PAT - Assignment 4

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Task 1

The paper shows examples of full-inversion, partial-inversion and semi-inversion. Below is four examples of inversions.

Multiplication to division:

This is a partial inversion with index sets: $I = \{2\}$ and $O = \{1\}$. The multiplication function mul(x, y) with inputs x and y and output z is partially inverted to the division function div(z, y).

Rules:

```
\begin{split} & \operatorname{mul}(0,y) \to 0 \\ & \operatorname{mul}(s(x),y) \to \operatorname{add}(y,\operatorname{mul}(x,y)) \\ & \operatorname{Inverted rules:} \\ & \operatorname{div}(0,y) \to 0 \\ & \operatorname{div}(z,y) \to s(x) \leftarrow \operatorname{sub}(z,y) \wedge \operatorname{div}(w,y) \end{split}
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Symmetric encrypter to decrypter This is a partial inversion with index sets: $I = \{2\}$ and $O = \{1\}$. The symmetric encrypter function is partially inverted to produce a decrypter function.

Rules:

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\begin{split} & \operatorname{encrypt}(\operatorname{nil}, \operatorname{key}) \to \operatorname{nil} \\ & \operatorname{encrypt}(x:xs, \operatorname{key}) \to z: zs \leftarrow \operatorname{mod4}(\operatorname{key}) \wedge \operatorname{add}(x,y) \wedge \operatorname{encrypt}(xs, \operatorname{key}) \\ & \operatorname{Inverted\ rules:} \\ & \operatorname{encrypt}^{\{2\}\{1\}}(\operatorname{key}, \operatorname{nil}) \to \operatorname{nil} \\ & \operatorname{encrypt}^{\{2\}\{1\}}(\operatorname{key}, z:zs) \to x: xs \leftarrow \operatorname{mod4}^{\{1\}\emptyset}(\operatorname{key}) \wedge \operatorname{add}^{\{2\}\{1\}}(y,z) \wedge \operatorname{encrypt}^{\{2\}\{1\}}(\operatorname{key}, zs) \end{split}
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Discrete simulation of free fall:

This is a full inversion with index sets: $I = \{0\}$ and $O = \{0\}$. The function for simulating free fall is fully inverted to determine initial conditions from final states.

Rules

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\begin{aligned} &\mathsf{fall}_0(v,h,0) \to \langle v,h \rangle \\ &\mathsf{fall}_0(v_0,h_0,s(t)) \to \langle v,h \rangle \leftarrow \mathsf{add}(v_0,s_5(0)) \land \mathsf{height}(h_0,v_n) \land \mathsf{fall}_0(v_n,h_n,t) \\ &\mathsf{Inverted rules:} \\ &\mathsf{fall}_1(v,h) \to \langle v_0,h_0,s(t) \rangle \leftarrow \mathsf{fall}_1(v_n,h_n) \land \mathsf{add}^{\{2\}\{1\}}(s_{10}(0),v_n) \land \mathsf{height}^{\{2\}\{1\}}(v_n,h_n) \end{aligned}
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Janus language inverter:

This is a full inversion with index sets: $I = \{\}$ and $O = \{1\}$. The Janus language inverter is fully inverted to produce the same procedure.

Rules:

$$\texttt{inv}(\texttt{proc}(u,v)) \to \langle \texttt{proc}(\texttt{name, progr}) \rangle \leftarrow \texttt{invName}(u) \wedge \texttt{inv}(v) \\ \texttt{inv}(+=x,y) \to \langle -=x,y \rangle$$

Inverted rules:

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\texttt{inv}(\texttt{proc}(name, progr)) \to \langle \texttt{proc}(u, v) \rangle \leftarrow \texttt{invName}(name) \land \texttt{inv}(\texttt{progr}) \\ \texttt{inv}(+ = x, y) \to \langle - = x, y \rangle
```

Task 2

Here is the CCS definition for unary subtraction:

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\begin{array}{l} \operatorname{sub}(z,0) \to z \\ \operatorname{sub}(s(z),s(x)) \to \operatorname{sub}(z,x) \end{array}
```

Subtracting zero from any unary number results in the number itself, hence the first case. Subtracting the successor of a number from the successor of another number reduces the problem to subtracting the numbers themselves, hence the second case.

To check that the system rewrites sub(s(s(0)), s(0)) into s(0), we will trace the rewrite steps. First we have

Then we apply the recursive rule to get

$$\mathrm{sub}(s(s(0)),s(0))\to\mathrm{sub}(s(0),0)$$

Then we apply the base case to get

$$\operatorname{sub}(s(0),0) \to s(0)$$

and thus we have that sub(s(s(0)), s(0)) correctly rewrites to s(0).

Task 3

We have chosen to perform semi-inversion on fall with $I = \{1\}$ and $O = \{2\}$. Below is the output of the program.

Inverting rules

```
add\{1,2\}\{\}(0, y) \rightarrow tp(y)
  add\{1,2\}\{\}(s(x), y) \rightarrow tp(s(z)) \mid add\{1,2\}\{\}(x, y) == tp(z)
  height{1,2}{ho, vn} \rightarrow tp(hn) \mid add{1,2}{ho, s(s(s(s(0)))))} == tp(ht)
             , sub{1,2}{}(vn, ht) == tp(hn)
  fall{1,2,3}{}(v, h, 0) \rightarrow tp(v, h)
  fall{1,2,3}{}(v0, h0, s(t)) \rightarrow tp(v, h)
             | add{1,2}{\{}(v0, s(s(s(s(s(s(s(s(s(0)))))))))) == tp(vn)
             , height{1,2}{}(h0, vn) == tp(hn)
             , fall{1,2,3}{}(vn, hn, t) == tp(v, h)
  sub{1,2}{}(0, y) \rightarrow tp(y)
  sub\{1,2\}\{\{s(x), s(z)\}\} -> tp(y) | sub\{1,2\}\{\{s(x, z)\}\} == tp(y)
)
(COMMENT)
-- PURE FULL RULE INVERTER --
(CONDITIONTYPE ORIENTED)
(VAR v h t v0 h0 vn hn ht x y z)
  (fall 3 2) (height 2 1) (add 2 1) (sub 2 1))
  fall{1}{2}(v, h) \rightarrow tp(h, 0, v)
  fall{1}{2}(v0, h) \rightarrow tp(h0, s(t), v) \mid fall{}{1,2}(v, h) == tp(vn, hn, t)
             , height{}{1}(hn) == tp(h0, vn)
             , add{\{1\}(vn) == tp(v0, s(s(s(s(s(s(s(s(s(s(0)))))))))))
  fall\{\{1,2\}(v, h) \rightarrow tp(v, h, 0)\}
  fall{\{1,2\}(v, h) \rightarrow tp(v0, h0, s(t)) \mid fall{\{1,2\}(v, h) == tp(vn, hn, t)\}}
             , height{}{1}(hn) == tp(h0, vn)
             , add{}{1}(vn) == tp(v0, s(s(s(s(s(s(s(s(s(s(s(s)))))))))))
  height{}{1}(hn) \rightarrow tp(h0, vn) \mid sub{}{1}(hn) == tp(vn, ht)
             , add{}{1}(ht) == tp(h0, s(s(s(s(0)))))
  add\{\}\{1\}(y) \rightarrow tp(0, y)
  add\{\{1\}(s(z)) \rightarrow tp(s(x), y) \mid add\{\{1\}(z) == tp(x, y)\}
  sub\{\}\{1\}(y) \rightarrow tp(0, y)
  sub\{\{1\}\{y\} \rightarrow tp(s(x), s(z)) \mid sub\{\{1\}\{y\} == tp(x, z)\}
(COMMENT )
-----
     PARTIAL RULE INVERTER --
_____
(CONDITIONTYPE ORIENTED)
(VAR v h t v0 h0 vn hn ht x y z)
  (fall 3 2) (height 2 1) (add 2 1) (sub 2 1))
(RULES
  fall{1}{2}(v, h) \rightarrow tp(h, 0, v)
  fall{1}{2}(v0, h) \rightarrow tp(h0, s(t), v) \mid fall{}{1,2}(v, h) == tp(vn, hn, t)
             , height{2}{1}(vn, hn) == tp(h0)
```

```
, add\{1,2\}\{1\}(v0, s(s(s(s(s(s(s(s(s(0))))))))), vn) == tp()
  fall\{\{1,2\}(v, h) \rightarrow tp(v, h, 0)\}
  fall\{\{1,2\}(v, h) \rightarrow tp(v0, h0, s(t)) \mid fall\{\{1,2\}(v, h) == tp(vn, hn, t)\}\}
              , height{2}{1}(vn, hn) == tp(h0)
              , add\{2\}\{1\}(s(s(s(s(s(s(s(s(s(0))))))))), vn) == tp(v0)
  height{2}{1}(vn, hn) \rightarrow tp(h0) \mid sub{1}{1}(vn, hn) == tp(ht)
              , add\{2\}\{1\}(s(s(s(s(s(0))))), ht) == tp(h0)
  add\{1,2\}\{1\}(0, y, y) \rightarrow tp()
  add\{1,2\}\{1\}(s(x), y, s(z)) \rightarrow tp() \mid add\{1,2\}\{1\}(x, y, z) == tp()
  add\{2\}\{1\}(y, y) \rightarrow tp(0)
  add\{2\}\{1\}(y, s(z)) \rightarrow tp(s(x)) \mid add\{2\}\{1\}(y, z) == tp(x)
  sub{1}{1}(0, y) \rightarrow tp(y)
  sub\{1\}\{1\}\{s(x), y\} \rightarrow tp(s(z)) \mid sub\{1\}\{1\}\{x, y\} == tp(z)
(COMMENT )
       SEMI RULE INVERTER
(CONDITIONTYPE ORIENTED)
(VAR v h t v0 h0 vn hn ht x y z)
(SIG
  (fall 3 2) (add 2 1) (height 2 1) (sub 2 1))
(RULES
  fall{1}{2}(v, h) \rightarrow tp(h, 0, v)
  fall{1}{2}(v0, h) \rightarrow tp(h0, s(t), v)
             add\{1,2\}\{\{(v0, s(s(s(s(s(s(s(s(s(s(0))))))))))\} == tp(vn)
              , fall{1}{2}(vn, h) == tp(hn, t, v)
              , height{2}{1}(vn, hn) == tp(h0)
  add\{1,2\}\{\}(0, y) \rightarrow tp(y)
  add\{1,2\}\{\}(s(x), y) \rightarrow tp(s(z)) \mid add\{1,2\}\{\}(x, y) == tp(z)
  height{2}{1}(vn, hn) \rightarrow tp(h0) \mid sub{1}{1}(vn, hn) == tp(ht)
              , add\{2\}\{1\}(s(s(s(s(s(0))))), ht) == tp(h0)
  sub{1}{1}(0, y) \rightarrow tp(y)
  sub\{1\}\{1\}\{s(x), y\} \rightarrow tp(s(z)) \mid sub\{1\}\{1\}\{x, y\} == tp(z)
  add{2}{1}(y, y) \rightarrow tp(0)
  add\{2\}\{1\}(y, s(z)) \rightarrow tp(s(x)) \mid add\{2\}\{1\}(y, z) == tp(x)
(COMMENT )
Comparisons
-----
         TRIVIAL RULE INVERTER
_____
- Comparing fall and its inversion
______
```

ORIG

TRIV

Number of functions Number of rules	4 7	5 9
<pre>Functional - Orthogonal . Non-overlapping . Left-linear - EV-Free - Left-to-right determ.</pre>	[x] [x] [x] [x] [x]	
Reversible - Functional - Output-orthogonal . Non-output-overlapping . Right-linear - Strictly non-erasing . Non-erasing . Weakly non-erasing	[x] [x] [x] [x] [x]	[x] [x]
- PURE FULL RULE INVERTER -		
- Comparing fall and its inversion Number of functions Number of rules	ORIG 4 7	FULL 5 9
Number of functions	4	5 9 []
Number of functions Number of rules Functional - Orthogonal . Non-overlapping . Left-linear - EV-Free	4 7 [x] [x] [x] [x]	5 9 [] [] [x] [x]

- Comparing fall and its inversion $% \left(1\right) =\left(1\right) \left(1$

	ORIG	PART
- SEMI RULE INVERTER -		
- Comparing fall and its inversion		
Number of functions Number of rules	ORIG 4 7	SEMI 5 9
<pre>Functional - Orthogonal . Non-overlapping . Left-linear - EV-Free - Left-to-right determ.</pre>	[x] [x] [x] [x]	[x] [x]
Reversible - Functional - Output-orthogonal . Non-output-overlapping . Right-linear - Strictly non-erasing . Non-erasing . Weakly non-erasing	[x] [x] [x] [x] [x]	[] [] [x] [] [] []

The trivial rule inversion function results in overlapping rules and loss of orthogonality, which undermines the functional and reversible properties of the system. The pure full rule inversion function introduces overlapping rules and does not preserve strict non-erasing properties, resulting in a non-functional and non-orthogonal system. Partial rule inversion does not fully address all aspects of the rules, resulting in potential gaps in functionality and orthogonality. The semi-rule inverter maintains left-to-right determinism but fails to preserve orthogonality and strict non-erasing properties, leading to a non-functional and non-reversible system.

Task 4

List reversal semi inversion:

The original program reverses the elements of a list. The semi-inversion of the program, reconstructs the original list given the head of the original list and the reversed list.

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Original program:  \texttt{reverse}([],[]) \\ \texttt{reverse}([H|T],R) \leftarrow \texttt{reverse}(T,\texttt{rev}_t) \land \texttt{append}(\texttt{rev}_t,[H],R) \\ \texttt{Semi-Inversion:} \\ \texttt{reverse\_si}(H,\texttt{rev}_t,r) \rightarrow L \\ \texttt{reverse\_si}(H,[],[H])
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 $\mathtt{reverse_si}(H,[H|\mathtt{rev}_t],R) \to [H|T] \leftarrow \mathtt{reverse}(T,\mathtt{rev}_t)$

The semi-inversion reverse_si reconstructs the original list L given its head H, the reversed tail rev_t and the fully reversed list R. This can be useful in a scenario where only partial knowledge of the input and output is available.