

use  $\Rightarrow$  sorting  $\rightarrow$  merge sort / quick sort  
trees / BST / Heaps / tries / segment trees  
DP (Dynamic Programming)  
Backtracking

### Observations:

- i) size keeps decreasing
- ii) similar dots only size changes
- iii) end dot (last dot to stop).

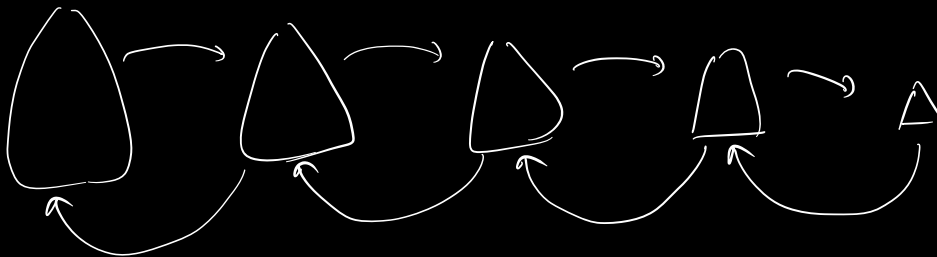
Recursion  $\Rightarrow$  a function calling itself.

technique of solving a problem using soln of

[smaller instances of same problem]



subproblem



$$\text{SUM}(N) = 1 + 2 + 3 + 4 + \dots + (N-1) + N$$

$\underbrace{\hspace{15em}}_{\text{sum}(N-1)}$

$$\Rightarrow \text{SUM}(N) = \underbrace{\text{SUM}(N-1)}_{\text{subproblem}} + N$$

$$\Rightarrow \text{SUM}(N-1) = 1 + 2 + 3 + 4 + \dots + (N-2) + (N-1)$$

$\underbrace{\hspace{15em}}_{\text{SUM}(N-2) + (N-1)}$

$$\Rightarrow \text{SUM}(N-1) = \text{SUM}(N-2) + (N-1)$$

$\Rightarrow$  learn how to write recursive code

$\Rightarrow$  How it works / dry run

$\Rightarrow$  TC / SC

# How do we write recursive code :-

Steps

i) Assumption

decide, what your func does?

ii) main logic

solving the main problem using subproblem

iii) base cond<sup>n</sup>

when should recursion stop.

# Find the sum of N natural nos. using recursion

```
int sum(N) { // assume :- sum of N natural nos.  
    // base case  
    if (N == 1)  
        return 1;  
    // main logic  
    return sum(N-1) + N;  
}
```

# Find the factorial of N.

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

$$N! = N \times \underbrace{N-1 \times N-2 \times N-3 \times \dots \times 1}_{(N-1)!}$$

$$\Rightarrow N! = N \times (N-1)!$$

$$\boxed{\begin{array}{l} 1! = 1 \\ 0! = 1 \end{array}}$$

code

```
int fact(N) { // assume :- that this fun calculates  
              fact(N)  
  // base condn  
  if (N == 0)  
    return 1;  
  
  // main logic  
  return fact(N-1) * N;  
}
```

# Given N, find  $N^{\text{th}}$  fibonacci number.

|                  |   |   |   |   |   |   |   |    |    |    |    |     |
|------------------|---|---|---|---|---|---|---|----|----|----|----|-----|
| N                | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7  | 8  | 9  | 10 | ... |
| fb $\rightarrow$ | 0 | 1 | 1 | 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 |     |

Fibonacci series  $\Rightarrow$  sum of last two nos.

[\* fibonacci series & golden ratio]

$$fb(8) = 21 \Rightarrow 8 + 13 \Rightarrow \underline{fb(6) + fb(7)}$$

$$fb(8) = fb(6) + fb(7)$$

$$\Rightarrow fb(N) = fb(N-1) + fb(N-2) \leftarrow$$

int fib(N) { // assume + calculate fibonacci for N<sup>th</sup> posn

if (N == 0)  
return 0

else if (N == 1)  
return 1

// main logic

return fib(N-1) + fib(N-2)

}

base case

N = 0

$$fib(0) = \cancel{fib(-1)} + \cancel{fib(-2)}$$

$$fib(1) = fib(0) + \cancel{fib(-1)}$$

$$fib(2) = \underbrace{fib(0)}_0 + \underbrace{fib(1)}_1$$

## # Working of recursion:-

→ Sum of N natural nos:-

```
int sum(N) {
```

```
    if (N == 1)
```

```
        return 1
```

```
    return sum(N-1) + N;
```

```
}
```

sum(5)

return ~~sum(4)~~ + 5

15

10

sum(4)

return ~~sum(3)~~ + 4

6

sum(3)

return ~~sum(2)~~ + 3

3

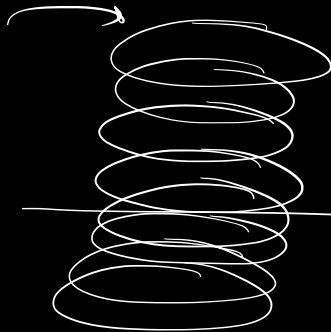
sum(2) ↓  
return ~~sum(1)~~ + 2

1

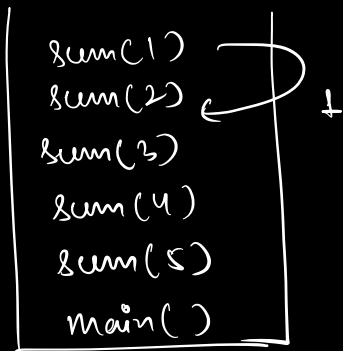
sum(1)

## # Call Stack

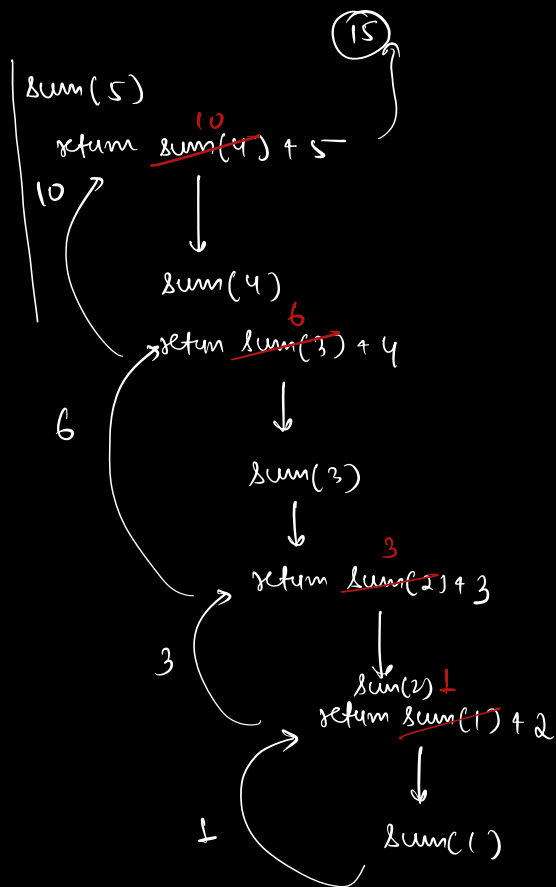
Stack :-



Stack → LIFO →  
↓ ↓ ↓ ↓  
last in first out



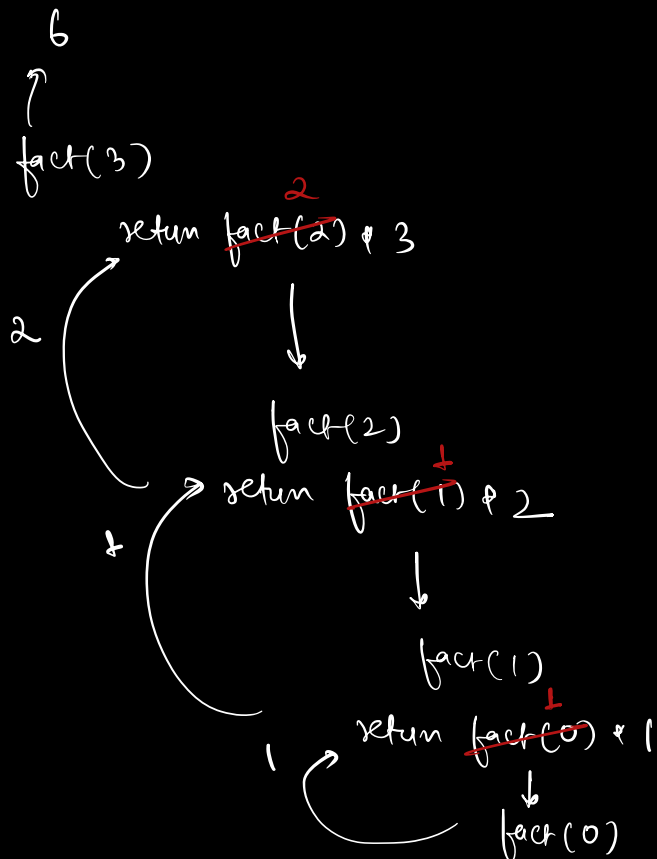
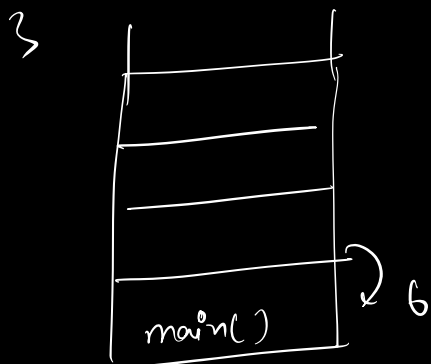
↓  
extra space  
required for  
call stack



# factorial dry run :-

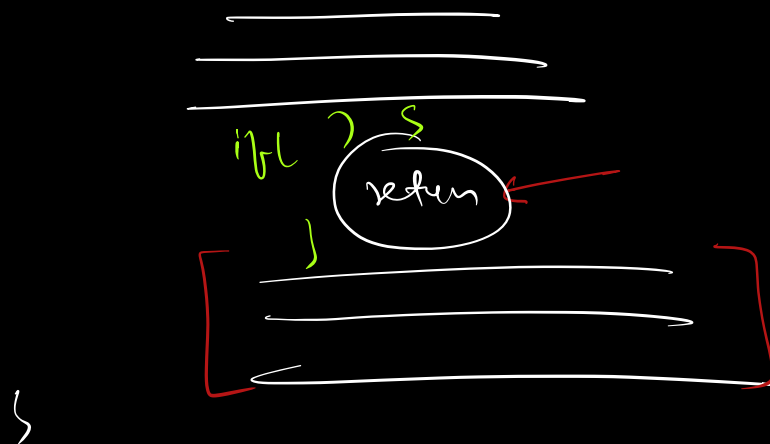
```

int fact(N) {
    if (N == 0)
        return 1
    return fact(N-1) * N
}
  
```



- return  $\Rightarrow$  keyword stops the execution of func

func() {



Q 1. Given a no.  $N$ , print all the nos. from 1 to  $N$ , increasing order using recursion.

If  $p = 5$ ,  $O(p) \Rightarrow 1 \ 2 \ 3 \ 4 \ 5$

void print( $N$ ) {

if ( $N == 0$ )  
return;

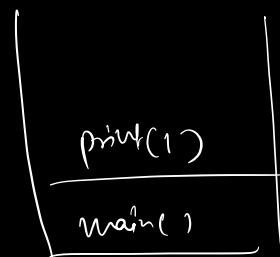
print( $N-1$ )

cout( $N$ )

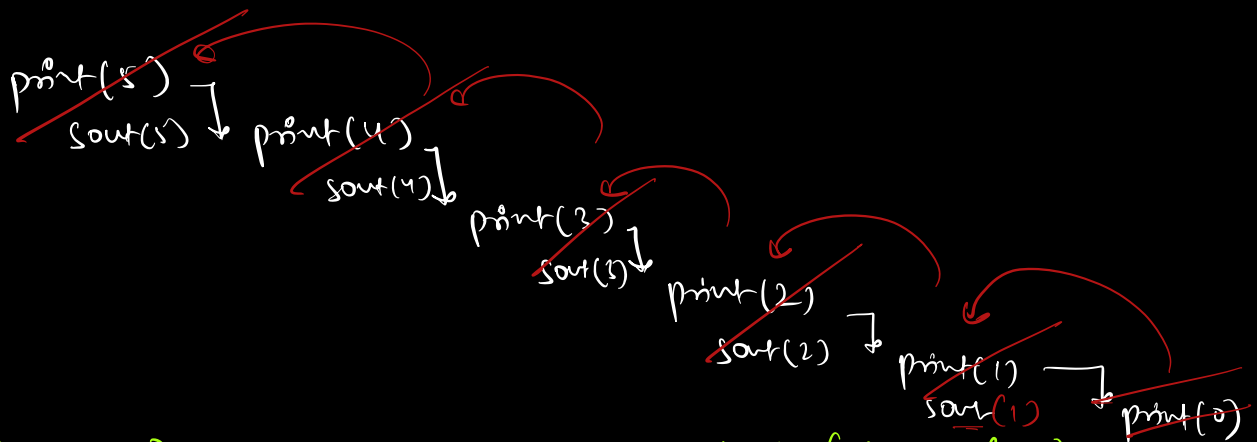
}

print(5)  $\Rightarrow 1 \ 2 \ 3 \ 4 \ 5$

print(1)







Q 2. Print all nos. from N to 1 (dec. order)

```
void print(N) {
```

```
    if (N == 0)
        return;
```

```
    sout(N)
```

```
    print(N-1)
```

```
}
```

↓

print(5)

↓  
5

print(4)

↓  
4

print(3)

↓  
3

print(2)

↓  
2

print(1)

↓  
1

print(0)

(4

print(5) → 5 4 3 2 1

print(4) → 4 3 2 1

print(5) → sout(5) + print(4)

Q 3. Given a string, check if it is a palindrome using recursion.

Palindrome  $\Rightarrow$  M A D A M  
L E V E L  
M A L A Y A L A M



$\uparrow \quad \uparrow$   
 $e \quad s$

Stopping cond<sup>n</sup>  $\Rightarrow \underline{s \geq e}$   
base case

```
boolean isPalindrome(str, s, e) {  
    if (s > e)  
        return true  
    if (str[s] == str[e]) {  
        return isPalindrome(str, s+1, e-1)  
    }  
    else { return false  
    }  
}
```

|   |   |   |   |   |
|---|---|---|---|---|
| 0 | 1 | 2 | 3 | 4 |
| M | A | D | A | M |
|   |   | ↑ |   |   |
|   |   | S |   |   |
|   |   |   | ↑ |   |
|   |   |   | e |   |

isPalin (str, 0, 4)

isPalin (str, 1, 3)

isPalin (str, 2, 2)

isPalin (str, 3, 1)

|   |   |   |   |
|---|---|---|---|
| 0 | 1 | 2 | 3 |
| A | B | A | A |
|   | ↑ |   |   |
|   | S |   |   |
|   |   | ↑ |   |
|   |   | e |   |

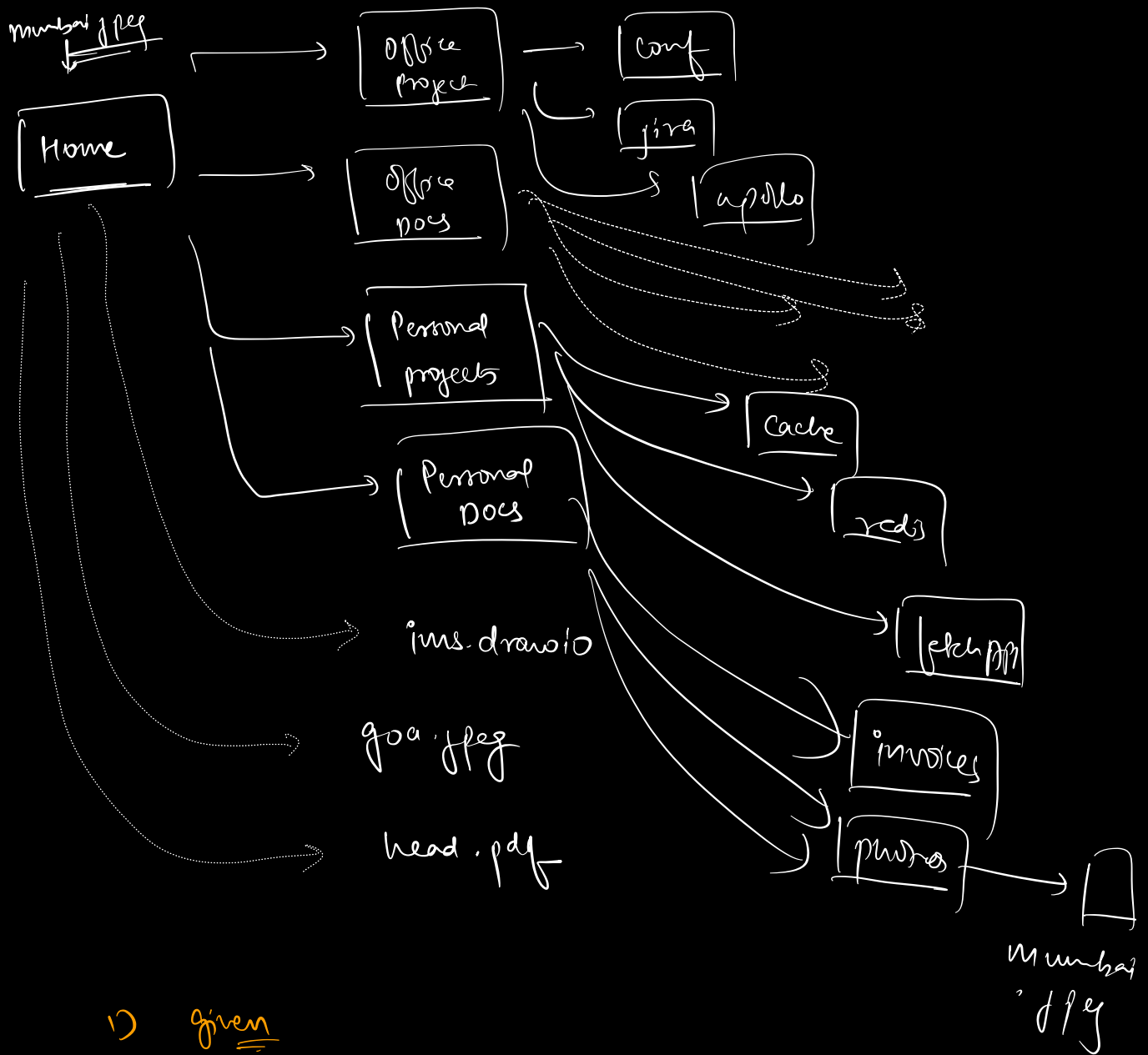
isPalin (str, 0, 3)

↳ isPalin (str, 1, 2) ↴

false

Facebook, Google

Q4, Given a directory structure and you have to search a file.



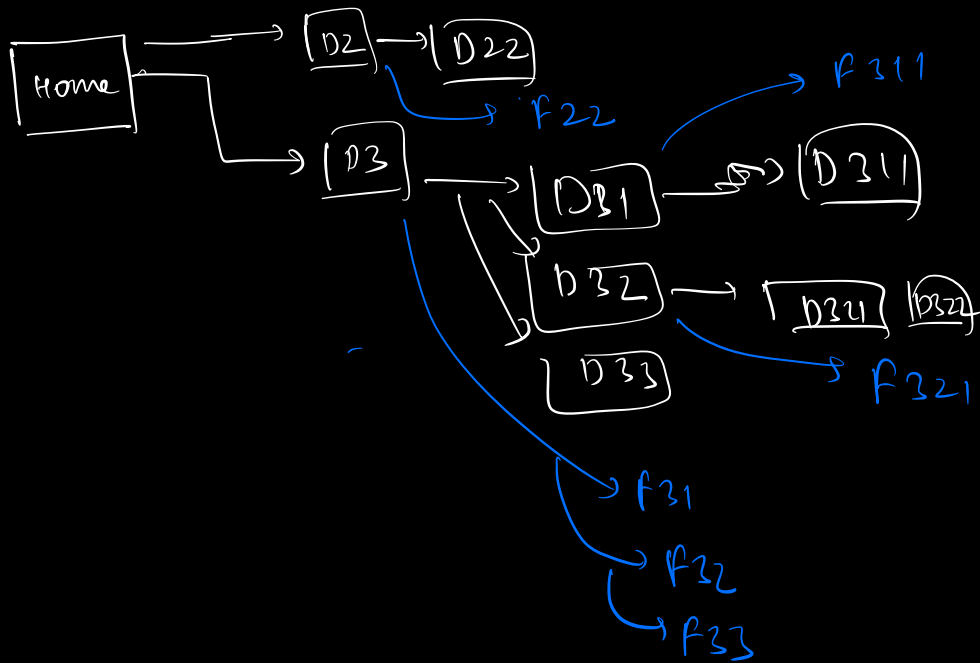
1) given

List < > get All Directory (directoryName)

List < > get All files (directoryName)

O/P ⇒ true / false ⇒ fileName

F311



```
boolean search (directoryName, fileName) {
```

```
    List<String> files = getAllFiles(directoryName);
```

```
    for (i=0; i < files.size(); i++) {
```

```
        if (files[i] == fileName)
```

```
            return true
```

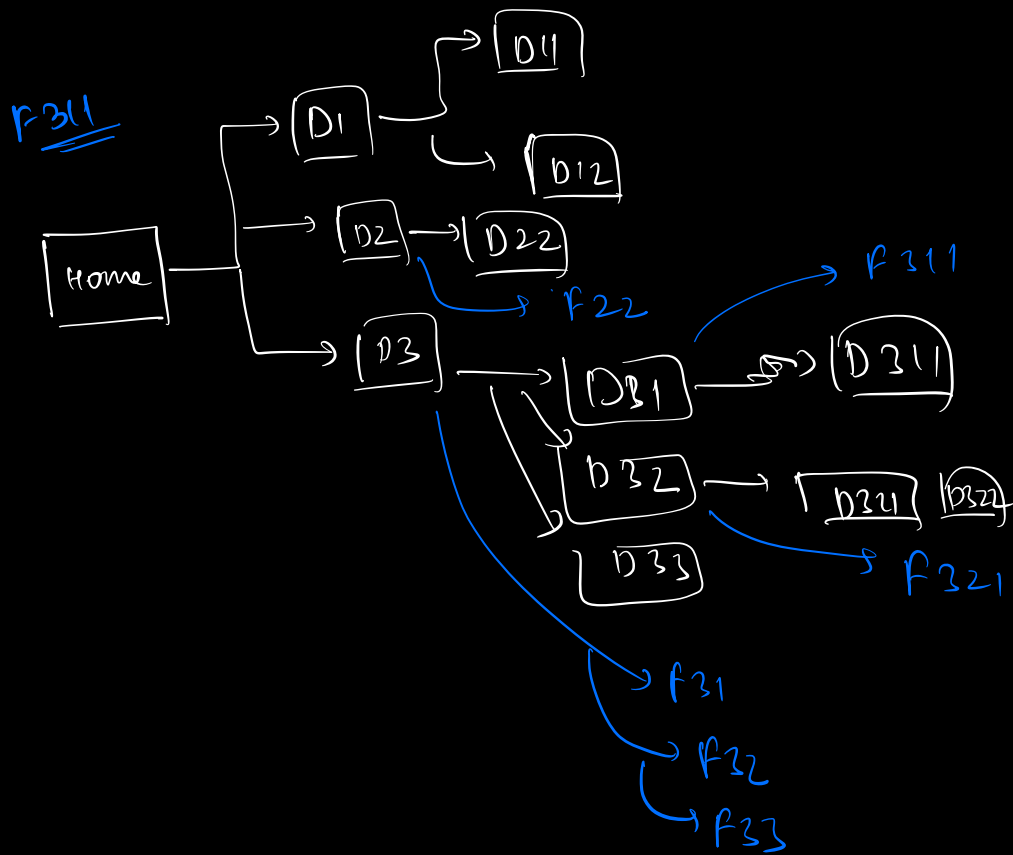
```
    }
```

```
    List<String> directories = getAllDirs(directoryName);
```

```

for( i=0; i < directories.size(); i++) {
    if( search( directories[i], fileName) )
        return true
}
return false
}

```



search( home, F311)

Search (D1, F311)

Rec-2

⇒ LinkedList ←

Queues

Trees 1

Trees 2

Someone else

Subsequence & Subsets

Advanced