

Example Write a function `all-numbers` such that:

If $L \Rightarrow$ a proper list, then

$(\text{all-numbers } L) \Rightarrow T$ if every element of the list is a number

$(\text{all-numbers } L) \Rightarrow \text{NIL}$ otherwise.

So: `(all-numbers '(6 2 6)) \Rightarrow T`; `(all-numbers '(7 1 DOG 9)) \Rightarrow NIL`

- **Final cleanup:**

Since `(if c T e) = (or c e)` if the value of `c` is always either T or NIL, we can simplify the above definition to:

```
(defun all-numbers (L)
  (or (null L)
      (and (numberp (car L)) (all-numbers (cdr L)))))
```

Common Mistake:

The following will not work:

```
(defun all-numbers (L)
  (and (numberp (car L)) (all-numbers (cdr L))))
```

This returns NIL whenever the argument is a proper list!

Example Write a function `safe-sum` such that:

- If $L \Rightarrow$ a proper list of numbers, then
(`safe-sum` L) \Rightarrow the sum of the elements of that list.
- If $L \Rightarrow$ a proper list whose elements are not all numbers, then
(`safe-sum` L) \Rightarrow the symbol **ERR!**.

So: (`safe-sum` '(7 2 4 0 9)) \Rightarrow 22; (`safe-sum` '(7 2 A 9)) \Rightarrow ERR!

```
(defun safe-sum (L)
  (if (null L)
```

0

```
    (let ((X (safe-sum (cdr L))))
```

an expression that \Rightarrow value of (`safe-sum` L)
and that involves X and, possibly, L

```
    )))
```

- Suppose $L \Rightarrow (7\ 2\ 4\ 9\ 3)$, so $(\text{cdr } L) \Rightarrow (2\ 4\ 9\ 3)$.
Then $X \Rightarrow 2+4+9+3 = 18$ and \dots should $\Rightarrow 7+2+4+9+3 = 25$.

○ We see $(+ (\text{car } L) X)$ is a good \dots expression for this L !

Q. For what non-null values of L is $(+ (\text{car } L) X)$ a good \dots ?

A. $(+ (\text{car } L) X)$ is a good \dots when $(\text{car } L)$ and $X \Rightarrow$ numbers
(equivalently, when $(\text{car } L) \Rightarrow$ a number and $X \not\Rightarrow$ ERR!).

Example Write a function **safe-sum** such that:

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```
(defun safe-sum (L)
  (if (null L)
      0
      (let ((X (safe-sum (cdr L))))
        (if (and (numberp X) (numberp (car L)))
            (+ (car L) X)
            'ERR!))))
```

Q. For what non-null values of L is $(+ (\text{car } L) X)$ a good ... ?

A. $(+ (\text{car } L) X)$ is a good ... when $(\text{car } L)$ and $X \Rightarrow$ numbers (equivalently, when $(\text{car } L) \Rightarrow$ a number and $X \not\Rightarrow$ ERR!).

Q. What is a good ... when $(\text{car } L) \not\Rightarrow$ a number or $X \Rightarrow$ ERR!?

A. A good ... expression in these cases is: 'ERR! _____

Example Write a function **safe-sum** such that:

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So: (safe-sum '(7 2 4 0 9)) \Rightarrow 22; (safe-sum '(7 2 A 9)) \Rightarrow ERR!

```
(defun safe-sum (L)
  (if (null L)
      0
      (let ((X (safe-sum (cdr L))))
        (if (and (numberp X) (numberp (car L)))
            (+ (car L) X)
            'ERR! )))))
```

Q. Should we eliminate the LET?

A. **No**, because **X** is used twice in the case where

(and (numberp X) (numberp (car L))) \Rightarrow T

- In this case **X** is used as the argument of (numberp X),
and used again as the argument of (+ (car L) X)!

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(safe-sum L) \Rightarrow the symbol **ERR!**.

So: (safe-sum '(7 2 4 0 9)) \Rightarrow 22; (safe-sum '(7 2 A 9)) \Rightarrow ERR!

```
(defun safe-sum (L)
  (if (null L)
      0
      (if (not (numberp (car L)))
          'ERR!
          (let ((X (safe-sum (cdr L))))
            (if (and (numberp X) (numberp (car L)))
                (+ (car L) X)
                'ERR! )))))
```

Q. Is there a case that should be moved outside the LET?

A. **Yes:** The case (numberp (car L)) \Rightarrow NIL should be moved out.
There's no need to use X in that case.

Example Write a function **safe-sum** such that:

- *If $L \Rightarrow$ a proper list of numbers, then
 $(\text{safe-sum } L) \Rightarrow$ the sum of the elements of that list.*
- *If $L \Rightarrow$ a proper list whose elements are not all numbers, then
 $(\text{safe-sum } L) \Rightarrow$ the symbol **ERR!**.*

```
(defun safe-sum (L)
  (cond ((null L) 0)
        ((not (numberp (car L))) 'ERR!)
        (t (let ((X (safe-sum (cdr L))))
              (cond ((numberp X) (+ (car L) X))
                    (t 'ERR!))))))
```

**2nd version
of the final
definition.**

- We didn't eliminate the LET, as its local variable X is used twice in the case where each of (car L) and $X \Rightarrow$ a number.
- Eliminating the LET would produce the function on the next slide, or an equivalent function that uses COND instead of nested IFs. Those functions would be extremely inefficient when L is a list of numbers: Their running time grows exponentially with the length of the list.

- Eliminating LET from the 1st version of the definition gives:

```
(defun safe-sum (L) ; very inefficient!
  (if (null L)
      0
      (if (not (numberp (car L)))
          'ERR!
          (if (numberp (safe-sum (cdr L)))
              (+ (car L) (safe-sum (cdr L)))
              'ERR!))))))
```

- Consider a call of `safe-sum` with argument value `(0 1 2 ... 49)`.
- It makes $2=2^1$ recursive calls with argument value `(1 2 3 ... 49)`.
- Each of those 2^1 calls makes 2 recursive calls with argument value `(2 3 4 ... 49)`, so there are a total of $2^1 \times 2 = 2^2$ recursive calls with argument value `(2 3 4 ... 49)`.
- Each of those 2^2 calls makes 2 recursive calls with argument value `(3 4 5 ... 49)`, so there are a total of $2^2 \times 2 = 2^3$ recursive calls with argument value `(3 4 5 ... 49)`.
- For $0 \leq d \leq 50$, there are 2^d calls with argument value `(d ... 49)`.

- Eliminating LET from the 1st version of the definition gives:

```
(defun safe-sum (L) ; very inefficient!
  (if (null L)
      0
      (if (not (numberp (car L)))
          'ERR!
          (if (numberp (safe-sum (cdr L)))
              (+ (car L) (safe-sum (cdr L)))
              'ERR!))))))
```

- Consider a call of **safe-sum** with argument value (0 1 2 ... 49).
- For $0 \leq d \leq 50$, there are 2^d calls with argument value ($d \dots 49$).
 \therefore the *total* no. of *recursive* calls is $2^1 + \dots + 2^{50} = 2^{51} - 2 > 2 \times 10^{15}$.
- **General Principle:** If a function f can make 2 or more direct recursive calls, then a single call of f might well produce 2^d or more recursive calls of f at recursion depth d .
- LET can be used to prevent a function from making 2 or more direct recursive calls with the very same argument values!

- **General Principle:** If a function `f` can make 2 or more direct recursive calls, then a single call of `f` might well produce 2^d or more recursive calls of `f` at recursion depth d .
- **LET can be used to prevent a function from making 2 or more direct recursive calls with the very same argument values!**
- The 1st and 2nd versions of `safe-sum` use `LET` in this way.

```
(defun safe-sum (L)
  (if (null L)
      0
      (if (not (numberp (car L)))
          'ERR!
          (let ((X (safe-sum (cdr L))))
            (if (numberp X)
                (+ (car L) X)
                'ERR!))))))
```

**1st version
of the final
definition.**

- These versions never make more than one direct recursive call, as a result of which `(safe-sum '(0 1 ... 49))` computes its result using **just 50** recursive calls rather than quadrillions!

Comments on Lisp Assignment 4

Problems 1–13 can be solved by starting with one of the templates below or a dual of the 2nd template in which the roles of e1 and e2 are switched. (These are just the templates presented earlier!)

```
(defun f (e)
  (if (null e) or (zerop e)
      value of (f nil) or (f 0)
      (let ((X (f (cdr e)) or (f (- e 1)))
            an expression that  $\Rightarrow$  value of (f e)
            and that involves X and, possibly, e )))
```

```
(defun f (e1 e2)
  (if (null e1) or (zerop e1)
      value of (f nil e2) or (f 0 e2)
      (let ((X (f (cdr e1) e2) or (f (- e1 1) e2))
            an expression that  $\Rightarrow$  value of (f e1 e2) and
            that involves X and, possibly, e1 and/or e2 )))
```

Comments on Lisp Assignment 4

Problems 1–13 can be solved by starting with one of the templates above or a dual of the 2nd template in which the roles of e1 and e2 are switched. (These are just the templates presented earlier!)

Recall that:

- If there is no case in which **X** is used more than once, then eliminate the LET.
- If the LET isn't eliminated, move any case in which **X** needn't be used out of the LET. If the LET is eliminated but there's a case where the recursive call's result isn't needed, deal with such cases as base cases--i.e., without making a recursive call.

In the recursive function definitions that were given above:

- In non-base cases the result is computed using just one recursive call, and it is the same recursive call in all non-base cases.
- The function has a formal parameter `e` for which it passes the value of `(cdr e)` or `(- e 1)` to the same parameter of the recursive call in non-base cases.
 - `e` may not be the only parameter, but the value of any other parameter is passed *without change* to the same parameter of the recursive call in non-base cases.

All 13 problems in section 2 of Lisp Assignment 4 can be solved using recursive functions of this simple kind, *but when doing Lisp Assignment 5 you must be prepared to write recursive functions that work differently!*

When a function makes a recursive call, there will often be a formal parameter `e` of the function for which the value passed to the same parameter of the recursive call is smaller in size than the value of `e`.

`(cdr e)` and `(- e 1)` may be used to produce the value of smaller size. Other expressions that can be used to do that include:

When a function makes a recursive call, there will often be a formal parameter `e` of the function for which the value passed to the same parameter of the recursive call is *smaller in size* than the value of `e`.

`(cdr e)` and `(- e 1)` may be used to produce the value of smaller size. *Other* expressions that can be used to do that include:

- `(cddr e)` if `e` \Rightarrow a nonempty list.
- `(- e 2)` if `e` \Rightarrow an integer ≥ 2 .
- `(floor e 2)` if `e` \Rightarrow an integer other than 0 or -1.
 - `(floor e 2) = $\lfloor e/2 \rfloor = e \gg 1$` in Java if `e` \Rightarrow an integer.
- `(/ e 2)` if `e` \Rightarrow an *even* integer other than 0.
- `(cdr L1)` if `e` \Rightarrow a nonempty list;
 - here `L1` \Rightarrow a list, obtained by *transforming* the list given by `e` in some way, whose length is \leq the length of that list.
- For Assignment 5, your function `SSORT` should use this kind of expression to produce the argument value for its recursive call.

Example of the Use of `(cddr L)` as a Recursive Call Argument

Recall from Assignment 4: If $L \Rightarrow$ a list then `(SPLIT-LIST L)` returns a list of two lists, in which the 1st list consists of the 1st, 3rd, 5th, ... elements of the list given by L , and the 2nd list consists of the 2nd, 4th, 6th, ... elements of the list given by L .

For example: `(SPLIT-LIST ()) => (NIL NIL)` `(SPLIT-LIST '(B)) => ((B) NIL)`
`(SPLIT-LIST '(A B C D 1 2 3 4 5)) => ((A C 1 3 5) (B D 2 4))`

```
(defun split-list (L)
  (if (null L)
      '(()())
      (let ((X (split-list (cddr L))))
        an expression that  $\Rightarrow$  value of (split-list L)
        and that involves X and, possibly, L.
      )))
```

- Let $L \Rightarrow (A B C D 1 2 3 4 5)$, so $(cddr L) \Rightarrow (C D 1 2 3 4 5)$.
Then $X \Rightarrow ((C 1 3 5) (D 2 4))$
and `...` should $\Rightarrow ((A C 1 3 5) (B D 2 4))$.
- Q. What is a good `...` expression in this case?
A. `(list (cons (car L) (car X)) (cons (cadr L) (cadr X)))`
- Q. For what non-null values of L is this not a good `...`?

Example of the Use of (cddr L) as a Recursive Call Argument

```
(defun split-list (L)
  (if (null L)
      '(()())
      (let ((X (split-list (cddr L))))
        an expression that  $\Rightarrow$  value of (split-list L)
        and that involves X and, possibly, L.))))
```

- Let $L \Rightarrow (A\ B\ C\ D\ 1\ 2\ 3\ 4\ 5)$, so $(cddr\ L) \Rightarrow (C\ D\ 1\ 2\ 3\ 4\ 5)$.
Then $X \Rightarrow ((C\ 1\ 3\ 5)\ (D\ 2\ 4))$
and \dots should $\Rightarrow ((A\ C\ 1\ 3\ 5)\ (B\ D\ 2\ 4))$.
- Q. What is a good \dots expression in this case?
A. $(list\ (cons\ (car\ L)\ (car\ X))\ (cons\ (cadr\ L)\ (cadr\ X)))$
- Q. For what non-null values of L is this not a good \dots ?
A. It's not good if $L \Rightarrow$ a list of length 1--e.g., $L \Rightarrow (B)$.
- Q. What is a good \dots expression in that case?
A. $(list\ L\ ())$

So \dots can
be written:

```
(cond ((null (cdr L)) (list L ()))
      (t (list (cons (car L) (car X))
                 (cons (cadr L) (cadr X))))))
```

Example of the Use of (cddr L) as a Recursive Call Argument

```
(defun split-list (L)
  (if (null L)
      '(()())
      (let ((X (split-list (cddr L))))
        (cond ((null (cdr L)) (list L ()))
              (t (list (cons (car L) (car X))
                        (cons (cadr L) (cadr X)))))))))
```

- As **X** is used twice in the **t** case, we must not eliminate the LET: The function would be very inefficient if it called (split-list (cddr L)) twice!
- As **X** is not used in the (null (cdr L)) case, it's good to move that case out of the LET.
- After that case is moved out of the LET, it can be combined with the (null L) base case, because (list L ()) is a good value to return in both cases.

Example of the Use of (cddr L) as a Recursive Call Argument

```
(defun split-list (L)
  (if (null L)
      '(()())
      (let ((X (split-list (cddr L))))
        (cond ((null (cdr L)) (list L ()))
              (t (list (cons (car L) (car X))
                        (cons (cadr L) (cadr X))))))))
```

- As **X** is not used in the (null (cdr L)) case, it's good to move that case out of the LET.
- After that case is moved out of the LET, it can be combined with the (null L) base case, because (list L ()) is a good value to return in both cases.

Final version: (defun split-list (L)

Note that calling (if (null (cdr L))
 (split-list (cddr L))
instead of
 (split-list (cdr L))
reduces the depth
of recursion.

```
                          (list L ())
                          (let ((X (split-list (cddr L))))
                            (list (cons (car L) (car X))
                                (cons (cadr L) (cadr X)))))
```