

# Database Management (4331603) - Winter 2023 Solution

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## Question 1(a) [3 marks]

Define the following terms: a). Data items b). Data dictionary c).Meta data

### Solution

Table:

**Table 1.** Database Terms

Term	Definition
<b>Data Items</b>	Basic units of data that cannot be subdivided further. Individual facts or values stored in database fields
<b>Data Dictionary</b>	Centralized repository containing metadata about database structure, tables, columns, and relationships
<b>Metadata</b>	Data about data that describes structure, constraints, and properties of database elements

- **Data Items:** Smallest unit of named data
- **Data Dictionary:** "System Catalog" storing schema definitions
- **Metadata:** Structural information describing the data

### Mnemonic

"DDM - Data Dictionary Manages"

## Question 1(b) [4 marks]

Explain disadvantages of File oriented system.

### Solution

Table:

**Table 2.** Disadvantages of File System

Disadvantage	Description
<b>Data Redundancy</b>	Same data stored in multiple files leading to storage waste
<b>Data Inconsistency</b>	Different versions of same data in different files
<b>Data Isolation</b>	Difficulty in accessing data scattered across multiple files
<b>Security Issues</b>	Limited access control and security mechanisms

- **Redundancy:** Duplication of data across files
- **Inconsistency:** Mismatched data due to poor synchronization
- **Isolation:** Lack of standard format hinders data sharing

- **Integrity:** Hard to enforce constraints across files

### Mnemonic

“RDIS - Really Difficult Information System”

## Question 1(c) [7 marks]

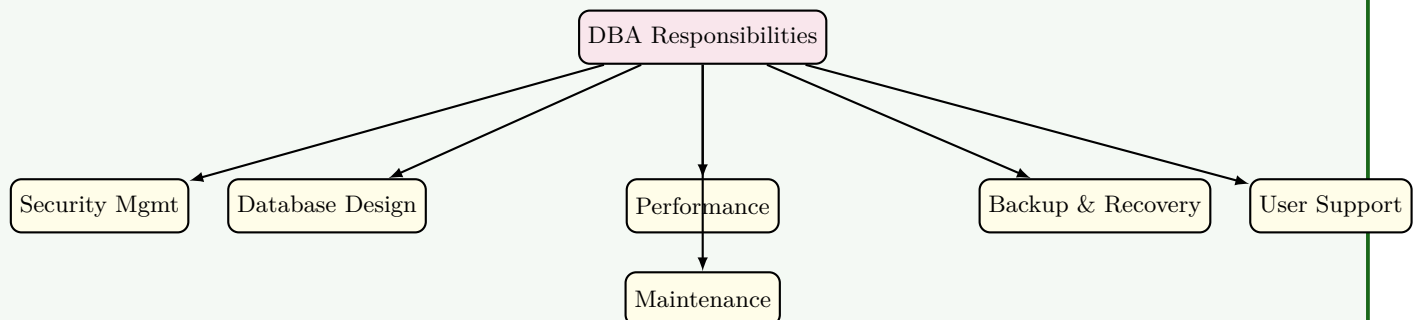
Describe the responsibilities of DBA in detail.

### Solution

Table:

**Table 3.** DBA Responsibilities

Responsibility	Details
<b>Database Design</b>	Creating logical and physical database structures
<b>Security Management</b>	Implementing user access controls and data protection
<b>Performance Monitoring</b>	Optimizing database performance and query execution
<b>Backup &amp; Recovery</b>	Ensuring data safety through regular backups
<b>User Support</b>	Providing technical assistance to database users
<b>System Maintenance</b>	Regular updates, patches, and system optimization



**Figure 1.** Key Responsibilities of Database Administrator

- **Design:** Schema definition and storage structure planning
- **Security:** Granting/revoking access and encryption
- **Recovery:** Disaster recovery planning
- **Tuning:** Indexing and query optimization

### Mnemonic

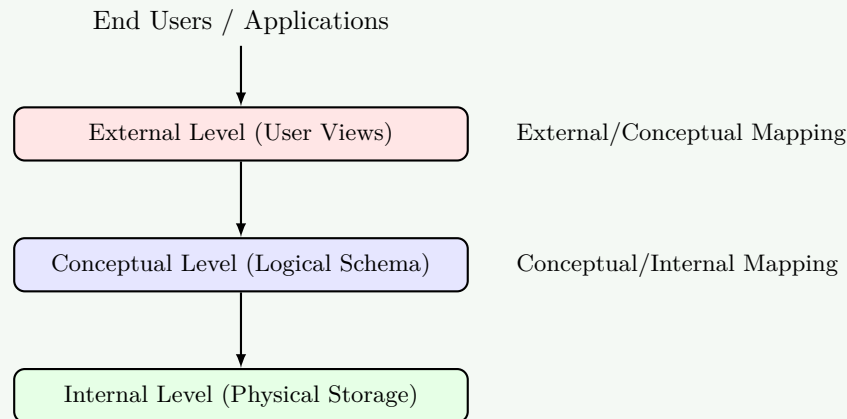
“DSPBUM - Database Specialists Provide Better User Management”

## Question 1(c OR) [7 marks]

Define data abstraction? Explain Three level Architecture of DBMS.

### Solution

**Data Abstraction:** Process of hiding complex implementation details while showing only essential features to users.



**Figure 2.** Three Level Architecture

Table:

**Table 4.** Architecture Levels

Level	Description	Purpose
External Level	User view of database	Individual user perspectives
Conceptual Level	Logical structure of entire database	Overall database organization
Internal Level	Physical storage details	How data is actually stored

- **View Level:** Describes only part of the database
- **Logical Level:** Describes what data is stored and relationships
- **Physical Level:** Describes complex low-level data structures

### Mnemonic

“ECI - Every Computer Industry”

## Question 2(a) [3 marks]

Define the Following Terms :a).Relationship set b).Participation c).Candidate key

### Solution

Table:

**Table 5.** Terminology

Term	Definition
Relationship Set	Collection of relationships of same type between entity sets
Participation	Constraint specifying whether entity occurrence is mandatory in relationship
Candidate Key	Minimal set of attributes that uniquely identifies each entity in entity set

- **Relationship:** Association among several entities
- **Total Participation:** Every entity must participate (double line)
- **Candidate Key:** Super key with no proper subset being a super key

**Mnemonic**

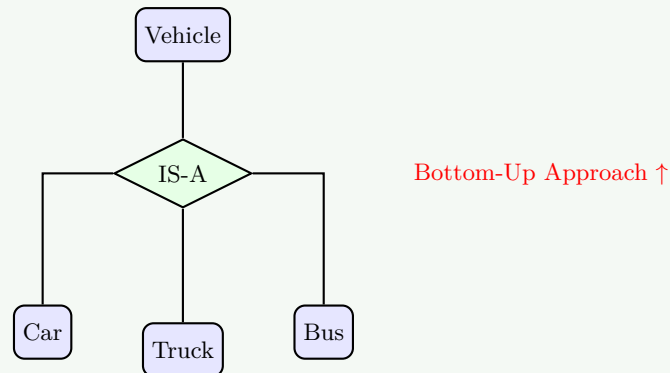
“RPC - Relationship Participation Candidate”

**Question 2(b) [4 marks]**

Explain Generalization with example.

**Solution**

**Generalization:** Bottom-up approach where common attributes of lower-level entities are combined into higher-level entity.



**Figure 3.** Generalization Example

**Table:**

**Table 6.** Generalization Concepts

Concept	Description
<b>Purpose</b>	Reduce redundancy by creating common superclass
<b>Direction</b>	Bottom-up (specific to general)
<b>Example</b>	Car, Truck, Bus → Vehicle

- **Superclass:** Generalized entity containing common properties
- **Subclass:** Specialized entities with unique properties
- **Attribute Inheritance:** Subclasses inherit attributes of superclass

**Mnemonic**

“GBU - Generalization Builds Up”

**Question 2(c) [7 marks]**

Define E-R diagram? Explain different symbols used in E-R diagram with example.

**Solution**

**E-R Diagram:** Graphical representation showing entities, attributes, and relationships in database design.

**Table:**

**Table 7.** ER Diagram Symbols

Symbol	Shape	Usage	Example
<b>Entity</b>	Rectangle	Represents objects	Student, Course
<b>Attribute</b>	Oval	Properties of entities	Name, Age, ID
<b>Relationship</b>	Diamond	Connections between entities	Enrolls, Teaches
<b>Primary Key</b>	Underlined oval	Unique identifier	Student_ID
<b>Multivalued</b>	Double oval	Multiple values	Phone_Numbers
<b>Derived</b>	Dashed oval	Calculated attributes	Age from DOB

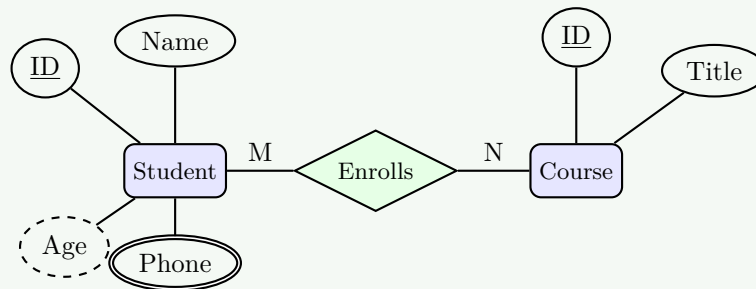


Figure 4. Sample ER Diagram

- **Entity Sets:** Things with independent existence
- **Relationship Sets:** Associations between entities
- **Attribute Types:** Key, Composite, Multivalued, Derived

**Mnemonic**

“EARPM - Every Attribute Represents Proper Meaning”

**Question 2(a OR) [3 marks]**

Define Relational Algebra? List out various operations in relational algebra?

**Solution**

**Relational Algebra:** Formal query language with operations for manipulating relational database tables.  
**Table:**

Table 8. Relational Algebra Operations

Operation Type	Operations
<b>Basic Operations</b>	Select ( $\sigma$ ), Project ( $\pi$ ), Union ( $\cup$ ), Set Difference ( $-$ ), Cartesian Product ( $\times$ )
<b>Additional Operations</b>	Intersection ( $\cap$ ), Join ( $\bowtie$ ), Division ( $\div$ ), Rename ( $\rho$ )

- **Procedural Language:** Specifies what to retrieve and how
- **Operators:** Takes relations as input and produces relation as output

**Mnemonic**

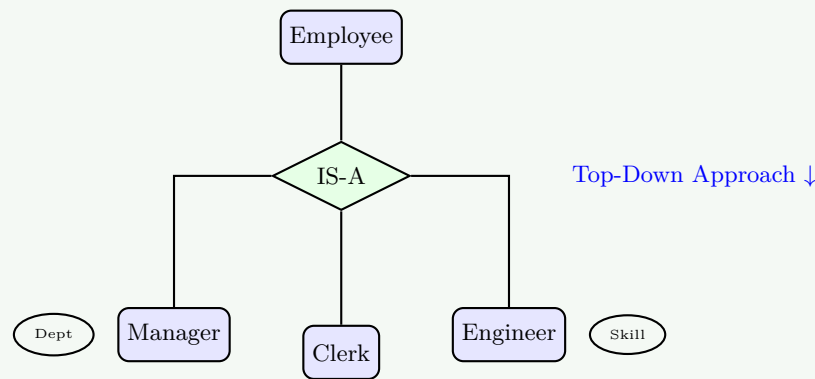
“SPUDC-IJDR - Simple People Use Database Concepts”

**Question 2(b OR) [4 marks]**

Explain Specialization with example.

**Solution**

**Specialization:** Top-down approach where higher-level entity is divided into specialized lower-level entities.



**Figure 5.** Specialization Example

**Table:**

**Table 9.** Specialization Concepts

Concept	Description
<b>Purpose</b>	Create specialized subclasses with unique attributes
<b>Direction</b>	Top-down (general to specific)
<b>Example</b>	Employee → Manager, Clerk, Engineer

- **Subgrouping:** Identifies subsets of entities with distinctive roles
- **Inheritance:** Lower-level entities inherit from higher-level

**Mnemonic**

“STD - Specialization Top Down”

**Question 2(c OR) [7 marks]**

**Define attribute? Explain different types of attributes with example.**

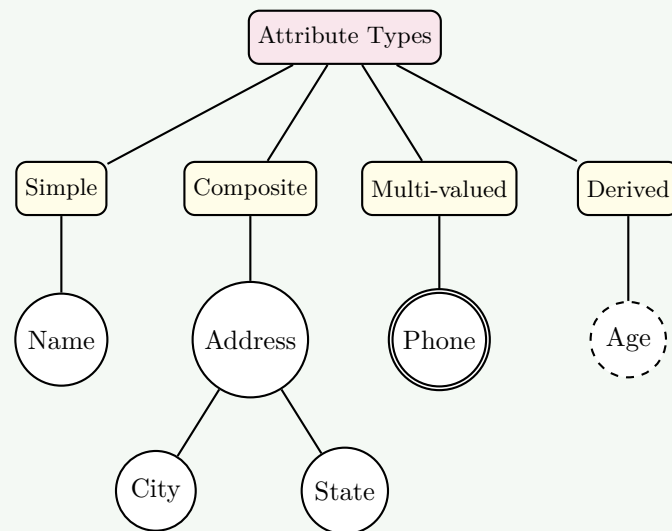
**Solution**

**Attribute:** Property or characteristic that describes an entity.

**Table:**

**Table 10.** Attribute Types

Attribute Type	Description	Example
<b>Simple</b>	Cannot be divided further	Age, Name
<b>Composite</b>	Can be subdivided	Address (Street, City)
<b>Single-valued</b>	Has one value	SSN, Employee_ID
<b>Multi-valued</b>	Can have multiple values	Phone_Numbers
<b>Derived</b>	Calculated from other attributes	Age from Birth_Date
<b>Key</b>	Uniquely identifies entity	Student_ID



**Figure 6.** Classification of Attributes

- **Domain:** Set of permitted values for each attribute
- **Null:** Value used when attribute is unknown or not applicable

#### Mnemonic

“SCSMDK - Simple Composite Single Multi Derived Key”

### Question 3(a) [3 marks]

Explain the GRANT and REVOKE statement in SQL.

#### Solution

Table:

**Table 11.** DCL Commands

Statement	Purpose	Syntax Example
<b>GRANT</b>	Provides privileges to users	GRANT SELECT ON table TO user
<b>REVOKE</b>	Removes privileges from users	REVOKE SELECT ON table FROM user

- **Privileges:** Permissions like SELECT, INSERT, UPDATE, DELETE
- **Control:** Manage which users can access or modify data
- **Security:** Fundamental for database security implementation

#### Mnemonic

“GR - Grant Removes (via REVOKE)”

### Question 3(b) [4 marks]

Explain following Character functions. 1) INSTR 2) LENGTH

## Solution

Table:

Table 12. SQL Character Functions

Function	Purpose	Syntax	Example
<b>INSTR</b>	Finds position of substring	INSTR(str, substr)	INSTR('Hello', 'e') → 2
<b>LENGTH</b>	Returns string length	LENGTH(str)	LENGTH('Hello') → 5

```

1 SELECT INSTR('Database', 'a') FROM dual; -- Returns 2
2 SELECT LENGTH('Database') FROM dual;      -- Returns 8

```

- **INSTR**: Case sensitive search for substring position
- **LENGTH**: Counts total characters including spaces

## Mnemonic

“IL - INSTR Locates, LENGTH measures”

## Question 3(c) [7 marks]

Write SQL statements for following table: Student(Enno,name,branch,sem,clgname,bdate)

## Solution

```

1  -- 1. Create a table Student
2  CREATE TABLE Student (
3      Enno VARCHAR(10) PRIMARY KEY,
4      name VARCHAR(50),
5      branch VARCHAR(20),
6      sem INT,
7      clgname VARCHAR(100),
8      bdate DATE
9  );
10
11 -- 2. Add a column mobno in Student table
12 ALTER TABLE Student ADD mobno VARCHAR(15);
13
14 -- 3. Insert one record in student table
15 INSERT INTO Student VALUES
16 ('E001', 'Raj Patel', 'IT', 3, 'GTU College', '2003-05-15', '9876543210');
17
18 -- 4. Find out list of students who have enrolled in "IT" branch
19 SELECT * FROM Student WHERE branch = 'IT';
20
21 -- 5. Retrieve all information about student where name begin with 'a'
22 SELECT * FROM Student WHERE name LIKE 'a%';
23
24 -- 6. Count the number of rows in student table
25 SELECT COUNT(*) FROM Student;
26
27 -- 7. Delete all record of student table
28 DELETE FROM Student;

```



**Mnemonic**

“CAIRSCD - Create Add Insert Retrieve Search Count Delete”

**Question 3(a OR) [3 marks]**

Explain equi join with example in SQL.

**Solution**

**Equi Join:** Join operation using equality condition to combine tables.

**Table:**

**Table 13.** Equi Join

Join Type	Condition	Result
Equi Join	Column1 = Column2	Matching rows from both tables

```

1  -- Example: List student names and their course names
2  SELECT s.name, c.course_name
3  FROM Student s, Course c
4  WHERE s.course_id = c.course_id;

```

- **Operator:** Uses equality operator (=)
- **Columns:** Typically compares primary key and foreign key

**Mnemonic**

“EE - Equi Equals”

**Question 3(b OR) [4 marks]**

Explain following Aggregate functions. 1) MAX 2) SUM

**Solution**

**Table:**

**Table 14.** SQL Aggregate Functions

Function	Purpose	Syntax	Example
MAX	Returns maximum value	MAX(column)	MAX(salary)
SUM	Returns total sum	SUM(column)	SUM(marks)

```

1  SELECT MAX(salary) FROM Employee; -- e.g., 50000
2  SELECT SUM(marks) FROM Student;   -- e.g., 450

```

- **Aggregation:** Performs calculation on a set of values to return single value
- **Usage:** Often used with GROUP BY clause

**Mnemonic**

“MS - MAX Sum”

### Question 3(c OR) [7 marks]

Write SQL statements for the following table: Employee(EmpID,Ename,DOB,Dept,Salary)

#### Solution

```

1  -- 1. Create a table Employee
2  CREATE TABLE Employee (
3      EmpID VARCHAR(10) PRIMARY KEY,
4      Ename VARCHAR(50),
5      DOB DATE,
6      Dept VARCHAR(30),
7      Salary DECIMAL(10,2)
8  );
9
10 -- 2. Find sum of salaries of all employee
11 SELECT SUM(Salary) FROM Employee;
12
13 -- 3. Insert one record in Employee table
14 INSERT INTO Employee VALUES
15 ('E001', 'John Doe', '1990-05-15', 'IT', 35000);
16
17 -- 4. Find names of employees who salary between 25000/- and 48000/-
18 SELECT Ename FROM Employee WHERE Salary BETWEEN 25000 AND 48000;
19
20 -- 5. Display detail of all employees in descending order of their DOB
21 SELECT * FROM Employee ORDER BY DOB DESC;
22
23 -- 6. List name of all employees whose name ends with 'a'
24 SELECT Ename FROM Employee WHERE Ename LIKE '%a';
25
26 -- 7. Find highest and least salaries of all employees
27 SELECT MAX(Salary) AS Highest, MIN(Salary) AS Lowest FROM Employee;

```

#### Mnemonic

“CSIDDHL - Create Sum Insert Display Display List HighLow”

### Question 4(a) [3 marks]

Consider a following relational schema & give Relational Algebra Expressions for the following queries. Student (Enrollment\_No,Name,DOB,SPI)

#### Solution

1. List all students whose SPI is greater than 7.0

$$\sigma_{SPI > 7.0}(Student)$$

2. Find Name of student whose Enrollment\_No is 007

$$\pi_{Name}(\sigma_{Enrollment\_No=007}(Student))$$

Table:

Table 15. Relational Algebra Symbols

Symbol	Operation	Purpose
$\sigma$	Selection	Filter rows based on condition
$\pi$	Projection	Select specific columns

- **Selection:** Returns subset of tuples
- **Projection:** Returns subset of attributes

#### Mnemonic

“SP - Select Project”

## Question 4(b) [4 marks]

Write a short note on partial functional dependency.

#### Solution

Table:

Table 16. Partial Dependency

Concept	Description
<b>Definition</b>	Non-prime attribute depends on part of composite primary key
<b>Occurs in</b>	Tables with composite primary keys
<b>Problem</b>	Causes redundancy and update anomalies
<b>Solution</b>	Decompose into 2NF

**Example:** In table StudentCourse(StudentID, CourseID, StudentName, CourseName):

- **Key:** (StudentID, CourseID)
- **Dependency:** StudentID  $\rightarrow$  StudentName (Partial Key)
- **Composite Key:** Primary key made of multiple attributes
- **Violation:** Breaks 2nd Normal Form rules

#### Mnemonic

“PDPR - Partial Dependency Problems Resolved”

## Question 4(c) [7 marks]

Explain need of Normalization? Discuss about 2NF with example.

#### Solution

**Need for Normalization:**

- Eliminates **Data Redundancy** (duplicate data)
- Prevents **Update Anomalies** (inconsistent updates)
- Prevents **Insert Anomalies** (inability to add data)
- Prevents **Delete Anomalies** (accidental data loss)

**Second Normal Form (2NF):**

1. Must be in 1NF
2. No partial functional dependencies allowed

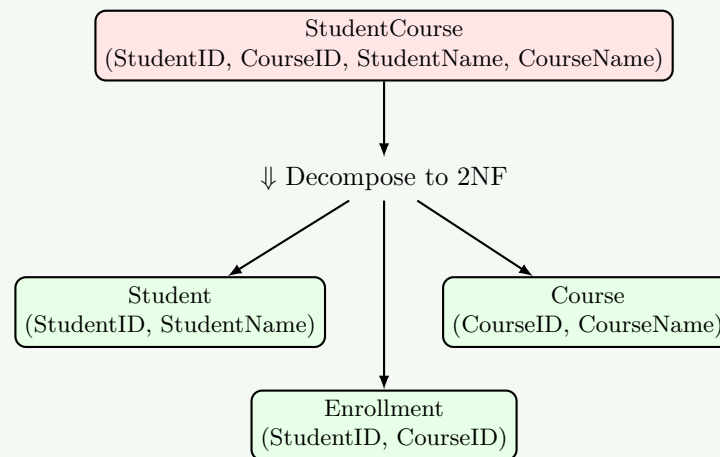


Figure 7. 2NF Decomposition

- **Decomposition:** Splitting table to separate themes
- **Full Dependency:** Non-keys must depend on whole primary key

### Mnemonic

“NUID2 - Normalization Unifies Important Data to 2NF”

## Question 4(a OR) [3 marks]

Consider a following relational schema & give Relational Algebra Expressions for the following queries. Student(Enno,name,age,address)

### Solution

1. Find names of students living in Surat

$$\pi_{name}(\sigma_{address='Surat'}(Student))$$

2. Find names of students older than 30

$$\pi_{name}(\sigma_{age>30}(Student))$$

Table:

Table 17. Query Operations

Query	Algebra Expression
Filter Rows	$\sigma_{condition}(Table)$
Select Columns	$\pi_{columns}(Result)$

- **Nesting:** Operations are nested inside each other
- **Order:** Selection usually performed before projection

## Question 4(b OR) [4 marks]

Define 1 NF? Explain 1NF with suitable example.

### Solution

**First Normal Form (1NF):** Database schema where each column contains atomic (indivisible) values, and there are no repeating groups.

**Table:**

**Table 18. 1NF Rules**

Rule	Description
<b>Atomic Values</b>	No multiple values in single cell
<b>No Repeating Groups</b>	No duplicate columns for same data
<b>Unique Rows</b>	Each row must be unique

**Example:**

Before 1NF	After 1NF
ID: 1, Subjects: Math, Science	ID: 1, Subject: Math
	ID: 1, Subject: Science

- **Atomicity:** Fundamental requirement for relational model
- **Scalability:** Makes querying and indexing efficient

### Mnemonic

“ANU - Atomic No-repeat Unique”

## Question 4(c OR) [7 marks]

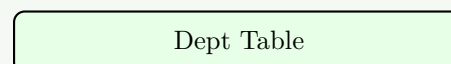
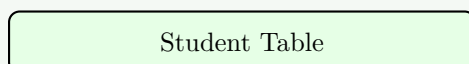
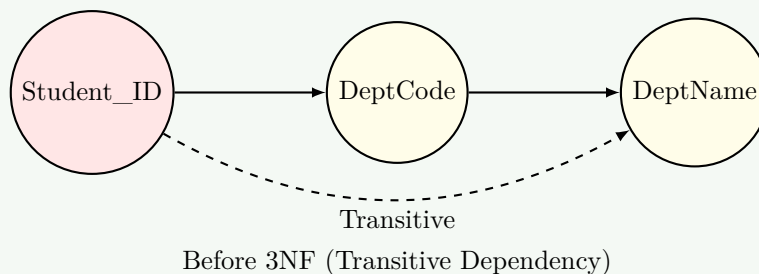
**Define Transitive Dependency? Explain 3NF with suitable example.**

### Solution

**Transitive Dependency:** Non-prime attribute depends on another non-prime attribute rather than directly on primary key.

**Third Normal Form (3NF):**

1. Must be in 2NF
2. No transitive dependencies exist



After 3NF (Split Tables)

**Figure 8. Transitive Dependency & 3NF**

- **Indirect Dependency:**  $A \rightarrow B \rightarrow C$  implies  $A \rightarrow C$
- **Solution:** Move transitive attributes to new table

**Mnemonic**

“T3ND - Transitive Third Normal Form No Dependencies”

**Question 5(a) [3 marks]**

Define Serializability? Explain rules of serializability?

**Solution**

**Serializability:** Property ensuring concurrent transaction execution produces same result as serial execution.  
**Table:**

**Table 19.** Serializability Rules

Rule	Description
<b>Conflict Serializability</b>	Schedule implies same order of conflicting operations as serial schedule
<b>View Serializability</b>	Schedule has same read-write patterns as serial schedule

**Example 2:** Consider a schedule with two transactions T1 and T2. If T1 reads a value, then T2 writes it, and then T1 writes it, this might not be serializable. A serial schedule means transactions are executed one after another.

- **Consistency:** Ensures database integrity during concurrency

**Mnemonic**

“SCV - Serial Conflict View”

**Question 5(b) [4 marks]**

Explain Attributes of Implicit Cursors.

**Solution**

**Table:**

**Table 20.** Implicit Cursor Attributes

Attribute	Description
<b>%FOUND</b>	TRUE if last SQL affected at least one row
<b>%NOTFOUND</b>	TRUE if last SQL affected no rows
<b>%ROWCOUNT</b>	Number of rows affected by last SQL
<b>%ISOPEN</b>	Always FALSE for implicit cursors (automatically closed)

- **SQL%Attribute:** Accessed using SQL prefix
- **Implicit:** Created automatically for DML statements

**Mnemonic**

“FNRI - Found NotFound RowCount IsOpen”

**Question 5(c) [7 marks]**

Explain two phase locking protocol with suitable example.

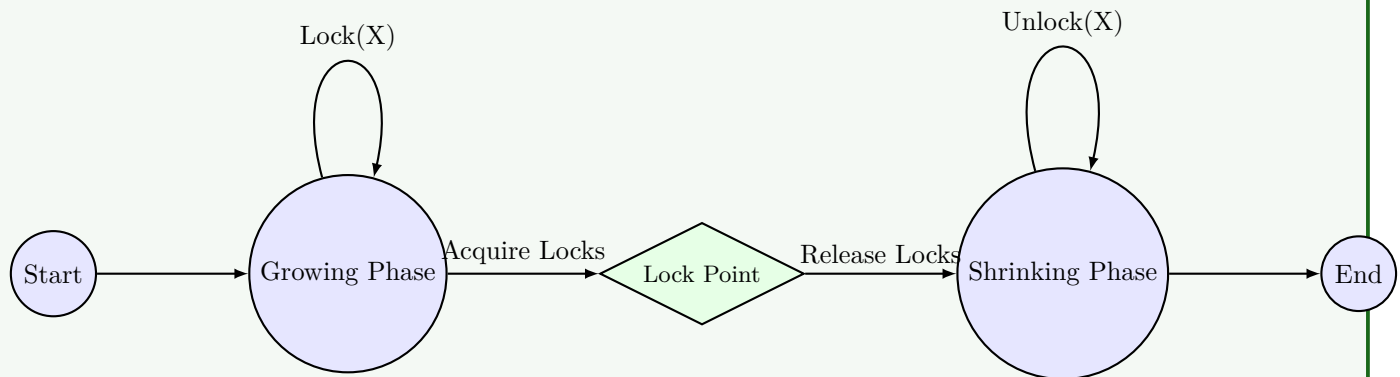
**Solution**

**Two Phase Locking (2PL):** Concurrency control protocol that ensures serializability by defining how transactions acquire and release locks.

**Table:**

**Table 21.** 2PL Phases

Phase	Description	Rule
<b>Growing Phase</b>	Transaction acquires locks	Cannot release any lock
<b>Lock Point</b>	Point where all locks are acquired	Maximum locks held
<b>Shrinking Phase</b>	Transaction releases locks	Cannot acquire new locks



**Figure 9.** Two Phase Locking Protocol

- **Serializability:** Guaranteed by 2PL
- **Deadlocks:** Possible in 2PL (unlike conservative 2PL)

**Mnemonic**

“2PGS - Two Phase Growing Shrinking”

**Question 5(a OR) [3 marks]**

**Explain ACID properties of transaction.**

**Solution**

**Table:**

**Table 22.** ACID Properties

Property	Description
<b>Atomicity</b>	Transaction is all-or-nothing unit of work
<b>Consistency</b>	Database transitions from one valid state to another
<b>Isolation</b>	Concurrent transactions do not interfere with each other
<b>Durability</b>	Committed changes are permanent despite failures

**Mnemonic**

“ACID - All Changes In Database”

## Question 5(b OR) [4 marks]

Define Triggers? Explain advantages of triggers.

### Solution

**Triggers:** Special stored procedures that automatically execute (fire) in response to specific events on a table or view.

**Table:**

**Table 23.** Advantages of Triggers

Advantage	Description
<b>Automatic Execution</b>	Runs automatically without manual invocation
<b>Data Integrity</b>	Enforces complex business rules and constraints
<b>Auditing</b>	Tracks history of data changes (who, when, what)
<b>Security</b>	Controls access and validates data modification

### Mnemonic

“ADAS - Automatic Data Auditing Security”

## Question 5(c OR) [7 marks]

List down problems of concurrency control. Explain any two with suitable example.

### Solution

**Problems of Concurrency Control:**

1. **Lost Update:** Overwriting uncommitted data
2. **Dirty Read:** Reading uncommitted data
3. **Non-repeatable Read:** Re-reading changed data
4. **Phantom Read:** New rows appearing in range query

**Table:**

**Table 24.** Concurrency Problems

Problem	Example Scenario
<b>Lost Update</b>	T1 reads X, T2 reads X. T1 updates X, T2 updates X. T1's update is lost.
<b>Dirty Read</b>	T1 updates X. T2 reads X. T1 fails/rollbacks. T2 has invalid data.

**Example 3:** Another example for Lost Update: T1: Read A(100) → A=A+50 T2: Read A(100) → A=A+30  
T1: Write A(150) T2: Write A(130) [**Overwrites T1's 150!**] Final should be 180, but is 130.

**Example 4 (Dirty Read):** T1: Update A=200 [**Not Committed**] T2: Read A(200) [**Dirty**] T1: Rollback to 100 T2 uses 200, which never officially existed.

### Mnemonic

“LDNP - Lost Dirty Non-repeatable Phantom”