

showing $\text{decode}(\text{code}(g(a), c))$ are Hpy equiv $\text{decode}(\text{code}(f(a), c))$
 remember $\text{code}(g(a), c) = (g(a) = c)$ and $\text{code}(f(a), c) = \sum_{a:A} (f(a) = f(a)) \times (c = g(a))$
 need to show decode is constant/natural

$$r : (g(a) = c)$$

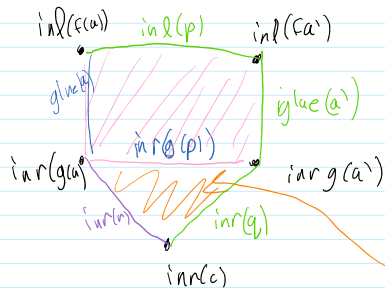
$$p : (f(a) = f(a))$$

Note, because f is injective, p implies $\exists p' : (a = a')$. b.c. B is a set, $ap + p' = p$. So we'll treat p as a path in A

$$q : (c = g(a'))$$

$$\text{decode}(r) = \text{ap}_{\text{inr}}(\text{inr}(r)) \quad \text{inr}(r) \text{ (switching up notations)}$$

$$\text{decode}(p, q) = \text{inr}(q) \circ \text{glue}(a') \circ \text{inl}(p)$$



glue(p)

we get this circle of things in C , but that's a set so it must be filled in

trying a new one ourselves

lets try $\text{ap}_{\text{decode}}(\text{code}(a), \text{inl}(b))$

Need to show $\text{decode}(\text{code}(f(a), b)) = \text{decode}(\text{code}(g(a), b))$
 $\text{code}(\text{inl}(f(a)), \text{inl}(b))$

$$\text{code}(f(a), b) = \text{P.O.} \quad \sum_{a:A} (f(a) = f(a)) \times (b = f(a)) \xrightarrow{\gamma} (f(a) = b)$$

$$\sum_{a, a':A} (f(a) = f(a')) \times (b = f(a')) \times (a = g(a')) \xrightarrow{\lambda} \text{code}(\text{inl}(f(a)), \text{inl}(b))$$

we almost decided that $p : \text{code}(f(a), b)$

must come from lower left corner
 which tells us which decode rule to use