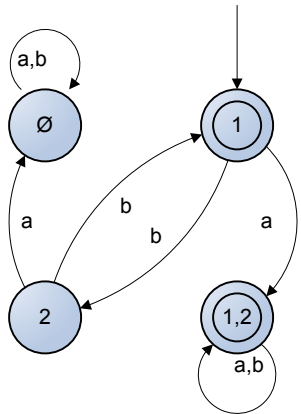


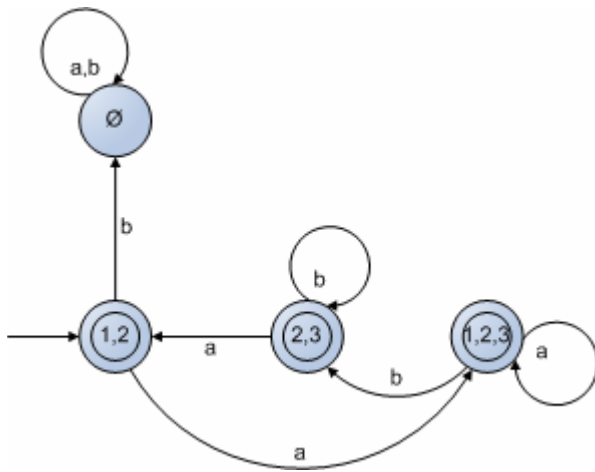
Homework #2 Solutions.

1.16

a.



b.



1.18

a. $1\Sigma^*0$

b. $\Sigma^*1\Sigma^*1\Sigma^*1\Sigma^*$

c. $\Sigma^*0101\Sigma^*$

d. $\Sigma\Sigma0\Sigma^*$

e. $(0(\Sigma\Sigma)^*) \cup (1\Sigma(\Sigma\Sigma)^*)$

f. $(0 \cup 10)^*1^*$

g. $\Sigma\Sigma\Sigma\Sigma\Sigma \cup \Sigma\Sigma\Sigma\Sigma \cup \Sigma\Sigma\Sigma \cup \Sigma\Sigma \cup \Sigma \cup \epsilon$

h. $\epsilon \cup 1 \cup 0 \cup 1(0 \cup 11(0 \cup 1))(0 \cup 1)^*$

i. $(1\Sigma)^*$

j. $000^* \cup (000^*1 \cup 100 \cup 010)0^*$

k. $\epsilon \cup 0$

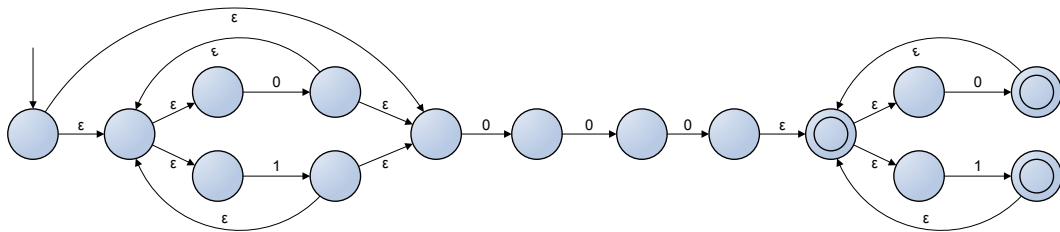
l. $(1^*01^*01^*)^*$

m. \emptyset

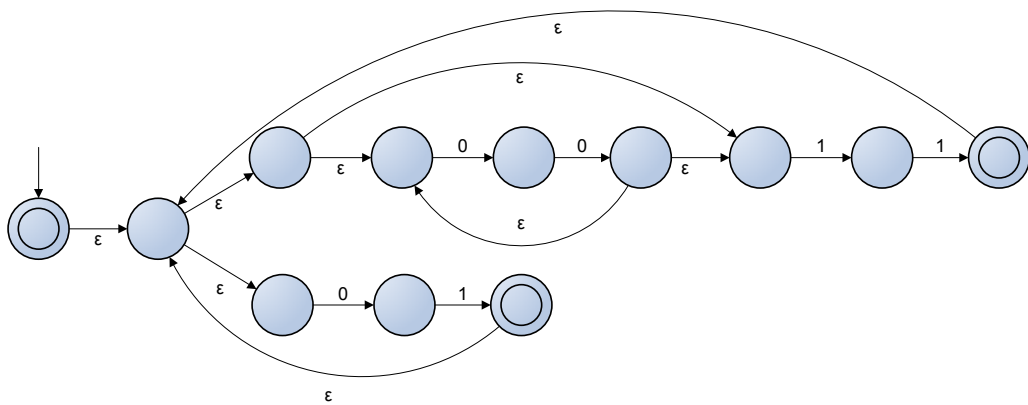
n. $(0 \cup 1)(0 \cup 1)^*$

1.19

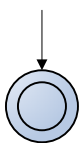
a.



b.



c.



1.20

c.

In: aaab
 In: b
 Not in: aba
 Not in: bbbbbbbba

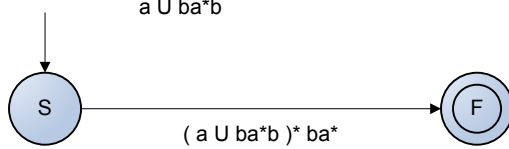
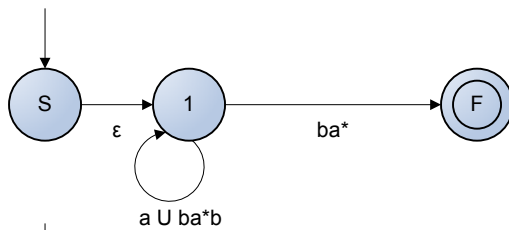
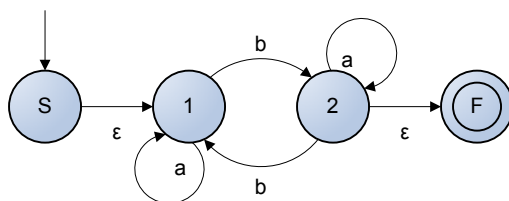
d.
 In: aaa
 In: aaaaaa
 Not in: b
 Not in: aaaaa

e.
 In: aaaaabbbbbbaaaaa
 In: aaba
 Not in: aaaaaabb
 Not in: bab

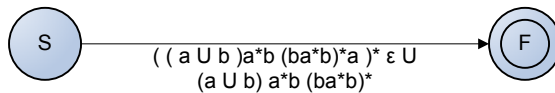
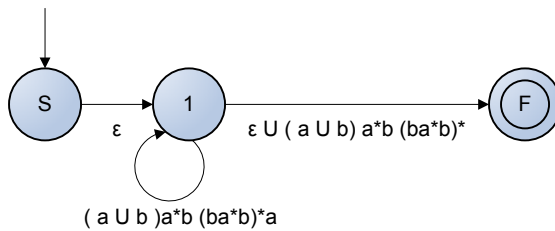
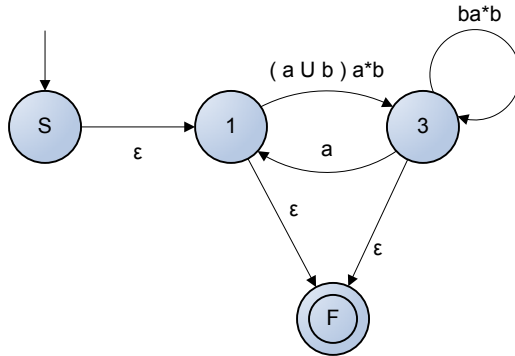
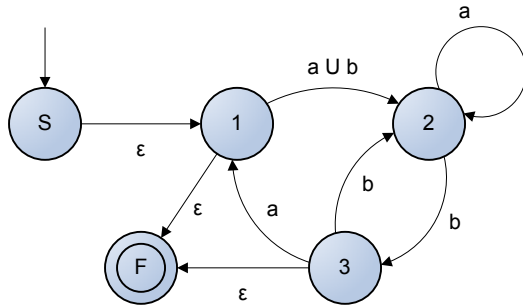
f.
 In: aba
 In: bab
 Not in: abab
 Not in: bb

1.21

a.



b.



1.29

a. text, page 96

b.

Let $w = ab^p$. Then $s = ab^p ab^p ab^p$. Now, with $|xy| \leq p$, then y must either be an a with some b 's, or all b 's. In the former case, pumping up will result in multiple a 's in the first w (of www) that are not repeated in the remaining two instances. Likewise, in the latter case, where y contains strictly b 's, then pumping up will result in too many b 's for the first instance that are, again, not echoed in the remaining instances of the substring.

c. text, page 96

1.46

a. To show a is not pumpable, let s be the string 0^p10^p . By $|xy| \leq p$, y must contain only 0s. Let y be 0^k , where $0 \leq k \leq p$. Now, if we pump up the string s , we get $0^{p+ki}10^p$ for each pump i , which is not in the accepted language.

b. text, page 97

c. Let s be the string 0^p10^p . By $|xy| \leq p$, y must contain only 0s. Pumping up will result in a string that is no longer a palindrome. Since regular languages are closed under complement, if non-palindromes were regular, then L would be too.

d. Let s be the string 0^p110^p1 . By $|xy| \leq p$, y must contain only 0s. Let y be 0^k , where $0 \leq k \leq p$. When we pump on this y , we will end up with $0^{p+k}110^p1$, which is not in the accepted language.

1.49

Let s be the string 0^p10^p . By $|xy| \leq p$, y must contain only 0s. If we pump down in this case, then the resulting string will be $0^{p-k}10^p$, which is not in the accepted language.

1.55

- a. 4
- b. 1
- c. 1
- d. 3
- e. 2