Milner

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Milner

Over the past three weeks, I've been creating my own programming language, Milner

```
external puts : (Cstring) -> ()

val main : () -> Int32
fun main() =
  puts("Hello, world!");
  0i32
```

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  - Execution is evaluation

A *reference implementation* is a specification for the language

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In practice, interpreters compile to a low-level *bytecode* which is then executed

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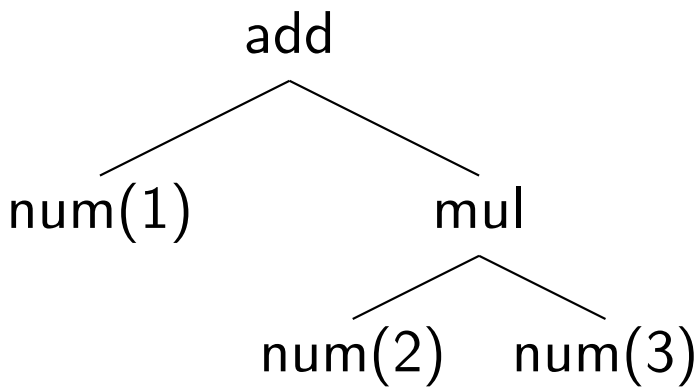
- 

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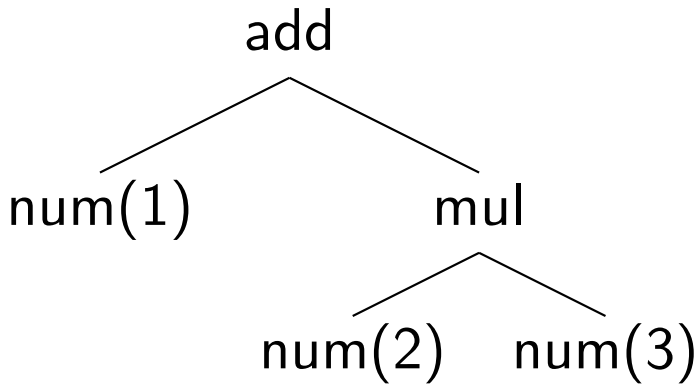
- 

- Sedlex lexer generator and Menhir parser generator

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  - $$\frac{e_1 : Int \quad e_2 : Int}{e_1 + e_2 : Int}$$
  - $$\frac{e_1 : A \to B \quad e_2 : A}{e_1\ e_2 : B}$$
- Hindley-Milner type system

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- LLVM = Low-Level Virtual Machine