Data Structures Chapter 3

- 1. Stack
- 2. Queue
- 3. Stack Applications
 - Arithmetic Expressions
 - Infix, Prefix, and Postfix
 - Arithmetic Expression Evaluation
 - Dijkstra's Two-Stack Algorithm
 - Postfix Evaluation

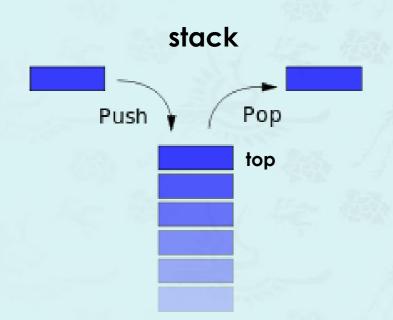


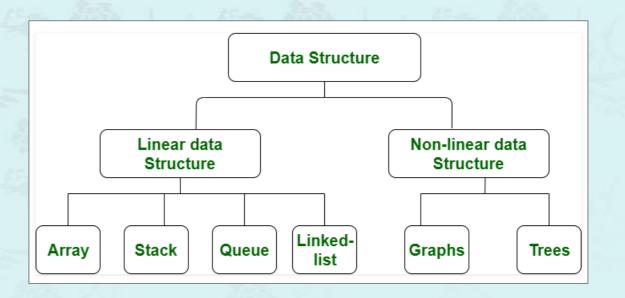
내 아들들을 먼 곳에서 이끌며 내 딸들을 땅 끝에서 오게 하며 내 이름으로 불려지는 모든 자 곧 내가 내 영광을 위하여 창조한 자를 오게 하라 그를 내가 지었고 그를 내가 만들었노라 (사 43:6-7)

Bring my sons from afar and my daughters from the end of the earth – everyone who is called by my name, whom I created for my glory, whom I formed and made.

Stack and Queue

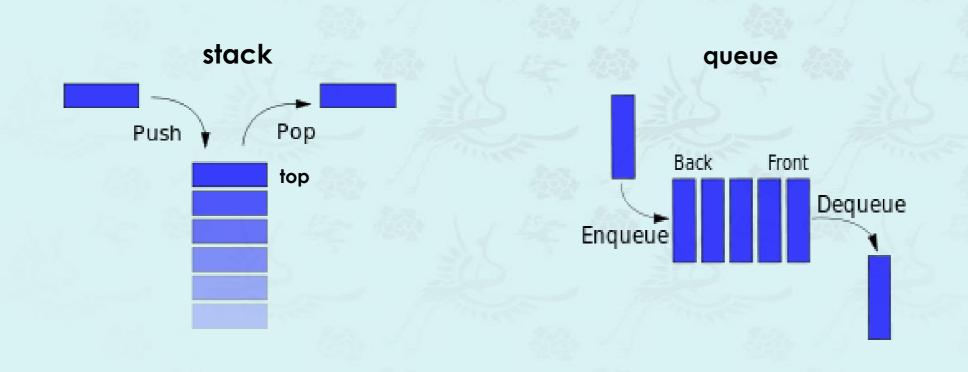
Stack is a linear data structure handled by a particular order of the operation is called LIFO(Last In First Out). It removes the item most recently added.





Stack and Queue

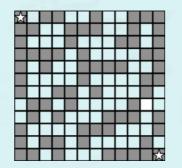
- Stack is a linear data structure handled by a particular order of the operation is called LIFO(Last In First Out). It removes the item most recently added.
- Queue is known as a Fist-in-first-out(FIFO) list since it removes the item least recently added.

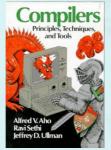


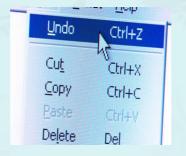
Stack Applications

- Parsing in a compiler.
- Undo in a word processor.
- Back button in a Web browser.
- PostScript language for printers.
- Backtracking as in a maze
- Implementing function calls in a compiler.

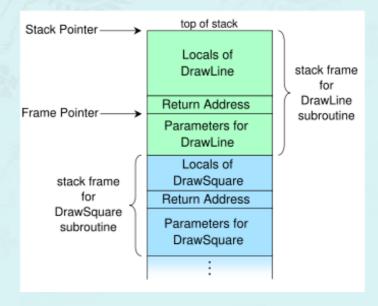
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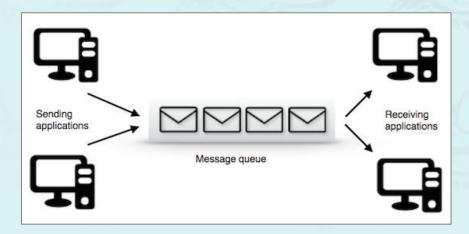


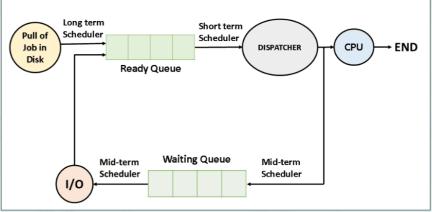




Queue Applications

- In a computer OS, requests for services come in unpredictable order and timing, sometimes faster than they can be serviced.
 - print a file
 - send an email
 - CPU scheduling
 - job scheduling in general







Infix, postfix and prefix expressions

Stacks can be used to implement algorithms involving Infix, postfix and prefix expressions.

Infix:

- An infix expression is a single letter, or an operator, proceeded by one infix string and followed by another infix string.
- A, A + B, (A + B) + (C D)

Prefix:

- A prefix expression is a single letter, or an operator, followed by two prefix strings. Every
 prefix string longer than a single variable contains an operator, first operand and second
 operand.
- A, + A B, + + A B C D

Postfix:

- A postfix expression (also called Reverse Polish Notation) is a single letter or an operator, preceded by two postfix strings. Every postfix string longer than a single variable contains first and second operands followed by an operator.
- A, A B +, A B + C D +

Infix, postfix and prefix expressions

- Prefix and postfix notations are methods of writing mathematical expressions without parenthesis.
- Why: Time to evaluate a postfix and prefix expression is O(n), where n is the number of elements in the array.

Infix	Prefix	Postfix
A + B	+ A B	A B +
A + B - C	- + A B C	A B + C -
(A + B) * C - D	- * + A B C D	A B + C * D -

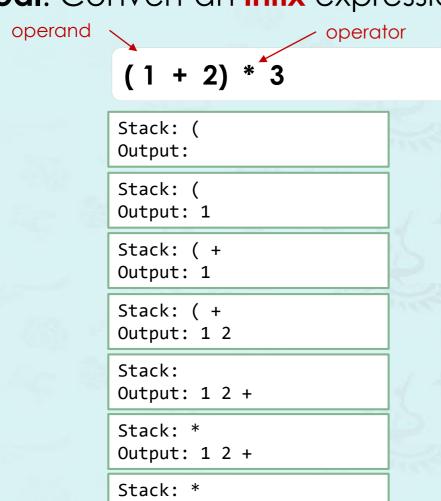
Infix, postfix and prefix expressions

- Prefix and postfix notations are methods of writing mathematical expressions without parenthesis.
- Why: Time to evaluate a postfix and prefix expression is O(n), where n is the number of elements in the array.

infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
(1 + 2) * 7	1 2 + 7 *
a * b / c	a b * c /
(a/(b-c+d))*(e-a)*c	a b c - d + / e a - * c *
a / b - c + d * e - a * c	a b / c - d e * + a c * -

Infix to Postfix Conversion

Goal: Convert an infix expression to a postfix expression using a stack.



Output: 12 + 3

Output: 1 2 + 3 *

Stack:

```
1. Operands are output immediately
                2. Push "(" always and operators in general.
                3. For ")", pop until "(". Discard "(" and ")".
                4. For operators, push it.
                   However, first pop any operators that have
                   higher or equal precedence.
  infix (1 + ((2 + 3) * (4 * 5)))
postfix 1 2 3 + 4 5 * * +
```

Infix to Postfix Conversion

Example 1: 3+4*5/6

in	stack(bottom to top)	postfix
3		3
+	+	
4		3 4
*	+ *	
5		3 4 5
/	+ /	3 4 5 *
6		3 4 5 * 6
		3 4 5 * 6 / +

- 1. Operands are output immediately
- 2. Push "(" always and operators in general.
- 3. For ")", pop until "(". Discard "(" and ")".
- For operators, push it.
 However, first pop any operators that have higher or equal precedence.

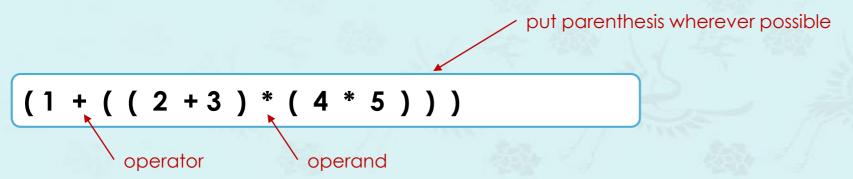
Infix to Postfix Conversion

Example 2: (1+3)*(4-2)/(5+7)

in	stack (bottom to top)	postfix	in	stack	postfix	
(((/ (1 3 + 4 2 - *	
1		1	5		1 3 + 4 2 - * 5	
+	(+		+	/ (+		
3		1 3	7		1 3 + 4 2 - * 5 7	
)		1 3 +)		1 3 + 4 2 - * 5 7 +	
*	*				1 3 + 4 2 - * 5 7 + /	
(* (
4		1 3 + 4				
-	* (-				 Operands are output immediately Push "(" always and operators in general. For ")", pop until "(". Discard "(" and ")". For operators, push it. However, first pop any operators that have higher or equal precedence. 	
2		1 3 + 4 2				
)	*	1 3 + 4 2 -				
/	1	1 3 + 4 2 - *				

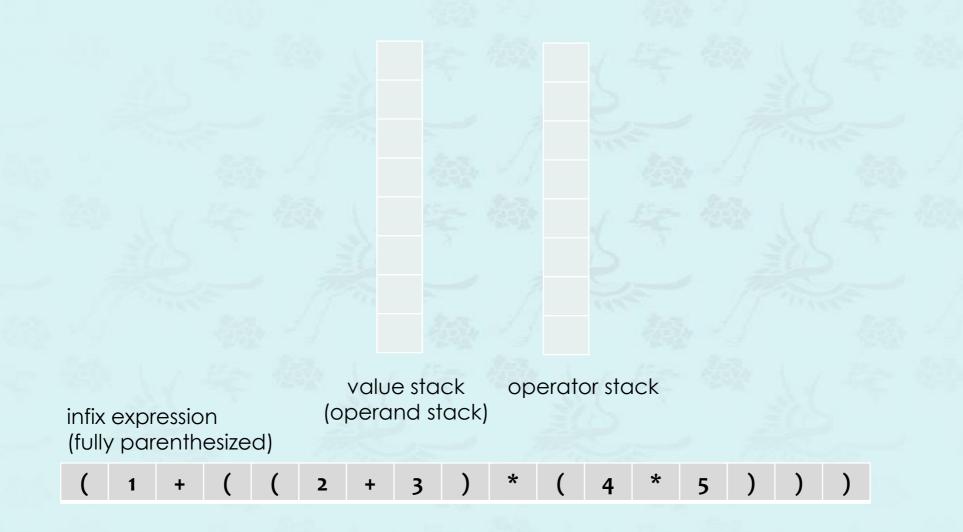
Infix Expression Evaluation

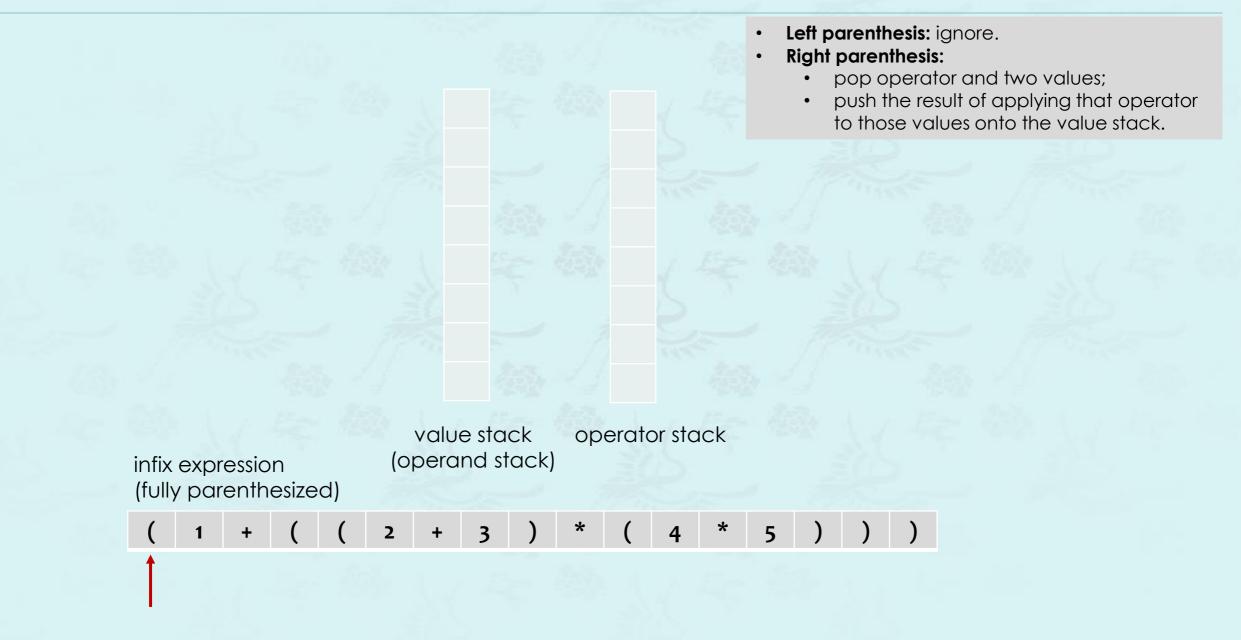
Goal: Evaluate infix expressions.

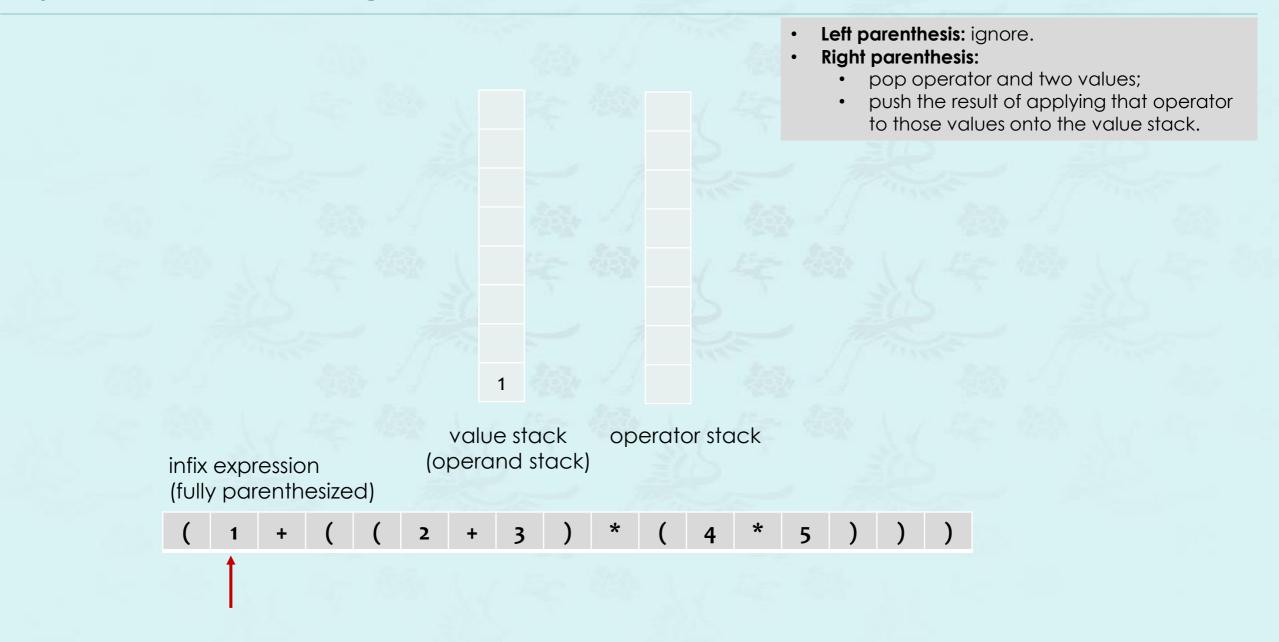


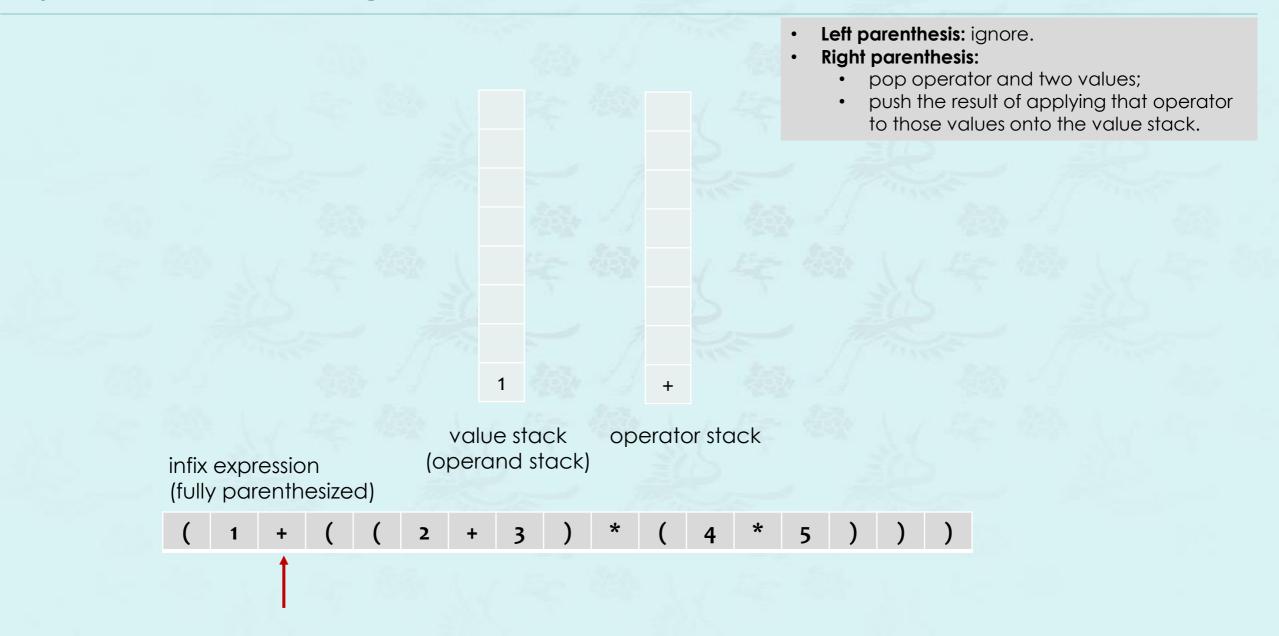
Two-stack algorithm. [E. W. Dijkstra]

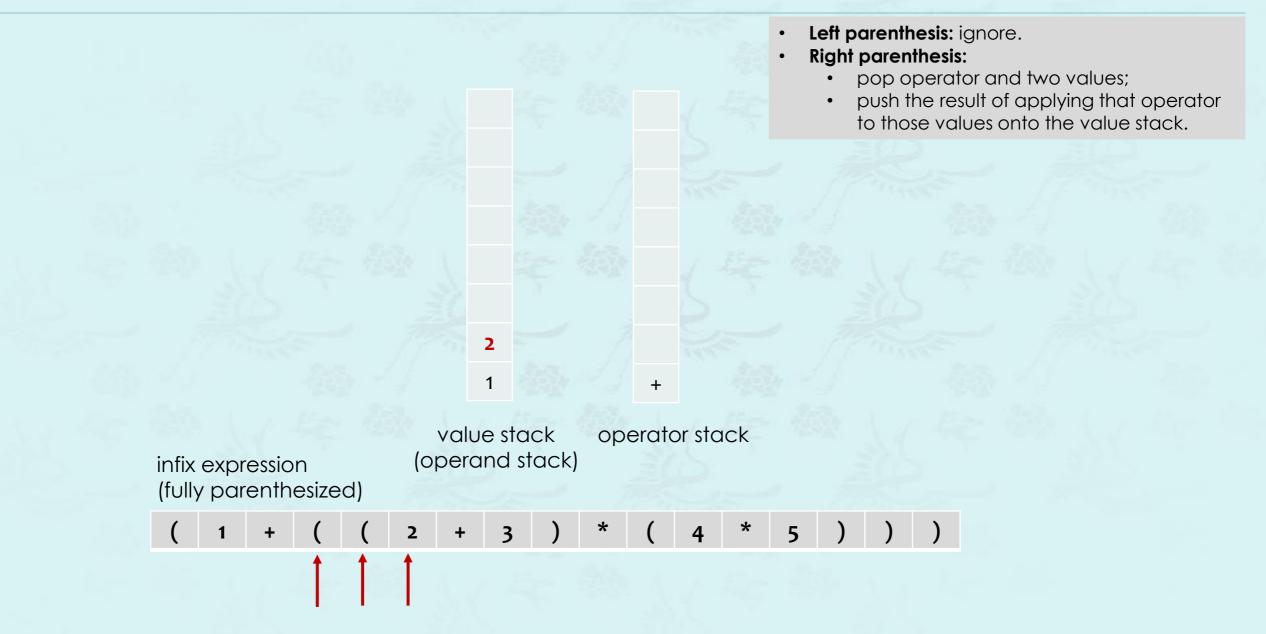
- Value: push onto the value stack.
- Operator: push onto the operator stack.
- Left parenthesis: ignore.
- Right parenthesis:
 - pop operator and two values;
 - push the result of applying that operator to those values onto the value stack.

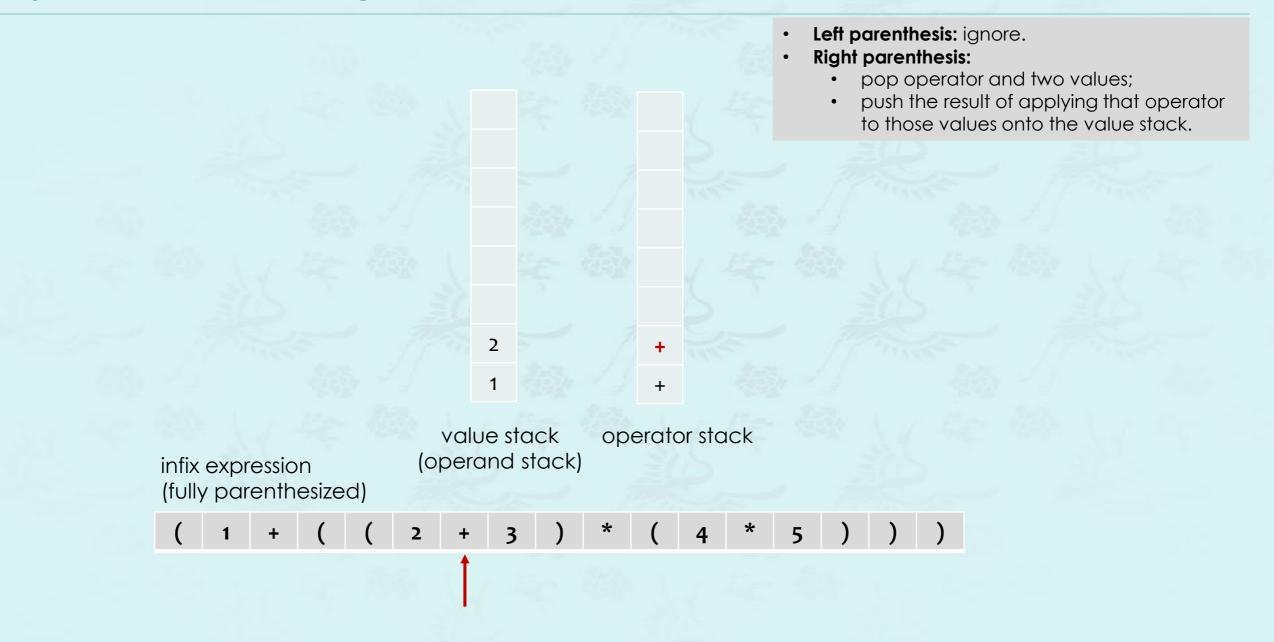


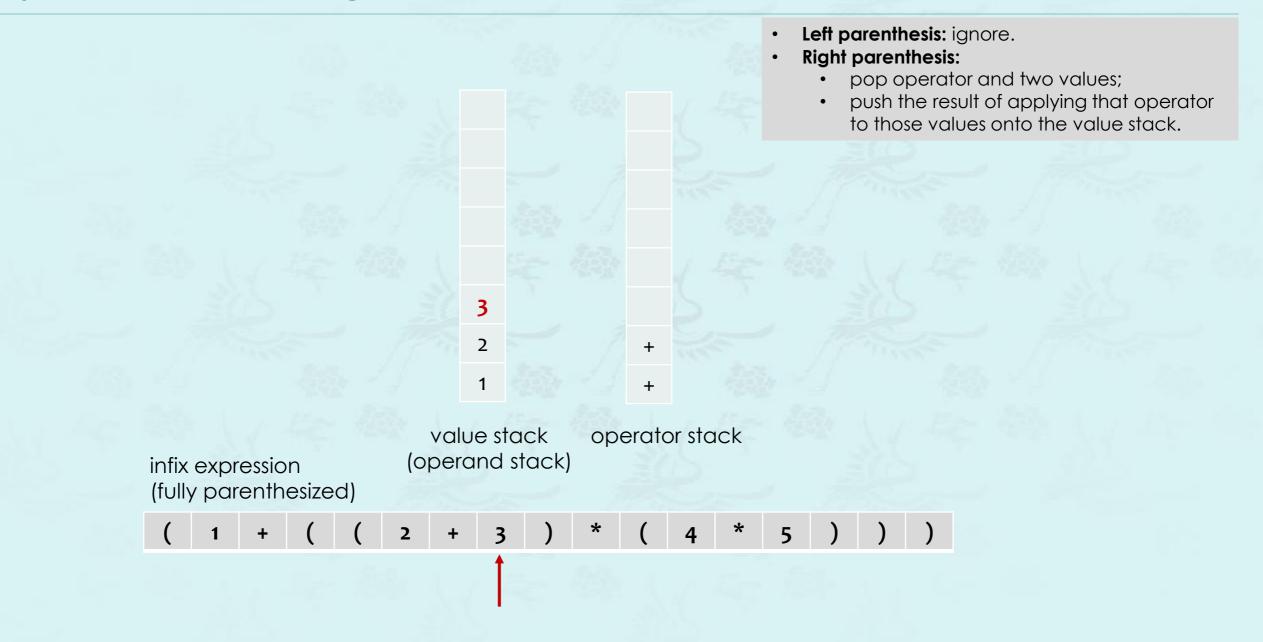


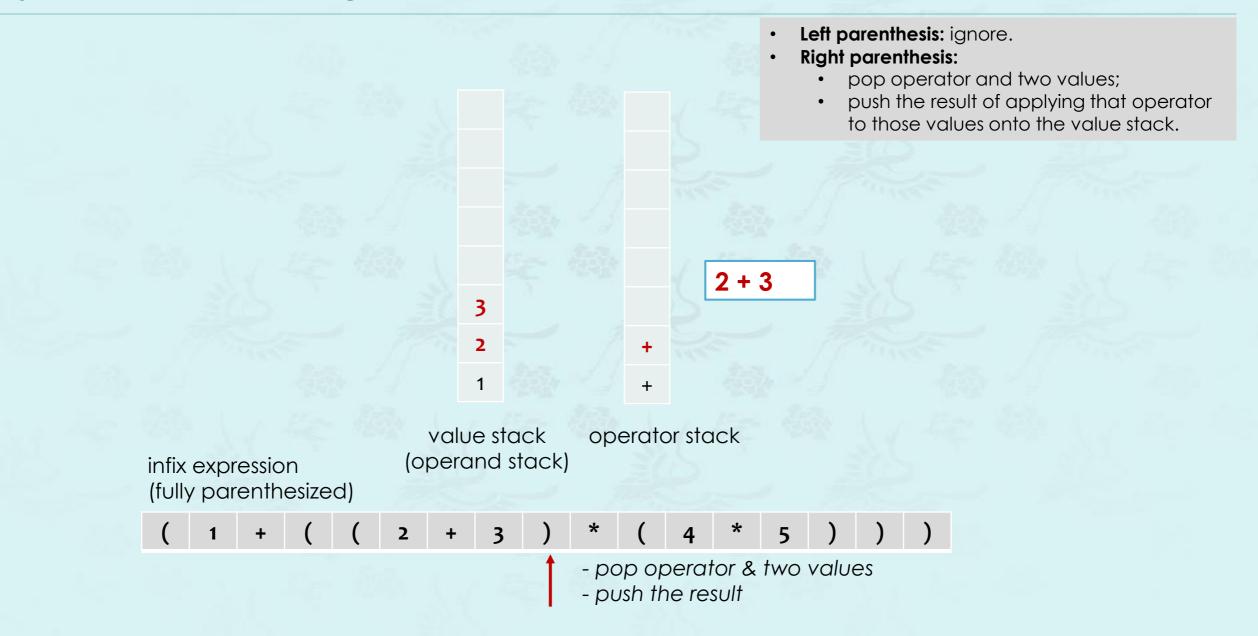


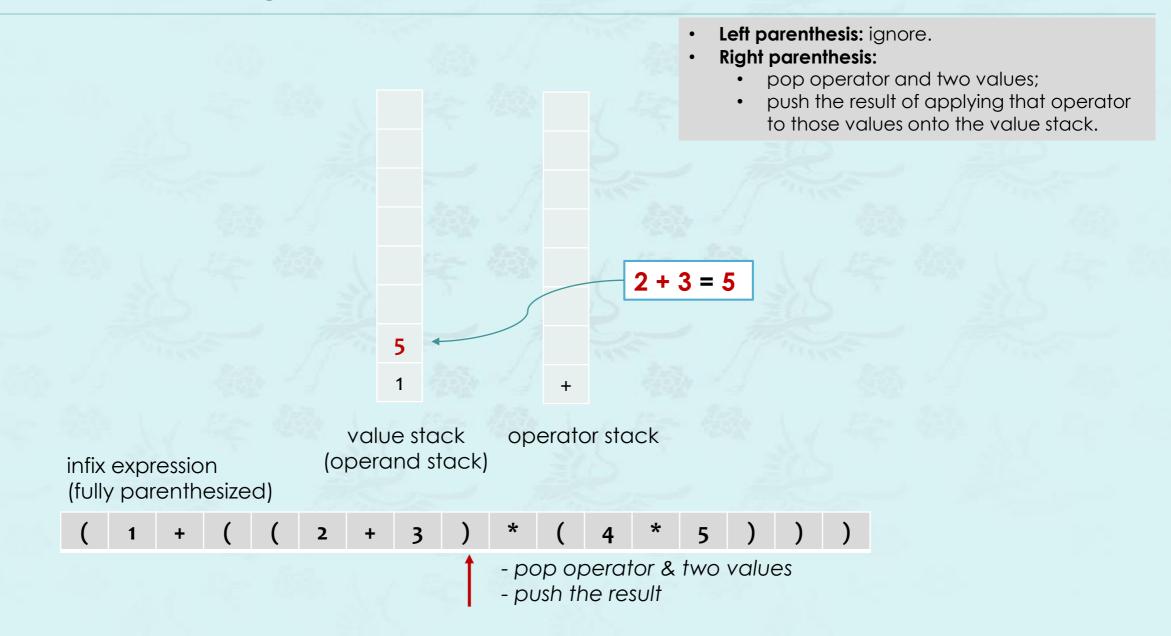


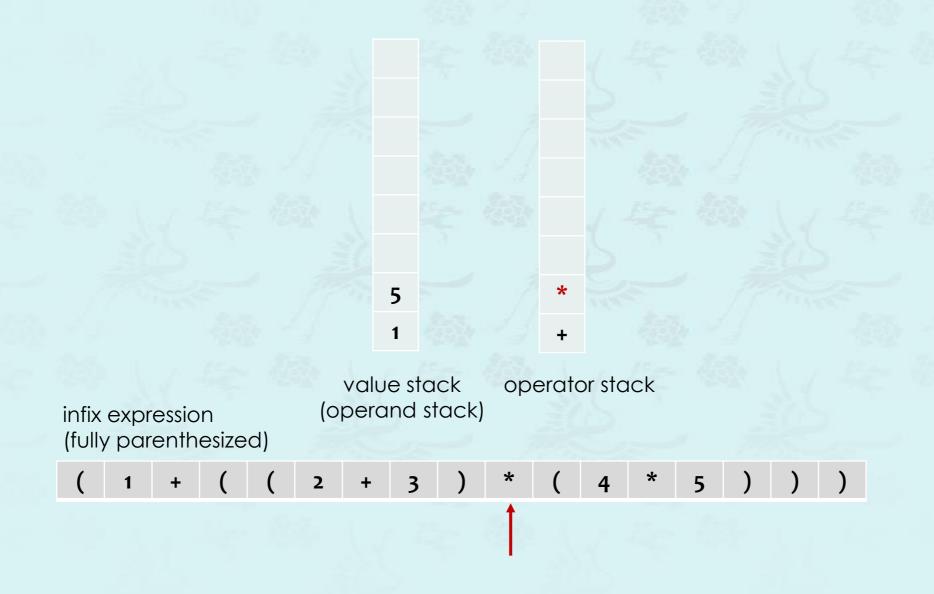


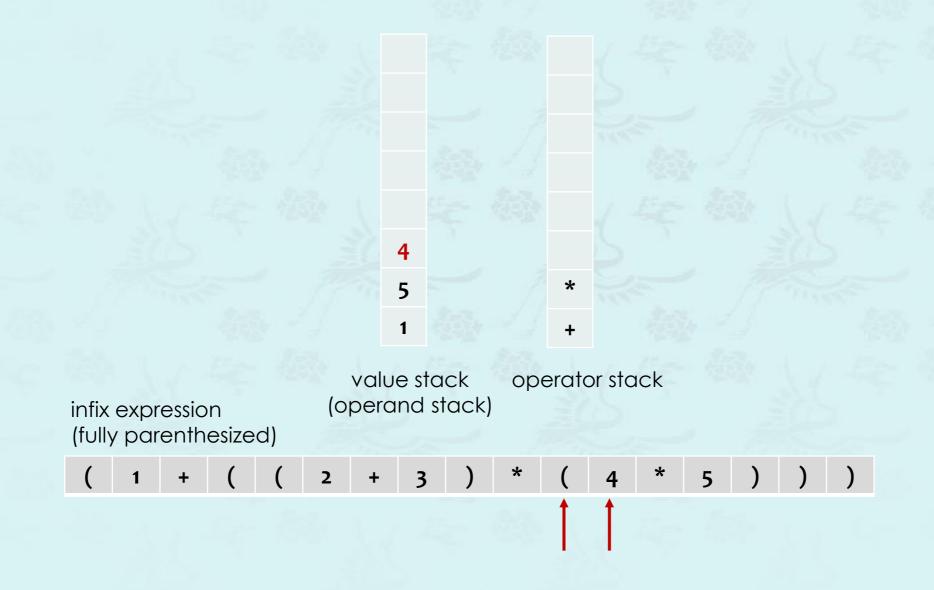


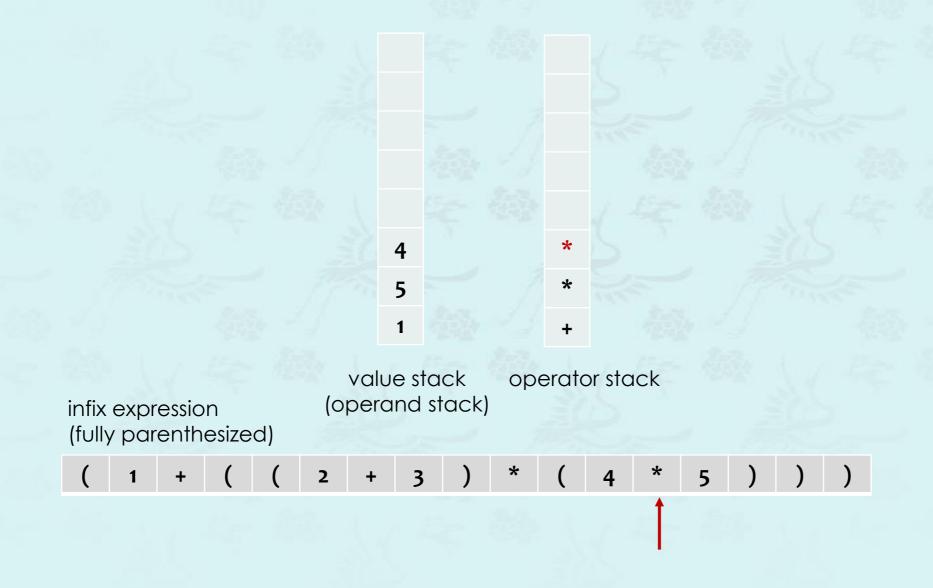


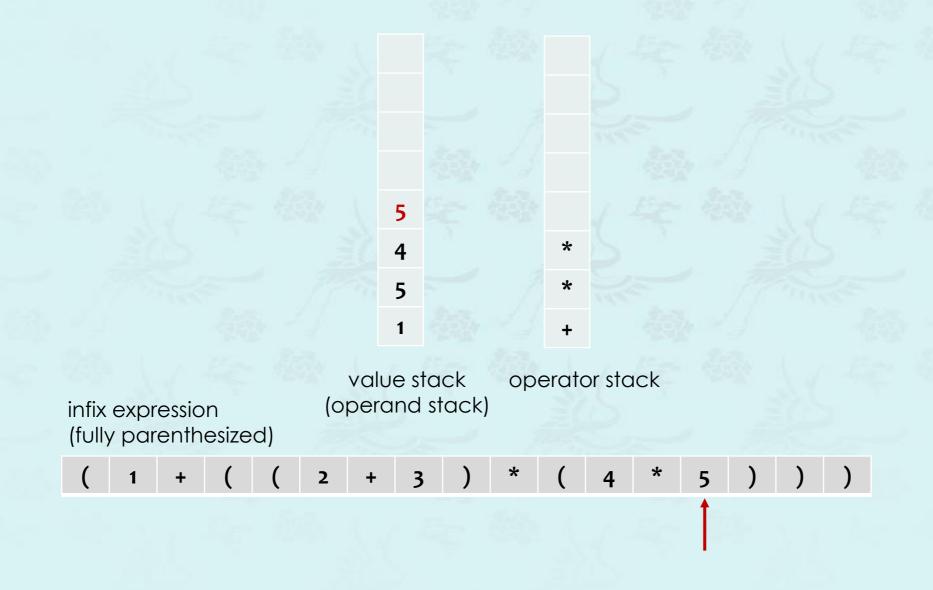


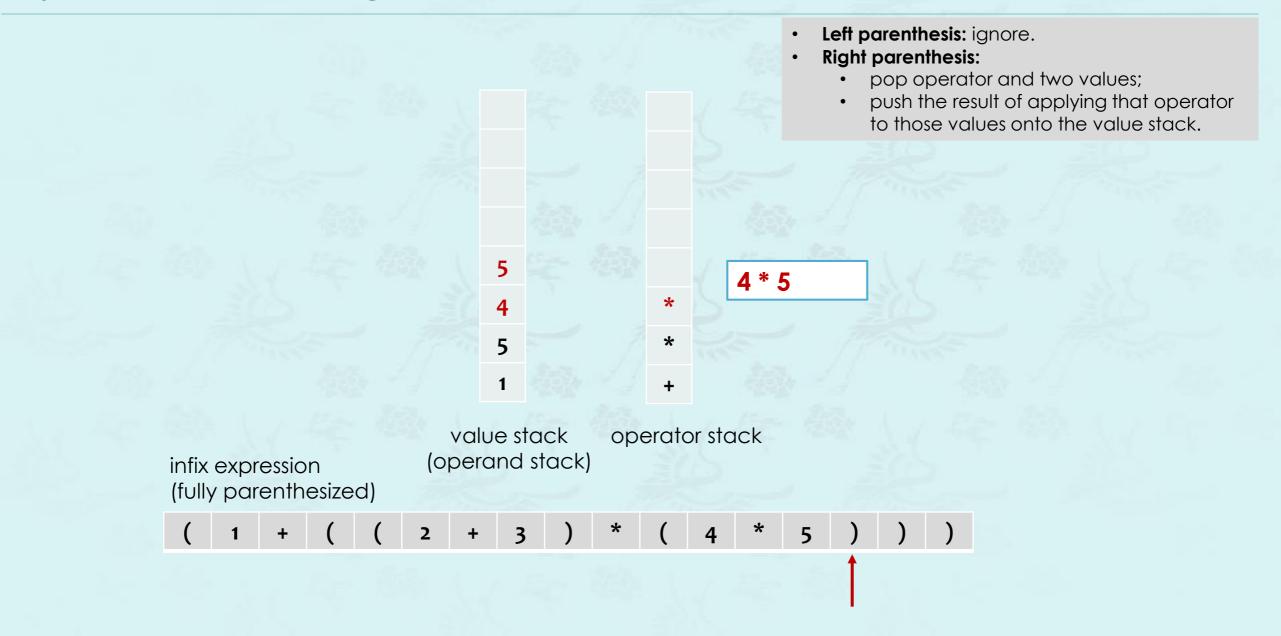


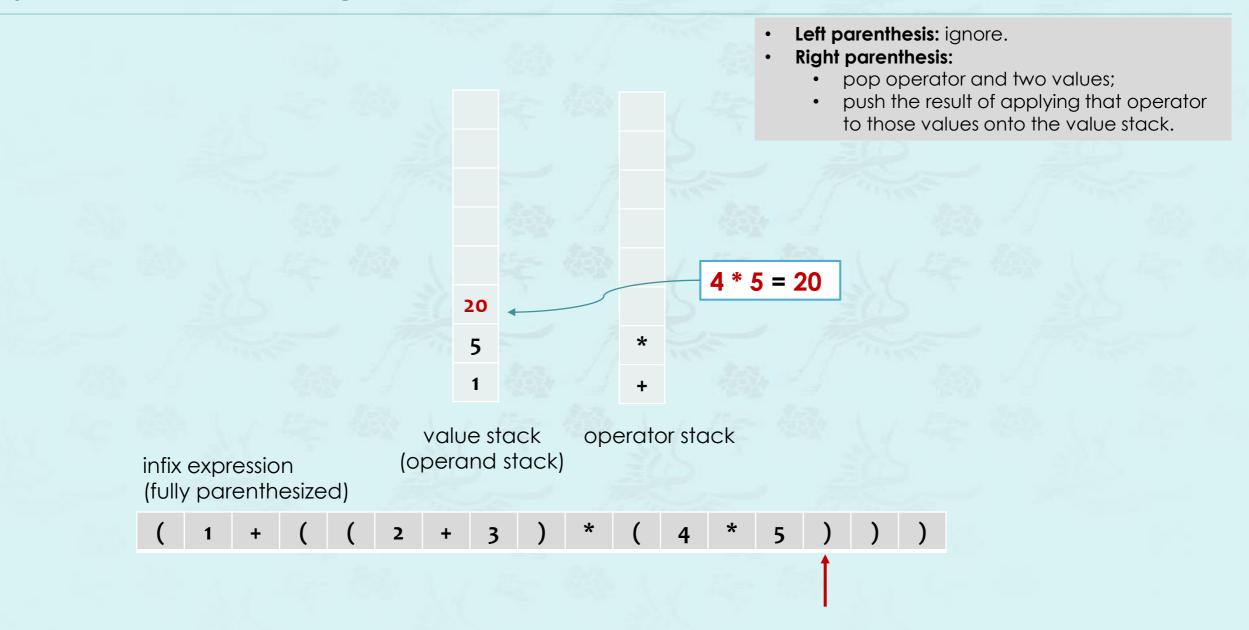


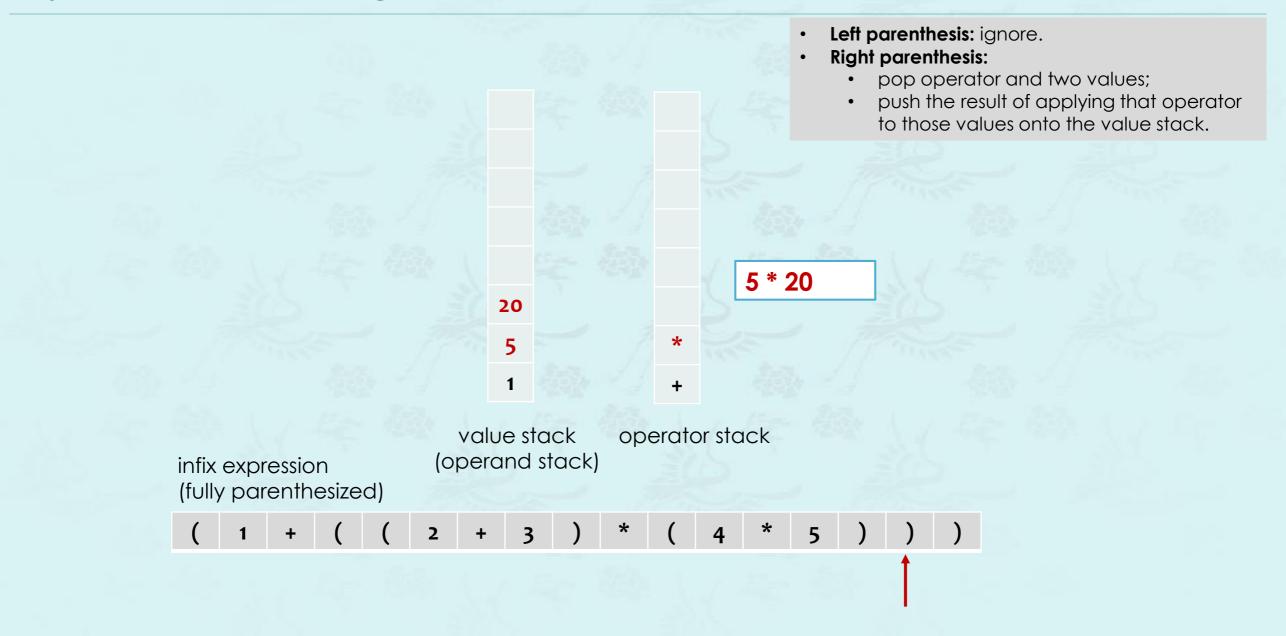


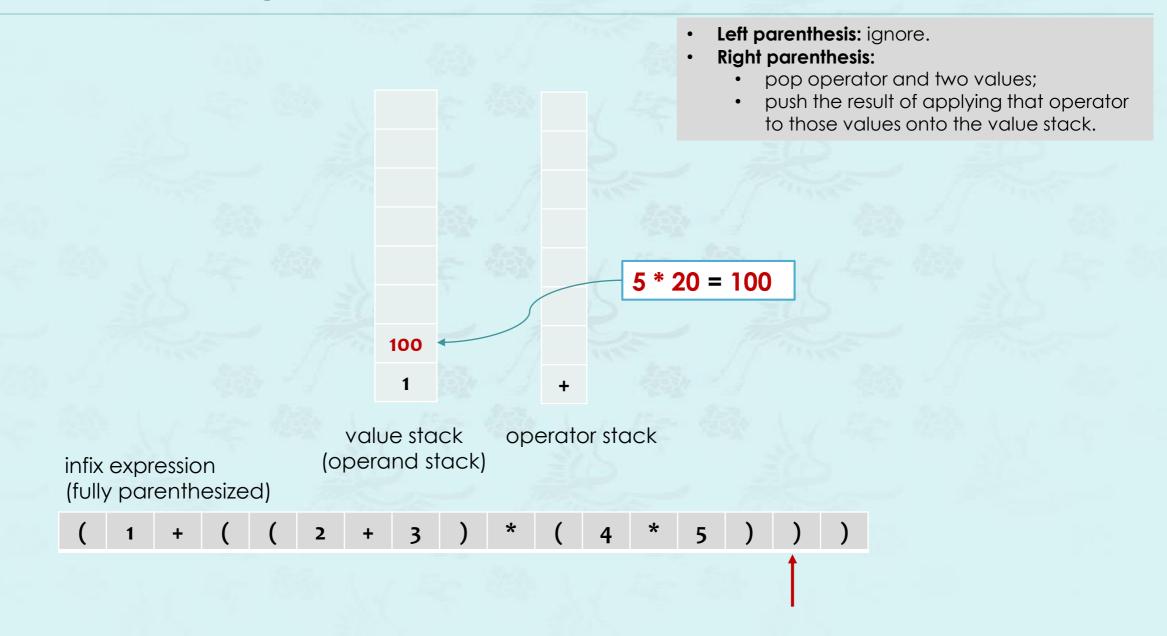


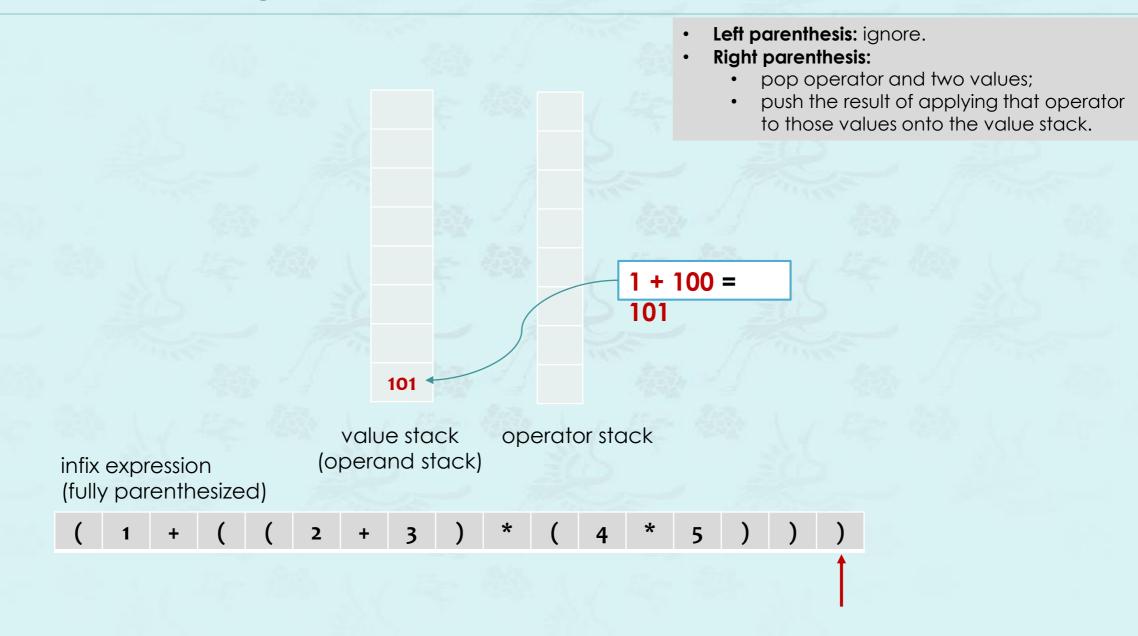


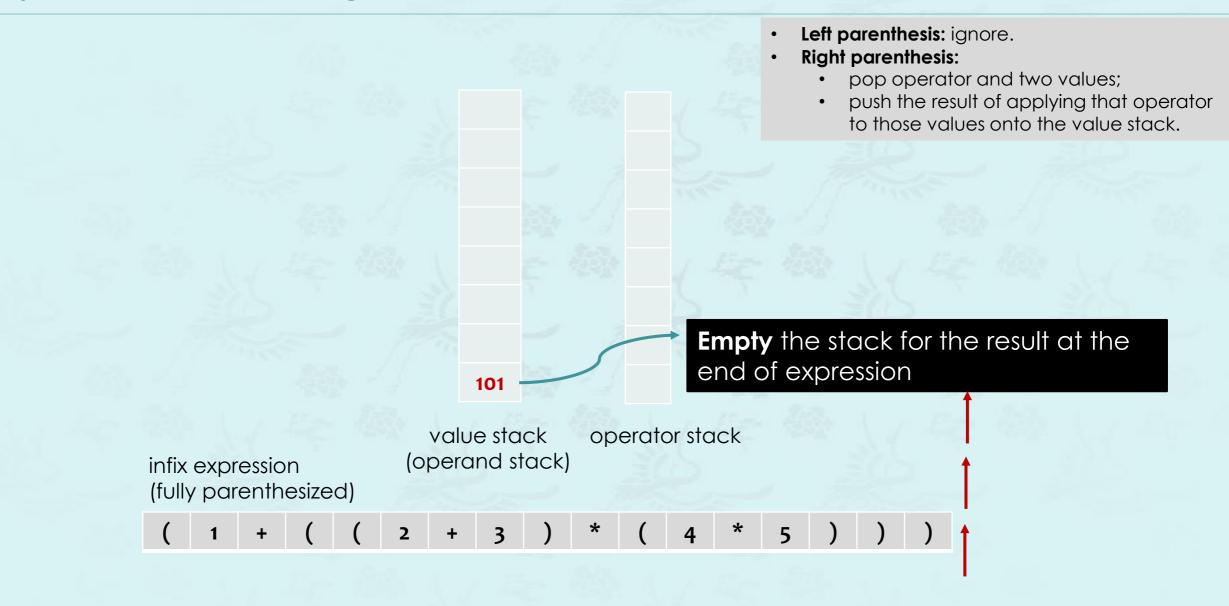












- Q: How does it work
- A: When algorithm encounters an operator surrounded by two values within parentheses, it leaves the result on the value stack.

```
(1 + ( ( 2 + 3 ) * ( 4 * 5 ) ) )
```

as if the original input were:

```
(1 + ( 5 * ( 4 * 5 ) ) )
```

Repeating the argument:

```
(1 + (5 * 20))
(1 + 100)
101
```

Extensions: More ops, precedence order, associativity.

 Observation 1. Dijkstra's two-stack algorithm computes the same value if the operator occurs after the two values.

Observation 2. All of the parentheses are redundant!

```
1 2 3 + 4 5 * * +
```

- Bottom line: Postfix or "reverse Polish" notation.
- Applications: Postscript, calculators, JVM,

Dijkstra's two-stack algorithm – version.1 (fully parenthesized)

- While there are still tokens to be read in,
 - 1.1 Get the next token.
 - 1.2 If the token is:
 - 1.2.1 A space: ignore it
 - 1.2.2 A left brace: ignore it
 - 1.2.3 A number:
 - 1.2.3.1 read the number (it could be a multiple digit.)
 - 1.2.3.2 push it onto the value stack
 - 1.2.4 A right parenthesis:
 - 1.2.4.1 Pop the operator from the operator stack.
 - 1.2.4.2 Pop the value stack twice, getting two operands.
 - 1.2.4.3 Apply the operator to the operands, in the correct order.
 - 1.2.4.4 Push the result onto the value stack.
 - 1.2.5 An operator
 - 1.2.5.1 Push the operator to the operator stack
- 2 (The whole expression has been parsed at this point.
 - Apply remaining operators in the op stack to remaining values in the value stack) While the operator stack is not empty,
 - 2.1 Pop the operator from the operator stack.
 - 2.2 Pop the value stack twice, getting two operands.
 - 2.3 Apply the operator to the operands, in the correct order.
 - 2.4 Push the result onto the value stack.
- 3 (At this point the operator stack should be empty, and the value stack should have only one value in it, which is the result.)
 Return the top item in the value stack.

Dijkstra's two-stack algorithm – version.1 (fully parenthesized)

```
int evaluate(string tokens) {
    stack<int> va stack;
                                         // stack to store operands or values
    stack<char> op stack;
                                          // stack to store operators.
    for (int i = 0; i < tokens.length(); i++) {</pre>
        if (tokens[i] == ' ') continue; // skip if token is whitespace or (
        if (tokens[i] == '(') continue;
        if (isdigit(tokens[i])) {
            // if token is a value, push it to va stack
            // add the code for multi-digits value(operand)
        else if (tokens[i] == ')') { // closing brace encountered
              // compute it and push the result to the value stack.
        else { // current token is an operator;
                 // push it to the operator stack.
    } // Parsing is over
   while (!op stack.empty()) {
         // apply remaining op stack to remaining va stack.
    return // va stack top contains the result, return it.
```

Dijkstra's two-stack algorithm – version.1 (fully parenthesized)

```
int evaluate(string tokens) {
    stack<int> va_stack;
                                         // stack to store operands or values
    stack<char> op stack;
                                          // stack to store operators.
    for (int i = 0; i < tokens.length(); i++) {
        if (tokens[i] == ' ') continue; // skip if token is whitespace or (
        if (tokens[i] == '(') continue;
        if (isdigit(tokens[i])) {
            // if token is a value, push it to va stack
            // add the code for multi-digits value(operand)
        else if (tokens[i] == ')') { // closing brace encountered
              // compute it and push the result to the value stack.
        else { // current token is an operator;
                 // push it to the operator stac int compute(stack<int>& va_stack, stack<char>& op_stack) {
                                                   int right = va_stack.top(); va_stack.pop();
    } // Parsing is over
                                                   int left = va stack.top(); va stack.pop();
                                                   char op = op_stack.top(); op_stack.pop();
   while (!op stack.empty()) {
                                                   int answer = apply_op(left, right, op);
         // apply remaining op stack to remainir
                                                   return answer;
    return // va stack top contains the result, return it.
```

Dijkstra's two-stack algorithm – version.2 (parenthesis are not required)

While there are still tokens to be read in, Get the next token. 1.2 If the token is: 1.2.1 A space: ignore it 1.2.2 A left brace: push it onto the operator stack. 1.2.3 A number: 1.2.3.1 read the number (it could be a multiple digit.) 1.2.3.2 push it onto the value stack 1.2.4 A right parenthesis: 1.2.4.1 While the item on top of the operator stack is not a left brace, Pop the operator from the operator stack. 1.2.4.1.2 Pop the value stack twice, getting two operands. 1.2.4.1.3 Apply the operator to the operands, in the correct order. 1.2.4.1.4 Push the result onto the value stack. 1.2.4.2 Pop the left brace from the operator stack and discard it. 1.2.5 An operator (let's call it thisOp) 1.2.5.1 While the operator stack is not empty, and the top item on the operator stack has the same or greater precedence as thisOp, Pop the operator from the operator stack 1.2.5.1.1 1.2.5.1.2 Pop the value stack twice, getting two values Apply the operator to two values in the correct order 1.2.5.1.3 Push the result on the value stack 1.2.5.1.4 1.2.5.2 Push the operator (thisOp) onto the operator stack (The whole expression has been parsed at this point. Apply remaining operators in the op stack to remaining values in the value stack) While the operator stack is not empty, Pop the operator from the operator stack. Pop the value stack twice, getting two values. Apply the operator to two values, in the correct order. 2.4 Push the result onto the value stack. (At this point the operator stack should be empty, and the value stack should have only one value in it, which is the result.)

Return the top item in the value stack.

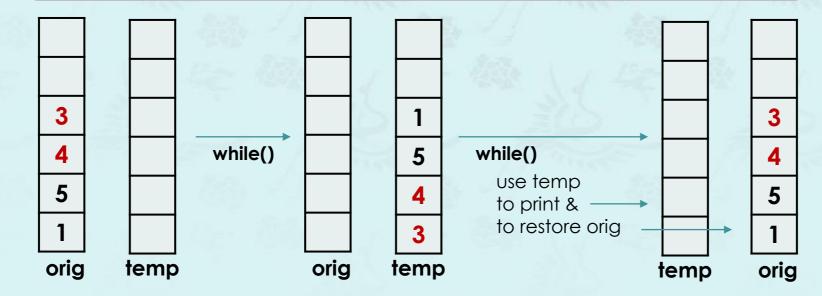
Dijkstra's two-stack algorithm – Implementation (iteration & template version)

- printStack() prints the contents of a stack from the bottom to top.
 - The stack contents should be the same as before after printing.

Algorithm:

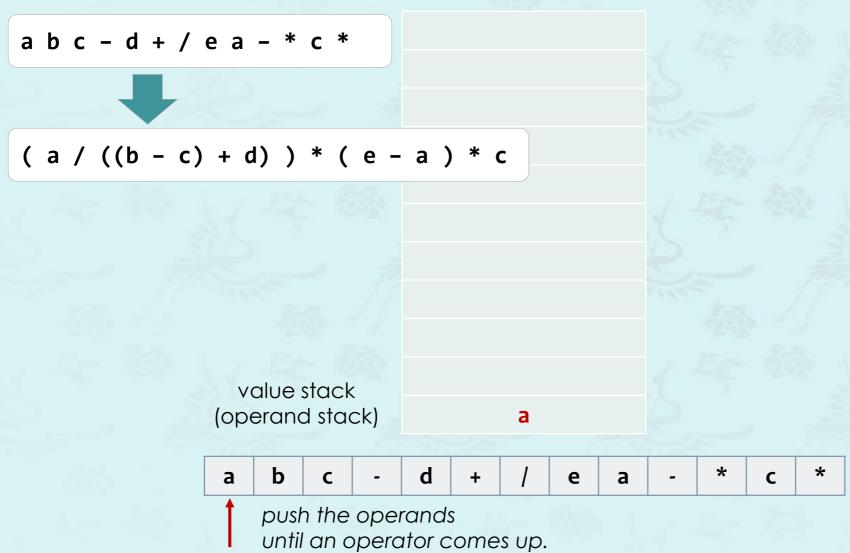
- Given a stack called orig.
- Create an empty stack called temp.
- While orig is not empty,
 - Top/Pop push an item from orig to temp.
- While temp is not empty,
 - Top/Pop an item from temp, print it and push it orig.

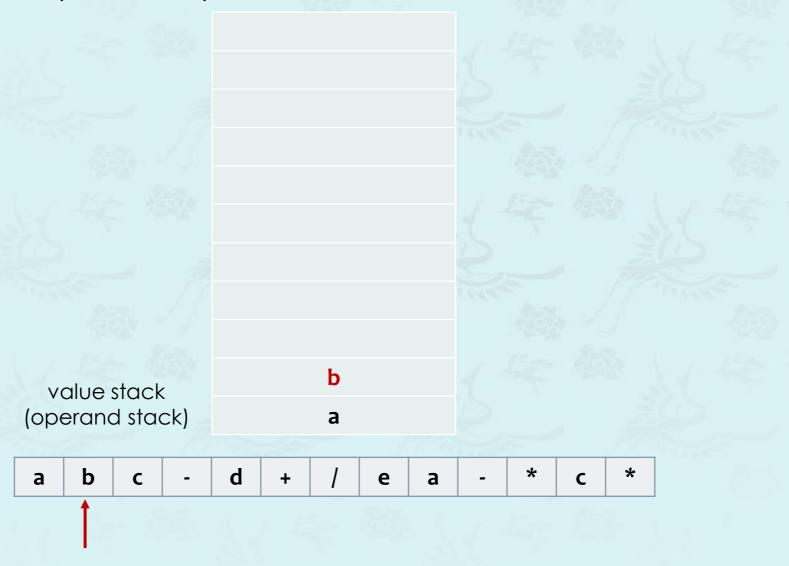
- Use both iteration and recursion version of stack in PSet.
- Use stack<double>
 to handle division
 operation in PSet.

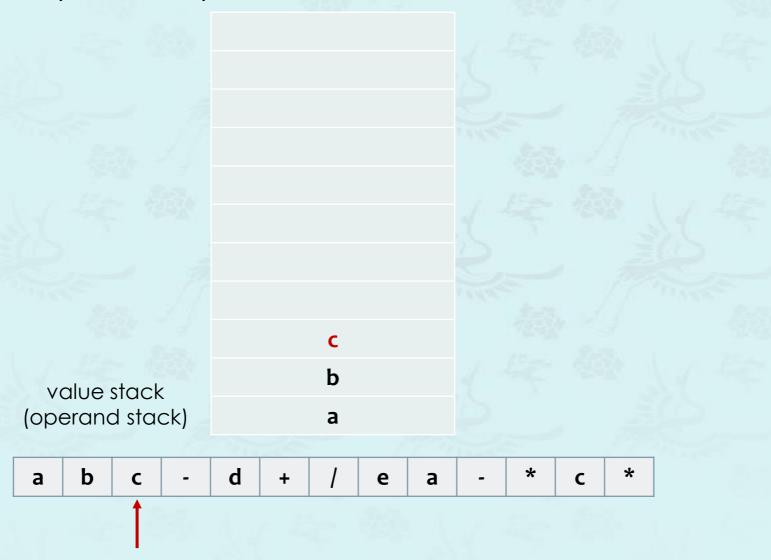


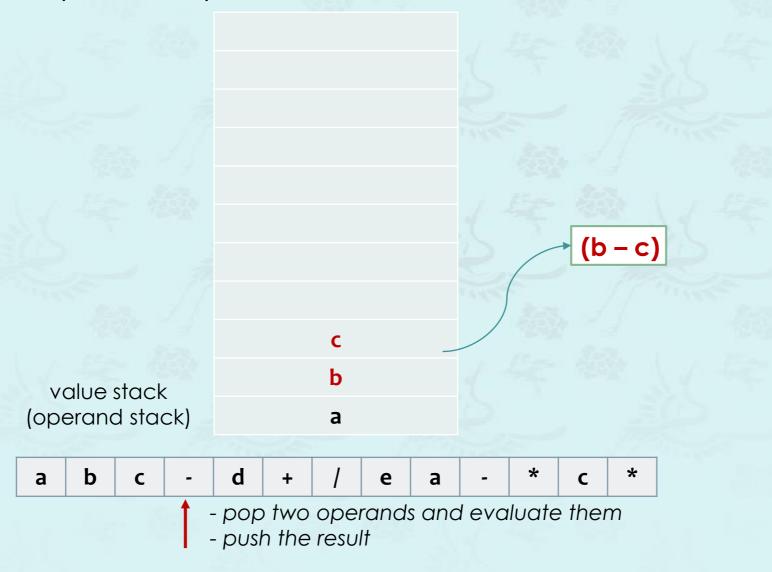
infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
(1 + 2) * 7	1 2 + 7 *
a * b / c	a b * c /
(a/(b-c+d))*(e-a)*c	a b c - d + / e a - * c *
a / b - c + d * e - a * c	a b / c - d e * + a c * -

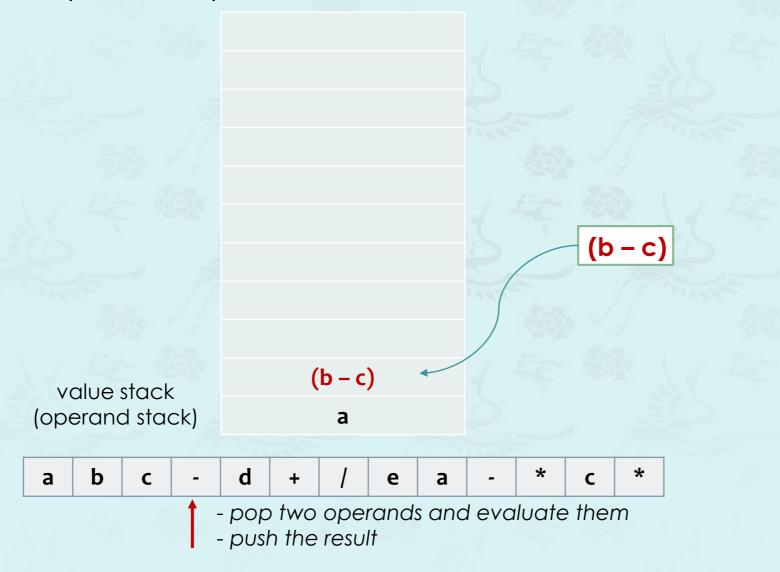


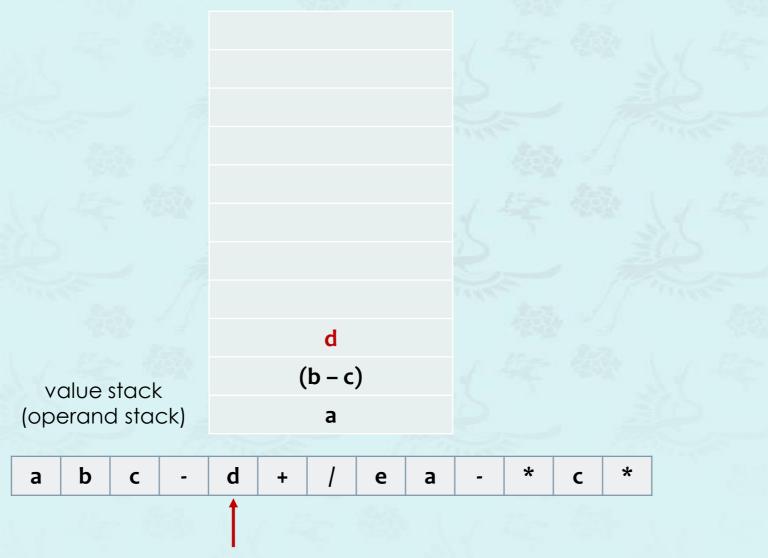


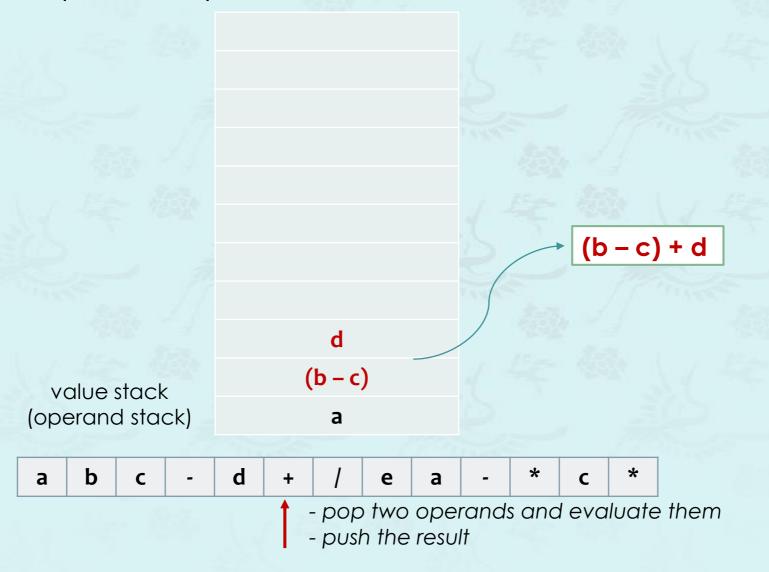


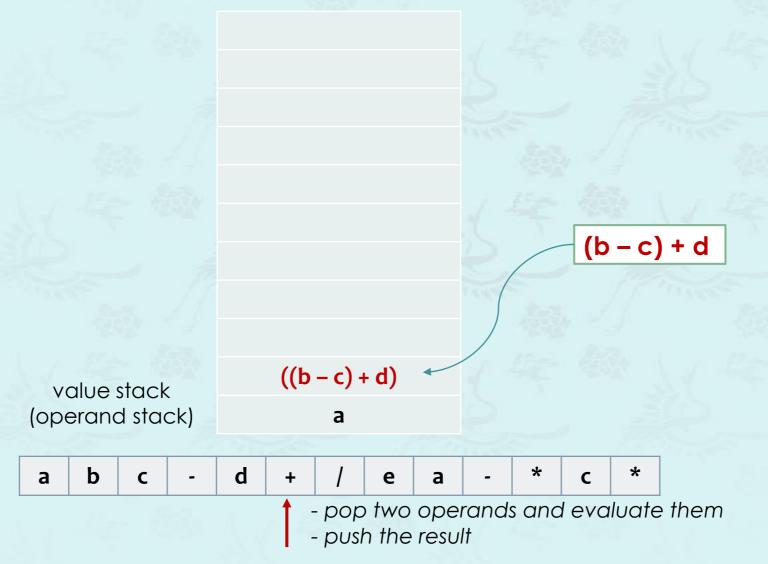


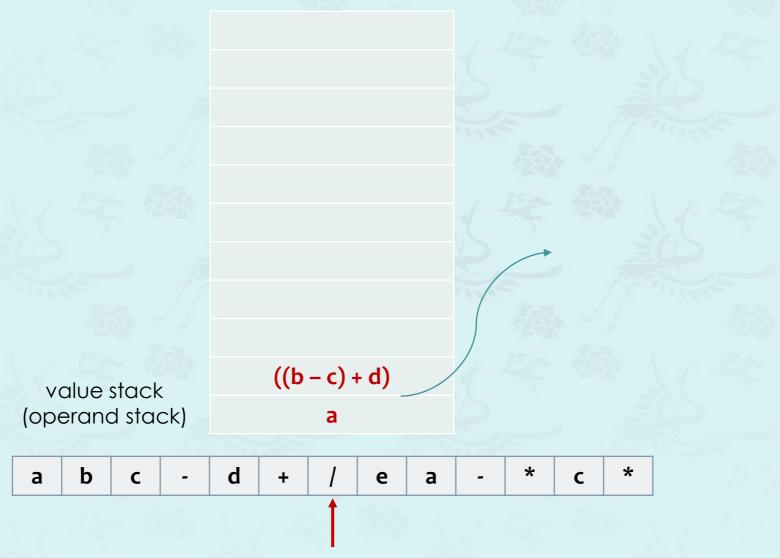


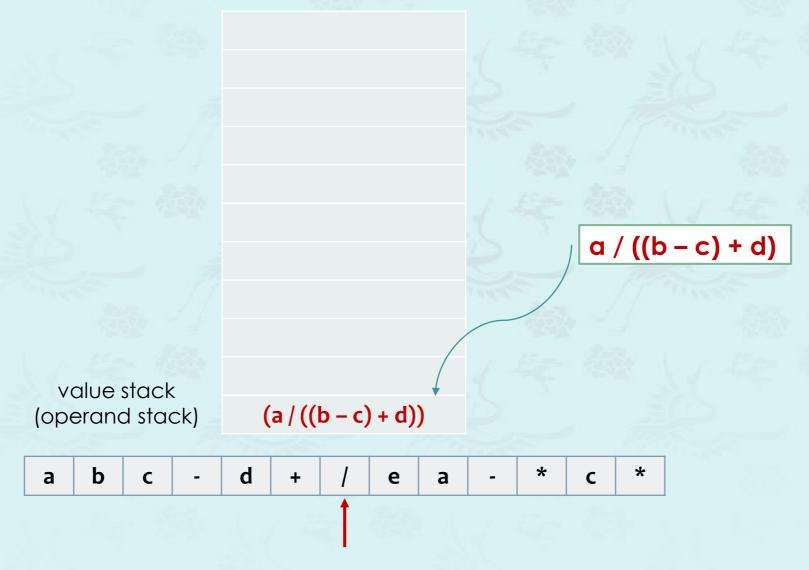


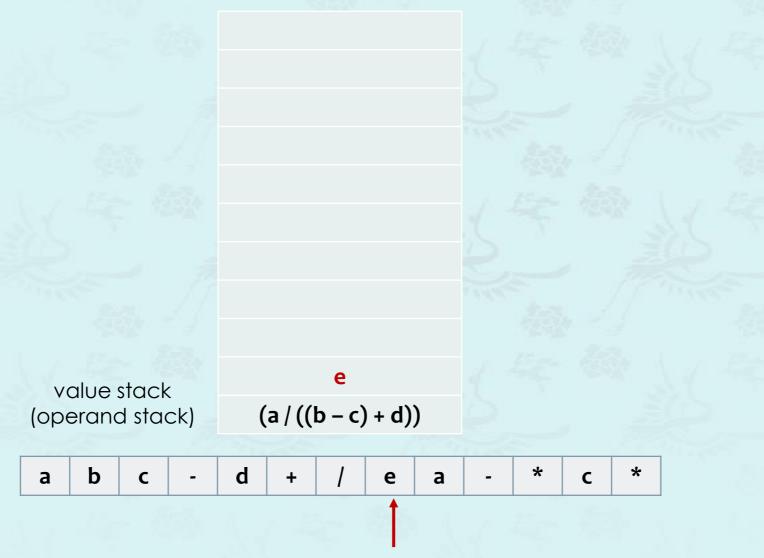


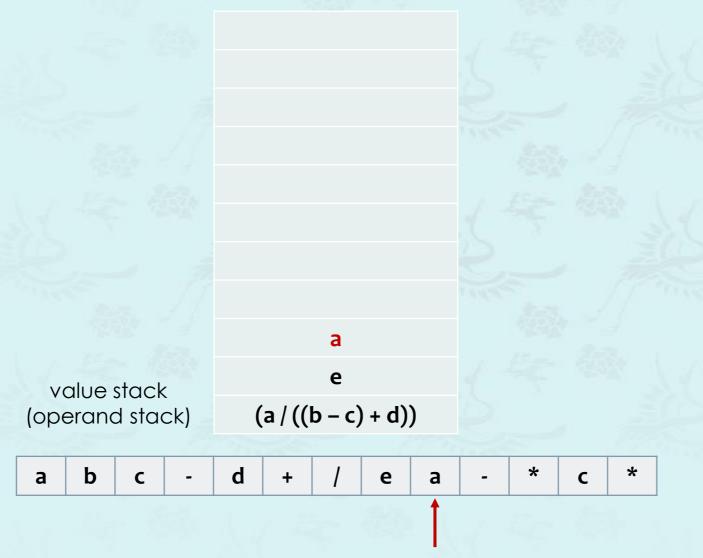


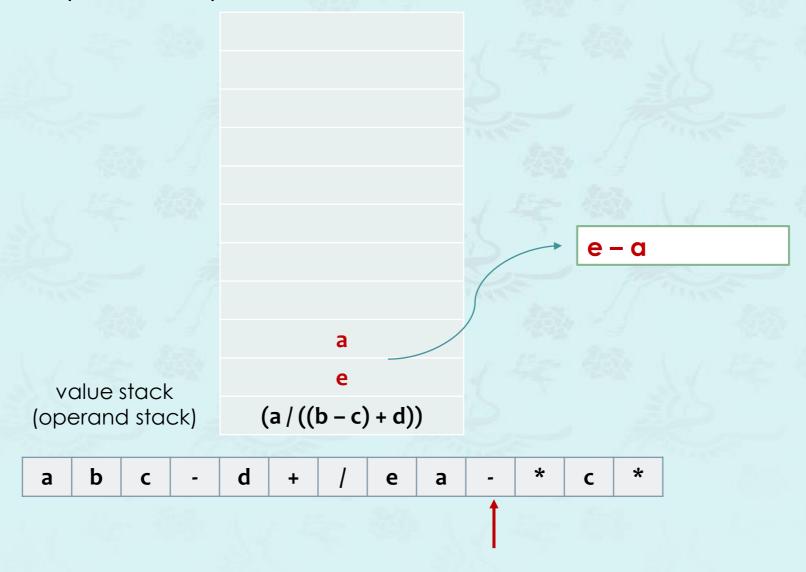


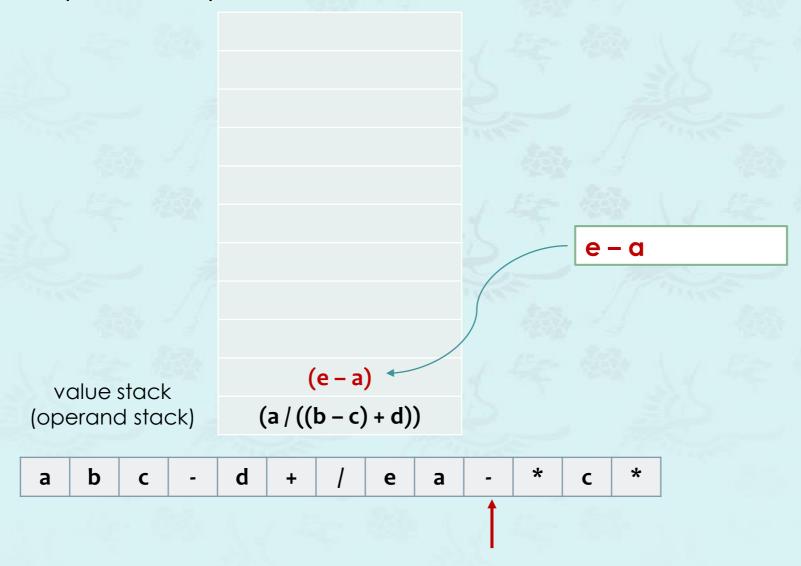


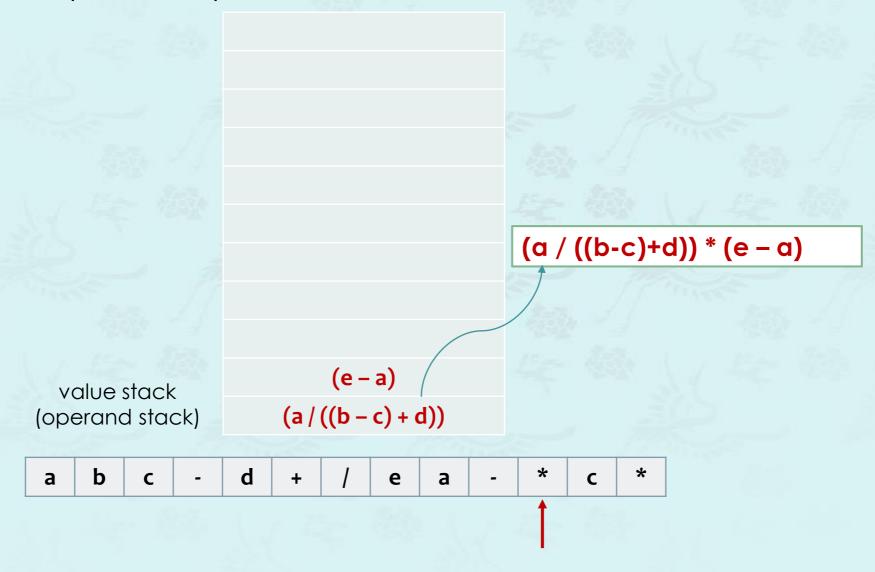


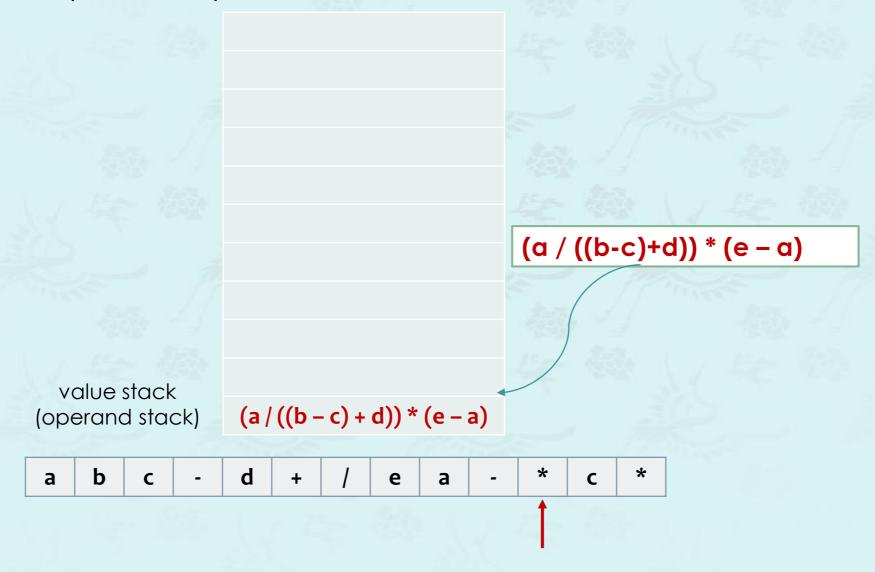


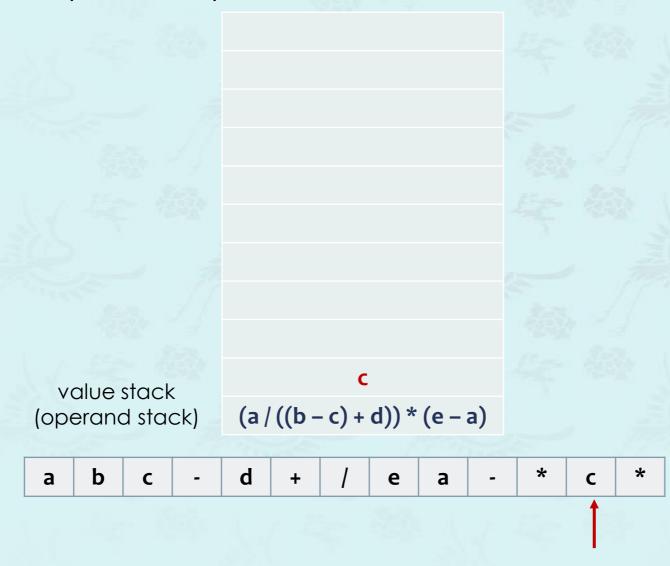


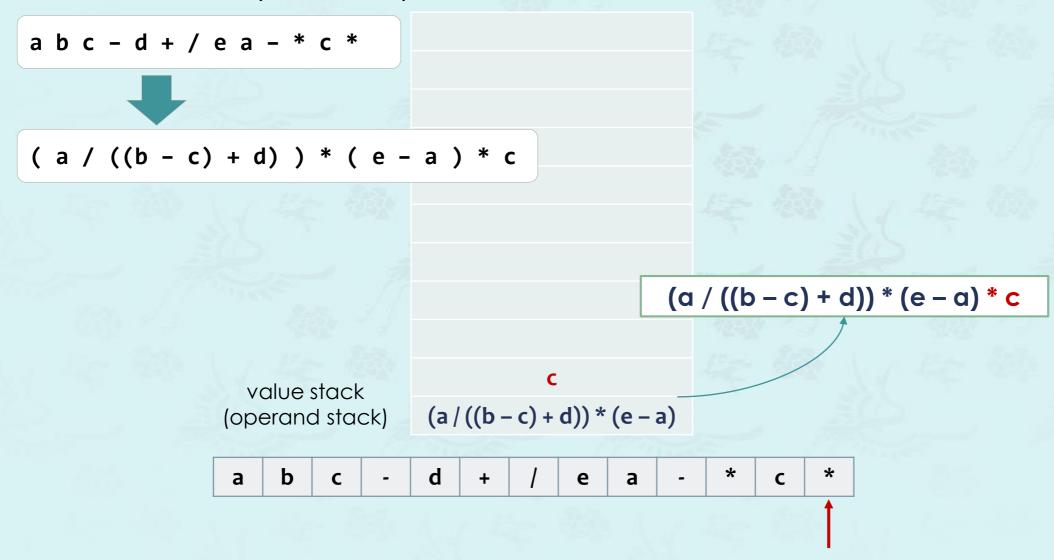


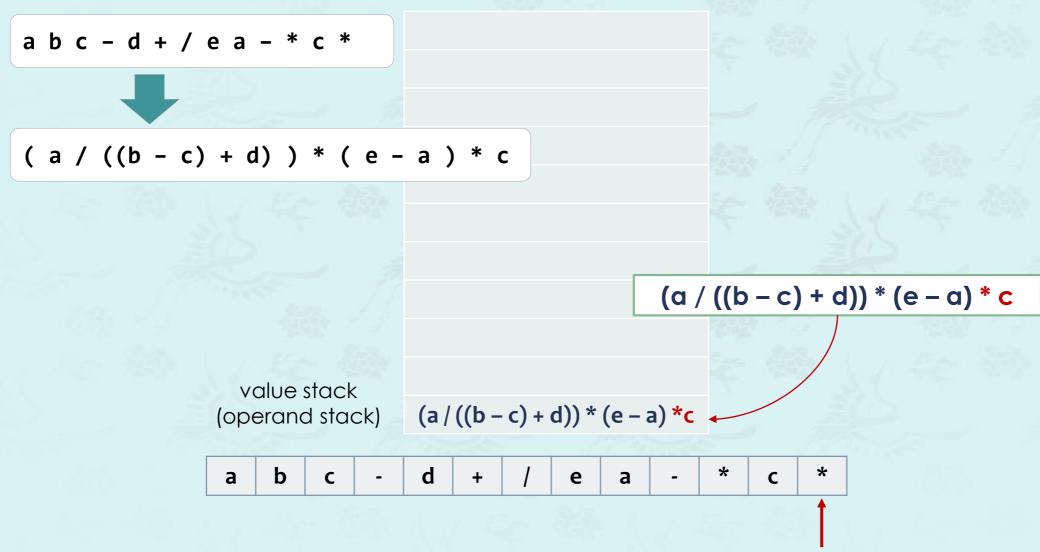


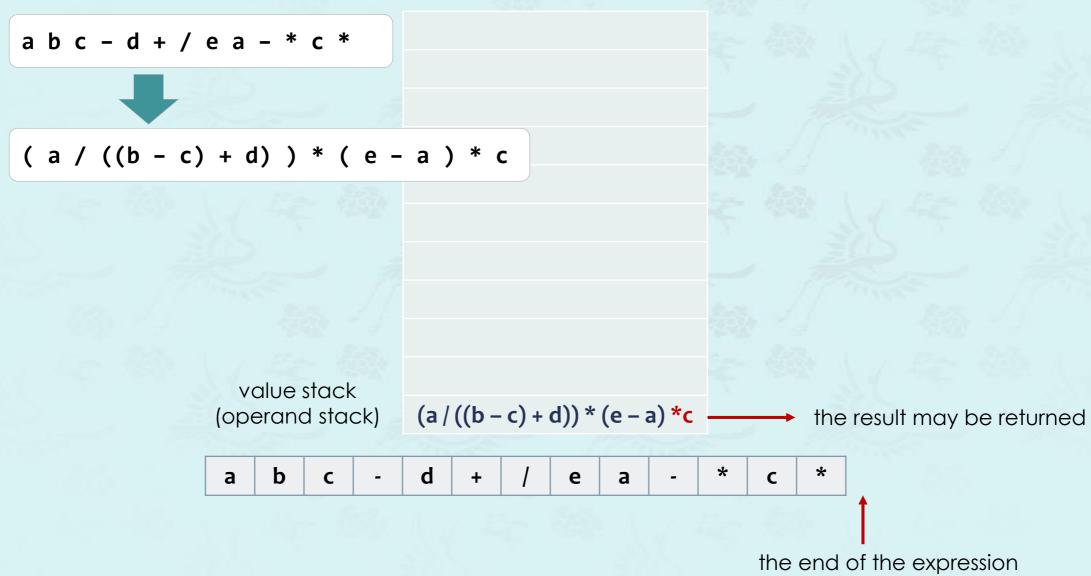






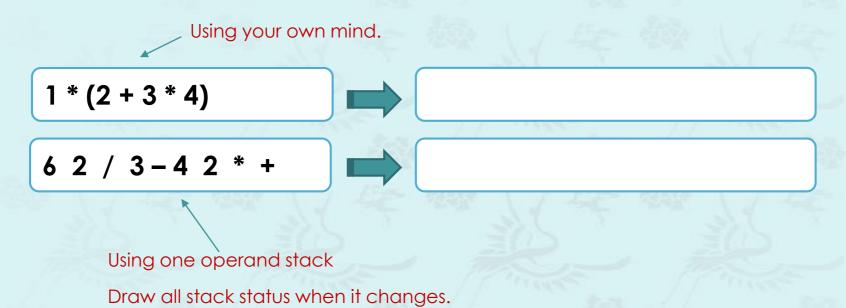






Exercise:

Infix to postfix, postfix to infix Conversion



Summary & quaestio quaesti quaest

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