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problem: knapsack of capacity  $c$   
 $n$  kinds of objects  $\rightarrow$  each object has a weight and profit :  $O_i (w_i, p_i)$   
there are unlimited copies of each  $O_i$   
maximise profit

step 1 formulate into subproblem

$\therefore$  'n' is not capped, we can't use  $\text{knapsack}(n)$   
we must use both  $n$  and  $c \rightarrow \text{knapsack}(n, c)$   
 $\left\{ \begin{array}{l} \text{if we don't choose } n \rightarrow \text{knapsack}(n-1, c) \\ \text{if we do choose } n \rightarrow \text{knapsack}(n, c-w_n) \end{array} \right.$   
 $\rightarrow$  until  $n=0$  or  $c=0$

$n^{\text{th}}$  object

$\rightarrow$  ① chosen :  $C_{\text{available}} = c - w_n$   
 $\text{profit} = p_n + \text{profit of rest}$   
 $\downarrow$   
 $\text{subproblem}(n, (c - w_n))$

$\rightarrow$  ② not chosen :  $C_{\text{available}} = c$   
 $\text{profit} = \text{profit of rest}$   
 $\downarrow$   
 $\text{subproblem}(n-1, c)$

$\therefore$  current profit depends  
on profit of subproblems.  
for current profit to be maximum,  
the solution of the subproblems should also be optimal.

$\uparrow$   
principle of optimality holds.

Step 2 recursive function definition

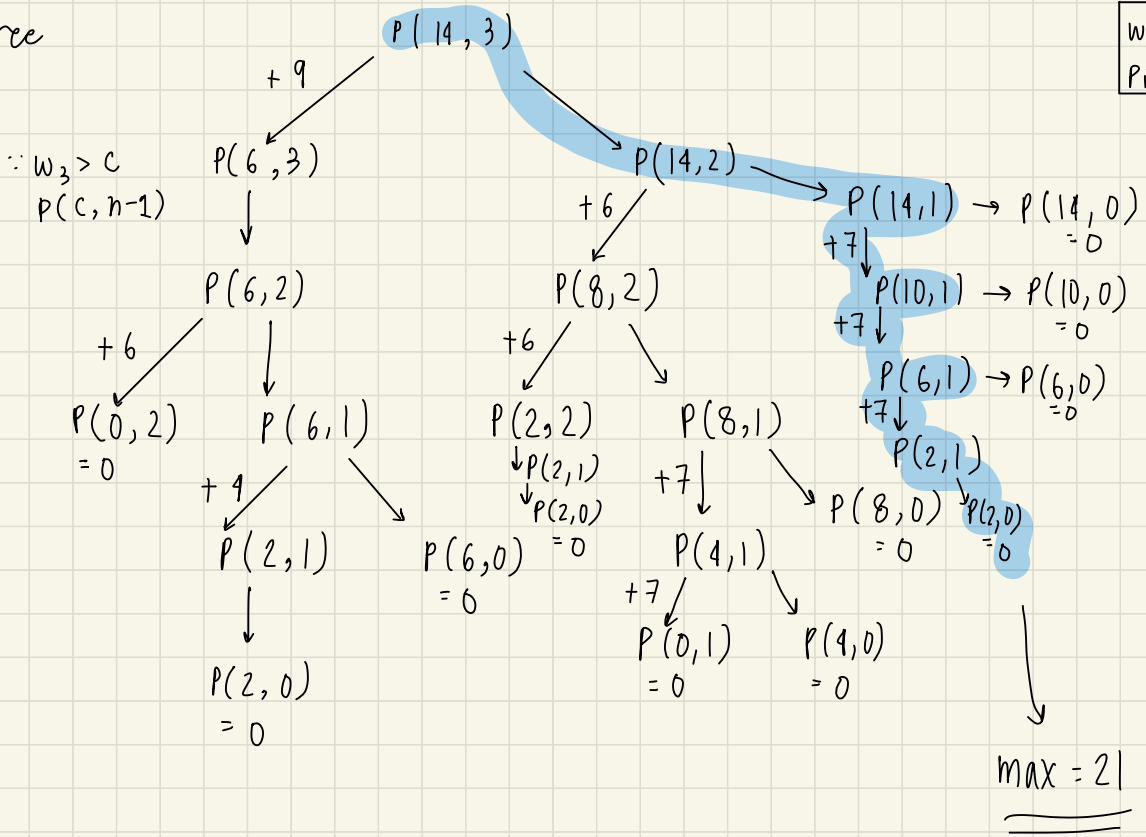
$$\text{Profit}(c, n) = \max \left( \begin{array}{l} P_n + \text{Profit}(c - W(n), n) \\ \text{Profit}(c, n - 1) \end{array} \right)$$

$$\text{Profit}(0, n) = \text{Profit}(c, 0) = 0$$

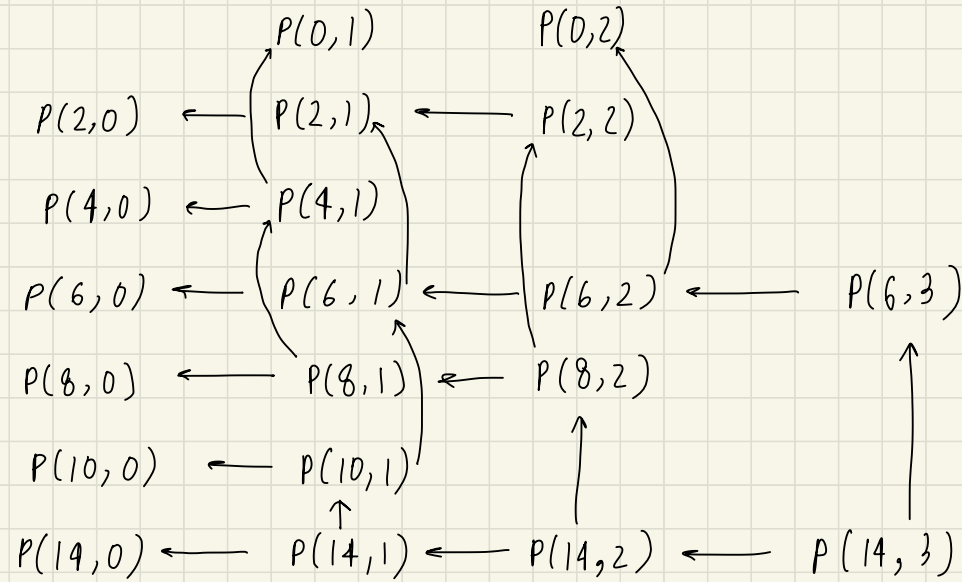
Step 3 subproblem graph  $\rightarrow$

Step 3  
subproblem tree

	1	2	3
$w_i$	4	6	8
$p_i$	7	6	9



# Subproblem graph



$P(a, b)$   
 $a$ : capacity  
 $b$ : object