

# **East West University Department of Computer Science and Engineering**

CSE 245(3) - LAB 02 [22 September 2016] Course Instructor: Dr. Mohammad Rezwanul Huq

# Implementation and Run-time Analysis of Merge and Quick Sort

### Lab Objective:

- To implement different sorting algorithms such as Quick sort, Merge sort.
- To compare and analyze their performance on different test cases.

#### Lab Outcome:

- Students are expected to be able to implement the aforesaid sorting algorithms and are also expected to solve real-life problems related to sorting.
- Furthermore, students are expected to understand the importance of the efficient algorithms by observing the performance of different sorting algorithms.

#### **Discussion:**

#### **Quick Sort:**

- Divide and conquer approach.
- Efficient than bubble, insertion and selection sort in most of the situations.
- Average case complexity  $\Theta$  (n log n); Worst case complexity  $\Theta$  (n<sup>2</sup>).

QUICKSORT $(A, p, r)$	PARTITION $(A, p, r)$
1 if $p < r$	1  x = A[r]
2   q = PARTITION(A, p, r)	2  i = p - 1
3 QUICKSORT $(A, p, q - 1)$	3 for $j = p \text{ to } r - 1$
4 QUICKSORT $(A, q + 1, r)$	4 if $A[j] \leq x$
	5   i = i + 1
	6 exchange $A[i]$ with $A[j]$
	7 exchange $A[i+1]$ with $A[r]$
	8 return $i+1$

#### **Merge Sort:**

- Divide and conquer approach.
- Time Complexity  $\Theta$  (n log n) but needs extra space.

```
MERGE-SORT(A, p, r)
                                 MERGE(A, p, q, r)
1 if p < r
                                     n_1 = q - p + 1
                                     n_2 = r - q
      q = \lfloor (p+r)/2 \rfloor
3
      MERGE-SORT(A, p, q)
                                  3 let L[1...n_1 + 1] and R[1...n_2 + 1] be new ar
      MERGE-SORT(A, q + 1, i
                                  4 for i = 1 to n_1
      MERGE(A, p, q, r)
                                  5
                                          L[i] = A[p+i-1]
                                  6 for j = 1 to n_2
                                  7
                                          R[j] = A[q+j]
                                     L[n_1+1]=\infty
                                  9 R[n_2 + 1] = \infty
                                 10 i = 1
                                      j = 1
                                 11
                                 12 for k = p to r
                                 13
                                          if L[i] \leq R[j]
                                              A[k] = L[i]
                                 14
                                              i = i + 1
                                 15
                                          else A[k] = R[j]
                                 16
                                 17
                                               j = j + 1
```

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## **Built-in Sort Function:**

```
#include <iostream>
#include <algorithm>
using namespace std;
const int SIZE = 7;
// function that determines the way to sort, e.g. ascending or descending
bool comp(int i, int j){
    if(i < j)
        return true;
    else
       return false;
int main()
    int intArray[SIZE] = \{5, 3, 32, -1, 1, 104, 53\};
    //Now we call the sort function
    sort(intArray, intArray + SIZE, comp);
    return 0;
}
```

#### **Lab Assignment (LAB 02):**

✓ Extend the program you have submitted in LAB01 by adding functions for merge sort, quick sort and built-in sort and then submitted the extended program.

Your program must be named as 245\_LAB02\_<your-student-id>.cpp and upload the file into <a href="https://dropitto.me/CSE245">https://dropitto.me/CSE245</a> by 26 September 2016 11:59:59 PM.

✓ Complete the evaluation sheet (soft copy posted in course page), reporting the performance of different sorting algorithms for different data size.

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