

Subject Name Solutions

4331601 – Summer 2024

Semester 1 Study Material

Detailed Solutions and Explanations

Question 1(a) [3 marks]

Differentiate between array and list.

Solution

Array

Fixed size at creation
Homogeneous data (same type)
Memory efficient - contiguous allocation
Faster access for calculations

List

Dynamic size - can grow/shrink
Heterogeneous data (mixed types)
Flexible but uses more memory
Built-in methods for operations

Mnemonic

“Arrays are Fixed Friends, Lists are Loose Leaders”

Question 1(b) [4 marks]

Explain the concept of class and object with the help of python program.

Solution

Class blueprint objects structure behavior define . **Object** class instance .

```
class Student:
    def __init__(self, name, age):
        self.name = name
        self.age = age

    def display(self):
        print(f"Name: {self.name}, Age: {self.age}")
```

```
\# Creating objects
s1 = Student("Ram", 20)
s2 = Student("Sita", 19)
s1.display()
```

- **Class:** Template
- **Object:** Real instance
- **Constructor:** Object initialize

Mnemonic

“Class Blueprints Create Object Instances”

Question 1(c) [7 marks]

Define constructor. Discuss different types of constructors with suitable python program.

Solution

Constructor special method object creation time automatically call . Python `__init__()` method constructor .

```
class Demo:
    \# Default Constructor
    def __init__(self):
        self.value = 0

    \# Parameterized Constructor
    def __init__(self, x, y=10):
        self.x = x
        self.y = y
```

```
\# Usage
d1 = Demo(5)    \#
x=5,
y=10 (default)

d2 = Demo(3, 7) \#
x=3,
y=7
```

Types of Constructors:

Type	Description	Usage
Default	No parameters	Object initialization
Parameterized	With parameters	Custom initialization
Copy	Creates copy of object	Object duplication

Mnemonic

“Default Parameters Copy Objects”

Question 1(c) OR [7 marks]

Define Polymorphism. Write a python program for polymorphism through inheritance.

Solution

Polymorphism same interface different objects different operations perform ability .

```
class Animal:
    def sound(self):
        pass

class Dog(Animal):
    def sound(self):
        return "Woof!"

class Cat(Animal):
    def sound(self):
        return "Meow!"

\# Polymorphic behavior
animals = [Dog(), Cat()]
```

```
for animal in animals:
    print(animal.sound())
```

- **Method Overriding:** Child class same method name
- **Dynamic Binding:** Runtime method selection
- **Code Reusability:** Same interface, different implementation

Mnemonic

“Many Objects, One Interface”

Question 2(a) [3 marks]

Explain Python specific data structure List, Tuple and Dictionary.

Solution

Data Structure	Properties	Example
List	Mutable, ordered, allows duplicates	[1, 2, 3, 2]
Tuple	Immutable, ordered, allows duplicates	(1, 2, 3, 2)
Dictionary	Mutable, key-value pairs, no duplicate keys	{'a': 1, 'b': 2}

Mnemonic

“Lists Change, Tuples Stay, Dictionaries Map”

Question 2(b) [4 marks]

Explain application of stack.

Solution

Stack Applications:

- **Function Calls:** Call stack management
- **Expression Evaluation:** Infix to postfix conversion
- **Undo Operations:** Text editors, browsers
- **Parentheses Matching:** Syntax checking

```
+{-{-}{-}+}
| 3 | {-} Top}
+{-{-}{-}+}
| 2 |
+{-{-}{-}+}
| 1 |
+{-{-}{-}+}
```

Mnemonic

“Functions Evaluate Undo Parentheses”

Question 2(c) [7 marks]

Define stack. Explain PUSH & POP operation with example. Write an algorithm for PUSH and POP operations of stack.

Solution

Stack LIFO (Last In First Out) principle follow linear data structure .

PUSH Algorithm:

1. Check if stack is full
2. If full, print "Stack Overflow"
3. Else, increment top
4. Add element at top position

POP Algorithm:

1. Check if stack is empty
2. If empty, print "Stack Underflow"
3. Else, remove element from top
4. Decrement top

Example:

```
stack = []
stack.append(10) \# PUSH
stack.append(20) \# PUSH
item = stack.pop() \# POP returns 20
```

Mnemonic

"Last In, First Out - Like Plates"

Question 2(a) OR [3 marks]

Define Following terms: I. Time Complexity II. Space Complexity III. Best case

Solution

Term	Definition	Example
Time Complexity	Algorithm execution time analysis	$O(n)$, $O(\log n)$
Space Complexity	Memory usage analysis	$O(1)$, $O(n)$
Best Case	Minimum time/space needed	Sorted array search

Mnemonic

"Time Space Best Performance"

Question 2(b) OR [4 marks]

Convert $A - (B / C + (D \% E * F) / G) * H$ into postfix expression

Solution

Step-by-step conversion:

Infix: $A - (B / C + (D \% E * F) / G) * H$

1. $A \ B \ C \ / \ D \ E \ \% \ F \ * \ G \ / \ + \ - \ H \ *$

Stack operations:

- Operators: $-$, $($, $/$, $+$, $($, $\%$, $*$, $)$, $/$, $)$, $*$

- Final: $A \ B \ C \ / \ D \ E \ \% \ F \ * \ G \ / \ + \ - \ H \ *$

Postfix Result: $A \ B \ C \ / \ D \ E \ \% \ F \ * \ G \ / \ + \ - \ H \ *$

Mnemonic

“Operands First, Operators Follow”

Mnemonic

“Operands First, Operators Follow”

Question 2(c) OR [7 marks]

Define circular queue. Explain INSERT and DELETE operations of circular queue with diagrams.

Question 2(c) OR [7 marks]

Define circular queue. Explain INSERT and DELETE operations of circular queue with diagrams.

Solution	
<p>Circular Queue queue modified version last position first position connected .</p> <pre> +{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+ 1 2 3 +{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+ \^{ }\^{}} front rear </pre>	
<p>INSERT Algorithm:</p> <ol style="list-style-type: none"> Check if queue is full rear = (rear + 1) % size queue[rear] = element If first element, set front = 0 	
<p>DELETE Algorithm:</p> <ol style="list-style-type: none"> Check if queue is empty element = queue[front] front = (front + 1) % size Return element <ul style="list-style-type: none"> Advantage: Memory efficiency Application: CPU scheduling, buffering 	

Solution	
<p>Circular Queue queue modified version last position first position connected .</p> <pre> +{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+ 1 2 3 +{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+ \^{ }\^{}} front rear </pre>	
<p>INSERT Algorithm:</p> <ol style="list-style-type: none"> Check if queue is full rear = (rear + 1) % size queue[rear] = element If first element, set front = 0 	
<p>DELETE Algorithm:</p> <ol style="list-style-type: none"> Check if queue is empty element = queue[front] front = (front + 1) % size Return element <ul style="list-style-type: none"> Advantage: Memory efficiency Application: CPU scheduling, buffering 	

Solution	
<p>Circular Queue queue modified version last position first position connected .</p> <pre> +{-{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+} 1 2 3 +{-{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+} \^{ }\^{} front rear </pre>	
<p>INSERT Algorithm:</p> <ol style="list-style-type: none"> Check if queue is full rear = (rear + 1) % size queue[rear] = element If first element, set front = 0 	
<p>DELETE Algorithm:</p> <ol style="list-style-type: none"> Check if queue is empty element = queue[front] front = (front + 1) % size Return element <ul style="list-style-type: none"> Advantage: Memory efficiency Application: CPU scheduling, buffering 	

Solution	
<p>Circular Queue queue modified version last position first position connected .</p> <pre> +{-{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+} 1 2 3 +{-{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+} \^{ }\^{}} front rear </pre>	
<p>INSERT Algorithm:</p> <ol style="list-style-type: none"> Check if queue is full rear = (rear + 1) % size queue[rear] = element If first element, set front = 0 	
<p>DELETE Algorithm:</p> <ol style="list-style-type: none"> Check if queue is empty element = queue[front] front = (front + 1) % size Return element <ul style="list-style-type: none"> Advantage: Memory efficiency Application: CPU scheduling, buffering 	

- Solution**

Circular Queue queue modified version last position first position connected .

```

+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
| 1 | 2 | 3 |   |
+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
  \^{      \^{}}
front      rear
  
```

INSERT Algorithm:

 1. Check if queue is full
 2. rear = (rear + 1) % size
 3. queue[rear] = element
 4. If first element, set front = 0

DELETE Algorithm:

 1. Check if queue is empty
 2. element = queue[front]
 3. front = (front + 1) % size
 4. Return element
 - **Advantage:** Memory efficiency
 - **Application:** CPU scheduling, buffering

Solution

Circular Queue queue modified version last position first position connected .

```

+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
| 1 | 2 | 3 |   |
+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
  \^{      \^{}}
front      rear
  
```

INSERT Algorithm:

1. Check if queue is full
2. rear = (rear + 1) % size
3. queue[rear] = element
4. If first element, set front = 0

DELETE Algorithm:

1. Check if queue is empty
2. element = queue[front]
3. front = (front + 1) % size
4. Return element

- **Advantage:** Memory efficiency
- **Application:** CPU scheduling, buffering

- Solution**

Circular Queue queue modified version last position first position connected .

```

+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
| 1 | 2 | 3 |   |
+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
  \^{      \^{}}
front      rear
  
```

INSERT Algorithm:

 1. Check if queue is full
 2. rear = (rear + 1) % size
 3. queue[rear] = element
 4. If first element, set front = 0

DELETE Algorithm:

 1. Check if queue is empty
 2. element = queue[front]
 3. front = (front + 1) % size
 4. Return element
 - **Advantage:** Memory efficiency
 - **Application:** CPU scheduling, buffering

- Solution**

Circular Queue queue modified version last position first position connected .

```

+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
| 1 | 2 | 3 |   |
+{-{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+{-}{-}{-}+}
  \^{      \^{}}
front      rear
  
```

INSERT Algorithm:

 1. Check if queue is full
 2. rear = (rear + 1) % size
 3. queue[rear] = element
 4. If first element, set front = 0

DELETE Algorithm:

 1. Check if queue is empty
 2. element = queue[front]
 3. front = (front + 1) % size
 4. Return element
 - **Advantage:** Memory efficiency
 - **Application:** CPU scheduling, buffering

Mnemonic
“Circle Back When Full”

Mnemonic
“Circle Back When Full”

Question 3(a) [3 marks]

Explain Implementation of Stack using List.

Question 3(a) [3 marks]

Explain Implementation of Stack using List.

Solution

Stack operations Python List :

```
stack = [] \# Empty stack
stack.append(10) \# PUSH
stack.append(20) \# PUSH
top = stack.pop() \# POP
```

- **PUSH:** append() method
- **POP:** pop() method
- **TOP:** stack[-1] for peek

Solution

Stack operations Python List :

```
stack = [] \# Empty stack
stack.append(10) \# PUSH
stack.append(20) \# PUSH
top = stack.pop() \# POP
```

- **PUSH:** append() method
- **POP:** pop() method
- **TOP:** stack[-1] for peek

Solution

Stack operations Python List :

```
stack = [] \# Empty stack
stack.append(10) \# PUSH
stack.append(20) \# PUSH
top = stack.pop() \# POP
```

- **PUSH:** append() method
- **POP:** pop() method
- **TOP:** stack[-1] for peek

- Solution

Stack operations Python List :

```
stack = [] \# Empty stack
stack.append(10) \# PUSH
stack.append(20) \# PUSH
top = stack.pop() \# POP
```

 - **PUSH:** append() method
 - **POP:** pop() method
 - **TOP:** stack[-1] for peek

Mnemonic
“Append Pushes, Pop Pulls”

Mnemonic
“Append Pushes, Pop Pulls”

Question 3(b) [4 marks]

Discuss different applications of linked list.

Solution

Linked List Applications:

- **Dynamic Memory:** Size varies at runtime
- **Insertion/Deletion:** Efficient at any position
- **Implementation:** Stacks, queues, graphs
- **Undo Functionality:** Browser history, text editors

Application	Advantage	Use Case
Music Playlist	Easy add/remove	Media players
Memory Management	Dynamic allocation	Operating systems
Polynomial Representation	Efficient storage	Mathematical operations

Mnemonic

“Dynamic Implementation Undo Memory”

Question 3(c) [7 marks]

Explain doubly linked list. Write an algorithm to delete a node from the beginning of doubly linked list

Solution

Doubly Linked List	node	data, next pointer	previous pointer	.
--------------------	------	--------------------	------------------	---

```
+{-}{-}{-}{-}{-}{-}{+}{-}{-}{-}{-}{-}{-}{+}{-}{-}{-}{-}{-}{-}{+}
| prev | data | next |
+{-}{-}{-}{-}{-}{-}{+}{-}{-}{-}{-}{-}{-}{+}{-}{-}{-}{-}{-}{-}{+}
      \~{           } \~{}
        NULL       points to next
```

Delete from Beginning Algorithm:

1. If list is empty, return
2. If only one node:
 - head = NULL
3. Else:
 - temp = head
 - head = head.next
 - head.prev = NULL
 - delete temp

```
def delete\_beginning(self):
    if self.head is None:
        return
    if self.head.next is None:
        self.head = None
    else:
        self.head = self.head.next
        self.head.prev = None
```

Mnemonic

“Two Way Navigation”

Question 3(a) OR [3 marks]

****Convert this Infix expression into Postfix expression: $A+B/C*D-E/F-G$ ****

Solution

Step-by-step conversion:

Infix: $A+B/C*D-E/F-G$

Postfix: A B C / D * + E F / - G -

Operator precedence: $\ast, / > +, -$

Left to right associativity

Mnemonic

“Multiply Divide Before Add Subtract”

Question 3(b) OR [4 marks]

Explain Circular Linked List with its disadvantages.

Solution

Circular Linked List last node next pointer first node point .

[illegible]

Disadvantages:

- **Infinite Loop Risk:** Improper traversal
- **Complex Implementation:** Extra care needed
- **Memory Overhead:** Additional pointer management
- **Debugging Difficulty:** Circular references

Mnemonic

“Circles Can Cause Confusion”

Question 3(c) OR [7 marks]

Write a Python Program to perform Insert operation in doubly Linked List. Explain with neat diagrams.

Solution

```
class Node:
    def __init__(self, data):
        self.data = data
        self.next = None
        self.prev = None

class DoublyLinkedList:
    def __init__(self):
        self.head = None

    def insert_beginning(self, data):
        new_node = Node(data)
        if self.head is None:
```

```

        self.head = new\_node
    else:
        new\_node.next = self.head
        self.head.prev = new\_node
        self.head = new\_node

```

Before: NULL {-} [10] {-} [20] {-} NULL}

After: NULL {-} [5] {-} [10] {-} [20] {-} NULL}
 \~{-}
 new head

Insert Operations:

- **Beginning:** Update head pointer
- **End:** Traverse to last node
- **Middle:** Update prev/next pointers

Mnemonic

“Begin End Middle Insertions”

Question 4(a) [3 marks]

Write an algorithm for Merge sort.

Solution

Merge Sort Algorithm:

1. If array size ≤ 1 , return
2. Divide array into two halves
3. Recursively sort both halves
4. Merge sorted halves

Time Complexity: $O(n \log n)$ **Space Complexity:** $O(n)$

Mnemonic

“Divide Conquer Merge”

Question 4(b) [4 marks]

Differentiate between Singly Linked List and Doubly Linked List.

Solution

Singly Linked List	Doubly Linked List
One pointer (next)	Two pointers (next, prev)
Forward traversal only	Bidirectional traversal
Less memory usage	More memory usage
Simple implementation	Complex implementation

Singly: [data|next] {- [data|next] {-} NULL}

Doubly: NULL {-} [prev|data|next] {-} [prev|data|next] {-} NULL}

Mnemonic

“Single Forward, Double Bidirectional”

Question 4(c) [7 marks]

Write an algorithm for selection sort. Give the trace to sort the given data using selection sort method.
Data are: 13, 2, 6, 54, 18, 42, 11

Solution

Selection Sort Algorithm:

1. For $i = 0$ to $n-2$:
2. Find minimum in array $[i..n-1]$
3. Swap minimum with array $[i]$

Trace for [13, 2, 6, 54, 18, 42, 11]:

Pass	Array State	Min Found	Swap
0	[13, 2, 6, 54, 18, 42, 11]	2	13 \leftrightarrow 2
1	[2, 13, 6, 54, 18, 42, 11]	6	13 \leftrightarrow 6
2	[2, 6, 13, 54, 18, 42, 11]	11	13 \leftrightarrow 11
3	[2, 6, 11, 54, 18, 42, 13]	13	54 \leftrightarrow 13
4	[2, 6, 11, 13, 18, 42, 54]	18	No swap
5	[2, 6, 11, 13, 18, 42, 54]	42	No swap

Final Result: [2, 6, 11, 13, 18, 42, 54]

Mnemonic

“Select Minimum, Swap Position”

Question 4(a) OR [3 marks]

Write an algorithm for Insertion sort.

Solution

Insertion Sort Algorithm:

1. For $i = 1$ to $n-1$:
2. $key = array[i]$
3. $j = i-1$
4. While $j \geq 0$ and $array[j] > key$:
5. $array[j+1] = array[j]$
6. $j = j-1$
7. $array[j+1] = key$

Time Complexity: $O(n^2)$ Best Case : $O(n)$ for sorted array

Mnemonic

“Insert In Right Position”

Question 4(b) OR [4 marks]

Write an algorithm to insert a new node at the end of circular linked list.

Solution

Algorithm:

1. Create new_node with data
2. If list is empty:
 - head = new_node

```

    - new_node.next = new_node
3. Else:
    - temp = head
    - While temp.next != head:
        - temp = temp.next
    - temp.next = new_node
    - new_node.next = head

def insert_end(self, data):
    new_node = Node(data)
    if self.head is None:
        self.head = new_node
        new_node.next = new_node
    else:
        temp = self.head
        while temp.next != self.head:
            temp = temp.next
        temp.next = new_node
        new_node.next = self.head

```

Mnemonic

“Circle Back To Head”

Question 4(c) OR [7 marks]

Write an algorithm for bubble sort. Give the trace to sort the given data using bubble sort method. Data are: 37, 22, 64, 84, 58, 52, 11

Solution

Bubble Sort Algorithm:

1. For i = 0 to n-2:
2. For j = 0 to n-2-i:
3. If array[j] > array[j+1]:
4. Swap array[j] and array[j+1]

Trace for [37, 22, 64, 84, 58, 52, 11]:

Pass	Comparisons & Swaps	Result
1	37 ↔ 22, 64 ↔ 84, 84 ↔ 58, 84 ↔ 52, 84 ↔ 11	[22, 37, 64, 58, 52, 11, 84]
2	37 ↔ 64, 64 ↔ 58, 64 ↔ 52, 64 ↔ 11	[22, 37, 58, 52, 11, 64, 84]
3	37 ↔ 58, 58 ↔ 52, 58 ↔ 11	[22, 37, 52, 11, 58, 64, 84]
4	37 ↔ 52, 52 ↔ 11	[22, 37, 11, 52, 58, 64, 84]
5	37 ↔ 11	[22, 11, 37, 52, 58, 64, 84]
6	22 ↔ 11	[11, 22, 37, 52, 58, 64, 84]

Final Result: [11, 22, 37, 52, 58, 64, 84]

Mnemonic

“Bubble Up The Largest”

Question 5(a) [3 marks]

Explain Binary search tree and application of it.

Solution

Binary Search Tree (BST) binary tree left subtree smaller values right subtree larger values

Properties:

- **Left child < Parent < Right child**
- **Inorder traversal** gives sorted sequence
- **Search time:** $O(\log n)$ average case

Applications:

Application	Benefit	Use Case
Database Indexing	Fast search	DBMS systems
Expression Trees	Evaluation	Compilers
Huffman Coding	Compression	Data compression

Mnemonic

“Binary Search Trees Organize Data”

Question 5(b) [4 marks]

Write Python Program for Linear Search and explain it with an example

Solution

```
def linear_search(arr, target):
    for i in range(len(arr)):
        if arr[i] == target:
            return i
    return -1

\# Example
numbers = [10, 25, 30, 45, 60]
result = linear_search(numbers, 30)
print(f"Element found at index: \{result\}") \# Output: 2
```

Working:

- **Sequential check:** Element by element
- **Time Complexity:** $O(n)$
- **Space Complexity:** $O(1)$
- **Works on:** Unsorted arrays

Step	Element	Found?
1	10	No
2	25	No
3	30	Yes!

Mnemonic

“Linear Line By Line”

Question 5(c) [7 marks]

Create a Binary Search Tree for the keys 45, 35, 12, 58, 5, 55, 58, 80, 35, 42 and write the Preorder, Inorder and Postorder traversal sequences.

Solution

BST Construction (duplicates ignored):

```
      45
     / {}
    35  58
   / {} / {}
  12 42 55 80
 /
5
```

Insertion Order: 45(root), 35(left), 12(left of 35), 58(right), 5(left of 12), 55(left of 58), 80(right of 58), 42(right of 12)

Traversals:

Traversal	Sequence	Rule
Preorder	45, 35, 12, 5, 42, 58, 55, 80	Root-Left-Right
Inorder	5, 12, 35, 42, 45, 55, 58, 80	Left-Root-Right
Postorder	5, 42, 12, 35, 55, 80, 58, 45	Left-Right-Root

Mnemonic

“Pre-Root First, In-Sorted, Post-Root Last”

Question 5(a) OR [3 marks]

Define following terms: I. Binary tree II. level number III. Leaf-node

Solution

Term	Definition	Example
Binary tree	Tree with max 2 children per node	Each node has ≤ 2 children
Level number	Distance from root (root = level 0)	Root=0, children=1, etc.
Leaf-node	Node with no children	Terminal nodes

```
      A    {{-} Level 0 (Root)}
     / {}
    B  C   {{-} Level 1}
   /
  D       {{-} Level 2 (Leaf)}
```

Mnemonic

“Binary Levels Lead To Leaves”

Question 5(b) OR [4 marks]

Differentiate between Linear Search and Binary search.

Solution

Linear Search	Binary Search
Works on unsorted arrays	Requires sorted array
Sequential checking	Divide and conquer
Time: $O(n)$	Time: $O(\log n)$
Simple implementation	Complex implementation

No preprocessing needed Sorting required

Linear: [1] [2] [3] [4] [5] {- Check each}

Binary: [1] [2] [3] [4] [5] {- Check middle, divide}

Mnemonic

“Linear Line, Binary Bisect”

Question 5(c) OR [7 marks]

Write an algorithm for insertion and deletion a node in Binary search tree.

Solution

Insertion Algorithm:

1. If root is NULL, create new node as root
2. If data < root.data, insert in left subtree
3. If data > root.data, insert in right subtree
4. If data == root.data, no insertion (duplicate)

Deletion Algorithm:

1. If node is leaf: Simply delete
2. If node has one child: Replace with child
3. If node has two children:
 - Find inorder successor
 - Replace data with successor's data
 - Delete successor

```
def insert(root, data):
    if root is None:
        return Node(data)
    if data < root.data:
        root.left = insert(root.left, data)
    elif data > root.data:
        root.right = insert(root.right, data)
    return root

def delete(root, data):
    if root is None:
        return root
    if data < root.data:
        root.left = delete(root.left, data)
    elif data > root.data:
        root.right = delete(root.right, data)
    else:
        \# Node to be deleted found
        if root.left is None:
            return root.right
        elif root.right is None:
            return root.left
        \# Node with two children
        temp = find_min(root.right)
        root.data = temp.data
        root.right = delete(root.right, temp.data)
    return root
```

Cases:

- Leaf deletion: Direct removal
- One child: Replace with child
- Two children: Replace with successor

Mnemonic

“Insert Compare, Delete Replace”