

# **Analysis of Type-Driven approach to systems programming: Implementation of OpenGL library for Rust**

(Analiza programowania systemowego z wykorzystaniem systemu typów:  
Implementacja biblioteki do OpenGL dla języka Rust)

Mikołaj Depta

Praca inżynierska

**Promotor:** dr Andrzej Łukaszewski

Uniwersytet Wrocławski  
Wydział Matematyki i Informatyki  
Instytut Informatyki

26 sierpnia 2024



## Abstract

For the past few years in the software development industry there has been a growing interest in strongly typed languages. It manifests itself in emergence of brand-new technologies in which strong type systems were one of the core founding principles or in changes introduced to existing languages. The most common examples of modern languages with powerful type systems are TypeScript as an alternative to JavaScript in the world of web development or Rust in domain of systems programming in place of C and C++. More mature languages also had their type systems revised for example in C# 8 explicit type nullability annotations were introduced, or even dynamically typed Python has seen major improvements to its type annotation system.

This study - an implementation of the OpenGL graphics API wrapper library for Rust - will attempt to demonstrate how Rust's type system can be utilized to improve low-level software safety and maintainability as well as how it affects API design and codebase structure.

---

For the past few years in the software development industry there has been a growing interest in strongly typed languages. It manifests itself in emergence of brand-new technologies in which strong type systems were one of the core founding principles or in changes introduced to existing languages. The most common examples of modern languages with powerful type systems are TypeScript as an alternative to JavaScript in the world of web development or Rust in domain of systems programming in place of C and C++. More mature languages also had their type systems revised for example in C# 8 explicit type nullability annotations were introduced, or even dynamically typed Python has seen major improvements to its type annotation system.

This study - an implementation of the OpenGL graphics API wrapper library for Rust - will attempt to demonstrate how Rust's type system can be utilized to improve low-level software safety and maintainability as well as how it affects API design and codebase structure.



# Contents

<b>1</b>	<b>The Rust programming language</b>	<b>7</b>
1.1	Data types . . . . .	9
1.1.1	User defined types . . . . .	9
1.2	Implementations . . . . .	10
1.2.1	Inherent implementations . . . . .	10
1.3	Generic types . . . . .	10
1.3.1	Type parameters . . . . .	11
1.3.2	Constant parameters . . . . .	11
1.3.3	Lifetime parameters . . . . .	11
1.4	Type level abstraction . . . . .	11
1.5	Implementation Coherence . . . . .	11
1.5.1	Orphan Rule . . . . .	11
<b>2</b>	<b>OpenGL and GLSL</b>	<b>13</b>
2.1	Introduction . . . . .	13
2.2	History . . . . .	13
2.3	OpenGL objects . . . . .	13
2.3.1	Buffer object . . . . .	14
2.3.2	Vertex Array object . . . . .	14
2.4	Graphics pipeline . . . . .	15
2.4.1	Vertex Specification . . . . .	15
2.4.2	Vertex Shader . . . . .	16
2.4.3	Tessellation . . . . .	16

2.4.4	Geometry shading . . . . .	17
2.4.5	Fixed function vertex post-processing . . . . .	17
2.4.6	Rasterization . . . . .	17
2.4.7	Fragment processing . . . . .	18
2.4.8	Fragment post processing . . . . .	18
2.5	GLSL . . . . .	18
2.5.1	Variables . . . . .	19
2.5.2	Variable types . . . . .	19
2.5.3	Variable storage qualifiers . . . . .	20
2.5.4	Variable layout qualifiers . . . . .	21
<b>3</b>	<b>OpenGL wrapper library</b>	<b>25</b>
3.1	External dependencies . . . . .	25
3.2	Auxiliary modules and crates . . . . .	26
3.3	Identified design patterns . . . . .	26
3.3.1	Compensation for language limitations . . . . .	26
3.3.2	Const generics in const expressions . . . . .	27
3.3.3	Effect system . . . . .	28
3.3.4	Application of existing features . . . . .	28
3.3.5	Markers . . . . .	28
3.3.6	Type State . . . . .	30
3.4	OpenGL wrapper . . . . .	30
	<b>Bibliography</b>	<b>31</b>

## Chapter 1

# The Rust programming language





# Introduction

Rust is an open-source, general purpose, multi paradigm, compiled, statically and strongly typed language maintained by Rust Foundation.

Although general purpose, language lends itself particularly well to systems programming, where software reliability, scalability and efficiency are paramount. These qualities can be largely attributed to powerful and expressive type system and ownership based memory management system which guarantees memory safety without a garbage collector.

## 1.1 Data types

### 1.1.1 User defined types

Rust provides 3 ways to construct user defined aggregate types which. These are: - **structs** - **enums** - **unions**

We focus on the first two of these.

**structs** Struct is a heterogeneous product of other types, they are analogous to struct types in C, the record types of the ML family, or the struct types of the Lisp family. They constitute the basic building blocks for any user defined types.

A struct type is a heterogeneous product of other types, called the fields of the type.

The memory layout of a struct is undefined by default to allow for compiler optimizations like field reordering, but it can be fixed with the repr attribute. In either case, fields may be given in any order in a corresponding struct expression; the resulting struct value will always have the same memory layout.

The fields of a struct may be qualified by visibility modifiers, to allow access to data in a struct outside a module.

A tuple struct type is just like a struct type, except that the fields are anonymous.

A unit-like struct type is like a struct type, except that it has no fields. The one value constructed by the associated struct expression is the only value that inhabits such a type.

1 struct types are analogous to struct types in C, the record types of the ML family, or the struct types of the Lisp family.

## 1.2 Implementations

Functionality of a type is not provided inline with its definition like in most C like languages. Instead it's associated with a type by so called *implementations*. Implementations for a type are contained within a block introduced by the `impl` keyword.

There are two types of implementations: - inherent implementations - trait implementations

All items within `impl` block are called associated items

Functions defined within `impl` blocks are called associated functions and can be accessed with qualification `<type-name>::<function-name>`.

Within an implementation both `self` and `Self` keywords can be used. `self` can be used in associated function definitions within an `impl` block as first parameter. Such functions are called methods and the `self` parameter denotes the receiver of method call. `self` can be additionally qualified with immutable or mutable reference `&` or `&mut`.

`Self` is a type alias that refers to implementing type.

### 1.2.1 Inherent implementations

We will shorthand implementation to `impl` which is common in Rust terminology.

Inherent `impls` associate contents of an `impl` block with specified nominal type. Such blocks can contain definition of a constants or functions.

## 1.3 Generic types

As of version 1.79 Rust provides 3 kinds of generic parameters types can use. These are: - type parameters - constant parameters - lifetime parameters Type which uses any generic parameters is said to be generic.

### 1.3.1 Type parameters

Type parameters represent an abstract type

### 1.3.2 Constant parameters

### 1.3.3 Lifetime parameters

## 1.4 Type level abstraction

Traits provide an ability to express shared behavior in abstract way [1]. We are mostly interested in their use in *trait bounds* on types and type parameters. Trait bounds declare contracts that types must fulfil or else the program will be rejected. We used it to enforce use of valid data formats and proper sequencing of operations.

## 1.5 Implementation Coherence

### 1.5.1 Orphan Rule



## Chapter 2

# OpenGL and GLSL

### 2.1 Introduction

OpenGL is an specification of an API for hardware accelerated computer graphics owned and maintained by the Khronos Group Inc.

Since it's inception and up until 2016 with release of Vulkan it has been the only widely supported cross platform graphics API.

### 2.2 History

IRIS GL, a proprietary graphics API, which later became OpenGL was initially developed by Silicon Graphics (SGI) during 1980's. SGI open sourced subset of their API as OpenGL due to mounting market pressure from SGI competitors (Sun Microsystems, Hewlett-Packard, IBM) who started providing their own APIs for hardware accelerated 3D graphics based on an existing open standard called PHIGS. In 1992 OpenGL Architectural Review Board (OpenGL ARB) was established and it was tasked with development and maintenance of the OpenGL specification. This task was passed on to Khronos Group in 2006 where it remained ever since.

### 2.3 OpenGL objects

OpenGL exposes an abstraction over GPU's resources called objects. These roughly correlate with object oriented design as they aggregate data for appropriate subset of operations albeit with certain unique caveats. In all but the latest opengl versions, to use given object it first must have been bound to a binding point in global in current OpenGL context. In OpenGL 4.6 the `ARB_direct_state_access` extension was made part of core specification which introduces duplicates of all object manipulating functions to accept as one of parameters the name of object to operate on.

Objects contain internal state which can be queried using introspection. Objects are identified by a *name* which is an unsigned 32 bit integer. There exists common object model which describes how most types of objects are managed.

Most types of objects can be created with a call to

```
void Gen*s(sizei n, uint *objects)
```

which will allocate the object's *name*. A subsequent call to

```
void Bind*(uint name, ...)
```

will bind the given object to the context. If the object has never been bound before, this will also allocate its internal state. Alternatively, one can use

```
void Create*s(sizei n, uint *objects)
```

which will allocate both the object's *name* and its state, but it will not set any context bindings. There exists a separate namespace for each object type.

Objects can be deleted with `void Delete*s(sizei n, uint *objects)`, bound with aforementioned `void Bind*(uint name, ...)` which usually accepts additional parameter that specifies binding point. The most notable outliers that do not conform to the rules above are program objects and shader objects.

OpenGL specification defines set of publicly available object parameters which can be queried using introspection with **GetInteger\*** family of functions. One notable usage is determining compilation and linking status for shaders and programs.

### 2.3.1 Buffer object

Buffer objects provide means to allocate data stores in GPU memory. They can contain data of different format and purpose depending on buffer's target. Primary usage for buffers is to provide geometric information which includes vertex attribute values and indices for indexed rendering.

### 2.3.2 Vertex Array object

Modern OpenGL is generic over vertex format and only poses limitation on the number of such attributes and limits their values to glsl's scalar or vector types. Each attribute is assigned a zero-based index. Vertex Array object (VAO) assigns each active attribute information on how and where from to source vertex data, as well as, what is the data type of provided attribute in glsl.

This can be viewed as two aspects: (1) memory layout and access frequency, and (2) data interpretation/conversion.

Originally all of that information was specified at one with

```
void VertexAttrib*Pointer(  
    )
```

## 2.4 Graphics pipeline

The modern OpenGL pipeline is a sequence of both programmable and fixed function stages that process geometric data to form discrete color values - pixels - that end up stored in a framebuffer.

### 2.4.1 Vertex Specification

Before rendering can begin, geometric information needs to be uploaded to GPU memory along with its description as generic vertex attributes.

Generic Vertex is an abstract composition of values (attributes) that is supposed to represent a vertex of the triangular mesh of an object. Generic stands from the fact that data associated with vertices has no intrinsic meaning. Semantics of data are decided by client provided vertex shader.

OpenGL sources data for each vertex attribute from a buffer. Each attribute is assigned a unique numeric index. Association between attribute with given index and a buffer, from which that attribute should be sourced, is established by Vertex Array Object (abbr. VAO).

Once all vertex attributes have their data sources assigned and properly configured, vertex specification can be considered finished and one could proceed with further pipeline configuration. In this instance vertices would be interpreted sequentially as appropriate geometric primitives. This forces vertex data to be specified redundantly for lines and especially raw triangles, since each triangle shares an edge with each neighboring triangle.

To better conserve memory one can use indexed rendering. This requires additional buffer filled with indices into main vertex buffer instead of inlined vertex data. In case of basic triangle rendering (without using compressed representations like triangle fan or triangle / line strip) will still cause repetition but now only few byte wide indices instead of whole attributes which are substantially larger.

### 2.4.2 Vertex Shader

Vertex shader is the first programmable stage of OpenGL Pipeline and is one of two required shaders to execute a draw call, the other being the fragment shader.

Most commonly vertex shader performs 3 translations. From initial model space, world space, view space to final clip space which we will now discuss briefly.

**model space** - when a 3D model is created in 3D modelling software its vertex positions are specified to some local coordinate system (commonly center of an object). These positions would commonly be loaded into gpu memory. Such objects can be easily placed in broader scene by providing a so called world transform performs transformation from model's local coordinate system to scene's coordinate system.

**world space** - world space refers to coordinate system of a scene that uses multiple models by translating, scaling or rotating them.

**view space** - its common for 3D rendering applications to provide means of interacting with the scene. Whether its a 3D computer game, CAD program or medical data visualization we would like to be able to control how scene is displayed by moving a virtual camera. This can be expressed as yet another transformation of the coordinate system - we would like to transform coordinate system to align with the position of our camera. This transformation is commonly called view transform.

**clip space** - having accounted for model position in a scene and user interactivity all that remains is to provide vertex data in form that subsequent fixed function pipeline stage - the rasterizer - expects. Namely once vertex shader is finished fixed function processing will clip all geometry then perform perspective divide to obtain vertices in normalized device coordinates (NDC). Output of vertex shader is a 4 component vector which corresponds to a 3D position in homogenous coordinate system used in computer graphics due to its ability to represent non linear transformations using matrices. The output position of vertex is divided by the forth component in order to introduce perspective

The main responsibility of vertex shader is to transform vertices to clip space, which will be discussed in future subsection.

### 2.4.3 Tessellation

Tessellation stages were added as graphics hardware compute capability grew. With raw compute throughput outperforming bus throughput GPUs were equipped with hardware tessellation unit which can subdivide a larger triangle into batch of smaller ones. This allows for efficient generation of geometric detail on chip alleviating the issue of limited PCI throughput. To drive the tessellation stage two new shaders were introduced:



- **Tessellation control shader** which configures how hardware tessellator should subdivide a triangle.
- **Tessellation evaluation shader** which performs transformations on vertices generated by the tessellator.

#### 2.4.4 Geometry shading

Geometry shader was introduced prior to tessellation stage. They operate on assembled geometric primitives and may even access primitives neighbors. Given primitive input geometry shaders output one or more primitive of the same type.

#### 2.4.5 Fixed function vertex post-processing

Once all programmable vertex processing has concluded, a series of fixed-function operations are applied to the vertices of the resulting primitives before rasterization. These operations include transform feedback, which captures processed vertex data, primitive queries to gather information about the primitives being processed and flat shading which applies a uniform attribute value to a whole primitive.

Primitives then get clipped against clip volume and client-defined half-spaces. The clip coordinates undergo perspective division, followed by viewport mapping to adjust for screen coordinates and depth range scaling.

#### 2.4.6 Rasterization

If neither tessellation stage nor geometry stage was used in vertex processing, primitive assembly takes place (presence of any of the aforementioned stages would necessitate early primitive assembly). OpenGL converts geometric primitives used in currently processed draw call into base primitives which are points, lines and triangles. Mathematical representation of primitives is used during rasterization to determine if given fragment falls inside of primitive being rasterized.

Process of rasterization requires determining if given pixel position falls inside of rendered primitive. This process needs to account for point and line thickness. Polygon rasterization is obviously the most complex of the three. Prior to insidiness test face culling is performed. This optimization culls a polygon based on the sign of surface normal computed based on edge ordering as specified in vertex array. This helps reduce overdraw which can be one of two main bottlenecks in modern rendering system, the latter being insufficient memory bandwidth.

Once pixel location was deemed inside a primitive a fragment is generated. A Fragment is a collection of data corresponding to specific pixel location. Most commonly its perspective corrected barycentric interpolation of vertex data across

the primitive's surface. Tough interpolation can be disabled from within vertex shader using `flat` qualifier on output variable declaration, as well as perspective correction with `noperspective` qualifier.

Once fragments are computed early per-fragment tests take place.

- **Ownership test** - determines if pixel at location (x, y) falls into the portion of the screen that active OpenGL context owns.
- **Scissor test** - checks if pixel at location (x, y) is contained within client provided list of axis aligned rectangles
- **Early Fragment tests** - stencil test, depth test and occlusion query which are normally performed after fragment processing can optionally be performed early. We discuss them in subsection on fragment post processing.

If all tests passed fragment is submitted for programable fragment processing.

### 2.4.7 Fragment processing

Programable fragment processing is performed by client provided fragment shader. The most essential task that fragment shader should perform is assign pixel a color. For that purpose data interpolated from rasterization is used. Most commonly fragment shaders perform texture mapping, lighting calculations, parallax mapping to emulate geometric detail and screen space effects like ambient occlusion, use signed distance functions and implicit surface equations to render otherwise complex scenes all by itself or create volumetric effects like clouds or visualize CT scan results.

### 2.4.8 Fragment post processing

## 2.5 GLSL

GLSL, which stands for OpenGL Shading Language, is a high level shading language with c like syntax developed by OpenGL Architecture Review Board to power programable processing stages in OpenGL pipeline. GLSL code is still relevant as it can be compiled into SPIR-V and used with Vulkan API.

### Shaders

Independent compilation units written in this language are called shaders. A program is a set of shaders that are compiled and linked together, completely creating one or more of the programmable stages of the API pipeline

In OpenGL 4.6 and GLSL 4.60 there exist 6 types of shaders: vertex, tessellation control and evaluation, geometry, fragment and compute. All shaders except compute shader control appropriate parts of OpenGL pipeline as described in subsections above.

Compute shaders operate completely outside of graphics pipeline. They can access same resources as fragment or vertex shader like textures, buffers, images and atomic counters but they are not expected to produce data with predetermined form or semantics. They offer general purpose compute capability on the GPU. They function similarly to other existing general purpose GPU compute APIs like CUDA or OpenCL.

### 2.5.1 Variables

The main purpose of shaders is to transform received data to some other form. The data that the shader expects is defined using global variables with appropriate qualifiers. During Program linking OpenGL matches outputs from previous stage with inputs of the next stage. In case of vertex shader `in` variables should match with vertex attribute definitions specified in vertex array object. Though in case of mismatch if attribute is disabled constant value can be provided, however that's rarely desired behavior. Similarly, `out` variables from fragment shader should match with framebuffer configuration. This process can be quite error prone and can lead to undefined behavior which can be difficult to diagnose and may have different consequences depending on actual hardware, OS or driver versions.

Under no circumstances erroneous pipeline configuration should be allowed. Program containing such malformed configuration should be rejected by static analysis, and that was one of most important aspects of this study. To achieve that we attempted to express both GLSL variable declarations along with full OpenGL pipeline in Rust's type system in such a way to force type errors for invalid pipeline configurations.

We determined that keeping track of three variable qualifiers is essential to achieve that.

### 2.5.2 Variable types

Expressing GLSL variable type in Rust types was the obvious first step. GLSL defines a set of built-in types along with ability to create aggregate data types with C-like array and struct definitions.

For this work we focused on builtin types and arrays and omitted structures due to Rust's inability to encode layout guarantees for arbitrary types. Rust's builtin numeric types and arrays have however have well defined memory layout and

create a close set of possible types which allowed us to enumerate them express their memory layout in type system using traits.

GLSL's built-in types are divided into two groups: transparent and opaque types. Transparent types represent numeric data (plain old data) whereas opaque types represent handles to different resources like texture image samplers.

In case of transparent types there are 5 base numeric types: `float`, `double`, `int`, `uint` and `bool`. Floating point types `float` and `double` are accordingly IEEE-754 single and double prevision numbers, integers are two's compliment 32bit values and `bool` undefined representation but it can take only two values `true` or `false`.

All base numeric types can be aggregated into 2, 3 or 4 component vector types. Each vector type is named `TvecN` where **T** depends on inner base type: `b` for `bool`, `d` for `double`, `i` for `int`, `u` for `uint` and for `float` nothing is prepended to `vecN`. **N** is the number of components vector should contain.

Finally, there are matrix types of form `TmatN`. Matrices can contain only `floats` (`matN`) or `doubles` (`dmatN`). The **N** depends on matrix dimensions it can be a single number 2, 3 or 4 for square matrices or can be arbitrarily combined pair of these numbers of form `NxM`, i.e. `mat2x4`, `mat4x2` or `dmat3x3`.

Data types used by GLSL have quite large memory footprint. Thats why OpenGL provides conversion mechanisms for data stored in buffers. Buffer data can be low bitwidth integer or float which will be normalized on access or even completely new OpenGL defined packed formats like `UNSIGNED_INT_2_10_10_10_REV` - a 32bit value which will be expanded to 4 `floats`.

This indirect mapping of OpenGL data to GLSL types is also essential to be statically verified just like shader input / output matching.

### 2.5.3 Variable storage qualifiers

The origin of data for a variable is encoded by a storage qualifier. We have already discussed that data within a shader can originate from previous stage / vertex buffers and that it can be saved as input to subsequent stage or framebuffer. These sources correspond to `in` and `out` qualifiers. Shaders can also declare uniform variables which are data associated with program itself. Value of these variables remains the same across the entire primitive being processed. All uniform variables are read-only and are initialized externally either at link time or through the API.

For bidirectional communication between shaders and API there exists `buffer` qualifier. Variables with such qualification are stored in buffer objects and can be both written to and read from by shaders and API.

Remaining qualifiers are ignored as they are irrelevant for the scope of this study.

### 2.5.4 Variable layout qualifiers

Statically asserting that the data passed through the pipeline is correct was the main goal of this study. .... Ability to statically assert that data flow through the pipeline is configured correctly depends on our ability to match neighboring stage inputs and outputs. In newer OpenGL versions one can use `location` layout qualifier to assign a variable a numeric index. This integer will be checked for uniqueness among all other variables that it shares storage qualifier with, and generate compilation error in case of overlap. When using locations based interface matching each `out` variable must have

It assigns an variable location and along with its storage qualifier it creates unique

TODO: tie this subsection together

“ An output variable is considered to match an input variable in the subsequent shader if:

- the two variables match in name, type, and qualification, and neither has a location qualifier, or
- the two variables are declared with the same location and component layout qualifiers and match in type and qualification.

For the purposes of interface matching, variables declared with a location layout qualifier but without a component layout qualifier are considered to have declared a component layout qualifier of zero. ”

/chapterExisting solutions

There are many qualities of any software library one could consider important. In this research we focused foremost on providing minimalistic wrapper and staying as faithful as possible to original specification of the API. By this we mean that appropriate GL functions take analogous parameters as in original spec and have their names and semantics preserved. Major benefit of this approach is that we could simply follow the OpenGL specification when creating type safe facades around procedures.

Starting from these minimalistic principles we focused on providing maximal level of type safety. The main goal was to enable rejection of as many ill-formed programs at compile time as possible.

There are many levels of safety guarantees we can expect from any software package. In this analysis we devise

Here we consider alternative ways of programming computer graphics with use of OpenGL as rendering backend.

We distinguish between a language of choice and any framework at use.

/sectionNative C / C++ bindings

The simplest way one can program with OpenGL is using platform provided C bindings contained within an os provided dynamic link library (.dll for MS Windows, .so ) along with appropriate function pointer loader. The requirement for the latter stems from common practice among OS vendors to officially guarantee distribution of very dated version of the specification (1.1 for Windows). This poses a requirement for manual function pointer loading at runtime, an approach that has two main benefits - abstracts away details of dynamic library loading for different platforms, - provides unified mechanism for using optional core standard extensions.

Additionally to function pointer loading one needs to initialize OpenGL context following platform defined protocol.

Most commonly there exists a library for each task, some examples for PC are GLEW for function loading and GLFW for window creation and context creation. Once these actions are accomplished one can use OpenGL in C or C++ provided appropriate attention to C interoperability.

Writing C application provides no auto

/sectionRust with unsafe bindings

Rust toolchain provides a utility for automatically generating Rust Foreign Function Interface bindings to C called **bindgen**. In this case all the setup needed for a Native C / C++ bindings application still applies. There exist appropriate counterparts to GLEW and GLFW. Once context is initialized and function pointers loaded one can call C functions but Rust will require one to use these functions inside unsafe context.

/chapterExisting solutions

There are many qualities of any software library one could consider important. In this research we focused foremost on providing minimalistic wrapper and staying as faithful as possible to original specification of the API. By this we mean that appropriate GL functions take analogous parameters as in original spec and have their names and semantics preserved. Major benefit of this approach is that we could simply follow the OpenGL specification when creating type safe facades around procedures.

Starting from these minimalistic principles we focused on providing maximal level of type safety. The main goal was to enable rejection of as many ill-formed programs at compile time as possible.

There are many levels of safety guarantees we can expect from any software package. In this analysis we devise

Here we consider alternative ways of programming computer graphics with use of OpenGL as rendering backend.

We distinguish between a language of choice and any framework at use.

/sectionNative C / C++ bindings

The simplest way one can program with OpenGL is using platform provided C bindings contained within an os provided dynamic link library (.dll for MS Windows, .so ) along with appropriate function pointer loader. The requirement for the latter stems from common practice among OS vendors to officially guarantee distribution of very dated version of the specification (1.1 for Windows). This poses a requirement for manual function pointer loading at runtime, an approach that has two main benefits - abstracts away details of dynamic library loading for different platforms, - provides unified mechanism for using optional core standard extensions.

Additionally to function pointer loading one needs to initialize OpenGL context following platform defined protocol.

Most commonly there exists a library for each task, some examples for PC are GLEW for function loading and GLFW for window creation and context creation. Once these actions are accomplished one can use OpenGL in C or C++ provided appropriate attention to C interoperability.

Writing C application provides no auto

/sectionRust with unsafe bindings

Rust toolchain provides a utility for automatically generating Rust Foreign Function Interface bindings to C called **bindgen**. In this case all the setup needed for a Native C / C++ bindings application still applies. There exist appropriate counterparts to GLEW and GLFW. Once context is initialized and function pointers loaded one can call C functions but Rust will require one to use these functions inside unsafe context.





## Chapter 3

# OpenGL wrapper library

In this chapter, we demonstrate how Rust’s type system can be harnessed to create a safe wrapper library for modern OpenGL, specifically targeting version 4.6. Our goal is to cover the most essential components of the OpenGL specification and staying as close to the original spec as possible. In many cases, we implement a minimal subset of functionality to demonstrate that, once a specific feature is in place, it can be readily extended to encompass a broader scope of the API.

Besides the wrapper library the purpose of this study was to identify common patterns that arise during type driven design.

### Overview

The resulting library was named *GPU bulwark* since it provides strong foundations for safer programming on the GPU, and could easily be extended to other GPU programming APIs.

Library at its root is logically divided into two halves: (1) main OpenGL wrapper and (2) general-purpose auxiliary modules which contain implementations of various patterns we have recognized.

### 3.1 External dependencies

Our library utilizes several publicly available crates from crates.io, we will briefly discuss their purposes below:

- `gl` - generates raw OpenGL bindings for Rust using build script. Additionally, it exposes a single function that loads function pointers using the provided routine. These bindings use C types and need to be invoked in `unsafe` context.

- `derive_move` - is a procedural macro crate that expands `derive` to support more built-in traits. It significantly reduces code boilerplate.
- `concat_ids` - provides singular procedural macro that allows to concatenate identifiers akin to C's `##` operator. We utilize this macro for identifier generation for certain OpenGL names that strictly follow a naming convention. This yet again helps to reduce boilerplate, makes code more succinct and minimizes risk of typos.
- `nalgebra` and `nalgebra-glm` define algorithms and types for linear algebra computations. They are not used directly in our library for their functionality but rather for optional integration with `gpu-bulwark`.

Remaining packages are imported for use in examples only.

- `thiserror` and `anyhow` - very popular crates that make error handling more ergonomic.
- `raw-window-handle`, `glutin` and `winit` allow for cross platform window creation and OpenGL context initialization.

## 3.2 Auxiliary modules and crates

All general purpose design patterns we encountered during development are implemented in these modules.

## 3.3 Identified design patterns

In our exploration we found that patterns which to emerge during programming with types can be broadly divided into two categories: (1) compensation for language limitations (2) validation of program structure at complete.

### 3.3.1 Compensation for language limitations

Rust is in continuous development. Some features have been work-in-progress for over years and are still nowhere near completion, and another have seen minimal-viable-product releases, and yet others are nothing more than subject of wishes and speculation. Features which we found to be useful for type-driven design fall into all of these categories. Some of them can be emulated with varying levels of complexity and user experience consequences stemming from often contrive usage of type system and language facilities.

#### Variadic Generics

**Problem** It is common practice among programming language developers to support variadic function arguments - functions which can accept arbitrarily

many arguments. This capability is great syntactic convenience, as well as, a tool for more complex abstractions.

It is substantially less common to support variadic type parameters in generic types, as a matter of fact Rust does not have such a language feature and discussion regarding it is very scarce. We identified one very desirable use case for such variadic generics - non homogenous collections.

**Solution** Rust have an excellent support for recursive types and since lists can be defined recursively like for example in LISP we derived a variadic generics emulation scheme from that. We used type level recursion on binary tuples (recursive step) and unit type as `Nil`. We call such type list **HList**. HLists can be wound to the left: first tuple component contains  $n - 1$  elements and second the  $n$ 'th type or in the reverse. Both these schemes are equivalent in terms of functionality but differ in terms of potential user experience. In our use case appending new types to the end of a HList was by far the most common use case and as we use almost exclusively such left wound HLists (LHLists).

Implementation of functionality for hlists needs to mirror their abstract and recursive nature. This can be achieved using two recursive implementations, one for unit in base case and another for a two tuple for recursive step.

These homogenous collections have been implemented as an independent module called `hlist` which can be found in the root of our crate.

**Use case** In `gpu-bulwark` we utilize hlists every time user configuration is required, most notably to represent shader inputs, outputs, used uniforms or external resources like textures. Almost always we create a facade marker trait which join together predefined pieces of functionality from `hlist` module and adds specialized requirements for the types depending on the usage which prohibits creation of invalid types.

### 3.3.2 Const generics in const expressions

**Problem** Const generics fall into the category of partially implemented features. Types can be dependant on values of limited subset of types, most notably for our purposes numeric types, `bool`, `unit`. Most crippling limitation is that these values must be expressions using only literals and in general two const parameters cannot be added or subtracted from each other. This prohibits us from performing arbitrary compile-time (CT) computations of these parameters for purposes of validation.

**Solution** There is one exception to that limitation however - associated constants. Associated constants can have their values computed using CT `const`

`fns` and themselves be used in such computations as parameters. These functions can panic with static error message (no formatting) and may cause compilation a error based on fully programable computations. As a consequence, different limitation was imposed: associated constants cannot be used as `const` parameters in types, they can only be used as values in code.

**Use case** Due to lack of negative reasoning, as of yet, in rust compiler we cannot express type inequality. The only viable solution would be to a blanket impl and stating that two types are different if they are not the same type, since such a blanket would apply to user defined types as well.

Such CT validation that types are all different is required to decide if glsl variable layout locations do not overlap. Such overlap technically is valid if layout components are used but that could be accounted for.

We use associated constants and conditionally panicking `const` function to check that location ranges do not overlap.

### 3.3.3 Effect system

**Problem** First class effect system is yet another non-existent feature that would be of immense value in this specific case. Ability to type check function invocation context in OpenGL would be especially useful as we could encode presence of appropriate object binding using an effect.

**Solution** We instead were forced to opt for more error prone and verbose approach. Objects like textures or buffers can generate binder objects which in their constructor bind object to context and in destructor unbind it. This lets us control binding using lexical scope but does not in any way prevent overriding of the binding.

**Usae Case** As already mentioned effect system would greatly improve the handling of context bindings in terms of statically verifiable correctness, as well as, user and developer experience.

### 3.3.4 Application of existing features

#### 3.3.5 Markers

**Problem** Enumeration types are a core component of almost all currently used programming languages. In recent years many languages even gained ability to store dynamic data in enum variants. Such enums provide simple mechanism for statically typed polymorphism with dynamic variants. However sometimes this dynamic-ness of enums is a hurdle causing constant match or

switch statements to crop you all over the code base, producing clutter and boilerplate. Sometime one simply wishes to encode static configuration based on a closed set of possible values.

**Solution** Markers are traits and types which don't provide any runtime behavior, but rather exist for purposes of conveying information and constrains on a type level. Marker traits provide no useful functionality but rather serve to impose relations and logical division on types.

Marker types don't hold any data and as such don't exist at runtime (they occupy zero bytes and are formally called Zero Sized Types - ZST). It is even possible to marker types to have type, lifetime and const parameters by using special compiler intrinsic datatype - `PhantomData` - which binds parameters but does not hold any value.

Marker traits along with marker types can be used as:

- compile-time enums – by limiting access to a marker using item visibility qualifiers we strictly control what types implement given functionality.
- marker trait based relations – we can express relations between types and make unsound parameter combinations a compile-time error.
- typing external resources – by using `PhantomData` we can attach type information to otherwise untyped parts of API.

**Use Case** We make heavy use of markers to implement entirety of glsl module which consists almost exclusively of ZSTs for purposes of modelling shader `in`, `out` and `uniform` variables. Types representing these variables aggregated into hlists are specified by the user with help of GLSL DSL implemented using lightweight declarative macros.

Marker traits in miscellaneous `_:valid` modules define relations between valid combinations of data types. Buffer in raw OpenGL, due to C's lack of generics, has its buffer populated using `*void` and the documentation enumerates valid types. To make things worse validity of data types changes depending on what's the buffer's target. It is illegal for index buffer to contain anything other than unsigned integers, pixel buffers can contain almost everything and vertex buffers, yet again, can contain only specific combinations of data. By associating a phantom type with a Buffer and using marker trait based validation relations on uploaded data we solve both of these issues.

This methodology can be extended to form **many modes** pattern, in which one uses marker types that implement trait containing generic associated types (GAT) to control behavior in more complex fashion than using non-generic associated types.

**many modes** We use many modes in `Variable<S, L, T, Store>` to abstract over kind of storage used for variable's type member - `Phantom` or `Inline`. `Phantom` uses `PhantomData` as its associated type and effectively discards value and `Inline` keeps it as is.

### 3.3.6 Type State

Type state is very powerful pattern that takes advantage of how rust understands generic types and allows for tracking runtime capabilities at compile-time.

## 3.4 OpenGL wrapper

# Bibliography

- [1] Steve Klabnik and Carol Nichols. *The Rust programming language*. No Starch Press, 2023.