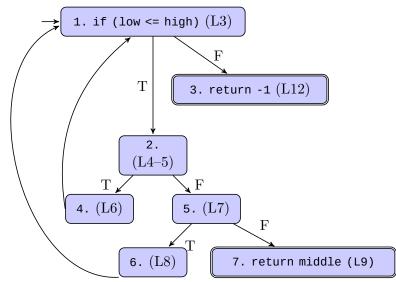
Software Testing, Quality Assurance and Maintenance	Winter 2019
Lecture 8 — January 23, 2019	
Patrick Lam	version 1

#### **Larger CFG example.** You can draw a 7-node CFG for this program:

```
1
      /** Binary search for target in sorted subarray a[low..high] */
2
      int binary_search(int[] a, int low, int high, int target) {
3
        while (low <= high) {</pre>
          int middle = low + (high-low)/2;
4
5
           if (target < a[middle)</pre>
6
            high = middle - 1;
7
           else if (target > a[middle])
            low = middle + 1;
9
           else
10
             return middle;
11
12
        return -1; /* not found in a[low..high] */
13
```



Here are more exercise programs that you can draw CFGs for.

```
/* effects: if x==null, throw NullPointerException
1
                otherwise, return number of elements in x that are odd, positive or both. */
2
    int oddOrPos(int[] x) {
3
4
      int count = 0;
5
      for (int i = 0; i < x.length; i++) {</pre>
6
         if (x[i]%2 == 1 || x[i] > 0) {
7
          count++:
9
      }
10
      return count;
11
12
    // example test case: input: x=[-3, -2, 0, 1, 4]; output: 3
```

Finally, we have a really poorly-designed API (I'd give it a D at most, maybe an F) because it's impossible to succinctly describe what it does. **Do not design functions with interfaces like this.** But we can still draw a CFG, no matter how bad the code is.

```
/** Returns the mean of the first maxSize numbers in the array,
 2
          if they are between min and max. Otherwise, skip the numbers. */
 3
       double computeMean(int[] value, int maxSize, int min, int max) {
 4
        int i, ti, tv, sum;
6
        i = 0; ti = 0; tv = 0; sum = 0;
7
        while (ti < maxSize) {</pre>
 8
9
          if (value[i] >= min && value[i] <= max) {</pre>
10
11
            sum += value[i];
12
          }
13
          i++;
14
15
        if (tv > 0)
16
          return (double)sum/tv;
17
18
           throw new IllegalArgumentException();
19
```

### Statement and Branch Coverage

We defined Control-Flow Graphs so that we can give principled definitions of statement and branch coverage. We can start with the definition of a test path:

**Definition 1** A test path is a path p (possibly of length 0) that starts at some initial node (i.e. in  $N_0$ ) and ends at some final node (i.e. in  $N_f$ ).

Here's a definition of coverage for graphs:

**Definition 2** Given a set of test requirements TR for a graph criterion C, a test set T satisfies C on graph G iff for every test requirement TR, at least one test path P in P in P in that P satisfies P in P satisfies P satisfies P in P satisfies P s

We'll use this notion to define a number of standard testing coverage criteria. But first, what are test paths?

**Test cases and test paths.** We connect test cases and test paths with a mapping path<sub>G</sub> from test cases to test paths; e.g. path<sub>G</sub>(t) is the set of test paths corresponding to test case t.

- usually we just write path since G is obvious from the context.
- we can lift the definition of path to test sets T by defining path $(T) = \{ path(t) | t \in T \}$ .
- each test case gives at least one test path. If the software is deterministic, then each test case gives exactly one test path; otherwise, multiple test cases may arise from one test path.

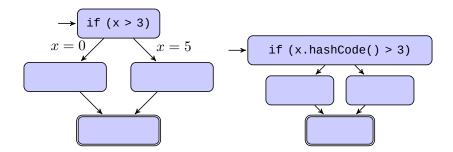
**Example.** Here is a short method, the associated control-flow graph, and some test cases and test paths.

```
1
   int foo(int x) {
                                                             (1) if (x < 5)
2
     if (x < 5) {
3
       x ++;
4
     } else {
                                                        (3) x - -
                                                                         (2) x++
5
       x --;
6
7
                                                              (4) return x
     return x;
8
  }
```

- Test case: x = 5; test path: [(1), (3), (4)].
- Test case: x = 2; test path: [(1), (2), (4)].

Note that (1) we can deduce properties of the test case from the test path; and (2) in this example, since our method is deterministic, the test case determines the test path.

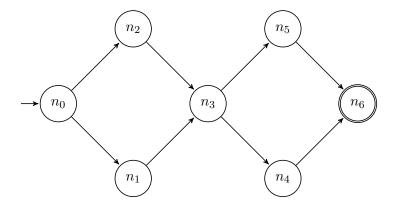
**Nondeterminism.** I mentioned the mapping between test cases and test paths above. The mapping is not one-to-one for nondeterministic code. Here's an example of deterministic and non-deterministic control-flow graphs:



Causes of nondeterminism include dependence on inputs; on the thread scheduler; and on memory addresses, for instance as seen in calls to the default Java hashCode() implementation.

Nondeterminism makes it hard to check test case output, since more than one output might be a valid result of a single test input.

As another (more abstract) example, consider the double-diamond graph D.



Here are the four test paths in D:

$$[n_0, n_1, n_3, n_4, n_6]$$

$$[n_0, n_1, n_3, n_5, n_6]$$

$$[n_0, n_2, n_3, n_4, n_6]$$

$$[n_0, n_2, n_3, n_5, n_6]$$

For the statement coverage criterion, we get the following test requirements:

$$\{n_0, n_1, n_2, n_3, n_4, n_5, n_6\}$$

That is, any test set T which satisfies statement coverage on D must include test cases t; the cases t give rise to test paths path(t), and some path must include each node from  $n_0$  to  $n_6$ . (No single path must include all of these nodes; the requirement applies to the set of test paths.)

Let's formally define statement coverage.

**Definition 3** Statement coverage: For each node  $n \in reach_G(N_0)$ , TR contains a requirement to visit node n.

For our example,

$$TR = \{n_0, n_1, n_2, n_3, n_4, n_5, n_6\}.$$

Let's consider an example of a test set which satisfies statement coverage on D.

Start with a test case  $t_1$ ; assume that executing  $t_1$  gives the test path

$$path(t_1) = p_1 = [n_0, n_1, n_3, n_4, n_6].$$

Then test set  $\{t_1\}$  does not give statement coverage on D, because no test case covers node  $n_2$  or  $n_5$ . If we can find a test case  $t_2$  with test path

$$path(t_2) = p_2 = [n_0, n_2, n_3, n_5, n_6],$$

then the test set  $T = \{t_1, t_2\}$  satisfies statement coverage on D.

What is another test set which satisfies statement coverage on D?

Here is a more verbose definition of statement coverage.

**Definition 4** Test set T satisfies statement coverage on graph G if and only if for every syntactically reachable node  $n \in N$ , there is some path p in path(T) such that p visits n.

A second standard criterion is that of branch coverage.

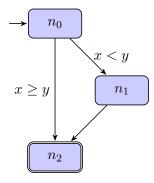
Criterion 1 Branch Coverage. TR contains each reachable path of length up to 1, inclusive, in G.

Here are some examples of paths of length  $\leq 1$ :

Note that since we're not talking about *test paths*, these reachable paths need not start in  $N_0$ . In general, paths of length  $\leq 1$  consist of nodes and edges. (Why not just say edges?)

Saying "edges" on the above graph would not be the same as saying "paths of length  $\leq 1$ ".

**Another example.** Here is a more involved example:



Let's define

$$path(t_1) = [n_0, n_1, n_2]$$
  
 $path(t_2) = [n_0, n_2]$ 

Then

$$T_1 = \langle ? \rangle$$
 satisfies statement coverage  $T_2 = \langle ? \rangle$  satisfies branch coverage

# Web Applications

Frontends HTML, CSS, JS (client-side) vs Backends

Web server, Database, PHP, Node.js (server-side)

## **Testing Frontend:**

- rendering
- cross-browser (frontend)
- wrong value errors
- malware
- performance
- cache issues

### Testing Backend:

- Dev Ops (simulations of failures)
- responsiveness to requests
- data persistence
- security (SQL injection prevention)
- race conditions
- performance testing