

Loosely speaking:

- **Syntax** means “form”.
- **Semantics** means “meaning”.

In the context of programming languages:

- The Java expressions  $a - 3 - b$  and  $u - 7 - v$  have the *same syntax*, but *different semantics*.
- The Java expressions  $a - 3 - b$  and  $((a - 3) - b)$  have the *same semantics*, but *different syntax*.
- The Lisp expression  $(- (- 2 3) 4)$  and the Java expression  $2 - 3 - 4$  have the *same semantics* but have *different syntax*.

## Operators, Arities, and Operands

An *operator of arity  $k$* , also called a  *$k$ -ary operator*, is a symbol that represents a function of  $k$  arguments.

- *Binary* means 2-ary.
- *Unary* means 1-ary.
- *Ternary* means 3-ary.

When a  $k$ -ary operator appears in an expression, each of the  $k$  subexpressions whose values are the  $k$  arguments of the function it represents is called an *operand* of the operator.

### Examples

In a Java or C++ expression  $x - y + z$

- $+$  is a binary operator whose operands are  $x - y$  and  $z$ .
- $-$  is a binary operator whose operands are  $x$  and  $y$ .

In the Lisp expression  $(+ (- x y) 3 (* z z))$

- $+$  is a ternary operator whose operands are  $(- x y)$ ,  $3$ , and  $(* z z)$ .

## Overview of Some Expression Notations

We will say that two expressions are equivalent if they have the same semantics.

Here is a Java expression: `f(g(h(1,2), f(3,4)), 5)`

We will now write equivalent expressions in:

- Infix notation, not making use of precedence classes
- Infix notation, making use of precedence classes
- Lisp notation (also called **Cambridge Polish Notation**)
- Prefix notation
- Postfix notation (also called **Reverse Polish Notation**)

### Infix Notation

In infix notation, we write binary operators *between* their operands--e.g., `f(3,4)` would be written `3 f 4`.

So the above Java expression is equivalent to the following infix expression: `((1 h 2) g (3 f 4)) f 5`

The designer of an infix notation will usually give *precedence and associativity rules* for the operators, to reduce the need for parentheses.

Suppose the notation designer specifies that:

- f and g belong to the *same* precedence class, and that precedence class is *left-associative*.
- h belongs to a *higher* precedence class than f and g.

Now the above infix expression  $((1\ h\ 2)\ g\ (3\ f\ 4))\ f\ 5$  is equivalent to:

$1\ h\ 2\ g\ (3\ f\ 4)\ f\ 5$

## An Analogous Example

In ordinary infix arithmetic notation:

- + and - belong to the same precedence class, and that precedence class is left-associative.
- \* belongs to a higher precedence class than + and -.

Thus the infix expression is equivalent to:

$((1\ *\ 2) - (3 + 4)) + 5$   
 $1\ *\ 2 - (3 + 4) + 5$

## Lisp Notation and Prefix Notation

The above Java expression `f(g(h(1,2), f(3,4)), 5)` is equivalent to the following Lisp expression: `(f (g (h 1 2) (f 3 4)) 5)`

Prefix notation is *Lisp notation without parentheses*.

Thus the above Java expression is equivalent to the following prefix expression: `f g h 1 2 f 3 4 5`

Prefix notation is not ambiguous if we know the arity of every operator: Then we can “put parentheses back” to produce an equivalent Lisp expression.

If we do not know the arities of operators, then prefix notation can be ambiguous.

For example, if `+` and `-` are binary then `+ 1 - 2 3` is equivalent to `(+ 1 (- 2 3))`, but if `+` is ternary and `-` is unary then `+ 1 - 2 3` is equivalent to `(+ 1 (- 2) 3)`.

## “RpnLisp” Notation and Postfix Notation

In Lisp, a function call is written as a list whose first element is the function name.

Now consider a notation we’ll call *rpnLisp* that’s the same as Lisp except in that a function call is written as a list whose last element is the function name.

Thus the Lisp expression `(+ (- 1 2) (* 3 4 5) 6)` is equivalent to the following *rpnLisp* expression: `((1 2 -) (3 4 5 *) 6 +)`

Just as prefix notation is “Lisp notation without parentheses”, postfix notation is “*rpnLisp notation without parentheses*”.

So the above expressions are equivalent to this postfix expression: `1 2 - 3 4 5 * 6 +`

In *rpnLisp*, *rpn* stands for “reverse polish notation”.

The above Java expression `f(g(h(1,2), f(3,4)), 5)`  
 is equivalent to the  
 Lisp expression `(f (g (h 1 2) (f 3 4)) 5)`  
 and this rpnLisp expression: `((1 2 h) (3 4 f) g) 5 f)`

Thus the expression is equivalent  
 to this postfix expression: `1 2 h 3 4 f g 5 f`

As is the case with prefix notation, postfix notation  
 isn't ambiguous if we know the arity of each operator:

If we do not know the arities of operators, then  
 postfix notation can be ambiguous.

**For example:**

- If + and - are binary, the postfix expression `1 2 3 + -`  
 is equivalent to the rpnLisp expression `(1 (2 3 +) -)`  
 and equivalent to the Lisp expression `(- 1 (+ 2 3)) => -4.`
- If + is ternary and - is unary, the expression `1 2 3 + -`  
 is equivalent to the rpnLisp expression `((1 2 3 +) -)`  
 and equivalent to the Lisp expression `(- (+ 1 2 3)) => -6.`

## More on Infix Notation

Infix notation allows unary and binary operators, but does not allow operators of arity  $> 2$ .

*Binary* operators are written between their operands.

The designer of an infix notation must specify, for each *unary* operator the notation allows, whether that unary operator is to be written as a ***prefix operator*** or is to be written as a ***postfix operator***:

- Prefix operators are written *before* their operands.
  - **Examples:** `-` in a Java or C++ expression `-x`  
`++` in a Java or C++ expression `++i`  
`*` in a C++ expression `*ptr`
- Postfix operators are written *after* their operands.
  - **Example:** `++` in a Java or C++ expression `i++`



## Syntactically Valid Infix Expressions

An expression  $e$  is a **syntactically valid infix expression** (**s.v.i.e.**) if one of the following is true:

1.  $e$  is a literal constant or an identifier.
2.  $e = ( e_1 )$ , where  $e_1$  is an s.v.i.e.
3.  $e = e_1 \text{ op } e_2$  where each of  $e_1$  and  $e_2$  is an s.v.i.e. and **op** is a binary operator.
4.  $e = \text{op } e_1$  where  $e_1$  is an s.v.i.e. and **op** is a prefix unary operator.
5.  $e = e_1 \text{ op}$  where  $e_1$  is an s.v.i.e. and **op** is a postfix unary operator.

Rules 2 – 5 give decompositions of  $e$  into two (rules 4 & 5) or three (rule 2 & 3) substructures, but some of these decompositions may **violate** the following important principle of syntax specification:

- *The semantics of a structure should be easily definable in terms of the semantics of its syntactic substructures.*

## Example of How the Principle May be Violated

Recall that the principle is:

- *The semantics of a structure should be easily definable in terms of the semantics of its syntactic substructures.*

Let  $e$  be this Java expression:  $-x - y * z + w$

Then 3.  $e = e_1 \text{ op } e_2$  where each of  $e_1$  and  $e_2$  is an s.v.i.e. and  $\text{op}$  is a binary operator.

and 4.  $e = \text{op } e_1$  where  $e_1$  is an s.v.i.e. and  $\text{op}$  is a prefix unary operator.

give the following decompositions of  $e$ :

- |       |                    |                       |                   |
|-------|--------------------|-----------------------|-------------------|
| (i)   | $e_1 = -x$         | $\text{op} = -$       | $e_2 = y * z + w$ |
| (ii)  | $e_1 = -x - y$     | $\text{op} = *$       | $e_2 = z + w$     |
| (iii) | $e_1 = -x - y * z$ | $\text{op} = +$       | $e_2 = w$         |
| (iv)  | $\text{op} = -$    | $e_1 = x - y * z + w$ |                   |

Decompositions (i), (ii), and (iv) violate the principle, as Java's semantics say  $e$  is equivalent to:  $(-x - y * z) + w$

- Sec. 2.5 of Sethi (assigned reading after Exam 1) gives another way to specify syntactically valid infix expressions *that does not have this drawback.*

## Semantics of an Infix Expression $e$

The semantics of  $e$  tells you how  $e$  can be evaluated.

Let  $e.value$  denote the value of  $e$ . Then:

1. If  $e$  is an identifier or a literal constant,  
 $e.value = \text{the value of the identifier / constant.}$
2. If  $e = (e_1)$ ,  $e.value = e_1.value$ .

Otherwise, *let  $op$  be the operator of  $e$  that should be applied last*. Then  $op$  can't be inside  $(...)$  — see rules 2.1–2.3 below — so  $e$  must be  $e_1 \text{ } op \text{ } e_2$ ,  $op \text{ } e_1$ , or  $e_1 \text{ } op$ .

- If  $e$  is  $e_1 \text{ } op \text{ } e_2$ ,  
 $e.value = \text{result of applying } op \text{ with } e_1.value \text{ and } e_2.value \text{ as the 1st and 2nd arguments.}$
- If  $e$  is  $op \text{ } e_1$  or  $e$  is  $e_1 \text{ } op$ ,  
 $e.value = \text{result of applying } op \text{ to } e_1.value$ .

**Key Question:** How can we determine which operator of  $e$  should be applied last?