#### GPU hardware

Cosmin Oancea and Troels Henriksen

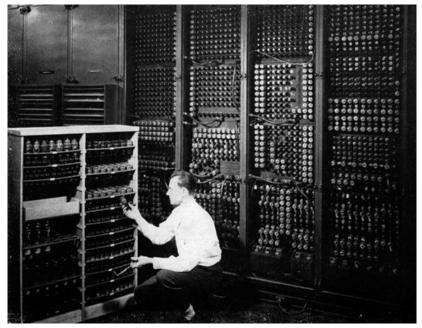
Spetember 2022

#### Hardware Trends

The GPU Architecture

The OpenCL Programming Model

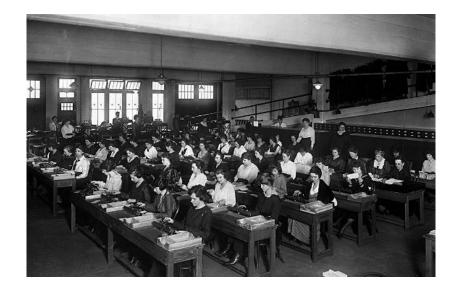
## The first computers were not this



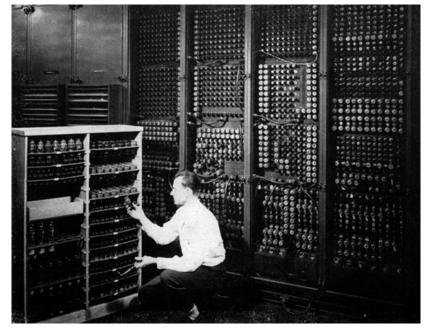
### **But this**



## And if you had a larger problem



## But then they started looking like this













## Then, from around 2005



## Then, from around 2005







### Then, from around 2005

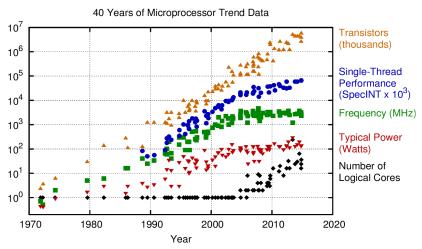


Improvements in *sequential performance* stalled, although computers still got smaller and faster.

### What Changed?

- ► Power complexity  $P_{dynamic} \sim Freq^3$ , preventing us from increasing processor frequency.
- Memory wall, ever-increasing performance gap between processor and memory (which means that memory becomes bottleneck, not processor speed).

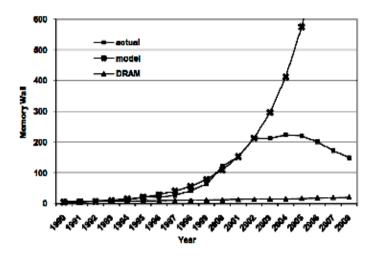
## **CPU** progress



Original data up to the year 2010 collected and plotted by M. Horowitz, F. Labonte, O. Shacham, K. Olukotun, L. Hammond, and C. Batten New plot and data collected for 2010-2015 by K. Rupp

Addressed with more cores.

## The Memory Wall



Memory Wall = processor cycles/memory cycles Addressed with caches (not scalable) and *latency hiding*.

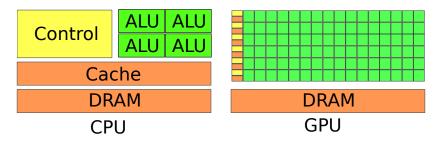
### This is why GPUs are useful

The design of GPUs directly attacks these two problems.

- ► Frequency scaling becomes less of an issue because we can instead use thousands of (slower) cores.
- ► The **memory wall** is partially circumvented by using faster and smaller memory, but mostly by *latency hiding*. With tens of thousands of threads, we can probably find something else to do while some threads are waiting for memory!

Ultimately, GPUs do *throughput processing*, and operations have (relatively) high latency.

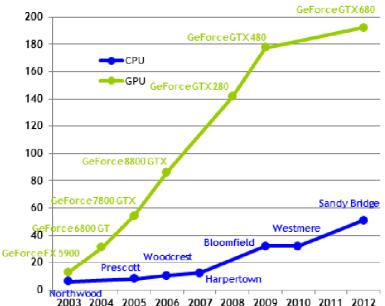
## **CPUs compared to CPUs**



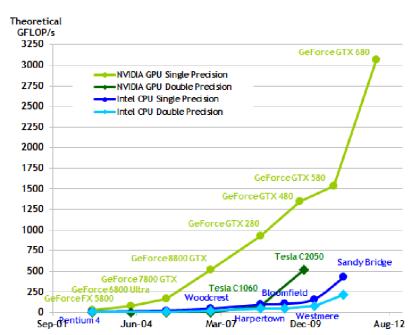
- GPUs have thousands of simple cores and taking full advantage of their compute power requires tens of thousands of threads.
- GPU threads are very restricted in what they can do: no stack, no allocation, limited control flow, etc.
- ► Potential *very high performance* and *lower power usage* compared to CPUs, but programming them is *hard*.

### **GPUs and Memory**





#### **GPUs and GFLOPS**



Hardware Trends

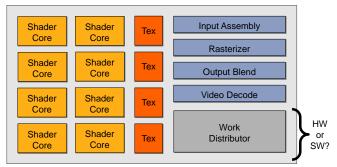
The GPU Architecture

The OpenCL Programming Model

The following slides are taken from the presentation *Introduction* to GPU Architecture by Ofer Rosenberg of AMD.

#### What's in a GPU?

A GPU is a heterogeneous chip multi-processor (highly tuned for graphics)



#### A diffuse reflectance shader

```
sampler mySamp;
Texture2D<float3> myTex;
float3 lightDir;

float4 diffuseShader(float3 norm, float2 uv)
{
   float3 kd;
   kd = myTex.Sample(mySamp, uv);
   kd *= clamp( dot(lightDir, norm), 0.0, 1.0);
   return float4(kd, 1.0);
}
```

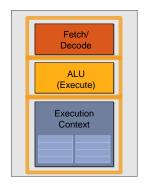
Shader programming model:

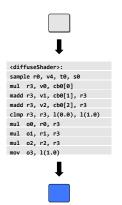
Fragments are processed independently, but there is no explicit parallel programming

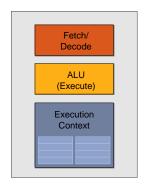
## Compile shader

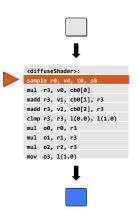
```
1 unshaded fragment input record
sampler mySamp;
Texture2D<float3> myTex;
float3 lightDir;
                                                                                 cdiffuseShader>.
                                                                                 sample r0, v4, t0, s0
                                                                                 mul r3, v0, cb0[0]
float4 diffuseShader(float3 norm, float2 uv)
                                                                                 madd r3, v1, cb0[1], r3
                                                                                 madd r3, v2, cb0[2], r3
                                                                                 clmp r3, r3, 1(0.0), 1(1.0)
  float3 kd:
                                                                                 mul 00, r0, r3
  kd = myTex.Sample(mySamp, uv);
                                                                                 mul o1, r1, r3
  kd *= clamp( dot(lightDir, norm), 0.0, 1.0);
                                                                                 mul o2, r2, r3
  return float4(kd, 1.0);
                                                                                 mov o3, 1(1.0)
```

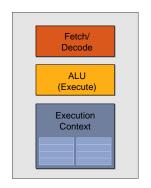
1 shaded fragment output record

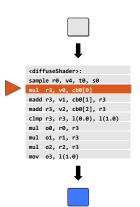


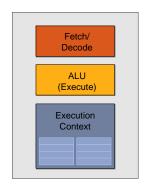


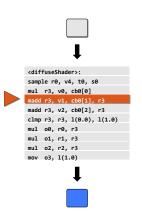


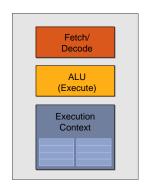


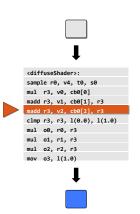


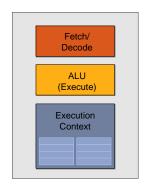


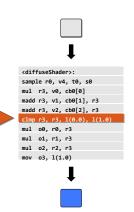


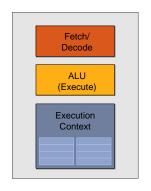


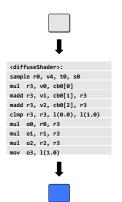




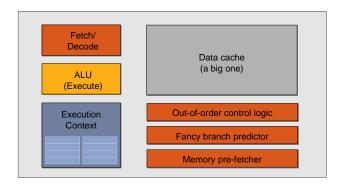




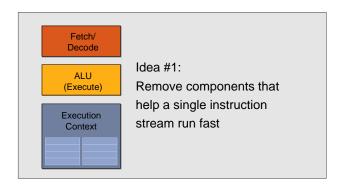




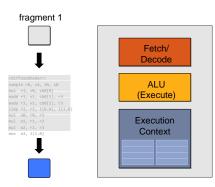
# "CPU-style" cores

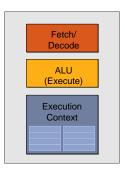


## Slimming down



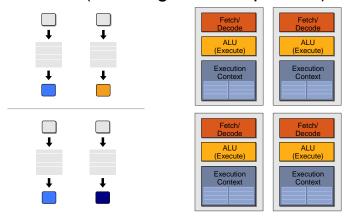
## Two cores (two fragments in parallel)



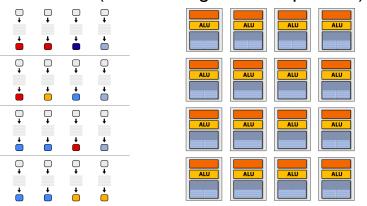




## Four cores (four fragments in parallel)

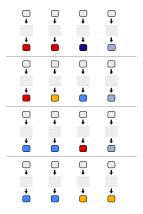


## Sixteen cores (sixteen fragments in parallel)



16 cores = 16 simultaneous instruction streams

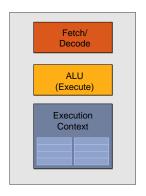
## Instruction stream sharing



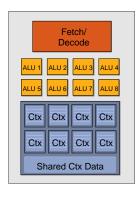
But ... many fragments should be able to share an instruction stream!

```
<diffuseShader>:
sample r0, v4, t0, s0
mul r1, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
cdd r73, v2, cb0[2], r3
cdd r0, r0, r3
mul o0, r0, r3
mul o1, r1, r3
mul o2, r2, r3
mov o3, 1(1.0)
```

# Recall: simple processing core



### Add ALUs

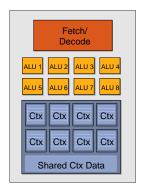


#### Idea #2:

Amortize cost/complexity of managing an instruction stream across many ALUs

SIMD processing

## Modifying the shader

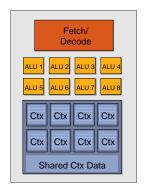


```
cdiffuseShader>:
sample r0, v4, t0, s0
mul r3, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
clmp r3, r3, l(0.0), l(1.0)
mul o0, r0, r3
mul o1, r1, r3
mul o2, r2, r3
mov o3, l(1.0)
```

Original compiled shader:

Processes one fragment using scalar ops on scalar registers

## Modifying the shader



```
VEC8_dif useShader>:

VEC8_sample vec_r0, vec_v4, t0, vec_s0

VEC8_mu1 vec_r3, vec_v4, cb0[0]

VEC8_madd vec_r3, vec_v1, cb0[1], vec_r3

VEC8_madd vec_r3, vec_r3, (10.0), 1(1.0)

VEC8_mu1 vec_00, vec_r0, vec_r3

VEC8_mu1 vec_01, vec_r1, vec_r3

VEC8_mu1 vec_02, vec_r1, vec_r3

VEC8_mu1 vec_02, vec_r2, vec_r3

VEC8_mu1 vec_02, vec_r2, vec_r3

VEC8_mu2 vec_03, vec_r2, vec_r3

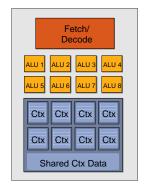
VEC8_mu2 vec_03, vec_r2, vec_r3

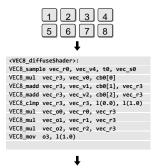
VEC8_mu2 vec_03, vec_r2, vec_r3
```

New compiled shader:

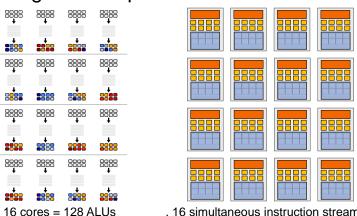
Processes eight fragments using vector ops on vector registers

## Modifying the shader

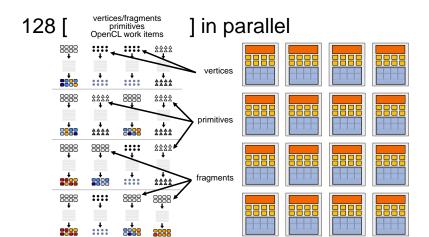


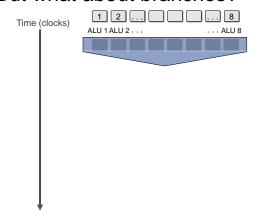


## 128 fragments in parallel

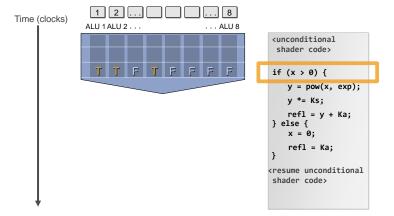


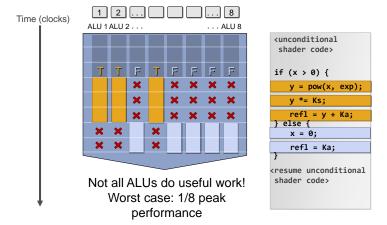
, 16 simultaneous instruction streams

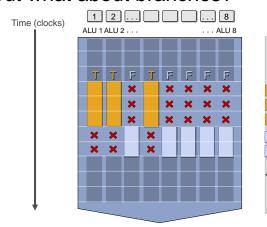




```
<unconditional</pre>
 shader code>
if (x > 0) {
    y = pow(x, exp);
    y *= Ks;
    refl = y + Ka;
} else {
    x = 0;
    refl = Ka;
kresume unconditional
shader code>
```







```
<unconditional
shader code>

if (x > 0) {
    y = pow(x, exp);
    y *= Ks;
    refl = y + Ka;
} else {
    x = 0;
    refl = Ka;
}
<resume unconditional
shader code>
```

### Clarification

### SIMD processing does not imply SIMD instructions

- · Option 1: explicit vector instructions
  - x86 SSE, AVX, Intel Larrabee
- · Option 2: scalar instructions, implicit HW vectorization
  - HW determines instruction stream sharing across ALUs (amount of sharing hidden from software)
  - NVIDIA GeForce ("SIMT" warps), ATI Radeon architectures ("wavefronts")







In practice: 16 to 64 fragments share an instruction stream.

# Stalls!

Stalls occur when a core cannot run the next instruction because of a dependency on a previous operation.

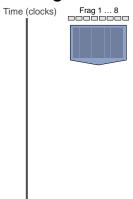
Texture access latency = 100's to 1000's of cycles

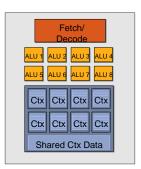
We've removed the fancy caches and logic that helps avoid stalls.

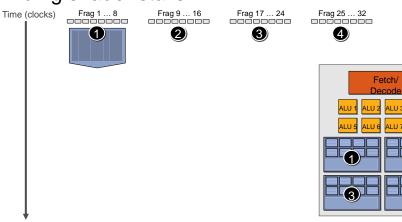
But we have LOTS of independent fragments.

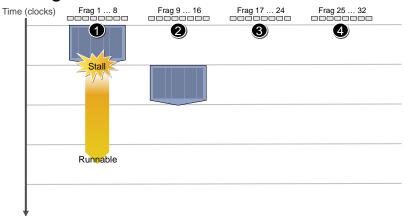
### Idea #3:

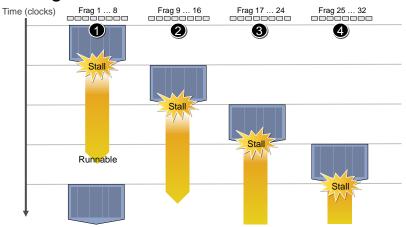
Interleave processing of many fragments on a single core to avoid stalls caused by high latency operations.



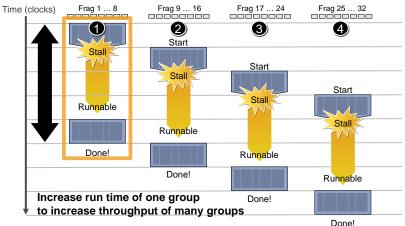




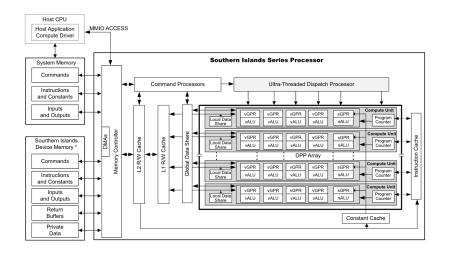




## Throughput!

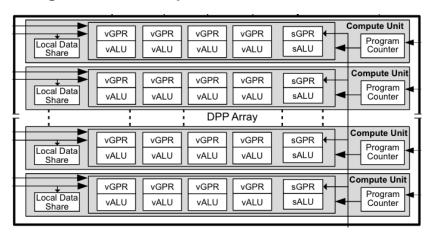


## The GPU we will be using: Radeon HD 7800<sup>1</sup>



<sup>&</sup>lt;sup>1</sup>https://developer.amd.com/wordpress/media/2012/12/ AMD\_Southern\_Islands\_Instruction\_Set\_Architecture.pdf

## **Zooming in on the Compute Units**



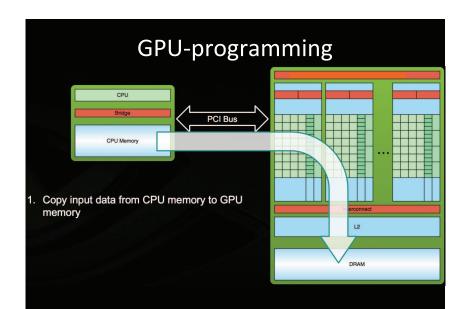
- Each vector-ALU executes a wavefront of 64 work-items over four clock cycles.
- Many wavefronts in flight at once to hide latency.

Hardware Trends

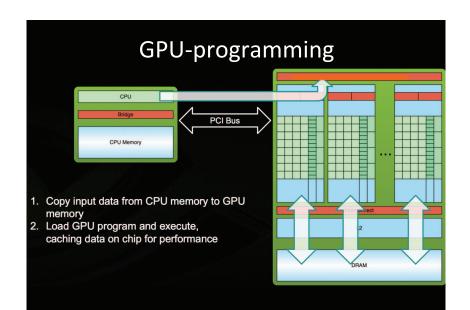
The GPU Architecture

The OpenCL Programming Model

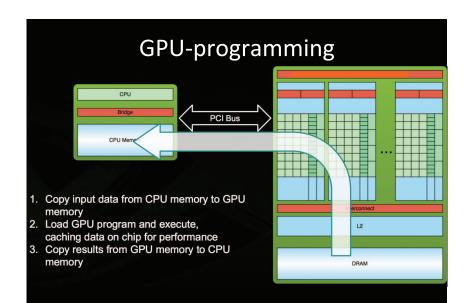
## **GPU** programming



## **GPU** programming



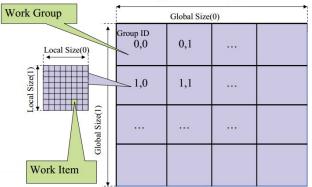
## **GPU** programming



### OpenCL is an SIMT model

Single Instruction Multiple Threads means we provide a sequential function that is executed in parallel by multiple threads ("work items" in OpenCL).

## OpenCL NDRange Configuration



Threads are arranged in *workgroups*, which form an *NDRange* (often called *grid*).