap_{\operatorname{code}}}(\operatorname{glue}(a),\operatorname{inr}c)\Big) : \left(\operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}f(a),\operatorname{inr}c\big)\Big) = \operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}g(a),\operatorname{inr}c\big)\Big) \right) \\ \operatorname{ap_{\operatorname{decode}}}\Big(\operatorname{ap_{\operatorname{code}}}(\operatorname{inr}c,\operatorname{glue}(a))\Big) : \left(\operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inr}c,\operatorname{inl}f(a)\big)\Big) = \operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inr}c,\operatorname{inr}g(a)\big)\Big) \right) \\ \operatorname{ap_{\operatorname{decode}}}\Big(\operatorname{ap_{\operatorname{code}}}(\operatorname{glue}(a),\operatorname{glue}(a))\Big) : \operatorname{isContr}\Big(\operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inl}f(a'),\operatorname{inl}f(a)\big)\Big) = \operatorname{decode}\Big(\operatorname{code}\big(\operatorname{inr}g(a'),\operatorname{inr}g(a)\big)\Big) \right) \end{split}$$