

# Homework Set 3

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## Problem 1 - a

Suppose  $\exists$  a maximal matching  $M$  and a maximum matching  $M^*$  in  $G=(V,E)$ .  $\rightarrow$  For each matched vertex in  $M$ , there must be at least two matched vertices in  $M^*$

AND

$\rightarrow$  There must be at least one matched vertex in  $M$  for which  $\exists$  4 matched vertices in  $M^*$

Contradiction: Take these four vertices  $v_1, v_2, v_3, v_4$  and at least 2 edges  $e_1$  and  $e_2$  which must connect either  $v_1 \rightarrow v_2$  and  $v_3 \rightarrow v_4$  or  $v_1 \rightarrow v_3$  and  $v_2 \rightarrow v_4$  in  $M^*$ . In  $M$ , only one of these 4 vertices is matched (thus using a third edge by necessity), leaving one of  $e_1$  or  $e_2$  as a free edge (with neither of its vertices included in the  $M$  matching).

$\therefore, |M|$  cannot be less than  $\frac{1}{2} |M^*|$ . Thus,  $|M| \geq \frac{1}{2} |M^*|$ .

Then, to compute a maximal matching, simply iterate over all edges in  $G$ , and match each free edge found, giving us a running time of  $O(E)$ .

Pseudocode:

for each edge  $e = (u,v)$  in  $E$ :  
if  $u$  and  $v$  are free, mark  $u$  and  $v$ , using  $e$ .

## Problem 2

If a graph is bipartite, then every path is an alternating path between odd and even vertices, and no vertex can be both even and odd (by definition).

Assume  $\exists$  some odd cycle  $v_1 \rightarrow v_2 \dots v_n$  where  $n$  is odd. Then, assign  $v_1$  to be odd and walk the cycle.  $v_2$  must be even,  $v_3$  odd, and so forth, with  $v_n$  being odd. then, walk from  $v_n$  to  $v_1$ . This path is not an alternating path since  $v_1$  and  $v_n$  are both odd.

$\therefore$ , if  $G$  is bipartite, there can be no odd cycle.

Now, any path that is not a cycle can be an alternating path. An even cycle can be drawn as an alternating path with  $v_1$  as odd,  $v_2$  even  $\dots v_n$  even since  $n$  is even (even cycle). Therefore,  $v_1 \rightarrow v_n \rightarrow v_1$  is an alternating path. so, if a graph is made up of only non-cyclical paths and/or even cycles then it is a bipartite graph.

$\therefore$ , if  $G$  has no odd cycles, it is bipartite

$G$  has no cycles  $\Leftrightarrow G$  is bipartite.