

A Mistake to Avoid When Doing TinyJ Assignment 1

A common mistake in writing recursive descent parsing code is to write

```
getCurrentToken() == X
```

or `accept(X)` [which performs a `getCurrentToken() == X` test]

using a Symbols constant *X* that represents a nonterminal. This is wrong, as `getCurrentToken()` returns a Symbols constant that represents a token. Here are two examples of this kind of mistake.

1. When writing the method `argumentList()`, which should be based on the EBNF rule

```
<argumentList> ::= '(' [<expr3>{,<expr3>}] ')'
```

it would be wrong to write:

```
accept(LPAREN);
if (getCurrentToken() == NExpr3) /* INCORRECT! */ {
    expr3();
    ... // a while loop that deals with {,<expr3>}
}
accept(RPAREN);
```

Here it would be correct to write code of the following form:

```
accept(LPAREN);
if (getCurrentToken() != RPAREN) /* CORRECT */ {
    expr3();
    ... // a while loop that deals with {,<expr3>}
}
accept(RPAREN);
```

2. When writing the method `expr1()`, one case you need to deal with relates to the following part of the EBNF rule that defines `<expr1>`:

```
IDENTIFIER ( . nextInt '(' ')' | [<argumentList>] { '[' <expr3> ']' } )
```

Here it would be wrong to write something like:

```
case IDENT:
    nextToken();
    if (getCurrentToken() != DOT) {
        if (getCurrentToken() == NTargumentList /* INCORRECT! */ ) argumentList();
        ... // a while loop that deals with { '[' <expr3> ']' }
    }
    else {
        ... // code to deal with . nextInt '(' ')'
    }
    break;
```

Instead, you can write something like:

```
case IDENT:
    nextToken();
    if (getCurrentToken() != DOT) {
        if (getCurrentToken() == LPAREN /* CORRECT */ ) argumentList();
        ... // a while loop that deals with { '[' <expr3> ']' }
    }
    else {
        ... // code to deal with . nextInt '(' ')'
    }
    break;
```

The use of `LPAREN` in the above code is correct because the first token of any instance of `<argumentList>` must be a left parenthesis, as we see from the EBNF rule

```
<argumentList> ::= '(' [<expr3>{,<expr3>}] ')'
```

An Old Exam Question

A student is debugging his current version of `Parser.java` for TinyJ Assignment 1. He compiles his file and then runs his program as follows:

```
java -cp . TJlasn.TJ X.java X.out
```

He also runs the solution that was provided, as follows:

```
java -cp TJ1solclasses:. TJlasn.TJ X.java X.sol
```

The first difference between the output files `X.out` and `X.sol` is that `X.sol` has a comma on **line 567**, but this is missing in `X.out`. Lines 556 – 568 of `X.sol` and `X.out` are reproduced below with line numbers. (Lines 556 – 566 are the same in both output files.)

```
Lines 556 - 568 of X.sol [Output produced by java -cp TJ1solclasses:. TJlasn.TJ ...]:
556      <expr1>
557      IDENTIFIER: leq
558      <argumentList>
559      (
560      <expr3>
561      <expr2>
562      <expr1>
563      IDENTIFIER: size
564      ... node has no more children
565      ... node has no more children
566      ... node has no more children
567      ,
568      <expr3>
```

```
Lines 556 - 568 of X.out [Output produced by java -cp . TJlasn.TJ ...]:
556      <expr1>
557      IDENTIFIER: leq
558      <argumentList>
559      (
560      <expr3>
561      <expr2>
562      <expr1>
563      IDENTIFIER: size
564      ... node has no more children
565      ... node has no more children
566      ... node has no more children
567      <expr3>
568      <expr2>
```

Hint: In reading this output, recall that the indentation levels of consecutive lines are either the same or differ by just 1; thus line 567 has the same indentation as line 559.

Now answer the following two questions. In each case, *circle the correct choice*. [The answers are given on the next page.]

- (i) The output files show there is probably an error in the student's version of the method
(a) `expr1()` (b) `expr2()` (c) `expr3()` (d) `argumentList()` (e) `ifStmt()`
[1 pt.]
- (ii) Which one of the following changes might well fix this error?
(a) Insert a missing call of `accept (COMMA)` or `nextToken()` in the student's `Parser.java`.
(b) Delete a call of `accept (COMMA)` from the student's `Parser.java`.
(c) Delete a call of `nextToken()` from the student's `Parser.java`.
(d) Insert a missing call of `expr3()` in the student's `Parser.java`.
(e) Delete a call of `expr2()` from the student's `Parser.java`.
[1 pt.]

Debugging Hints for TinyJ Assignment 1

1. It is a very common mistake to omit a call of `accept(...)` or `nextToken()`: For *each* token in the EBNF definition of a non-terminal $\langle N \rangle$, the body of the corresponding parsing method $N()$ should contain a call of `accept(...)` or `nextToken()` whose execution may cause that token to be output as a parse tree node. Another common mistake is to call `nextToken()` when `accept(...)` should be called; this often produces the following error message:

Internal error in parser: Token discarded without being inspected

A third common mistake is to pass a `Symbols` object that represents a *non*-terminal as an argument to `accept(...)`, as in `accept(NTexpr7)`;—see [A Mistake to Avoid When Doing TinyJ Assignment 1](#) above.

2. The sideways parse tree in the output file can be regarded as an *execution trace* of your program, and can be useful when debugging your code! If your program is not working correctly, and you have produced both *k.sol* and *k.out* for some *k* (as described on page 4 of the assignment document), then the first line in *k.sol* that isn't in *k.out* shows "something my solution did that your program didn't do". (You can find that line from the output of `diff -c [on euclid/venus] or fc.exe /n [on a PC].`) When reading the output file for debugging purposes, bear the following in mind:
 - A. In a sideways parse tree, the parent of a node appears on the most recent previous line that has lower indentation. (Note that adjacent lines of the tree either have the same indentation or have indentation levels that differ by just 1.) For example, in the Old Exam Question, the parent of the comma on line 567 of *X.sol*, and of $\langle \text{expr3} \rangle$ on line 568, is $\langle \text{argumentList} \rangle$ on line 558.
 - B. Each non-terminal $\langle N \rangle$ in the output file is written when the corresponding parsing method $N()$ is called (by the call of `TJ.output.printSymbol(...)` at the beginning of N 's body). The value of `getCurrentToken()` at that time is shown by the first token in the output file *after* $\langle N \rangle$'s line. $\langle N \rangle$'s parent in the parse tree shows the caller of $N()$. For example, in the Old Exam Question, $\langle \text{expr3} \rangle$ on line 560 of *X.sol* was written when `expr3()` was called. The value of `getCurrentToken()` was `IDENT` at the time of the call (as shown by line 563); `expr3()` was called by the method corresponding to the parent of the $\langle \text{expr3} \rangle$ node on line 560—i.e., by `argumentList()`, as we see from line 558.
 - C. Each token in the output file is written during execution of a call of `accept(T)` or `nextToken()` in some non-terminal's parsing method, at a time when the value of `getCurrentToken()` is *T*; here *T* is the `Symbols` object that represents the token. The parsing method in question is shown by the token's parent in the parse tree. For example, in the Old Exam Question, the comma on line 567 of *X.sol* was written during execution of a call of `accept(COMMA)` or `nextToken()` in a non-terminal's parsing method; the value of `getCurrentToken()` was `COMMA` at the time of the call, and we see from line 558 that the parsing method in question was `argumentList()`.
 - D. The `... node has no more children` line that is a child of a node $\langle N \rangle$ of the tree is written *just before* the corresponding call of method $N()$ returns control to its caller. The value of `getCurrentToken()` at that time is shown by the first token in the output file *after* the line `... node has no more children`. For example, in the Old Exam Question, the line `... node has no more children` on line 565 of *X.sol* is a child of the node $\langle \text{expr2} \rangle$ on line 561, and was therefore written just before the corresponding call of `expr2()` returned control to its caller. The caller was `expr3()`, since the parent of $\langle \text{expr2} \rangle$ is the node $\langle \text{expr3} \rangle$ on line 560. Line 567 of *X.sol* shows that the value of `getCurrentToken()` was `COMMA` when `expr2()` returned control to `expr3()`.

The correct answers to the [Old Exam Question](#) are (i)—(d) and (ii)—(a). This follows from 2A, 2B, and 2C above.