

# Advanced Topics in Programming Languages (ATPL)

## Synthesizing Futhark Quantum Circuit Simulation

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December 3, 2024

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## Context: Data-Parallel Programming

- High-performance demands have lead to SIMD architectures, which are well-suited for data-parallel programming.
- Parallel solutions to many problems benefit from support for *nested parallelism* (e.g., maps within maps) and even for *irregular nested parallelism* (i.e., performing uneven amounts of parallel work in parallel).

## Gap: A large performance penalty (open research problem)

- General approaches to supporting irregular parallelism exist (i.e., Blelloch's flattening transformation, as implemented in NESL) but it has a **significant overhead in practice**, due to the administration of segment descriptors (i.e., flag vectors).

## Futhark – 101

- Futhark is a strict data-parallel pure functional programming language. It is also a compiler aimed at running Futhark programs efficiently on massively-parallel hardware (e.g., GPUs).
- Futhark shows impressive performance numbers! In some cases it beats even hand-optimized OpenCL and CUDA programs.
- Futhark incrementally turns nested regular data-parallelism into flat data-parallelism and creates multiple versions of code that utilises different levels of parallelism.
- Futhark provides a series of *zero-cost abstraction* features including **higher-order modules**, **type polymorphism**, a restricted form of **higher-order functions**, and **size polymorphism**.



## Futhark – the Language

- Syntactically and conceptually similar to Haskell and ML.
- Simple functional language with conditionals and support for looping.
- Datatypes are booleans, basic numeric types, such as `i32` and `f64`, and the types of multi-dimensional arrays (`[] $\tau$` , where  $\tau$  is any type.)
- Futhark is augmented with a set of **parallel** SOACs for operating on arrays.

*-- Simple looping*

```
let fact(n: i32): i32 =  
  loop res = 1  
    for i < n do (i+1) * res
```

*-- Looping*

```
let fib(n: i32): i32 =  
  fst(loop (x,y) = (1,1)  
    for i<n do (y,x+y))
```

*-- Initialize and fill*

```
let fib (n: i32): []i32 =  
  loop arr = replicate n 1  
    for i < n-2 do  
      let arr[i+2] =  
        arr[i] + arr[i+1]  
    in arr
```

## Futhark Second-Order Array Combinators (SOACs)

`map`  $f$   $a$  Probably the simplest parallel SOAC. Maps  $f$  over the outer dimension of  $a$ .

`map2`  $f$   $a_1$   $a_2$  Applies the binary function  $f$  pair-wise to the outer elements of  $a_1$  and  $a_2$ .

`reduce`  $\oplus$   $ne$   $a$  Similar to Haskell's `foldl` and `foldr`, but assumes  $A \oplus (B \oplus C) = (A \oplus B) \oplus C$  and  $ne \oplus A = A \oplus ne = A$ .

`scan`  $\oplus$   $ne$   $a$  Parallel inclusive prefix scan. Returns an array of the same shape as  $a$  but with each element being identical to the result of reducing the prefix of the array  $a$ .  
The `scan` SOAC is critical for implementing parallel algorithms.<sup>1</sup>

Other built-in data-parallel operations include `filter` and `scatter`.

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<sup>1</sup>Blelloch, Guy E. 1989. "Scans as Primitive Parallel Operations." IEEE Transactions on Computers 38(11), pp. 1526–1538.

## Size Polymorphism and Statically-Enforced Size-Constraints

Programmers may specify that functions are parametric in array sizes:

```
def dotprod [n] (xs: [n]f64) (ys: [n]f64) : f64 =  
  f64.sum (map2 (*) xs ys)
```

Here is the type of `map2`:

```
val map2 'a 'b 'c [n] : (a → b → c) → [n]a → [n]b → [n]c
```

We could also have written:

```
def dotp (n:i64) (x: [n]f64) (y: [n]f64): f64 =  
  f64.sum (map2 (*) x y)
```

**Notice:** Size polymorphism makes the size argument implicit; instances are inserted by the compiler.

## Dependent Sizes

```
val iota : (n:i64) : [n]i64
```

```
iota (x+y) : [x+y]i64
```

Like in dependantly typed-languages:

```
val concat 'a [n][m] : [n]a → [m]a → [n+m]a
```

```
concat (iota (x+y)) (iota z) : [x+y+z]i64
```

## Existential Size Types

```
val filter 'a [n] : (a→bool) → [n]a → ?[m].[m]a
```

Existential size-types are “opened implicitly” in `let`-bindings making it possible for type inference to track sizes and report size errors statically.



## Size Expressions are Syntactic

**val** *flatten*     'a [n][m] : [n][m]a → [n\*m]a

More interestingly:

**val** *unflatten*   'a [n][m] : [n\*m]a → [n][m]a

**val** *split*        'a [n][m] : [n+m]a → ([n]a, [m]a)

## Size Constraints

- A size constraint  $x \geq [n*2]f64$  has type  $[n*2]f64$ .
- It is checked statically that  $x$  is an  $f64$ -array and dynamically that the array has size  $n*2$ .

## Static Properties and Restrictions

- It is a static error if a size-argument cannot be determined by the context:

```
def iota [n] () = iota n      -- invalid!
```

- For an array of arrays, all elements must be of the same size.
- It is enforced that coherent sizes are given to functions like `matmul`:

```
def matmul [n][p][k] (A: [n][p]f64) (B: [p][k]f64)
  : [n][k]f64 =
  map (\Arow → map (\Bcol → dotprod Arow Bcol)
      (transpose B))
      A
```

## Dynamic Checks

- Array indexing is checked dynamically, but may be discharged statically.
- `iota n` may fail dynamically (if  $n < 0$ ).

## Generating Futhark Programs from Standard ML

- Futhark is **not** meant to be a general programming language.
- For instance, it does not directly support recursive data types.

## Use a Host Language (e.g., Standard ML) for Generating Futhark Code!

- It is straightforward to embed a simple Futhark expression language in a functional language using algebraic data types.
- Use a monad to keep track of local let-bindings.
- A simple pretty-printer outputs well-formed Futhark code.

## Futhark Matrix Utility Library

```

module type MAT = {
  type t
  val tensor [m][n][p][q] : [m][n]t → [p][q]t → [m*p][n*q]t
  val control      [m] : [m][m]t → [2*m][2*m]t
  val matmul       [m][n][k] : [m][n]t → [n][k]t → [m][k]t
  ...
}

```

## Implementation

```

module type NUM = { type t
  val i64 : i64 → t
  val *   : t → t → t
  val +   : t → t → t
}

module mk_mat (T : NUM) : MAT with t = T.t = {
  type t = T.t
  def scale [m][n] (s:t) (a:[m][n]t) : [m][n]t =
    map (map (T.((s*))) a
  def tensor [m][n][p][q] (a:[m][n]t) (b:[p][q]t) : [m*p][n*q]t =
    map (map (\s → scale s b)) a
    |> map transpose |> flatten |> map flatten
  ...
}

```

## Futhark Expressions and Bindings in Standard ML

```
structure Futhark :> sig
  type var = string
  type ty = string
  datatype exp = VAR of var
               | BINOP of string * exp * exp
               | CONST of string
               | APP of string * exp list
               | ARR of exp list
               | TYPED of exp * ty
               | SEL of int * exp

  type  $\alpha$  M
  val ret      :  $\alpha \rightarrow \alpha$  M
  val >>=     :  $\alpha$  M * ( $\alpha \rightarrow \beta$  M)  $\rightarrow \beta$  M
  val LetNamed : var  $\rightarrow$  exp  $\rightarrow$  var M
  val FunNamed : var  $\rightarrow$  (exp  $\rightarrow$  exp M)  $\rightarrow$  ty  $\rightarrow$  ty  $\rightarrow$  var M
  val Let      : exp  $\rightarrow$  var M
  val Fun      : (exp  $\rightarrow$  exp M)  $\rightarrow$  ty  $\rightarrow$  ty  $\rightarrow$  var M
  val run      : exp M  $\rightarrow$  string
  val runBinds : unit M  $\rightarrow$  string
end
```

## Compiling Circuits To Futhark Matrix Expression

```

fun comp (t:Circuit.t) : Futhark.exp Futhark.M =
  let open Circuit
    open Futhark infix >>=
  in case t of
    I ⇒ ret (APP("I",[ ]))
  | X ⇒ ret (APP("X",[ ]))
  | Y ⇒ ret (APP("Y",[ ]))
  | Z ⇒ ret (APP("Z",[ ]))
  | H ⇒ ret (APP("H",[ ]))
  | SW ⇒ ret (APP("SW",[ ]))
  | Seq(t1,t2) ⇒
    comp t1 >>= (fn e1 ⇒
    comp t2 >>= (fn e2 ⇒
      let val n = Int.toString (pow2 (dim t))
        val ty = "[" ^ n ^ "]"[" ^ n ^ "]"C.complex"
      in ret (TYPED(APP("matmul", [e2,e1]),ty))
      end))
  | Tensor(t1,t2) ⇒
    comp t1 >>= (fn e1 ⇒
    comp t2 >>= (fn e2 ⇒
      let val n = Int.toString (pow2 (dim t))
        val ty = "[" ^ n ^ "]"[" ^ n ^ "]"C.complex"
      in ret (TYPED(APP("tensor", [e1,e2]),ty))
      end))
  | C t' ⇒ comp t' >>= (fn e ⇒
    let val n = Int.toString (pow2 (dim t))
      val ty = "[" ^ n ^ "]"[" ^ n ^ "]"C.complex"
    in ret (TYPED(APP("control",[e]),ty))
    end)
  end

```

## Prelude with Futhark Quantum Gates

```

import "lib/github.com/diku-dk/complex/complex"
import "futlib"
module C = mk_complex(f64)
module mat = mk_mat(C)
open mat

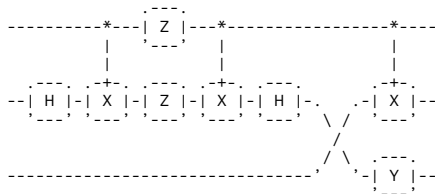
def I () : [2][2]C.complex =
  [[C.i64 1,C.i64 0], [C.i64 0,C.i64 1]]
def H () : [2][2]C.complex =
  [[C.mk_re (1.0 / f64.sqrt(2.0)),C.mk_re (1.0 / f64.sqrt(2.0))],
   [C.mk_re (1.0 / f64.sqrt(2.0)),C.mk_re ((-1.0) / f64.sqrt(2.0))]]
def X () : [2][2]C.complex =
  [[C.i64 0,C.i64 1], [C.i64 1,C.i64 0]]
def Y () : [2][2]C.complex =
  [[C.i64 0,C.mk_im (-1)], [C.mk_im 1,C.i64 0]]
def Z () : [2][2]C.complex =
  [[C.i64 1,C.i64 0], [C.i64 0,C.i64 (-1)]]
def SW () : [4][4]C.complex =
  [[C.i64 1,C.i64 0,C.i64 0,C.i64 0], [C.i64 0,C.i64 0,C.i64 1,C.i64 0],
   [C.i64 0,C.i64 1,C.i64 0,C.i64 0], [C.i64 0,C.i64 0,C.i64 0,C.i64 1]]

$ futhark repl prelude.fut
[0]> tensor (matmul (X()) (H())) (Y())
[[ (0.0, 0.0), (0.0, -0.7071067811865475), (-0.0, 0.0), (0.0, 0.7071067811865475)],
  [(0.0, 0.7071067811865475), (0.0, 0.0), (-0.0, -0.7071067811865475), (-0.0, 0.0)],
  [(0.0, 0.0), (0.0, -0.7071067811865475), (0.0, 0.0), (0.0, -0.7071067811865475)],
  [(0.0, 0.7071067811865475), (0.0, 0.0), (0.0, 0.7071067811865475), (0.0, 0.0)]]
[1]>

```

## Example Generated Futhark Matrix Expression

Circuit: (I \*\* H oo C X oo Z \*\* Z oo C X oo I \*\* H) \*\* I oo I \*\* SW oo C X \*\* Y



$$\begin{bmatrix} -i & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & i & 0 & 0 & 0 & 0 & 0 \\ 0 & -i & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & i & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & -i & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & i \\ 0 & 0 & 0 & 0 & -i & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & i & 0 \end{bmatrix}$$

## Generated Futhark Matrix Expression:

```
def m = (matmul ((tensor ((control (X()))):>[4][4]C.complex)
(Y())):>[8][8]C.complex) ((matmul ((tensor (I()) (SW())):>[8][8]C.complex)
((tensor ((matmul ((tensor (I()) (H())):>[4][4]C.complex) ((matmul ((control
(X())):>[4][4]C.complex) ((matmul ((tensor (Z()) (Z())):>[4][4]C.complex)
((matmul ((control (X())):>[4][4]C.complex) ((tensor (I())
(H())):>[4][4]C.complex)):>[4][4]C.complex)):>[4][4]C.complex))
:>[4][4]C.complex)):>[4][4]C.complex)
(I())):>[8][8]C.complex)):>[8][8]C.complex)):>[8][8]C.complex)
```



## Can We Simulate Quantum Circuits Without Tensor-Explosions?

Instead of first generating a matrix from a circuit and then applying the matrix to a state vector, perhaps “apply the circuit to the state vector”:

```
type vec = C.complex vector
val interp : Circuit.t → vec → vec
```

### Use the “Vec Trick”<sup>2</sup> and the Definition of Matrix Multiply:

- If  $A : [m][m]\mathbb{C}$  and  $B : [n][n]\mathbb{C}$ , then  $(A \otimes B) : [mn][mn]\mathbb{C}$ .
- $(A \otimes B) v = \text{vec}(BVA^T) = \text{vec}((A(BV)^T)^T)$ , where  $V = \text{unvec } v$
- $\text{vec } V : [nm]\mathbb{C}$  is the result of “stacking” all columns in  $V : [n][m]\mathbb{C}$ .
- $\text{unvec } v$  is the matrix  $V : [n][m]\mathbb{C}$  such that  $\text{vec } V = v$ .
- $\llbracket C_A \rrbracket V = (\text{map } (\text{interp } C_A) V^T)^T$ , where  $\llbracket C_A \rrbracket = A$

---

<sup>2</sup>See [https://en.wikipedia.org/wiki/Kronecker\\_product](https://en.wikipedia.org/wiki/Kronecker_product)

## Utility Functions for Kronecker-Free Quantum Simulation

```

fun flatten (a:mat) : vec =
  let val (rows,cols) = M.dimensions a
  in Vector.tabulate(rows * cols,
                    fn i  $\Rightarrow$  M.sub(a,i div cols,i mod cols))
  end

fun unflatten (r,c) (v:vec) : mat =
  if Vector.length v <> r * c then raise Fail "unflatten"
  else M.tabulate(r,c,fn (i,j)  $\Rightarrow$  Vector.sub(v,i*c+j))

fun unvec (r:int,c:int) (v:vec) : mat =
  unflatten (r,c) v |> M.transpose

fun vec (a:mat) : vec =
  M.transpose a |> flatten

fun vecSplit (v:vec) : vec * vec =
  let val n = Vector.length v div 2
  in (VectorSlice.vector(VectorSlice.slice(v,0,SOME n)),
    VectorSlice.vector(VectorSlice.slice(v,n,NONE)))
  end

fun mapRows (f:vec $\rightarrow$ vec) (a:mat) : mat =
  List.tabulate(M.nRows a,
                fn i  $\Rightarrow$  f(M.row i a))
  |> M.fromVectorList

```

## The Kronecker-Free Quantum Simulation Function

```
fun interp (t:Circuit.t) (v:vec) : vec =
  case t of
    Circuit.Seq(t1,t2)  $\Rightarrow$  interp t2 (interp t1 v)
  | Circuit.Tensor(A,B)  $\Rightarrow$ 
    let val V = unvec (pow2(Circuit.dim A),
                      pow2(Circuit.dim B)) v
        val W = mapRows (interp B) (M.transpose V)
        val Y = mapRows (interp A) (M.transpose W)
    in vec Y
    end
  | Circuit.C A  $\Rightarrow$ 
    let val (v1,v2) = vecSplit v
    in Vector.concat[v1,interp A v2]
    end
  | Circuit.I  $\Rightarrow$  v
  | _  $\Rightarrow$  eval t v
```

## Derivation for the Tensor Case

$$\begin{aligned}
 \text{interp } (A \otimes B) v &= \text{vec}((A(BV)^T)^T) & BV &= (\text{map } (\text{interp } B) (V^T))^T \\
 &= \text{vec}((A((\text{map } (\text{interp } B) (V^T))^T)^T)^T) \\
 &= \text{vec}((A(\text{map } (\text{interp } B) (V^T)))^T) \\
 &= \text{let } W = \text{map } (\text{interp } B) (V^T) \\
 &\quad \text{in } \text{vec}((AW)^T) & AW &= (\text{map } (\text{interp } A) (W^T))^T \\
 &= \text{let } W = \text{map } (\text{interp } B) (V^T) \\
 &\quad \text{in } \text{vec}(((\text{map } (\text{interp } A) (W^T))^T)^T) \\
 &= \text{let } W = \text{map } (\text{interp } B) (V^T) \\
 &\quad \text{in } \text{vec}(\text{map } (\text{interp } A) (W^T)) \\
 &= \text{let } W = \text{map } (\text{interp } B) (V^T) \\
 &\quad \text{let } Y = \text{map } (\text{interp } A) (W^T) \\
 &\quad \text{in } \text{vec}(Y)
 \end{aligned}$$

qed.

## Utilities for Futhark Quantum Simulation Synthesis

```

fun allI (t:Circuit.t) : bool =
  let open Circuit
  in case t of I  $\Rightarrow$  true
    | Tensor(a,b)  $\Rightarrow$  allI a andalso allI b
    | Seq(a,b)  $\Rightarrow$  allI a andalso allI b
    | _  $\Rightarrow$  false
end

fun vecTyFromDim d =
  "[" ^ Int.toString(pow2 d) ^ "]C.complex"

fun FunC A (f:F.exp  $\rightarrow$  F.exp F.M) : F.var option F.M =
  if allI A then ret NONE
  else let val ty = vecTyFromDim (Circuit.dim A)
    in Fun f ty ty >=> (ret o SOME)
    end

fun splitF d v =
  let val ty = "[" ^ Int.toString (pow2 d) ^ "+" ^
    Int.toString (pow2 d) ^ "]C.complex"
  in APP("split",[TYPED(v,ty)])
  end

fun concatF d a b = TYPED(APP("concat",[a,b]),vecTyFromDim d)
fun unvecF (e,ty) = APP("unvec", [TYPED(e,ty)])
fun vecF e = APP("vec",[e])
fun mapF f e = APP("map", [VAR f,e])
fun matvecmulF m v = APP("matvecmul", [m,v])
fun transposeF m = APP("transpose", [m])
fun mapF' NONE e = e
  | mapF' (SOME f) e = mapF f e

```

## Futhark Quantum Simulation Synthesis

```

fun icomp (t:Circuit.t) (v:exp) : exp M =
  case t of
    Circuit.I ⇒ ret v
  | Circuit.Seq(t1,t2) ⇒ icomp t1 v >>= (icomp t2)
  | Circuit.C t' ⇒
    Let (splitF (Circuit.dim t') v) >>= (fn p ⇒
      icomp t' (SEL(1,VAR p)) >>= (fn v1 ⇒
        ret (concatF (Circuit.dim t) (SEL(0,VAR p)) v1)))
  | Circuit.Tensor(A,B) ⇒
    FunC A (icomp A) >>= (fn Af ⇒
      FunC B (icomp B) >>= (fn Bf ⇒
        let val dA = pow2(Circuit.dim A)
          val dB = pow2(Circuit.dim B)
          val ty = "[" ^ Int.toString dA ^ "*" ^
            Int.toString dB ^ "]"C.complex"
        in Let (unvecF(v,ty)) >>= (fn V ⇒
          Let (mapF' Bf (transposeF (VAR V))) >>= (fn W ⇒
            Let (mapF' Af (transposeF (VAR W))) >>= (fn Y ⇒
              ret (TYPED(vecF (VAR Y),vecTyFromDim (Circuit.dim t))))))
        end))
  | Circuit.H ⇒ ret (matvecmulF (APP("H",[])) v)
  | Circuit.X ⇒ ret (matvecmulF (APP("X",[])) v)
  | Circuit.Y ⇒ ret (matvecmulF (APP("Y",[])) v)
  | Circuit.Z ⇒ ret (matvecmulF (APP("Z",[])) v)
  | Circuit.SW ⇒ ret (matvecmulF (APP("SW",[])) v)

```

## Example Simulator

```

def f (v6:[8]C.complex) : [8]C.complex =
  let f0 (v7:[4]C.complex) : [4]C.complex =
    let f1 (v8:[2]C.complex) : [2]C.complex = matvecmul (H()) v8
    let v9 = unvec (v7:>[2*2]C.complex)
    let v10 = map f1 (transpose v9)
    let v11 = transpose v10
    let v12 = split (((vec v11):>[4]C.complex):>[2+2]C.complex)
    let f2 (v13:[2]C.complex) : [2]C.complex = matvecmul (Z()) v13
    let f3 (v14:[2]C.complex) : [2]C.complex = matvecmul (Z()) v14
    let v15 = unvec(((concat(v12.0)(matvecmul (X())(v12.1))):>[4]C.complex):>[2*2]C.complex)
    let v16 = map f3 (transpose v15)
    let v17 = map f2 (transpose v16)
    let v18 = split (((vec v17):>[4]C.complex):>[2+2]C.complex)
    let f4 (v19:[2]C.complex) : [2]C.complex = matvecmul (H()) v19
    let v20 = unvec(((concat (v18.0)(matvecmul(X())(v18.1))):>[4]C.complex):>[2*2]C.complex)
    let v21 = map f4 (transpose v20)
    let v22 = transpose v21
  in (vec v22):>[4]C.complex
let v23 = unvec (v6:>[4*2]C.complex)
let v24 = transpose v23
let v25 = map f0 (transpose v24)
let f5 (v26:[4]C.complex) : [4]C.complex = matvecmul (SW()) v26
let v27 = unvec (((vec v25):>[8]C.complex):>[2*4]C.complex)
let v28 = map f5 (transpose v27)
let v29 = transpose v28
let f6 (v30:[4]C.complex) : [4]C.complex =
  let v31 = split (v30:>[2+2]C.complex)
  in (concat (v31.0) (matvecmul (X()) (v31.1))):>[4]C.complex
let f7 (v32:[2]C.complex) : [2]C.complex = matvecmul (Y()) v32
let v33 = unvec (((vec v29):>[8]C.complex):>[4*2]C.complex)
let v34 = map f7 (transpose v33)
let v35 = map f6 (transpose v34)
in (vec v35):>[8]C.complex

```

## Exercises

See <http://github.com/diku-dk/atpl-sml-quantum>

- Clone the repository and read the README.md file.

## Project Ideas

- The README.md file also mentions a series of project ideas.