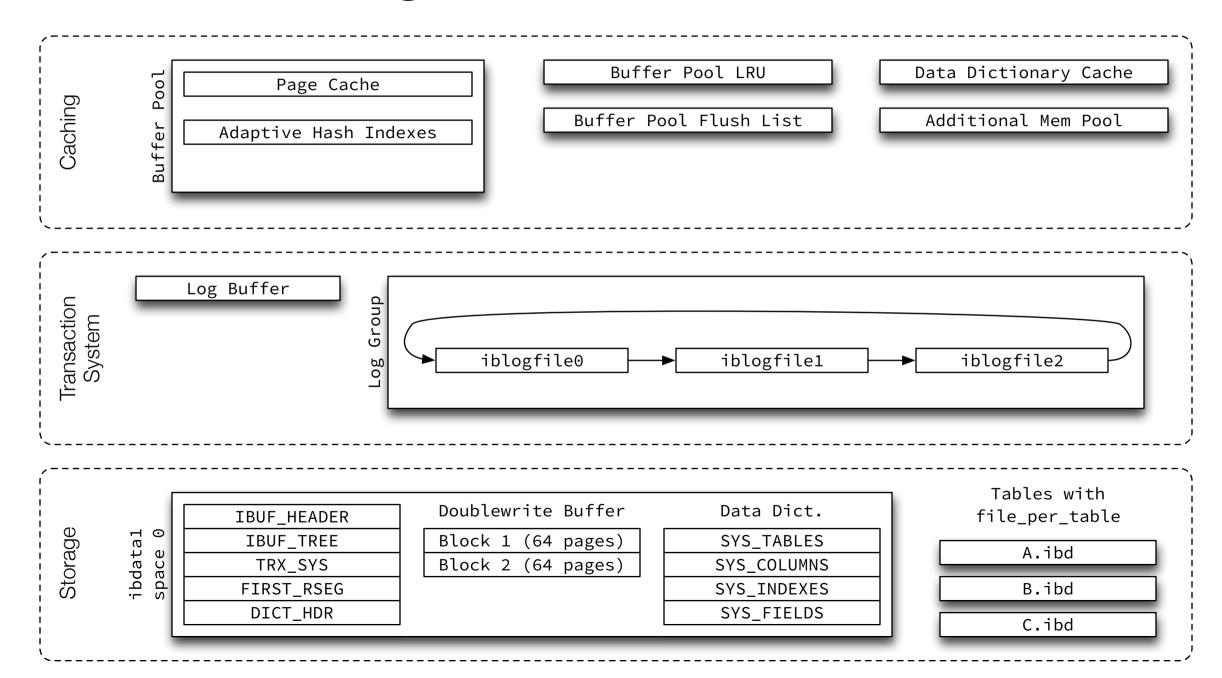
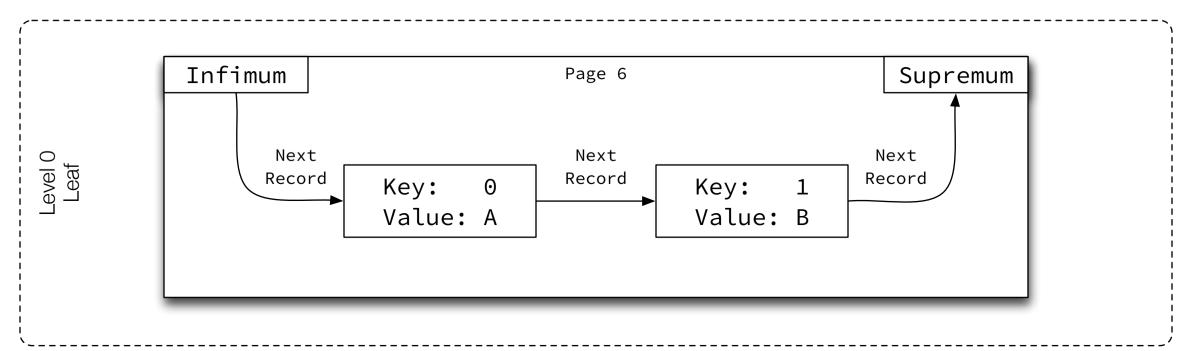
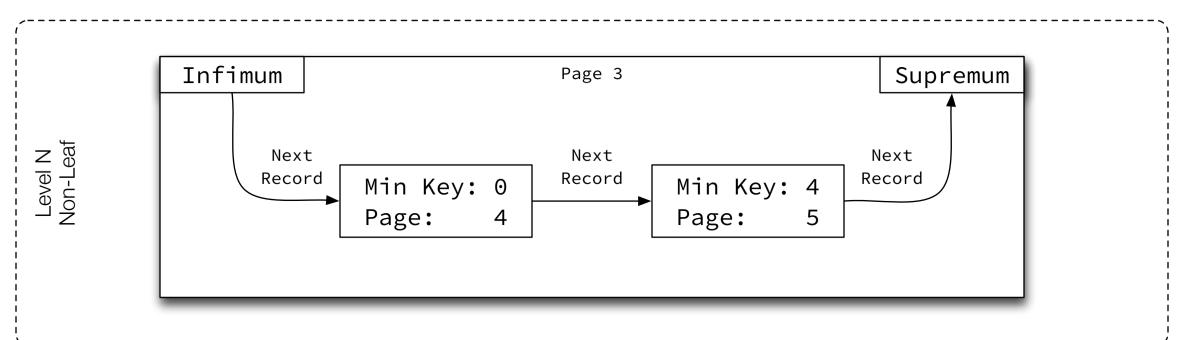
## High-level Overview



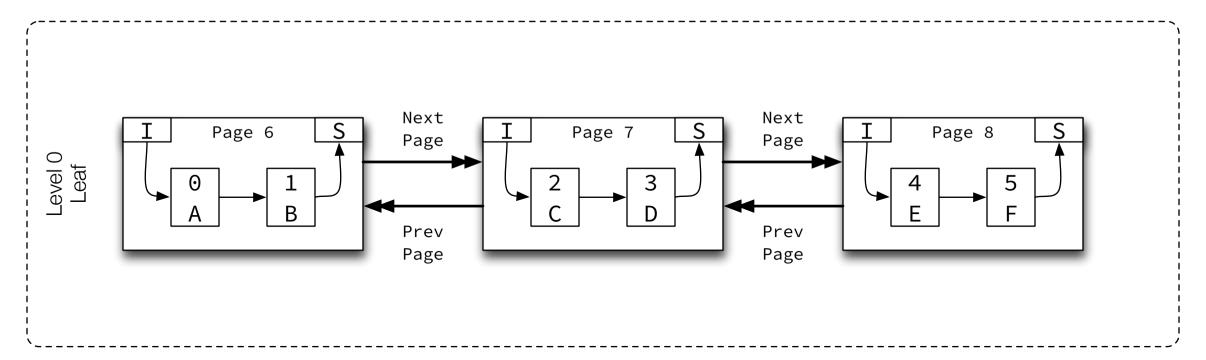
# B+Tree Simplified Leaf Page



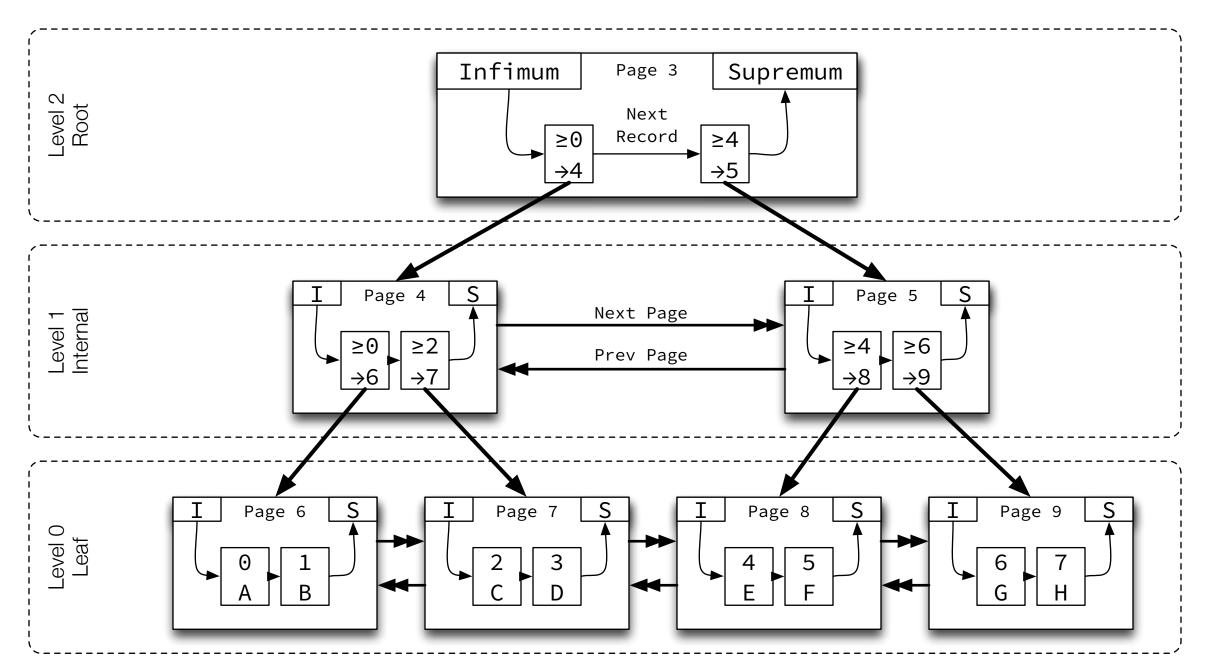
# B+Tree Simplified Non-Leaf Page



# B+Tree Simplified Level



### B+Tree Structure



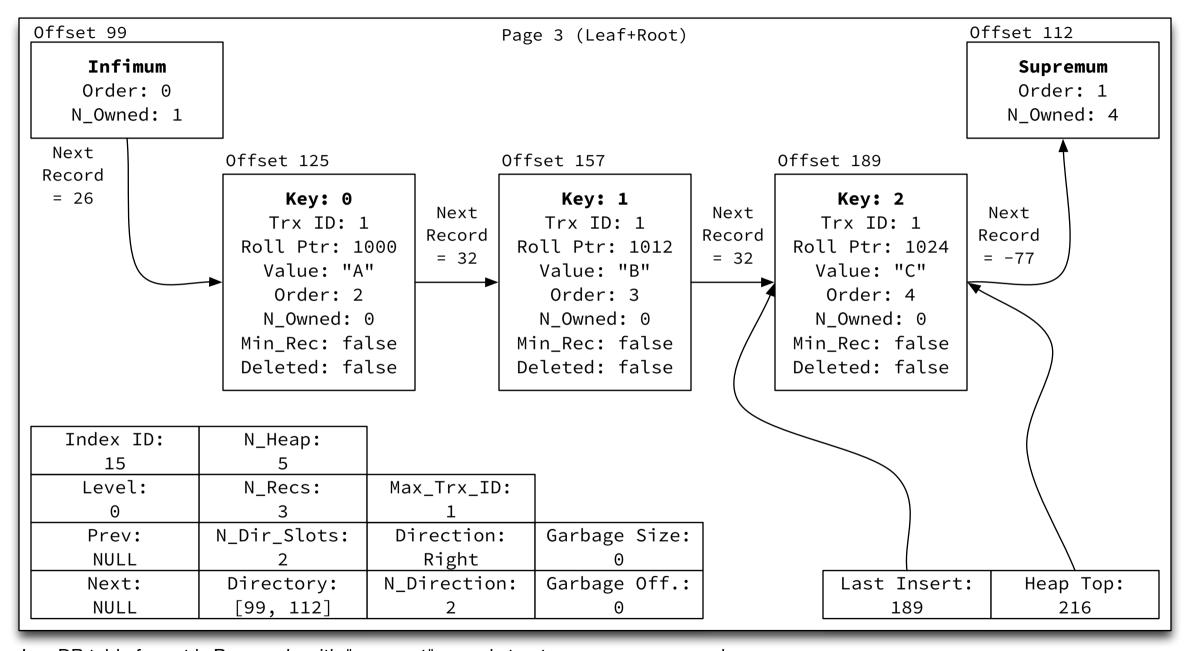
Levels are numbered starting from 0 at the leaf pages, incrementing up the tree.

Pages on each level are doubly-linked with previous and next pointers in ascending order by key.

Records within a page are singly-linked with a next pointer in ascending order by key.

Infimum represents a value lower than any key on the page, and is always the first record in the singly-linked list of records. Supremum represents a value higher than any key on the page, and is always the last record in the singly-linked list of records. Non-leaf pages contain the minimum key of the child page and the child page number, called a "node pointer".

# B+Tree Detailed Page Structure



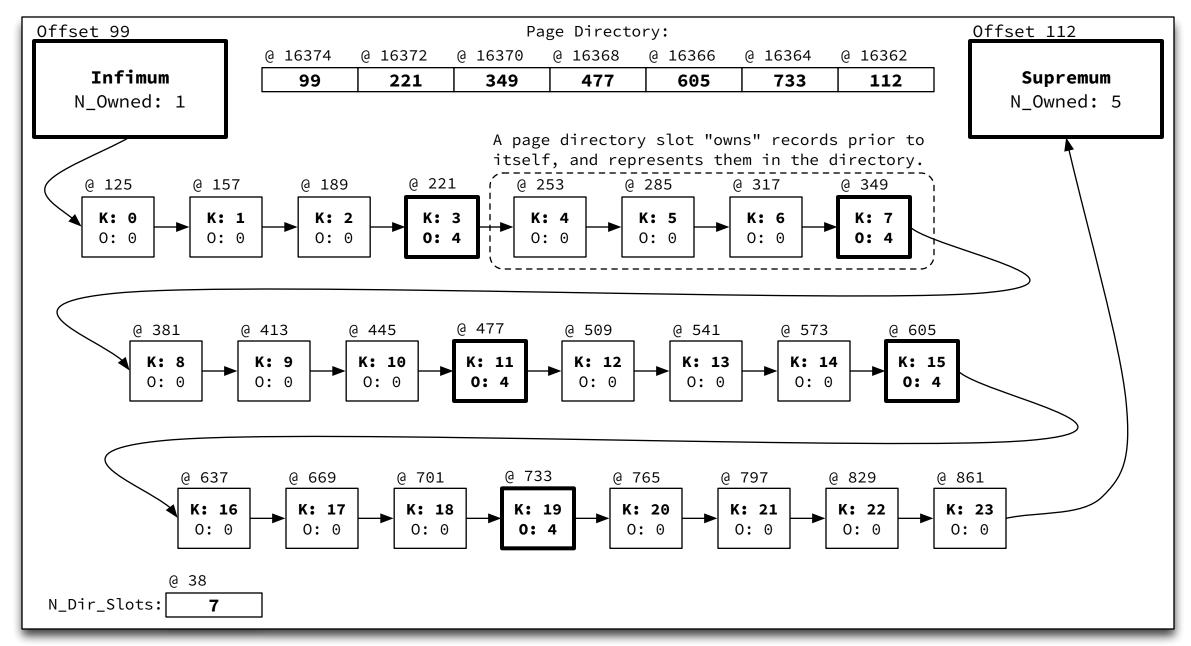
InnoDB table format is Barracuda with "compact" record structure, non-compressed.

Table created with: CREATE TABLE t (i INT NOT NULL, s CHAR(10) NOT NULL, PRIMARY KEY(i)) ENGINE=InnoDB;

Table populated with: INSERT INTO t (i, s) VALUES (0, "A"), (1, "B"), (2, "C");

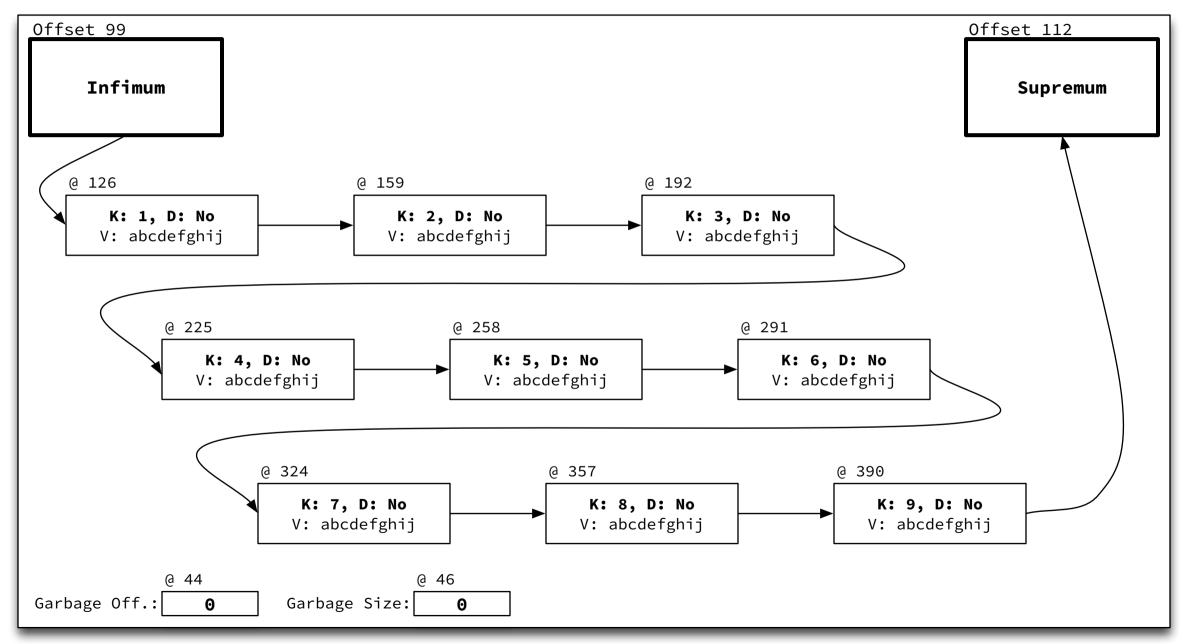
Record size: 5 (header) + 4 (PK) + 6 (TRX\_ID) + 7 (ROLL\_PTR) + 10 (non-key fields) = 32 bytes

# B+Tree Page Directory Structure



Infimum always owns only itself, so will always have a slot in the page directory with N\_Owned = 1. Supremum always owns the last few records in the page, and is allowed to own less than 4 records (if the page has fewer). All directory slots will own a minimum of 4 and maximum of 8 records, except supremum, which may own fewer. The page directory grows "downwards" from offset 16376, the beginning of the FIL trailer; the first directory entry starts at 16374.

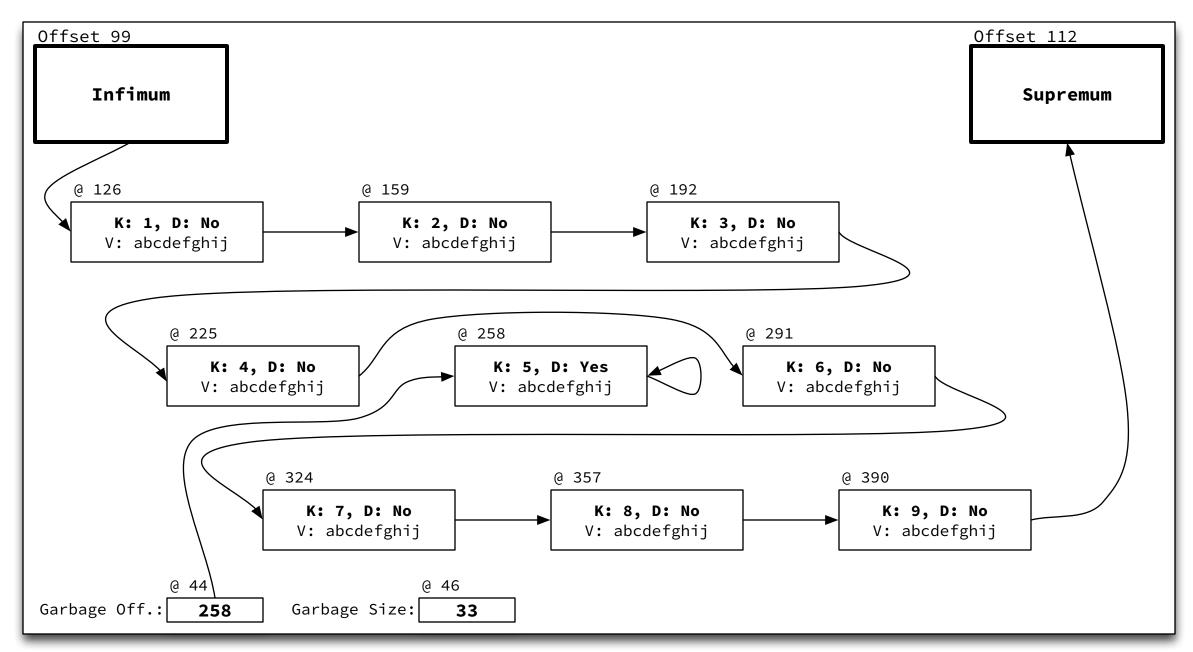
### B+Tree Record Initial State



SQL: create table t (i int not null, s varchar(100) not null, primary key(i)) engine=innodb;

SQL: insert into t (i, s) values (1, "abcdefghij"); for i in 1..9

### B+Tree Record Delete 1



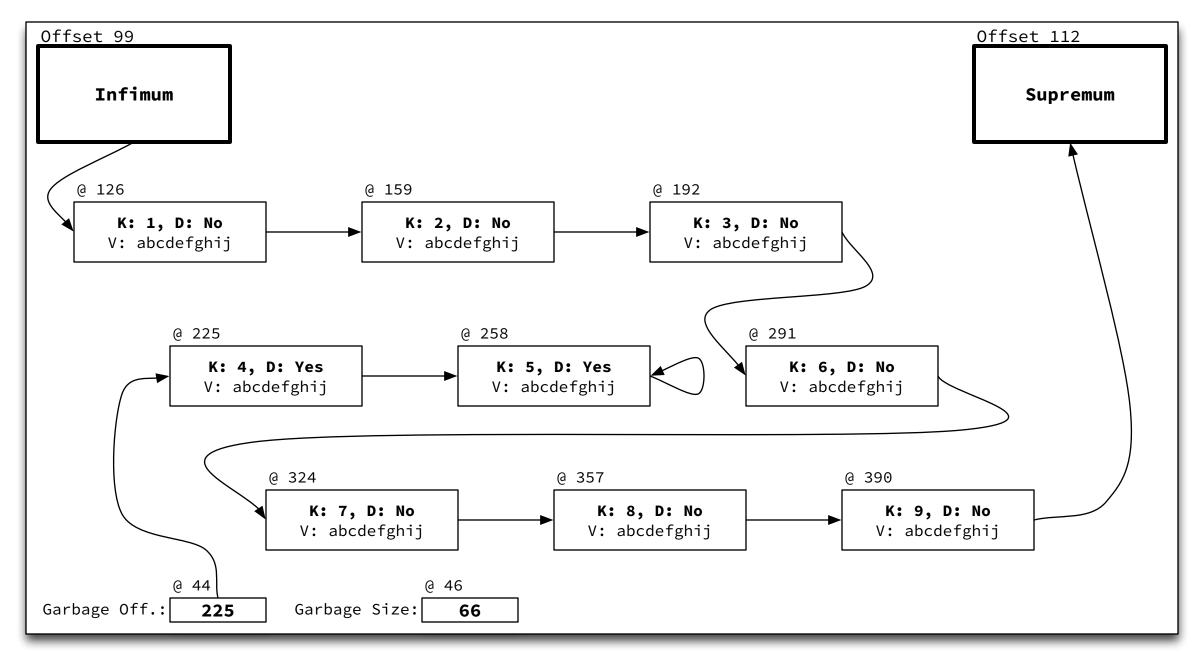
SQL: delete from t where i = 5;

Row is marked as deleted.

Garbage size is incremented by total row size.

Garbage offset is pointed to row, and row next pointer is pointed back to self.

## B+Tree Record Delete 2



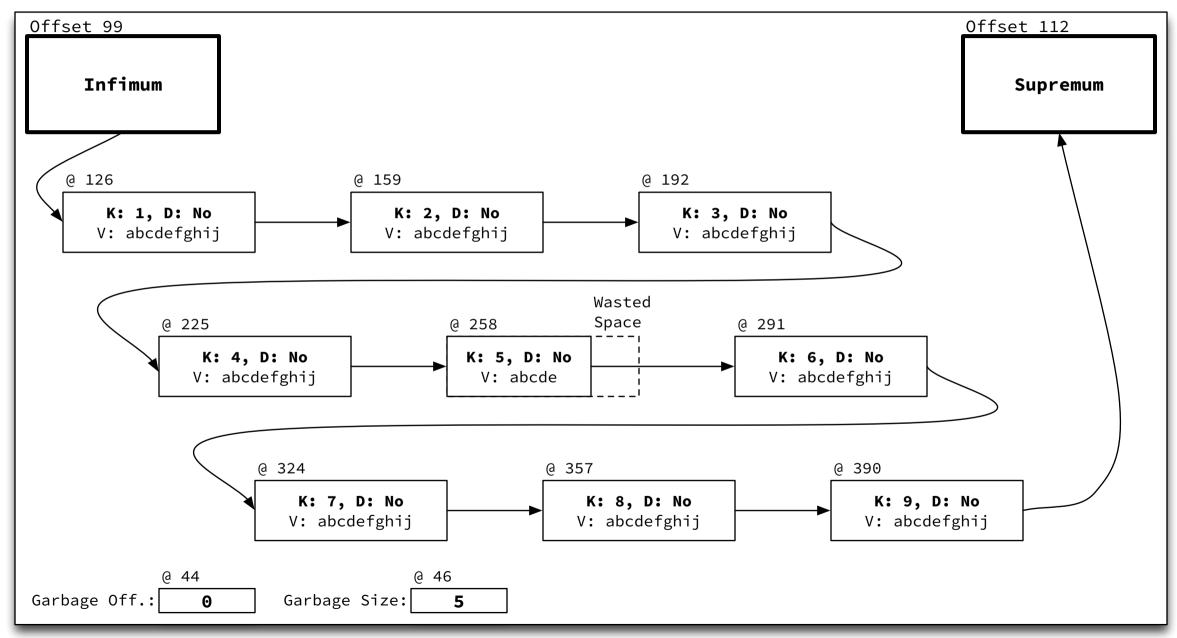
SQL: delete from t where i = 5; delete from t where i = 4;

Garbage size is incremented by total row size for each delete.

Garbage offset is pointed to row @ 258 initially, and row next pointer is pointed back to self.

Garbage offset is updated to row @ 225, and row next pointer is pointed to previous garbage offset (garbage is added to head of list).

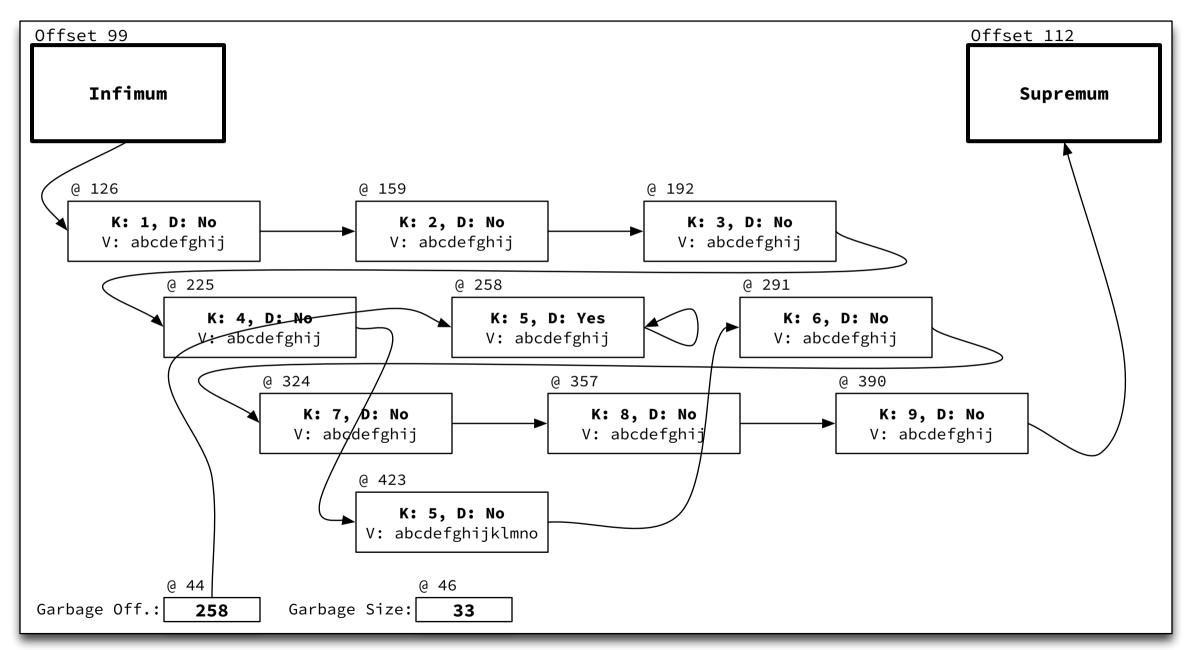
## B+Tree Record Update - Smaller



SQL: update t set s="abcde" where i = 5;

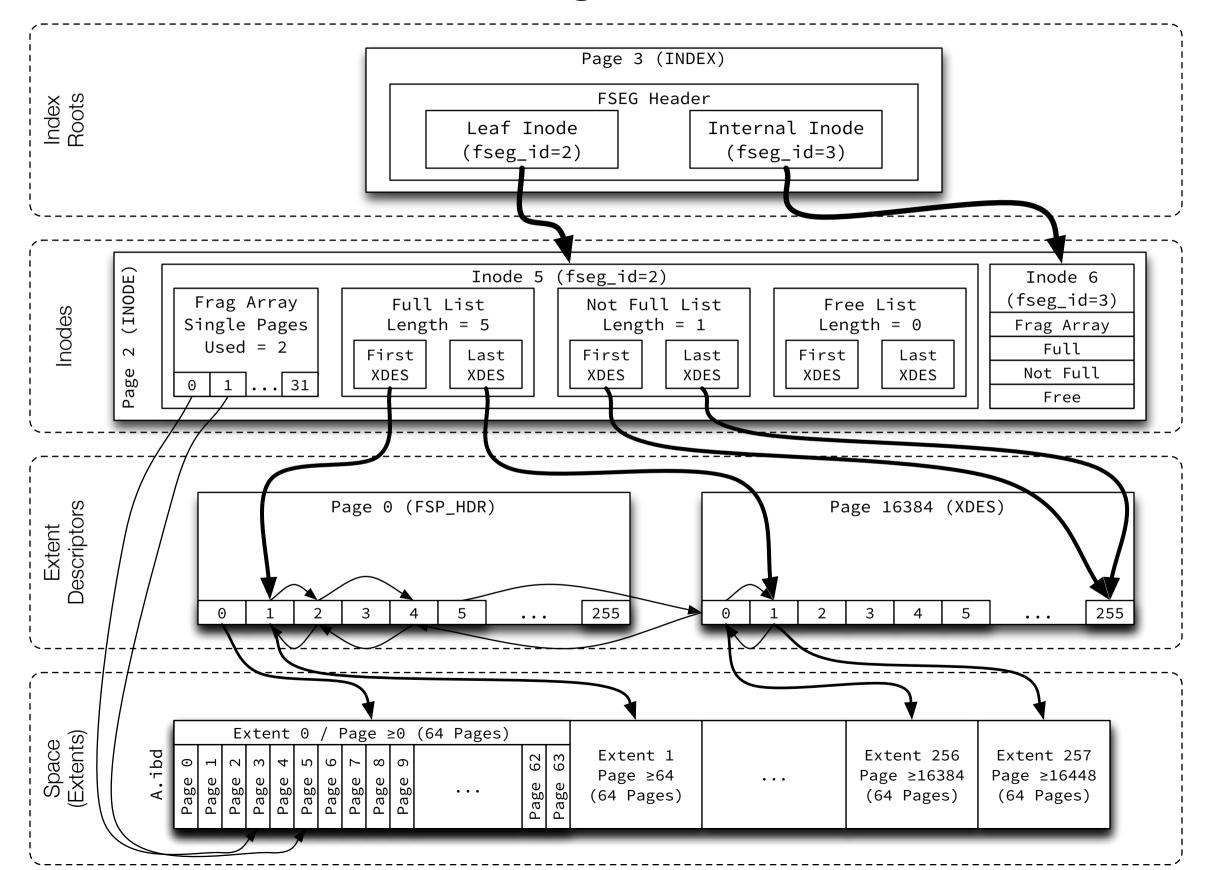
Garbage size is incremented by size of row shrinkage, but wasted space is not tracked in garbage list.

# B+Tree Record Update - Larger

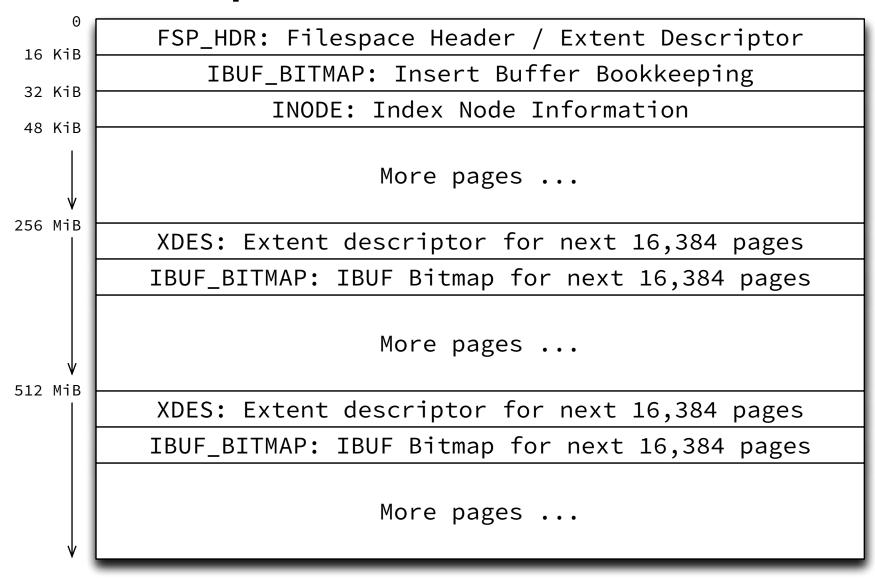


SQL: update t set s="abcdefghijklmno" where i = 5; Row is deleted, and a new row is inserted into the heap.

# Index File Segment Structure



## Space File Overview



## ibdata1 File Overview

0 16 KiB	FSP_HDR: Filespace Header / Extent Descriptor
	IBUF_BITMAP: Insert Buffer Bookkeeping
32 KiB -	INODE: Index Node Information
48 KiB	SYS: Insert Buffer Header
64 KiB  - 	INDEX: Insert Buffer Root
	TRX_SYS: Transaction System Header
	SYS: First Rollback Segment
	SYS: Data Dictionary Header
	More pages
Page 64	Double Write Buffer Block 1 (64 pages)
Page 128	Double Write Buffer Block 2 (64 pages)
Page 192	More pages

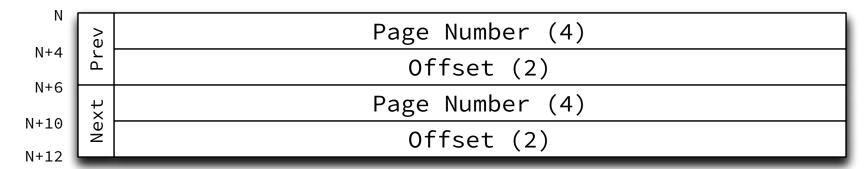
### **IBD File Overview**

```
FSP_HDR: Filespace Header / Extent Descriptor
16 KiB
             IBUF BITMAP: Insert Buffer Bookkeeping
32 KiB
                  INODE: Index Node Information
48 KiB
                 INDEX: Root page of first index
64 KiB
                INDEX: Root page of second index
                      INDEX: Node pages ...
                      INDEX: Leaf pages ...
            ALLOCATED: Reserved but unused pages ...
                          More pages ...
```

## List Base Node

N.I.		
N		List Length (4)
N+4	st	Page Number (4)
N+8	Fir	Offset (2)
N+10 N+14	st	Page Number (4)
N+14 N+16	La	Offset (2)
INT IO		

## List Node



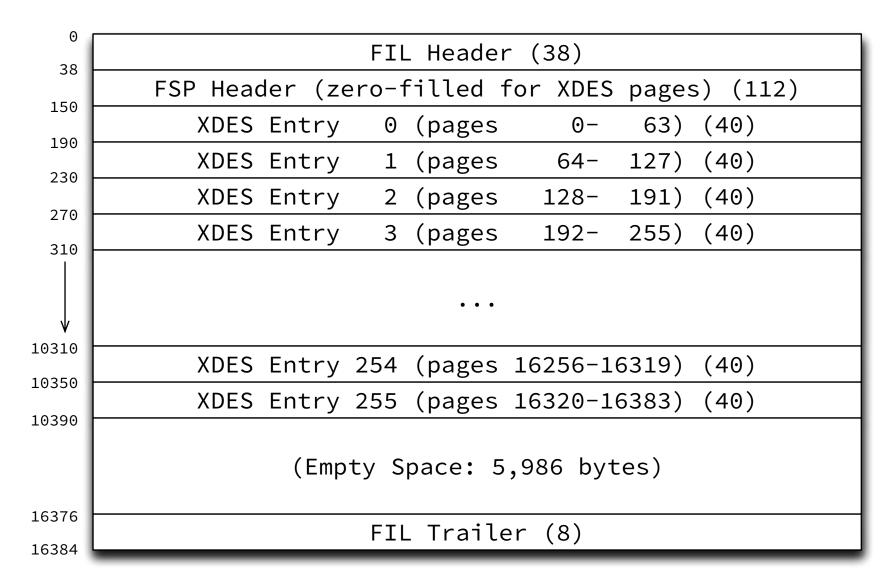
## Basic Page Overview

0 38	FIL Header (38)
	Other headers and page data, depending on page type.
	Total usable space: 16,338 bytes.
16376 16384	FIL Trailer (8)

## FIL Header/Trailer

_	
0	Checksum (4)
4	Offset (Page Number) (4)
8	Previous Page (4)
12	Next Page (4)
16	LSN for last page modification (8)
24	Page Type (2)
26	Flush LSN (0 except space 0 page 0) (8)
34	Space ID (4)
38	
	• • •
16376	
	Old-style Checksum (4)
16380	Low 32 bits of LSN (4)
16384	

#### FSP\_HDR/XDES Overview



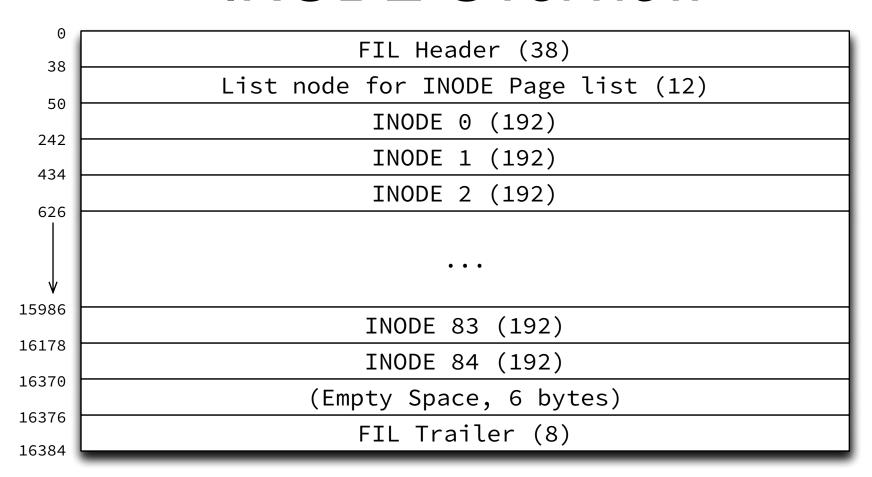
## **FSP Header**

~ ~	
38	Space ID (4)
42	(Unused) (4)
46	Highest page number in file (size) (4)
50	Highest page number initialized (free limit) (4)
54	Flags (4)
58	Number of pages used in "FREE_FRAG" list (4)
62	List base node for "FREE" list (16)
78	List base node for "FREE_FRAG" list (16)
94	List base node for "FULL_FRAG" list (16)
110	Next Unused Segment ID (8)
118	List base node for "FULL_INODES" list (16)
134 150	List base node for "FREE_INODES" list (16)
TO0	

## **XDES Entry**

N.I.	
N	File Segment ID (8)
N+8	List node for XDES list (12)
N+20	State (4)
N+24	Page State Bitmap (16)
N+40	2 bits per page, 1=free, 2=clean

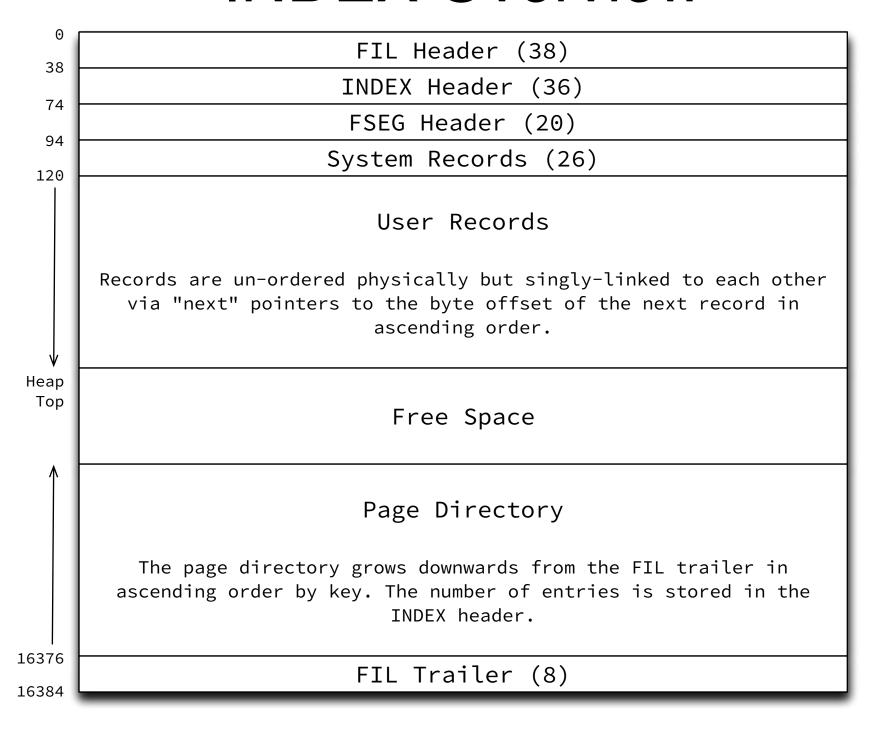
#### **INODE Overview**



# **INODE** Entry

N	FSEG ID (8)
N+8	Number of used pages in "NOT_FULL" list (4)
N+12	List base node for "FREE" list (16)
N+28	List base node for "NOT_FULL" list (16)
N+44	List base node for "FULL" list (16)
N+60	Magic Number = 97937874 (4)
N+64	Fragment Array Entry 0 (4)
N+68	
N+188 N+192	Fragment Array Entry 31 (4)
14 · ±52	

### **INDEX Overview**



## **INDEX Header**

38	
30	Number of Directory Slots (2)
40	Heap Top Position (2)
42	Heap Top Position (2)
	Number of Heap Records / Format Flag (2)
44	First Garbage Record Offset (2)
46	
	Garbage Space (2)
48	Last Incort Dosition (2)
	Last Insert Position (2)
50	Page Direction (2)
52	Number of Inserts in Page Direction (2)
54	Number of Inserts in rage birection (2)
54	Number of Records (2)
56	` ,
C 4	Maximum Transaction ID (8)
64	Page Level (2)
66	rage Level (2)
	Index ID (4)
74	· · · /

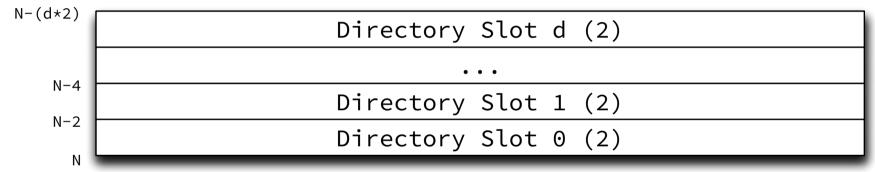
## FSEG Header

7 4	
74 78	Leaf Pages Inode Space ID (4)
82	Leaf Pages Inode Page Number (4)
84	Leaf Pages Inode Offset (2)
	Internal (non-leaf) Inode Space ID (4)
88	Internal (non-leaf) Inode Page Number (4)
92	Internal (non-leaf) Inode Offset (2)
94	, , ,

## **INDEX System Records**

0.4	
94	Info Flags (4 bits)
0.5	Number of Records Owned (4 bits)
95	Order (13 bits)
	Record Type (3 bits)
97	Next Record Offset (2)
99	"infimum\0" (8)
107	Info Flags (4 bits)
108	Number of Records Owned (4 bits)
100	Order (13 bits)
110	Record Type (3 bits)
112	Next Record Offset (2)
120	"supremum" (8)

## **INDEX Page Directory**



#### Record Format - Header

Record Type (3 bits)

Next Record Offset (2)

- 1	Variable field lengths (1-2 bytes per var. field)
N-5	Nullable field bitmap (1 bit per nullable field)
	Info Flags (4 bits)
N-4	Number of Records Owned (4 bits)
	Order (13 bits)

N-2

#### Record Format - Clustered Key - Leaf Pages

N =	Variable field lengths (1-2 bytes per var. field)
N-5	Info Flags (4 bits)
N-4	Number of Records Owned (4 bits)
	Order (13 bits)
N-2	Record Type (3 bits)
N	Next Record Offset (2)
N+k	Cluster Key Fields (k)
N+k+6	Transaction ID (6)
N+k+13	Roll Pointer (7)
l+k+13+j	Non-Key Fields (j)

#### Record Format - Clustered Key - Non-Leaf Pages

N E	Variable field lengths (1-2 bytes per var. field)
N-5	Info Flags (4 bits)
N-4	Number of Records Owned (4 bits)
	Order (13 bits)
N-2	Record Type (3 bits)
N	Next Record Offset (2)
N+k	Cluster Key Min. Key on Child Page (k)
N+k+4	Child Page Number (4)

#### Record Format - Secondary Key - Leaf Pages

N-5

N-4

N-2

N+k

N+k+j

Ν

/ariable field lengths (1-2 bytes per var. field)
Nullable field bitmap (1 bit per nullable field)
Info Flags (4 bits)
Number of Records Owned (4 bits)
Order (13 bits)
Record Type (3 bits)
Next Record Offset (2)
Secondary Key Fields (k)
Cluster Key Fields (j)

#### Record Format - Secondary Key - Non-Leaf Pages

N-5

N-4

N-2

N+k

N+k+j

N+k+j+4

Ν

Variable field lengths (1-2 bytes per var. field)
Nullable field bitmap (1 bit per nullable field)
Info Flags (4 bits)
Number of Records Owned (4 bits)
Order (13 bits)
Record Type (3 bits)
Next Record Offset (2)
Secondary Key Min. Key on Child Page (k)
Cluster Key Fields (j)
Child Page Number (4)