

SUBJECT REDUCTION FOR PURE TYPE SYSTEMS

Charlotte Marchal | 261031516 Dashiell Rich | 261002837
Milton Rosenbaum | 260972050 Zhaoshen Zhai | 261003108

ABSTRACT. Following [GN91] and [Bar91], we study the basics of *pure type systems*, which abstract many of the constructs found in the eight systems of the λ -cube. We start with a brief introduction to the systems of the λ -cube, discuss their expressive power, and introduce pure type systems as a unifying framework in which they can be studied. We then give a detailed proof of subject reduction for arbitrary pure type systems.

Introduction. Subject reduction is a crucial property of a type system that guarantees its ‘computational consistency’ by ensuring that reductions of a well-typed expression remains well-typed, and which supports the slogan that ‘well-typed programs do not go wrong’. It is thus desirable that we can prove it uniformly across many different type systems, and this is the goal of the present note.

To this end, **TODO:** review λ -cube and the dependencies.

Definition 1. A *pure type system* is a tuple $\sigma := (\mathcal{C}, \mathcal{V}, \mathcal{S}, \mathcal{A}, \mathcal{R})$ consisting of a set \mathcal{C} of *constants*, a set \mathcal{V} of *variables*, a set $\mathcal{S} \subseteq \mathcal{C}$ of *sorts*, a set $\mathcal{A} \subseteq \mathcal{C}^2$ of *axioms*, and a set $\mathcal{R} \subseteq \mathcal{S}^3$ of *rules*.

Notation 2. Throughout, let σ be denote an arbitrary pure type system.

TODO: link this with the λ -cube by interpreting λ_{\rightarrow} as a PTS.

Definition 3. The collection of σ -*pseudoterms* is defined by $T := \mathcal{V} \mid \mathcal{C} \mid (TT) \mid (\lambda\mathcal{V}:T.T) \mid (\Pi\mathcal{V}:T.T)$. Pairs $(A, B) \in T^2$ are called σ -*assignments*, written $A:B$, and a finite sequence thereof is called a σ -*pseudocontext*.

Definition 4. The β -*reduction* relation is the least relation on σ -terms satisfying the following for all σ -terms A, A', A'' : the *principal reduction rule* $(\lambda x:A.A')A'' \rightarrow_{\beta} A'[A''/x]$, and the *congruence rules* $AA' \succ A A''$, $A' A \succ A'' A$, $\lambda x:A.A' \succ \lambda x:A.A''$, $\lambda x:A'.A \succ \lambda x:A''.A$, $\Pi x:A.A' \succ \Pi x:A.A''$, and $\Pi x:A'.A \succ \Pi x:A''.A$.

Notation 5. We write \rightarrow_{β} for the reflexive and transitive closure of \rightarrow_{β} , and $=_{\beta}$ for the equivalence relation generated by \rightarrow_{β} . A σ -term of the form $(\lambda x:A.A')A''$ is called a β -*redex*.

Definition 6. Let Γ be a σ -pseudocontext and let M, N be σ -pseudoterms. We say that Γ *proves* $M:N$, and write $\Gamma \vdash M:N$, if there is a finite well-founded tree \mathcal{D} , called a *derivation*, such that the following hold.

1. Vertices of \mathcal{D} are of the form $\Delta \vdash A:B$, where A and B are σ -pseudoterms and Δ is a σ -pseudocontext.
2. The root of \mathcal{D} is $\Gamma \vdash M:N$ and the leaves of \mathcal{D} are instances of $\vdash c:c'$, where $(c, c') \in \mathcal{A}$.
3. Each interior vertex of \mathcal{D} is a conclusion of an *inference rule*, whose successors are exactly the premises.

The inference rules of σ are as follows. Below, $s \in \mathcal{S}$, $x \in \mathcal{V} \setminus \text{dom } \Gamma$, $(s_1, s_2, s_3) \in \mathcal{R}$, and $C =_{\beta} C'$.

$$\begin{array}{c} \frac{\Gamma \vdash A:s}{\Gamma, x:A \vdash x:A} \text{ INIT} \quad \frac{\Gamma \vdash A:s \quad \Gamma \vdash B:C}{\Gamma, x:A \vdash B:C} \text{ WEAK} \quad \frac{\Gamma \vdash B:C \quad \Gamma \vdash C':s}{\Gamma \vdash B:C'} \text{ CONV} \quad \frac{\Gamma \vdash B_1:s_1 \quad \Gamma, x:B_1 \vdash B_2:s_2}{\Gamma \vdash (\Pi x:B_1.B_2):s_3} \text{ PI-RULE} \\ \frac{\Gamma \vdash B_1:s_1 \quad \Gamma, x:B_1 \vdash B_2:s_2 \quad \Gamma, x:B_1 \vdash C:B_2}{\Gamma \vdash (\lambda x:B_1.C):(\Pi x:B_1.B_2)} \text{ L-RULE} \quad \frac{\Gamma \vdash B_1:(\Pi x:C_1.C_2) \quad \Gamma \vdash B_2:C_1}{\Gamma \vdash B_1 B_2:C_2[B_2/x]} \text{ APP} \end{array}$$

Definition 7. If $\Gamma \vdash A:B$, then Γ is a σ -*context* and A, B are σ -*terms*.

Lemma 8 (Substitution Lemma; [GN91, Lemma 17]). *Let Γ and $\Gamma_1, y:A, \Gamma_2$ be σ -contexts and let A, M, N, P be σ -terms. If $\Gamma_1, y:A, \Gamma_2 \vdash M:N$ and $\Gamma \vdash P:A$, then $(\Gamma_1, \Gamma_2)[P/y] \vdash M[P/y]:N[P/y]$.*

Date: April 22, 2025.

Extended abstracted for the final project for COMP527: LOGIC AND COMPUTATION taught by Professor Brigitte Pientka.

Link to presentation: **TODO**

Lemma 9 (Stripping Lemma; [GN91, Lemma 19]). *Let Γ be a σ -context and let M, N, P be σ -terms.*

1. *If $\Gamma \vdash c : P$ where $c \in \mathcal{C}$, then $P =_\beta c'$ and $(c, c') \in \mathcal{A}$ for some $c' \in \mathcal{C}$.*
2. *If $\Gamma \vdash x : P$ where $x \in \mathcal{V}$, then $P =_\beta Q$ for some σ -term Q such that $(x : Q) \in \Gamma$.*
3. *If $\Gamma \vdash (\Pi x : M.N) : P$, then $\Gamma \vdash M : s_1$, $\Gamma, x : M \vdash N : s_2$, and $P =_\beta s_3$ for some $(s_1, s_2, s_3) \in \mathcal{R}$.*
4. *If $\Gamma \vdash (\lambda x : M.N) : P$, then $\Gamma \vdash M : s_1$, $\Gamma, x : M \vdash Q : s_2$, $\Gamma, x : M \vdash N : Q$, $\Gamma \vdash P : s_3$, and $P =_\beta \Pi x : M.Q$ for some $(s_1, s_2, s_3) \in \mathcal{R}$ and σ -term Q .*
5. *If $\Gamma \vdash M N : P$, then $\Gamma \vdash M : (\Pi x : A.B)$, $\Gamma \vdash N : A$, and $P =_\beta B[N/x]$ for some σ -terms A and B .*

Theorem 10 (Subject Reduction; [GN91, Lemma 22]). *Let Γ, Γ' be σ -contexts and let M, M', N be σ -terms.*

1. *If $\Gamma \vdash M : N$ and $M \rightarrow_\beta M'$, then $\Gamma \vdash M' : N$.*
2. *If $\Gamma \vdash M : N$ and $\Gamma \rightarrow_\beta \Gamma'$, then $\Gamma' \vdash M : N$.*

Proof. We proceed by simultaneous induction on the derivation $\mathcal{D} : \Gamma \vdash M : N$ when $M \rightarrow_\beta M'$ and $\Gamma \rightarrow_\beta \Gamma'$; the general case follows by iteration. We first prove (1), and split into cases with similar proofs.

- If \mathcal{D} ends with INIT, then there is no redex in M . If \mathcal{D} ends with CONV, then there are derivations $\mathcal{D}_1 : \Gamma \vdash M : N'$ and $\mathcal{D}_2 : \Gamma \vdash N' : s$ for some $s \in \mathcal{S}$ and some σ -term N' such that $N' =_\beta N$. By IH₁, we have $\Gamma \vdash M' : N'$, on which CONV with \mathcal{D}_2 gives $\Gamma \vdash M' : N$. The case when \mathcal{D} ends with WEAK is similar.
- If \mathcal{D} ends with Π -RULE, say with $M = \Pi x : B_1.B_2$, then the Stripping Lemma furnish some $(s_1, s_2, s_3) \in \mathcal{R}$ and derivations $\mathcal{D}_1 : \Gamma \vdash B_1 : s_1$ and $\mathcal{D}_2 : \Gamma, x : B_1 \vdash B_2 : s_2$ such that $N =_\beta s_3$. By definition of \rightarrow_β , two cases occur: if there is a σ -term B'_1 such that $B_1 \rightarrow_\beta B'_1$, then by IH₁ on \mathcal{D}_1 , we have $\mathcal{D}'_1 : \Gamma \vdash B'_1 : s_1$. Moreover, IH₂ on \mathcal{D}_2 gives $\mathcal{D}'_2 : \Gamma, x : B'_1 \vdash B_2 : s_2$, so applying Π -RULE on \mathcal{D}'_1 and \mathcal{D}'_2 gives $\Gamma \vdash (\Pi x : B'_1.B_2) : s_3$, on which CONV gives $\Gamma \vdash (\Pi x : B'_1.B_2) : N$. The second case when $B_2 \rightarrow_\beta B'_2$ for some σ -term B'_2 is the same (in fact, easier). The case when \mathcal{D} ends with λ -RULE is similar (and again has two subcases).
- If \mathcal{D} ends with APP, say with $M = B_1 B_2$, then reductions within either B_1 or B_2 are trivial. Thus, we can take $x \in \mathcal{V} \setminus \text{dom } \Gamma$ such that $B_1 = \lambda x : A_1.A_2$, and assume $M = (\lambda x : A_1.A_2)B_2 \rightarrow_\beta A_2[B_2/x]$. The Stripping Lemma then furnish σ -terms C_1 and C_2 such that $N =_\beta C_2[B_2/x]$ and derivations $\mathcal{D}_1 : \Gamma \vdash (\lambda x : A_1.A_2) : (\Pi x : C_1.C_2)$ and $\mathcal{D}_2 : \Gamma \vdash B_2 : C_1$. Again, the Stripping Lemma applied to \mathcal{D}_1 then furnish $(s_1, s_2, s_3) \in \mathcal{R}$, a σ -term C'_2 such that $\Pi x : C_1.C_2 =_\beta \Pi x : A_1.C'_2$, and derivations $\mathcal{E}_1 : \Gamma \vdash A_1 : s_1$, $\mathcal{E}_2 : \Gamma, x : A_1 \vdash C'_2 : s_2$, and $\mathcal{E}_3 : \Gamma, x : A_1 \vdash A_2 : C'_2$. Observe that $A_1 =_\beta C_1$, so CONV on \mathcal{D}_2 and \mathcal{E}_1 gives $\mathcal{D}_0 : \Gamma \vdash B_2 : A_1$, and using the Substitution Lemma with $(\mathcal{D}_0, \mathcal{E}_2)$ and $(\mathcal{D}_0, \mathcal{E}_3)$ give $\mathcal{E}'_2 : \Gamma \vdash C'_2[B_2/x] : s_2$ and $\mathcal{E}'_3 : \Gamma \vdash A_2[B_2/x] : C'_2[B_2/x]$; note that $\Gamma[B_2/x] = \Gamma$ since $x \notin \text{dom } \Gamma$. Finally, since $C_2 =_\beta C'_2$ and $N =_\beta C_2[B_2/x]$, applying CONV on \mathcal{E}'_2 and \mathcal{E}'_3 gives $\Gamma \vdash A_2[B_2/x] : N$. ■

REFERENCES

- [GN91] H. Geuvers and M. Nederhof, *Modular proof of strong normalization for the calculus of constructions*, Journal of Functional Programming **1** (1991), no. 2, 155-189.
- [Bar91] H. Barendregt, *Introduction to Generalized Type Systems*, Journal of Functional Programming **1** (1991), no. 2, 125-154.