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## CYBER-PHYSICAL SYSTEMS

physical/engineered systems whose operations are monitored, coordinated and controlled by a reliable computing and communication core

- automotive systems (autonomous driving, parking assistance, airbag controls)
- avionics (flight navigation and control)
- manufacturing systems (robotics, process controls)
- medical systems (robotic surgery, devices)

• RTOS was needed because :

① Proliferation of sensor devices.

these devices include embedded processors which ofc need an OS

② Systems being integrated as system of systems

• CPS need .

• Real time processing. The applications must process inputs and generate outputs correctly within a given deadline

• logical / functional + temporal correctness

• Applications : pacemaker, collision detection + avoidance systems

What should the deadline be?

① functionality?

- collision avoidance in an automotive (milliseconds)
- pacemaker (second)
- robotic surgery (depends on target)

↑  
longer timings are taken as design-time planning and not real-time requirements

② environment constraints?

- available computing and communication resources
- timing characteristics of sensors/actuators/operations

③ failure-mitigation strategies? ↔ design consideration

- time to detect and recover from failures
- consider timing overhead of strategies for detecting and recovering from failures  
eg. replication of computing for redundancy.

THIS 'DEADLINE' IS THE WORST CASE TIME TAKEN FOR THE APPLICATION TO RESPOND

## REAL TIME CPS / REAL TIME OS

def · system whose correctness depends not only on logical / functional aspects, but also temporal aspects

key performance measures :

- ① timeliness / Predictability on timing constraints (deadlines)
- ② significance of worst-case over average-case

how to model

→ RTOS Process

def · Real Time Process :  $\langle R, C, D \rangle$

R : process release time

C : execution requirement

D : relative deadline

meaning : C time units in the CPU are needed between time  $[R, R+D)$

for simplicity, we assume : 1 CPU burst (duration = C)  
no I/O bursts

## RECURRENT RTS

- most RTS are recurrent (pacemaker, collision avoidance)
- each instance of execution is a RTP  $\langle R, C, D \rangle$

## PERIODIC RTP

- repeats periodically
- processes generated by a time triggered phenomena (eg. sensor sending data periodically)
- defined as  $\langle T, C, D \rangle$ 
  - T : process period  $\rightarrow$  release every T time units
  - C : CPU time
  - D : deadline

releases as  $R = 0, T, 2T, \dots$

## SPORADIC REAL TIME PROCESS

- repeats sporadically with a minimum gap between releases
- processes generated by an event triggered phenomena (eg. anti-lock braking function)  $\hookrightarrow$  induced by humans
- defined as  $\langle T, C, D \rangle$ 
  - T : minimum release - separation time

short term scheduler / ready queue scheduler

↖ RT CPU Scheduling

• classic algorithms like FCFS, SJF, RR fail as they don't prioritize deadlines

### ① Fixed Priority Scheduling

• all instances of the recurrent processes always have the same priority

← A. Rate Monotonic Scheduler

- Priority is assigned based on process periods / minimum release separation time  $\rightarrow \underline{T}$

shorter  $T \Rightarrow$  larger priority  
Ties are broken arbitrarily

- gives a predictable schedule for high-priority periodic sporadic process
- $\therefore$  RM ignores  $D$ , it can still miss a deadline

VERY VERY  
POPULAR

## B. Deadline Monotonic (DM) Scheduler

- Assigns priorities based on process deadlines (D)

Shorter D  $\Rightarrow$  greater probability

Ties are broken arbitrarily

- $\therefore$  Since DM can't change priorities at the level of process instances, it can miss deadlines for certain process sets.

## 2. DYNAMIC PRIORITY SCHEDULER

### Earliest Deadline First Scheduler

- dynamic priority scheduler assigns priorities based on process instance deadlines
- instances with shorter deadlines are given higher priority
- ties are broken arbitrarily

RM/DM	EDF
<ul style="list-style-type: none"><li>Simpler implementation (separate queue for each recurrent process)</li></ul>	<ul style="list-style-type: none"><li>Harder implementation (online sorting of queue based on instance deadlines)</li></ul>
<ul style="list-style-type: none"><li>Predictability for high priority process, even under overload</li></ul>	<ul style="list-style-type: none"><li>misbehavior during overload</li></ul>