

# Link Prediction in Undirected Networks

## A Probabilistic Foundation

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# Motivation

Graphs (Networks) are collections of nodes (vertices) and edges (connections) that can be used to model relationships.

- Social Networks (Facebook, Twitter, LinkedIn, ego-nets)
- Protein-protein interaction networks
- Telecommunication Networks
- Buying/Selling Networks (Amazon, Ebay, etc.)



Problems:

- Community Detection, Clustering, Partitioning
- Centrality Measures
- Link Prediction
- Graph Drawing/Visualization
- Diffusion Analysis

We look at Link Prediction today, using an Effective Resistance based metric

## Graph

*Graph  $G(V, E)$  is a set of  $n$  nodes (vertices)  $V = \{1, 2, \dots, n\}$ , with pairs of nodes connected by edges (links) in the set  $E$ .*

## Adjacency Matrix

$$A_{ij} = \begin{cases} 1 & \text{if there exists an edge in } E \text{ from node } i \text{ to node } j, \\ 0 & \text{otherwise.} \end{cases}$$

Note we can replace 1 by weight  $w_e$  for the weight of edge  $e = (i, j)$ .

# Graph Basics (cont.)

## Degree Matrix

$$D = \text{diag}(d_1, d_2, \dots, d_n)$$

where  $d_i = \deg(i) = \# \text{ of nodes that node } i \text{ connects to.}$

## Graph Laplacian

Given the adjacency matrix  $A$  and degree matrix  $D$ ,

$$L = D - A$$

Other Graph Laplacians:

- $L_{rw} = I - D^{-1}A$
- $L_{sym} = I - D^{-1/2}AD^{-1/2}$

# Directed vs Undirected

Edges  $\implies$  relationships or flow of information.

Not all relationships in the world are mutual, or bidirectional.

## Undirected Graph

*A graph  $G$  is **undirected** if all of the connections are bidirectional. This is equivalent to  $A$  and  $L$  being symmetric.*

*Otherwise, the graph is **directed** (digraph), with  $A$  and  $L$  not symmetric.*

Note:  $\lambda = 0 \in \sigma(L)$ , with eigenvector  $\mathbf{e} = \mathbb{1}$ .

# Undirected Example

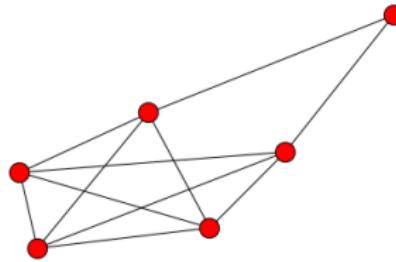
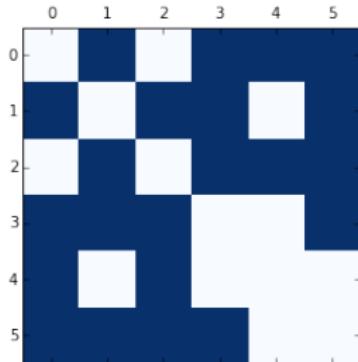


Figure: Undirected Adjacency Matrix and Corresponding Graph

# Focus Today : Undirected Networks

We focus only on *undirected* networks, so we have symmetry.

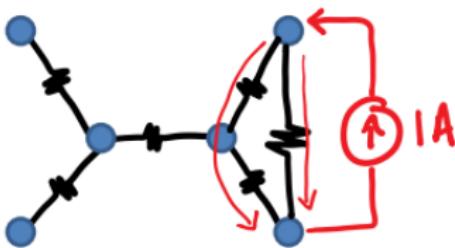
# Link Prediction Problem

Given an observed, undirected network  $G(V, E)$ , what is the most likely *unobserved* edge  $e \notin E$  that should be in  $E$ , or is likely to be in  $E$  in the future?

## Problems

- *ill-posed problem*
- *how to measure quality of link prediction?*
- *complex nature of networks, underlying dynamics*

# Effective Resistances



## Effective Resistance

The effective resistance between nodes  $i, j \in V$  is the energy dissipation when a unit current is injected at node  $i$  and removed at node  $j$ . It can be calculated as the potential difference

$$R_{\text{eff}}(i, j) = v(i) - v(j)$$

photo credit: Nikhil Srivastava, Graph Sparsification I: Sparsification via Effective Resistances

# Effective Resistances

Model our network as an electrical resistor network:

It can be shown using Kirchoff and Ohm's laws that  $R_{\text{eff}}(i, j)$  can be found via:

$$R_{\text{eff}}(i, j) = (\mathbf{e}_i - \mathbf{e}_j)^T L^\dagger (\mathbf{e}_i - \mathbf{e}_j) = L_{ii}^\dagger - 2L_{ij}^\dagger + L_{jj}^\dagger$$

- $L^\dagger$  : Moore-Penrose Pseudoinverse of the Graph Laplacian (symmetric)
- $\mathbf{e}_i, \mathbf{e}_j$  :  $i^{\text{th}}$  and  $j^{\text{th}}$  standard  $\mathbb{R}^n$  basis vectors

# Sparsification via Effective Resistances, Daniel Spielman and Nikhil Srivastava

Spielman, Srivastava (2009)

*Sparsify dense graphs via random sampling of edges based on the effective resistances across edges.*

$$\begin{array}{ccc} \text{dense } G(V, E) & \xrightarrow{\quad} & \text{sparse } H(V, E_s) \\ & R_{\text{eff}}(G) & \end{array}$$

Sparsified graph  $H$  retains certain “spectral” properties of  $G$ :

- eigenvalues and eigenvectors are “close”
- graph cuts
- clustering

If  $L_G, L_H$  are the corresponding Graph Laplacians of  $G$  and  $H$ , respectively:

$$(1 - \epsilon)\mathbf{x}^T L_G \mathbf{x} \leq \mathbf{x}^T L_H \mathbf{x} \leq (1 + \epsilon)\mathbf{x}^T L_G \mathbf{x}$$

$\forall \mathbf{x} \in \mathbb{R}^n$  with high probability.

# Connecting Sparsification and Link Prediction

Sparsification:

$$\text{dense } G(V, E) \implies \text{sparse } H(V, E_s)$$

Link Prediction:

$$\text{"sparse" } H(V, E_s) \implies \text{"dense" } G(V, E)$$

# Link Prediction via Effective Resistances

With  $e = (i, j) \in E$ , we have that

$$R_{\text{eff}}(e) : E \rightarrow [0, \infty)$$

defines a metric on the edge set,  $E$ .

- Effective Resistances  $\implies$  “distance”
- more short paths  $\implies$  lower  $R_{\text{eff}}()$   $\implies$  “closer” electrically

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*Extend this metric to all pairs of nodes.*

## Link Prediction Routine

Given an observed, undirected graph  $G(V, E)$ , we predict the link  $\hat{e} \notin E$  s.t.

$$\hat{e} = \operatorname{argmin}_{e \notin E} R_{\text{eff}}(e) = \operatorname{argmin}_{(i,j) \notin E} L_{ii}^\dagger - 2L_{ij}^\dagger + L_{jj}^\dagger$$

where  $L^\dagger$  is the Moore-Penrose Pseudoinverse of the Graph Laplacian.

# Quick Example

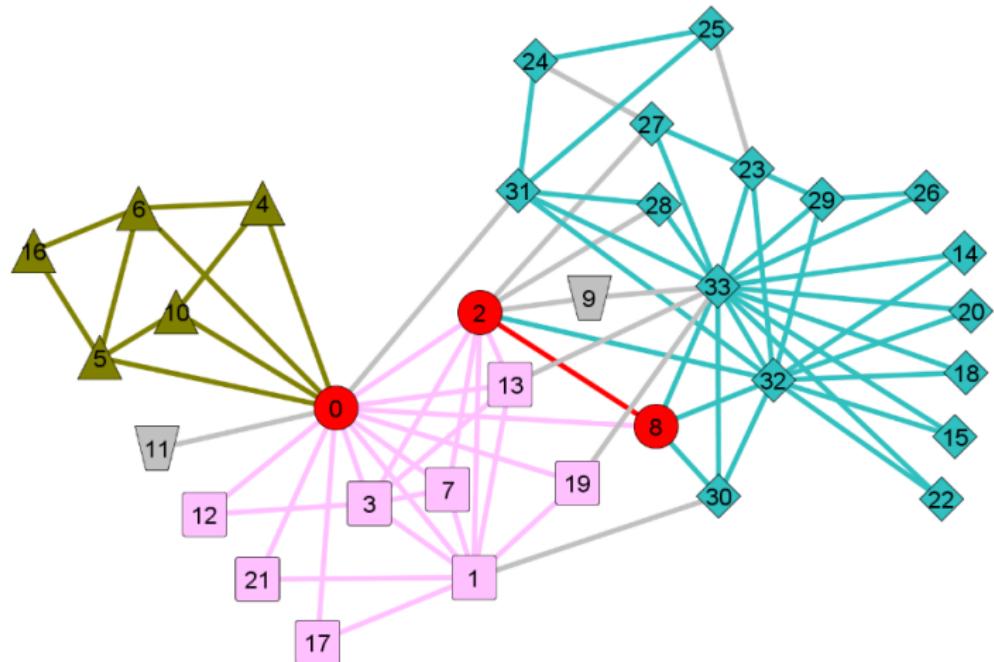


Figure: Zachary's Karate Club Network

# Justification for this Method?

Questions:

- Empirically good, but justified?
- In what sense is this predicted link, “the best” or “most likely”?
- Different metrics, different results? Which is best?
- Computationally efficient?

# Justification for this Method?

We show that Link Prediction via Effective Resistances yields the “most likely” link in a probabilistic sense, when we view the observed graph as a draw from the probability distribution across edges as defined for Sparsification via Effective Resistances.

# Formalization of Link Prediction

Consider an observed, undirected graph  $G_o(V, E_o)$  with edge weights  $\{w_e\}_{e \in E_o}$  then we define:

- *plus-one graph* = a graph  $G_1(V, E_1)$  s.t.  $E_1 = E_o \cup \{e_1\}$ , with ( $e_1 \notin E_o$ )
- $\mathbb{G} = \{G_1(V, E_1) : G_1 \text{ is a plus-one graph of } G_o(V, E_o)\}$
- $\mathbb{E} = \{E_1 : E_1 \text{ is a plus-one edge set of } E_o\}$
- $R_{eff_{E_o}}(e)$  = effective resistance of the edge  $e$  in the edge set  $E_o$

# Probabilistic Foundation for Link Prediction via Effective Resistances

## Theorem

Given an undirected, observed graph  $G_o(V, E_o)$  and a prior on all edge weights  $\{w_e\}_{e \notin E_o}$ , the edge  $\hat{e} \notin E_o$  s.t.

$$\hat{e} = \operatorname{argmin}_{e \notin E_o} w_e \operatorname{Reff}_{E_o}(e)$$

then  $\hat{G}(V, E_o \cup \{\hat{e}\})$  is most-likely plus-one graph to have produced  $G_o(V, E_o)$ .

# Bibliography

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