

Example

What TinyJ VM code would the TinyJ compiler generate for the declaration **static int x, y = 10;** in the program on p. 7 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> ?

Solution

We use the static memory allocation rule on p. 3 and Code Generation Rules 1 and 2 on p. 9 of:

<https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>

x is allocated data memory address 0.

y is allocated data memory address 1.

[*Blue* items are on EXPRSTACK after

The generated instructions are: *execution of each instruction.*]

0: PUSHSTATADDR 1	Pushes pointer to y.
1: PUSHNUM 10	Pushes 10.
2: SAVETOADDR	Pops 10. Pops ptr to y. Stores 10 into y's location.

Example

What is the 1st VM instruction generated by the TinyJ compiler for each of the methods **main**, **f**, and **g** in the program on p. 7 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>?

Solution

From Code Generation Rule 4, the instruction is:

INITSTKFRM <no. of stackframe locations given to local vars declared in the method's body>

- **main**'s stackframe has **no locations for local vars** declared in main's body--there are no such vars *in this example!*

From the stack-dynamic memory allocation rules on p. 3 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>:

- The 3 local variables declared in **f**'s body are given 3 different locations in each stackframe of **f**.
- The only local variable declared in **g**'s body is given 1 location in each stackframe of **g**.

Hence **main**'s code begins with: 3: **INITSTKFRM** 0
f's code begins with: ?: **INITSTKFRM** 3
g's code begins with: ?: **INITSTKFRM** 1

Code memory address is not known until main and f have been translated.

Code memory address is not known until main has been translated.

Example

What TinyJ VM code would the TinyJ compiler generate for the 1st two statements of main method in the program on p. 7 of <https://euclid.cs.qc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>?

The two statements are: **System.out.print("Enter num: ");**
x = input.nextInt();

Solution

The static memory allocation rules on p. 3 of <https://euclid.cs.qc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> imply:

- x is allocated data memory address 0.
- y is allocated data memory address 1.
- The 11 characters of "Enter num: " are allocated data memory addresses 2, 3, ..., 11, 12.

[Blue items are on EXPSTACK after

Hence the generated instructions are: *execution of each instruction.*]

4: WRITESTRING 2 12	Writes "Enter num: " to the screen.
5: PUSHSTATADDR 0	Pushes pointer to x.
6: READINT	Reads an int from kbd; pushes its value.
7: SAVETOADDR	Pops the int; pops the pointer to x; stores the int into x's location.

What TinyJ VM code does the TinyJ compiler generate for this method?

Solution

The static and stack-dynamic memory allocation rules imply:

- y is given data memory address 1.
- z is given the location with offset +1 in a stackframe of g.
- e is given the location with offset -2 in a stackframe of g.
- d is given the location with offset -3 in a stackframe of g.

```
static void g (int d, int e)
{
    int z;
    y = d / e;
}
```

Generated Code:

[Blue items are on EXPSTACK after execution

INITSTKFRM 1	From Code Generation Rule 4.	of each instruction.]
PUSHSTATADDR 1	Pushes ptr to y.	
PUSHLOCADDR -3	Pushes ptr to d.	
LOADFROMADDR	Replaces ptr to d with d's value on top of stack.	
PUSHLOCADDR -2	Pushes ptr to e.	
LOADFROMADDR	Replaces ptr to e with e's value on top of stack.	
DIV	Pops e's value; pops d's value; pushes (d/e)'s value.	
SAVETOADDR	Pops (d/e)'s value;	
	pops ptr to y;	
	stores (d/e)'s value into y's location.	
RETURN 2	From Code Generation Rule 6, as g has 2 parameters.	

What TinyJ VM code does the TinyJ compiler generate for the statement `return y - a % u;` in method `f` in the program on p. 7 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>?

Solution

The static and stack-dynamic memory allocation rules imply:

- `y` is given the data memory address 1.
- `a` is given the location with offset -4 in a stackframe of `f`.
- `u` is given the location with offset +3 in a stackframe of `f`.

In view of Code Generation Rule 7, the generated code is:

PUSHSTATADDR 1	Pushes ptr to <code>y</code>.
LOADFROMADDR	Replaces ptr to <code>y</code> with <code>y</code>'s value on top of stack.
PUSHLOCADDR -4	Pushes ptr to <code>a</code>.
LOADFROMADDR	Replaces ptr to <code>a</code> with <code>a</code>'s value on top of stack.
PUSHLOCADDR +3	Pushes ptr to <code>u</code>.
LOADFROMADDR	Replaces ptr to <code>u</code> with <code>u</code>'s value on top of stack.
MOD	Pops <code>u</code>'s and <code>a</code>'s values; pushes <code>(a%u)</code>'s value.
SUB	Pops <code>(a%u)</code>'s and <code>y</code>'s values; pushes <code>(y-a%u)</code>'s value.
RETURN 3	From Code Generation Rule 7, as <code>f</code> has 3 parameters.

Code Generation Rule 4 implies that the first TinyJ VM instruction generated for the method **f** in the program on p. 7 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> is **INITSTKFRM 3**. Hand-translation of **main** shows that this instruction is placed in code memory at address 34.

What code does the compiler generate so **f(21,22,23)**'s value will be on top of EXPRSTACK when the next instruction in main's code is executed?

Solution

In view of Code Generation Rule 8, the generated code is:

PUSHNUM 21	Pushes 21.
PASSPARAM	Pops 21. Stores 21 in 1 st param's loc in f's stackframe.
PUSHNUM 22	Pushes 22.
PASSPARAM	Pops 22. Stores 22 in 2 nd param's loc in f's stackframe.
PUSHNUM 23	Pushes 23.
PASSPARAM	Pops 23. Stores 23 in 3 rd param's loc in f's stackframe.
CALLSTATMETHOD 34	Next instr. to be executed will be: 34 INITSTKFRM 3 f's execution will leave f(21,22,23) 's value on stack.

What TinyJ VM code does the TinyJ compiler generate for the statement `System.out.println(y + f(21,22,23));` in `main` in the program on p. 7 of

<https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> ?

Solution

Recall that `y` is given the data memory address 1.

Generated Code:

`PUSHSTATADDR 1` Pushes ptr to `y`.
`LOADFROMADDR` Replaces ptr to `y` with `y`'s value on top of stack.

`PUSHNUM 21`

`PASSPARAM`

`PUSHNUM 22`

`PASSPARAM`

`PUSHNUM 23`

`PASSPARAM`

`CALLSTATMETHOD 34`

`ADD` Pops value returned by `f(21,22,23)`; pops `y`'s value; pushes `(y+f(21,22,23))`'s value.

`WRITEINT`

`WRITELNOP`

[SEE
EARLIER
SLIDE.](#)

Execution of method `f` puts the value returned by `f(21,22,23)` on top of stack.

Pops `(y+f(21,22,23))`'s value; writes it to the screen.
Writes a newline to the screen.

What TinyJ VM code does the TinyJ compiler generate for the statement `f(17,y,x-y);` in `main` in the program on p. 7 of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> ?

Solution

Recall: `x` is given data mem. addr. `0`; `y` is given data mem. addr. `1`.
`f`'s code begins with `34 INITSTKFRM 3`

Generated Code:

<code>PUSHNUM 17</code>	Pushes 17.
<code>PASSPARAM</code>	Pops 17 & stores it in 1 st param's loc in <code>f</code> 's stackfrm.
<code>PUSHSTATADDR 1</code>	Pushes ptr to <code>y</code> .
<code>LOADFROMADDR</code>	Replaces ptr to <code>y</code> with <code>y</code> 's value on top of stack.
<code>PASSPARAM</code>	Pops <code>y</code> 's value & stores it in 2 nd param's loc in <code>f</code> 's stackfrm.
<code>PUSHSTATADDR 0</code>	Pushes ptr to <code>x</code> .
<code>LOADFROMADDR</code>	Replaces ptr to <code>x</code> with <code>x</code> 's value on top of stack.
<code>PUSHSTATADDR 1</code>	Pushes ptr to <code>y</code> .
<code>LOADFROMADDR</code>	Replaces ptr to <code>y</code> with <code>y</code> 's value on top of stack.
<code>SUB</code>	Pops <code>y</code> 's and <code>x</code> 's values; pushes <code>(x-y)</code> 's value.
<code>PASSPARAM</code>	Pops <code>(x-y)</code> 's value & stores it in 3 rd param's loc in <code>f</code> 's stackfrm.
<code>CALLSTATMETHOD 34</code>	Next instr. to be executed will be: <code>34 INITSTKFRM 3</code> <code>f</code> 's execution will leave <code>f(17,y,x-y)</code> 's value on stack.
<u><code>DISCARDVALUE</code></u>	Pops <code>f(17,y,x-y)</code> 's value as per Code Generation Rule 9.

Hand-translation of the program on p. 7 of

<https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>

shows that the first instruction of method **g** is placed in code memory at address 60.

*What code is generated for the statement **g(c,b+u);** in method **f**?*

Solution

- **b** is given the location with offset -3 in **f**'s stackframe.
- **c** is given the location with offset -2 in **f**'s stackframe.
- **u** is given the location with offset +3 in **f**'s stackframe.

Generated Code:

PUSHLOCADDR -2	Pushes ptr to c .
LOADFROMADDR	Replaces ptr to c with c 's value on top of stack.
PASSPARAM	Pops c 's value & stores it in 1 st param's loc in f 's stackfrm.
PUSHLOCADDR -3	Pushes ptr to b .
LOADFROMADDR	Replaces ptr to b with b 's value on top of stack.
PUSHLOCADDR +3	Pushes ptr to u .
LOADFROMADDR	Replaces ptr to u with u 's value on top of stack.
ADD	Pops u 's and b 's values; pushes (b+u) 's value.
PASSPARAM	Pops (b+u) 's value & stores it in 2 nd param's loc in f 's stackfrm.
CALLSTATMETHOD 60	Next instr. to be executed will be g 's 1 st instr.
NOP	Does nothing. See Code Generation Rule 9 .

EXECUTION OF VARIOUS TINYJ VM INSTRUCTIONS

BEFORE execution of: WRITESTRING 3 9

s t k o f f r a m e t	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)	a d d r e s s	DATA MEMORY	a d d r e s s	HEAP (Part of Data Memory)	a d d r e s s	CODE MEMORY
		0					
		1					
		2					
		3	'T'				
		4	'h'				
		5	'e'				
		6	' '				
		7	'C'				
		8	'a'				
		9	't'				
		10					
		11					
		⋮					

EXPRSTACK

	WRITESTRING 3 9
--	-----------------

AFTER execution of: **WRITESTRING 3 9**

s t k o f f r a m e t	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

EXPRSTACK

a d d r e s s	DATA MEMORY
0	
1	
2	
3	'T'
4	'h'
5	'e'
6	' '
7	'C'
8	'a'
9	't'
10	
11	
:	

a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
	WRITESTRING 3 9

NOTE: In this example, execution of **WRITESTRING 3 9** writes the string The Cat to the screen.

BEFORE execution of: **PUSHNUM 23**

s t k o f f r a m e t	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

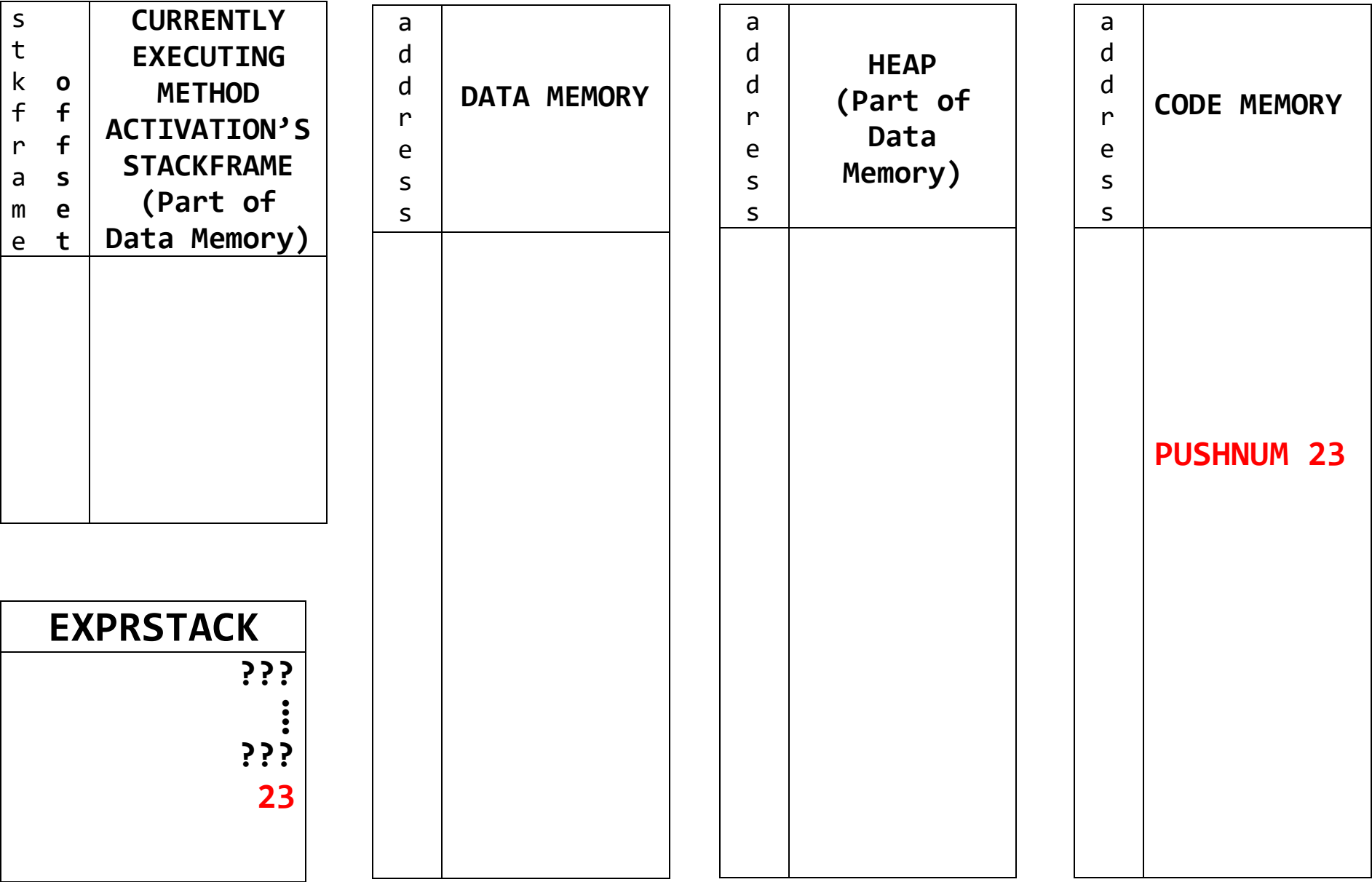
EXPRSTACK
???
⋮
???

a d d r e s s	DATA MEMORY

a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
	PUSHNUM 23

AFTER *execution of:* **PUSHNUM 23**



BEFORE execution of: PUSHSTATADDR 17

Stack frame	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

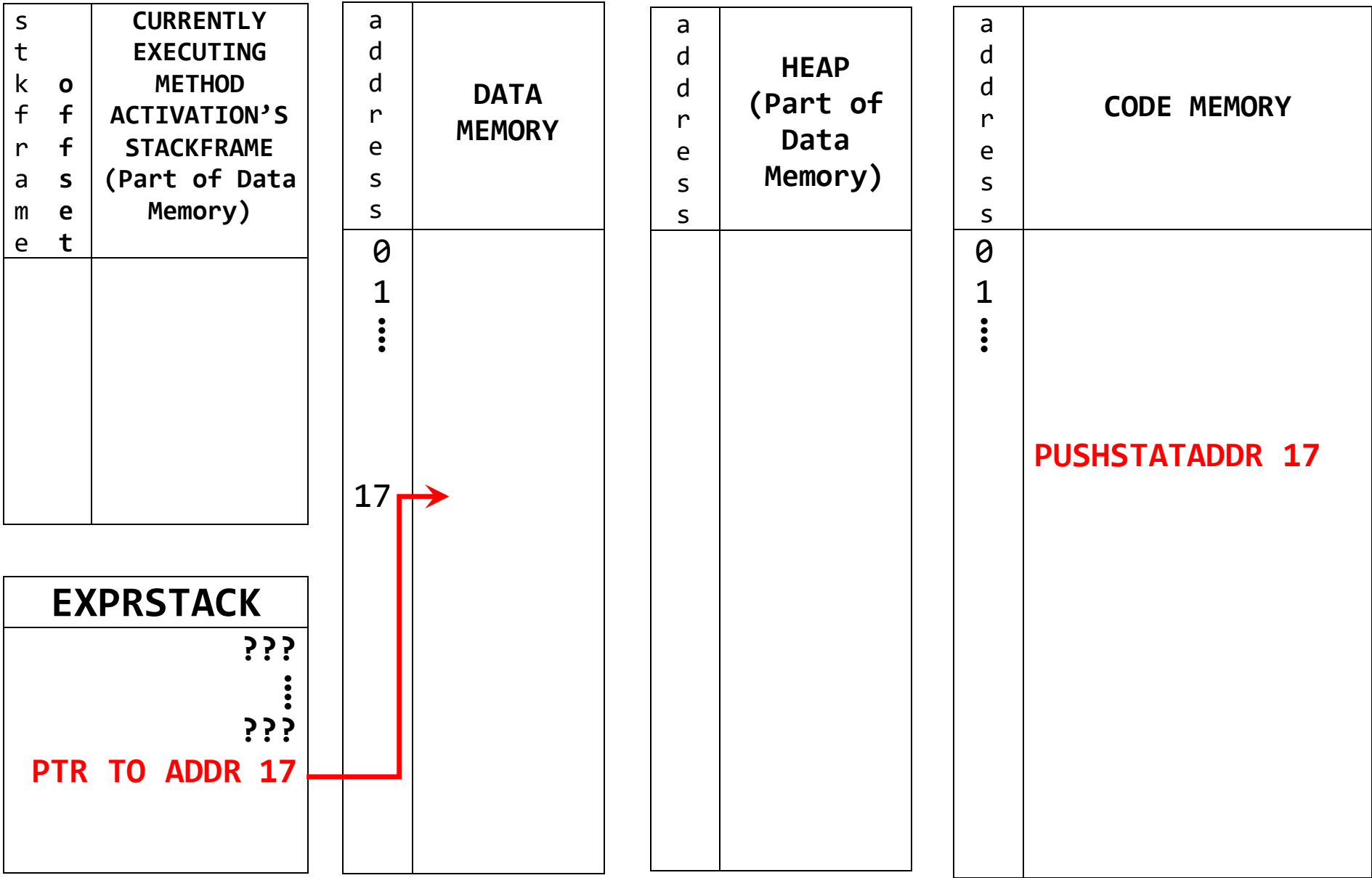
EXPRSTACK
???
:
???

a d d r e s s	DATA MEMORY
0 1 ⋮ 17	

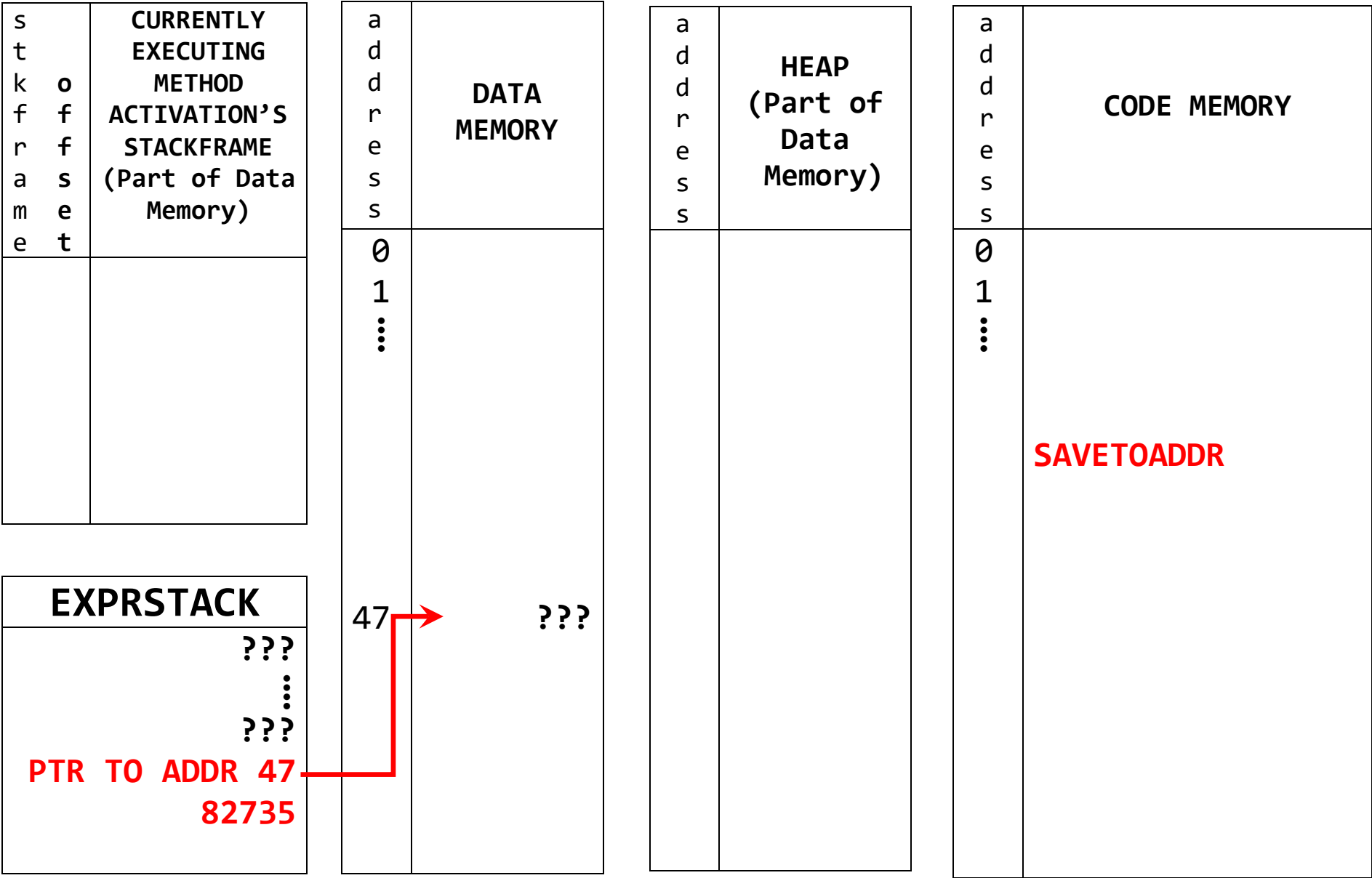
a d d r e s s	HEAP (Part of Data Memory)

addresses	CODE MEMORY
0 1 ⋮	PUSHSTATADDR 17

AFTER execution of: **PUSHSTATADDR 17**



BEFORE execution of SAVETOADDR



AFTER *execution of SAVETOADDR*

Stack offset address memory entry	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

EXPRSTACK
???
⋮
???

a d d r e s s	DATA MEMORY
0 1 ⋮	
47	82735

a d d r e s s	HEAP (Part of Data Memory)

address	CODE MEMORY
0 1 ⋮	SAVETOADDR

BEFORE *execution of* READINT

s t o c k f r a m e s	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

EXPRSTACK
???
⋮
???

a d d r e s s	DATA MEMORY
0	
1	
⋮	

a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0	
1	
⋮	
	READINT

AFTER *execution of READINT*

s t a c k o f f r a m e s	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)

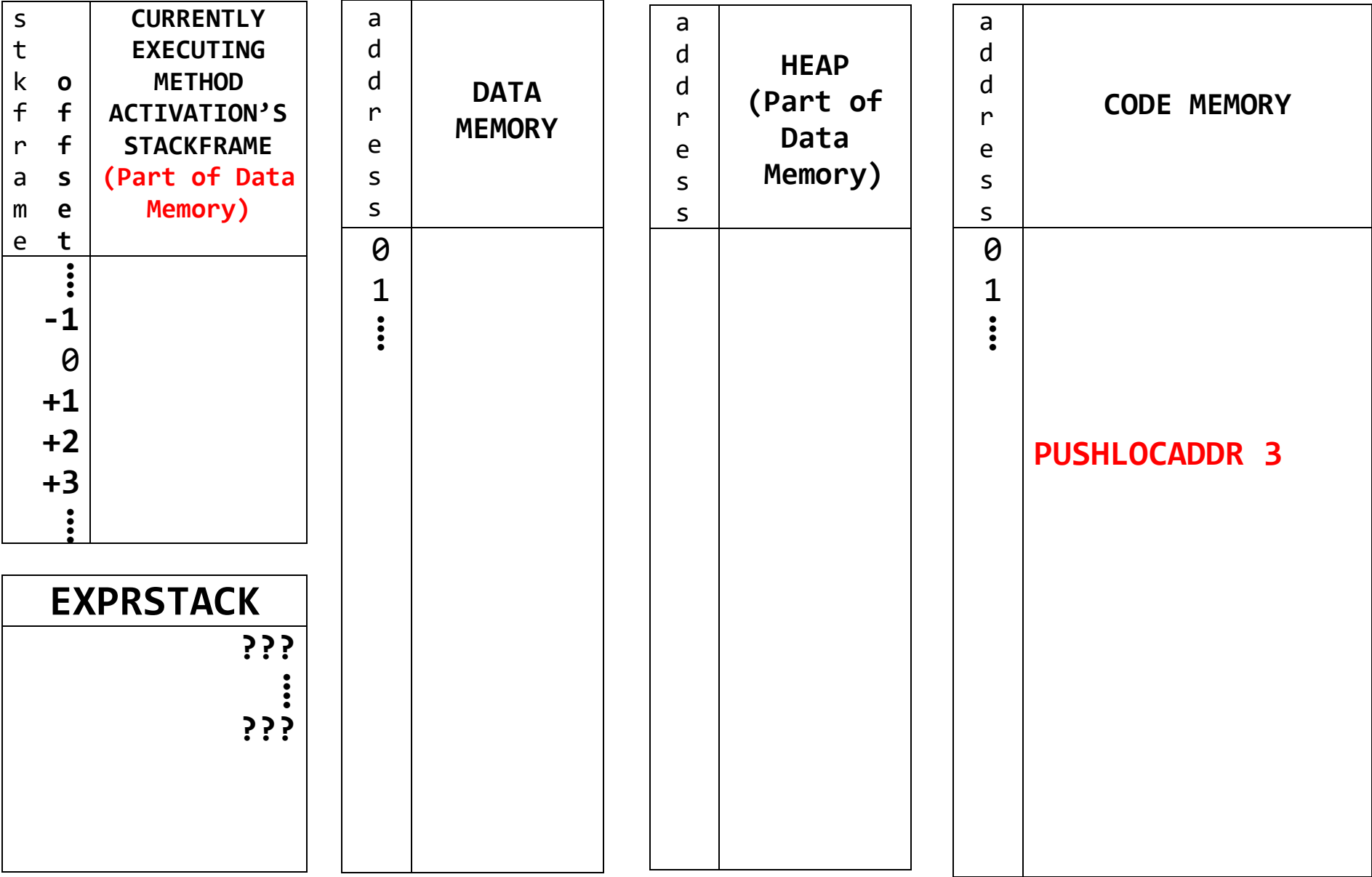
EXPRSTACK
???
⋮
???
int entered on kbd

a d d r e s s	DATA MEMORY
0	
1	
⋮	

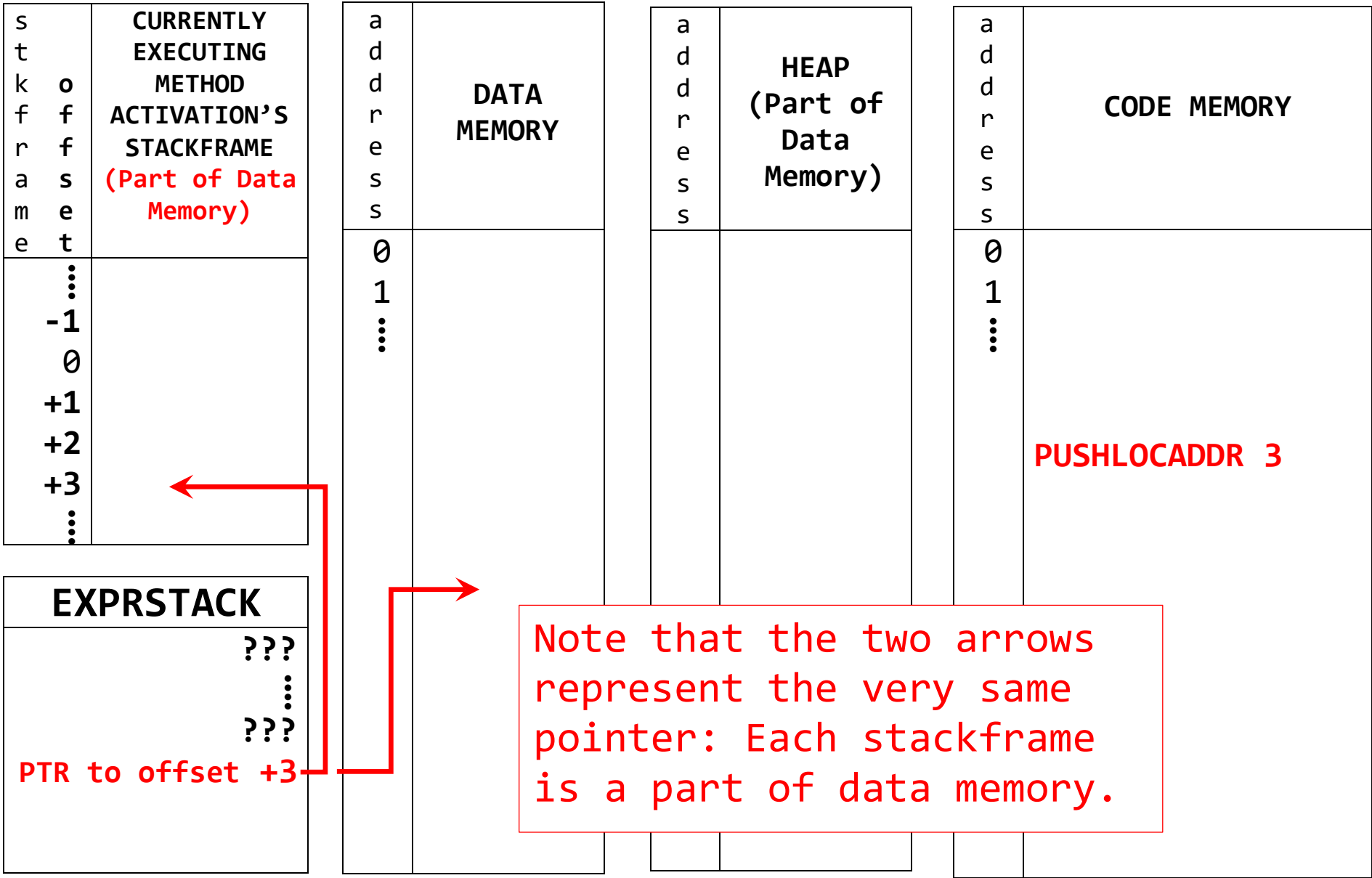
a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0	
1	
⋮	
	READINT

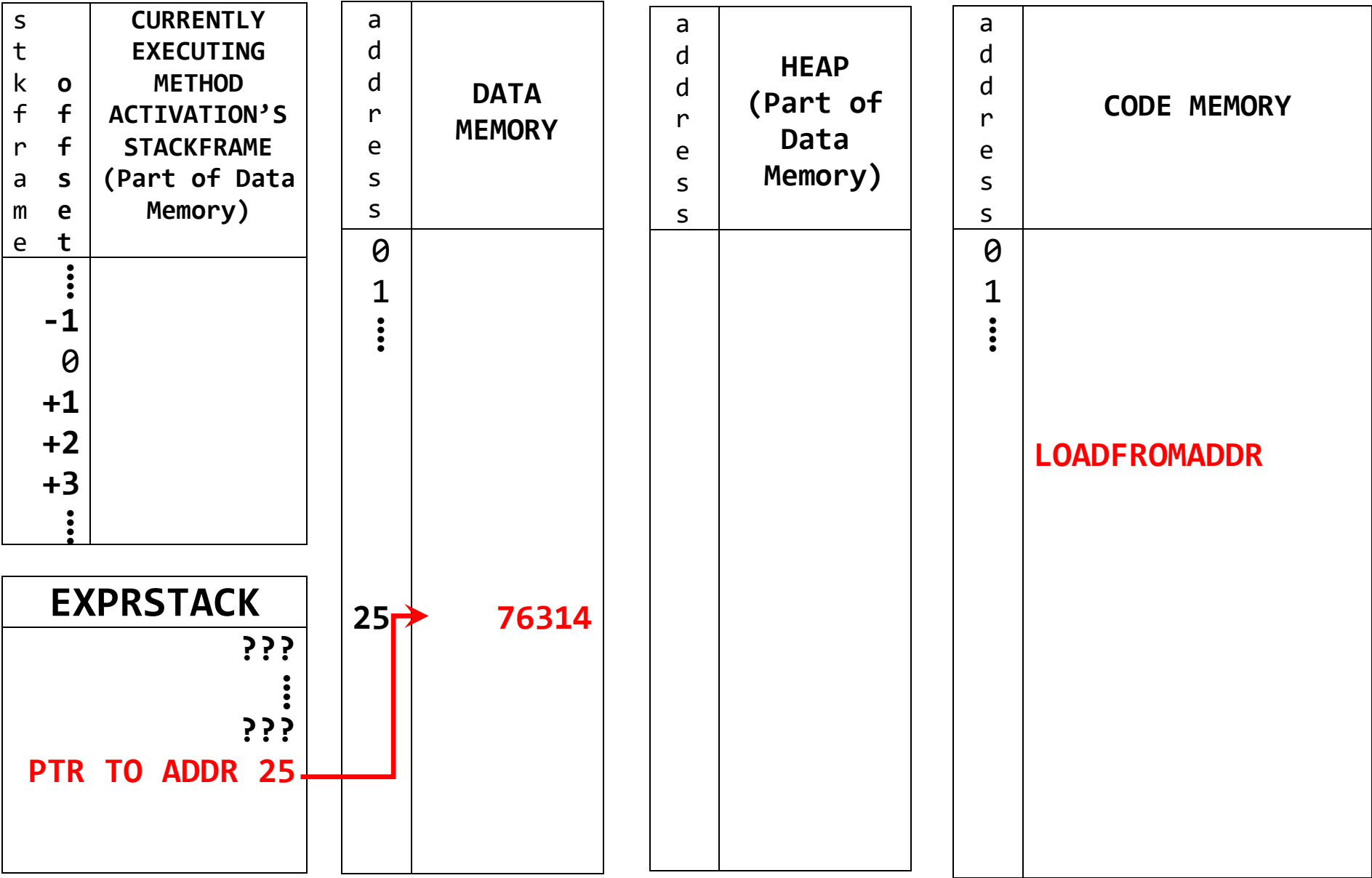
BEFORE execution of: **PUSHLOCADDR 3**



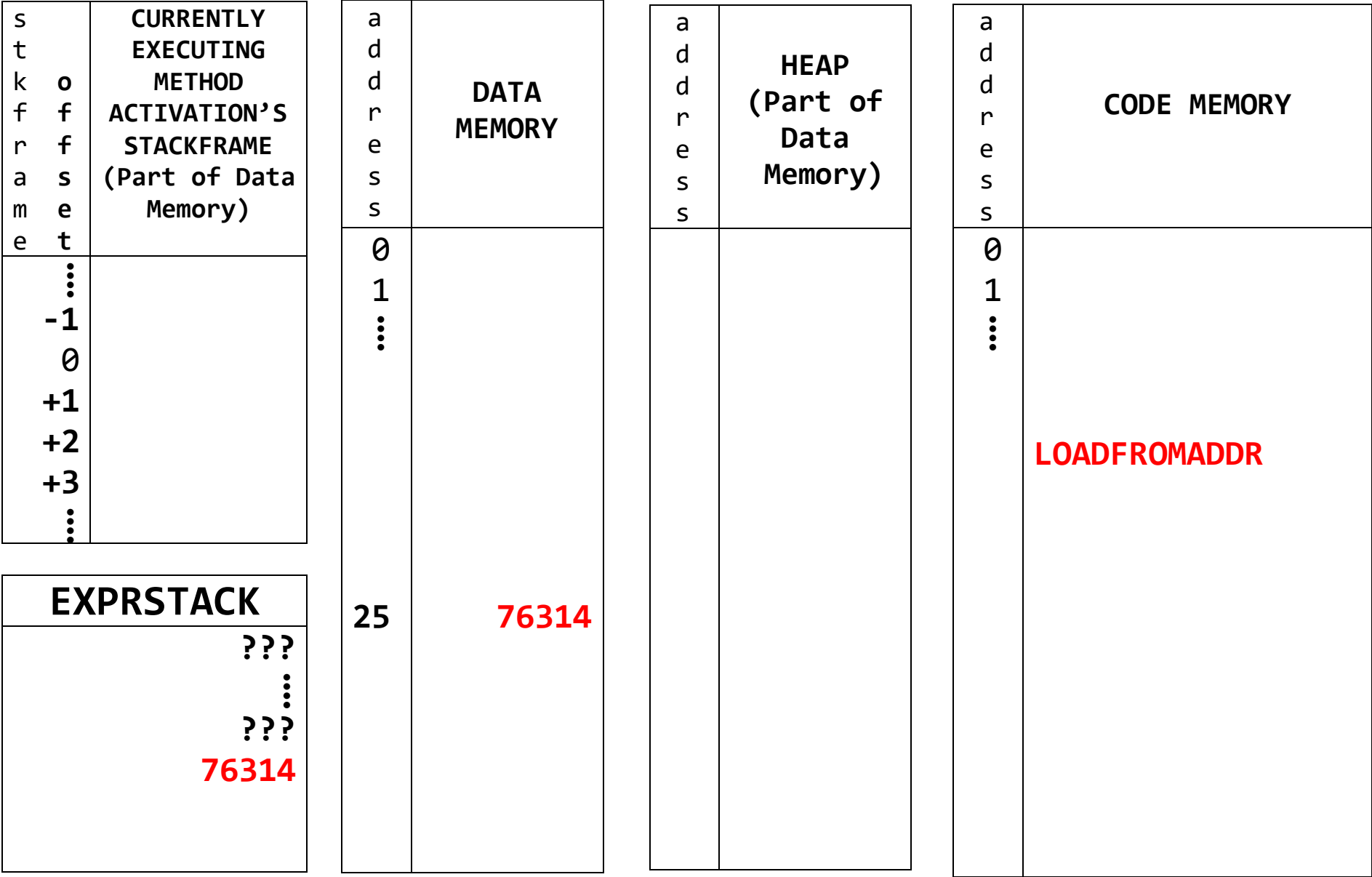
AFTER execution of: **PUSHLOCADDR 3**



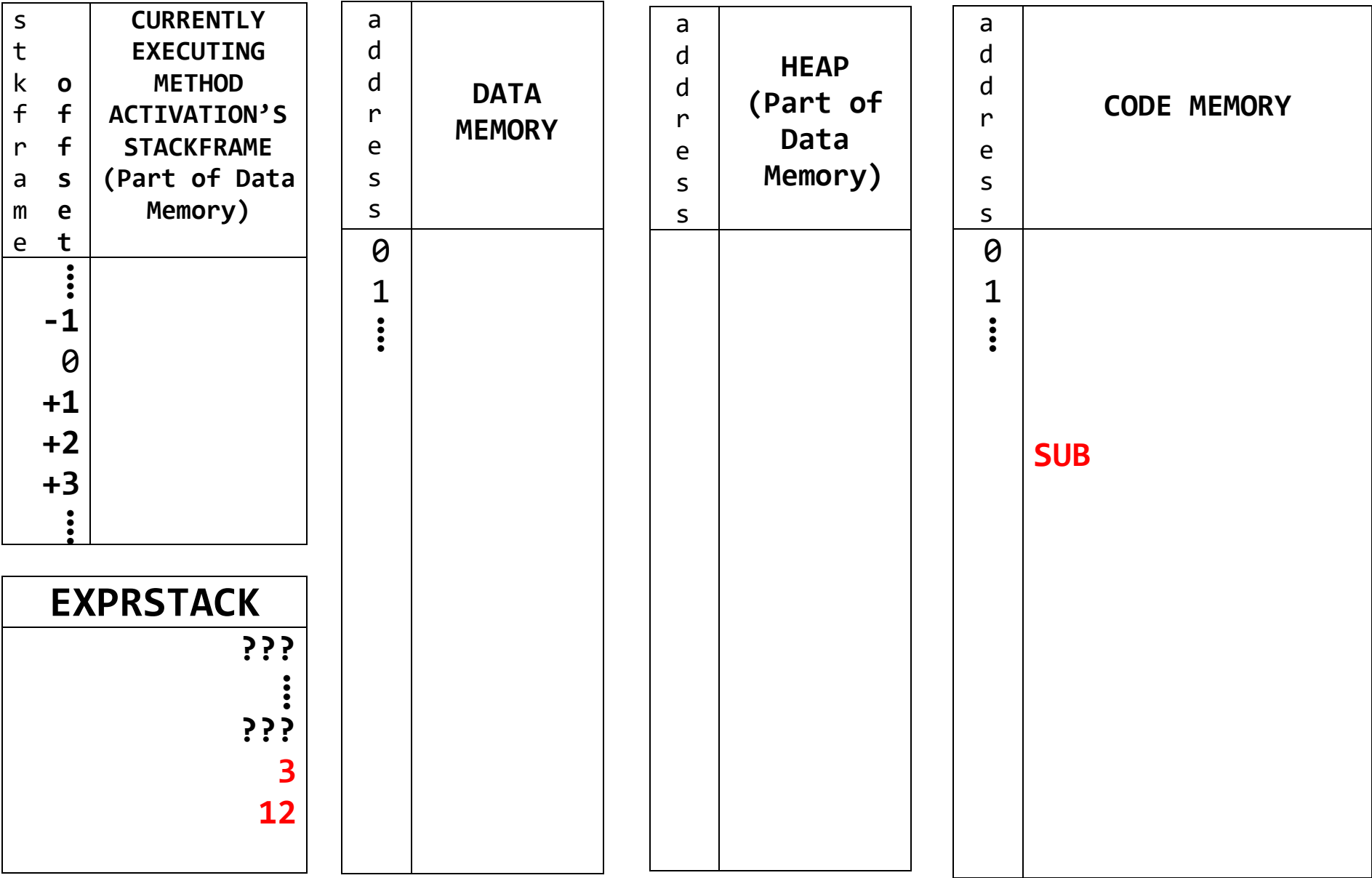
BEFORE execution of: **LOADFROMADDR**



AFTER execution of: **LOADFROMADDR**



BEFORE execution of: **SUB**



AFTER *execution of:* **SUB**

s t k o f f r a m e s t	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)
⋮ -1 0 +1 +2 +3 ⋮	

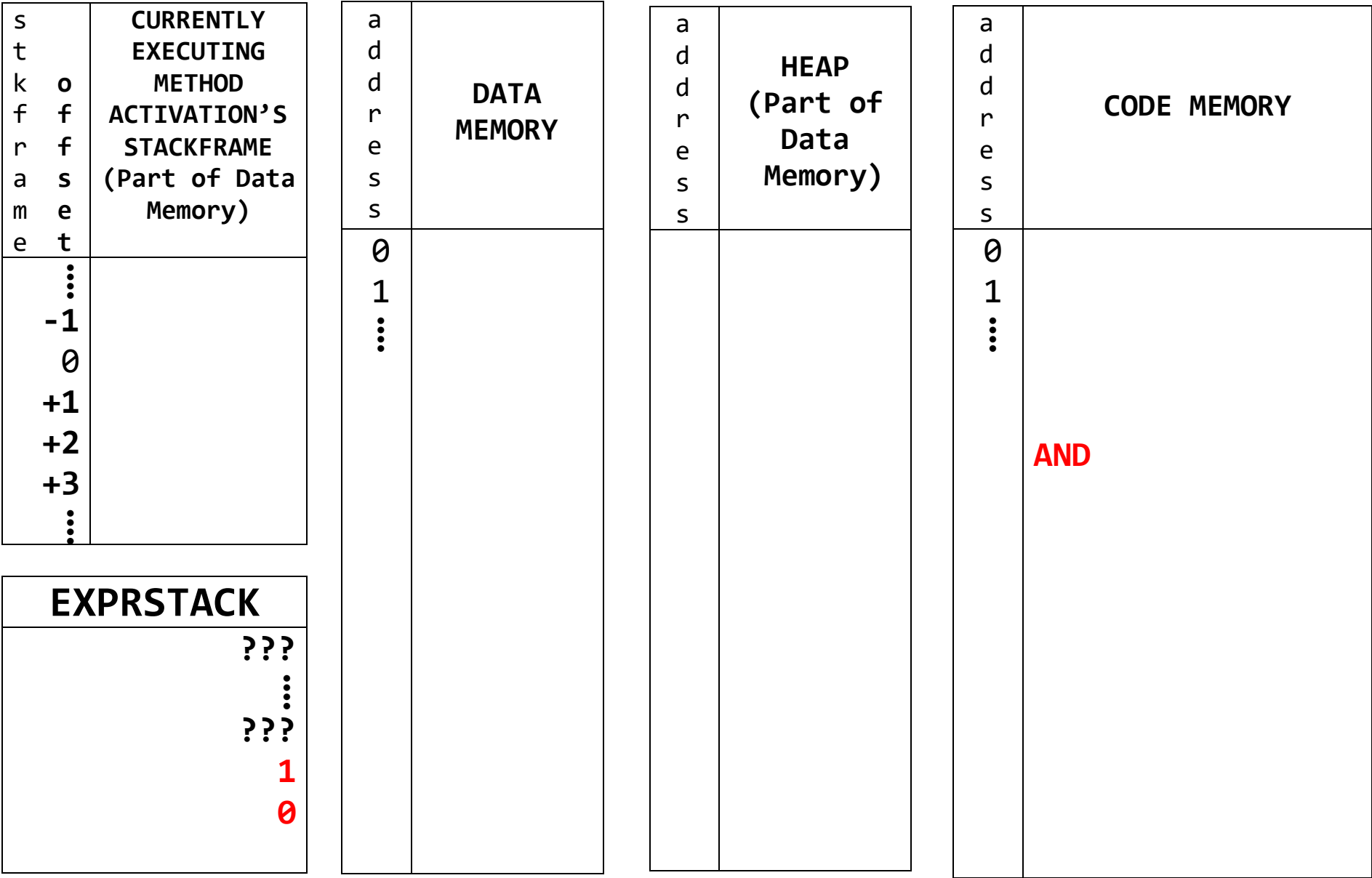
EXPRSTACK
??? ⋮ ??? -9

a d d r e s s	DATA MEMORY
0 1 ⋮	

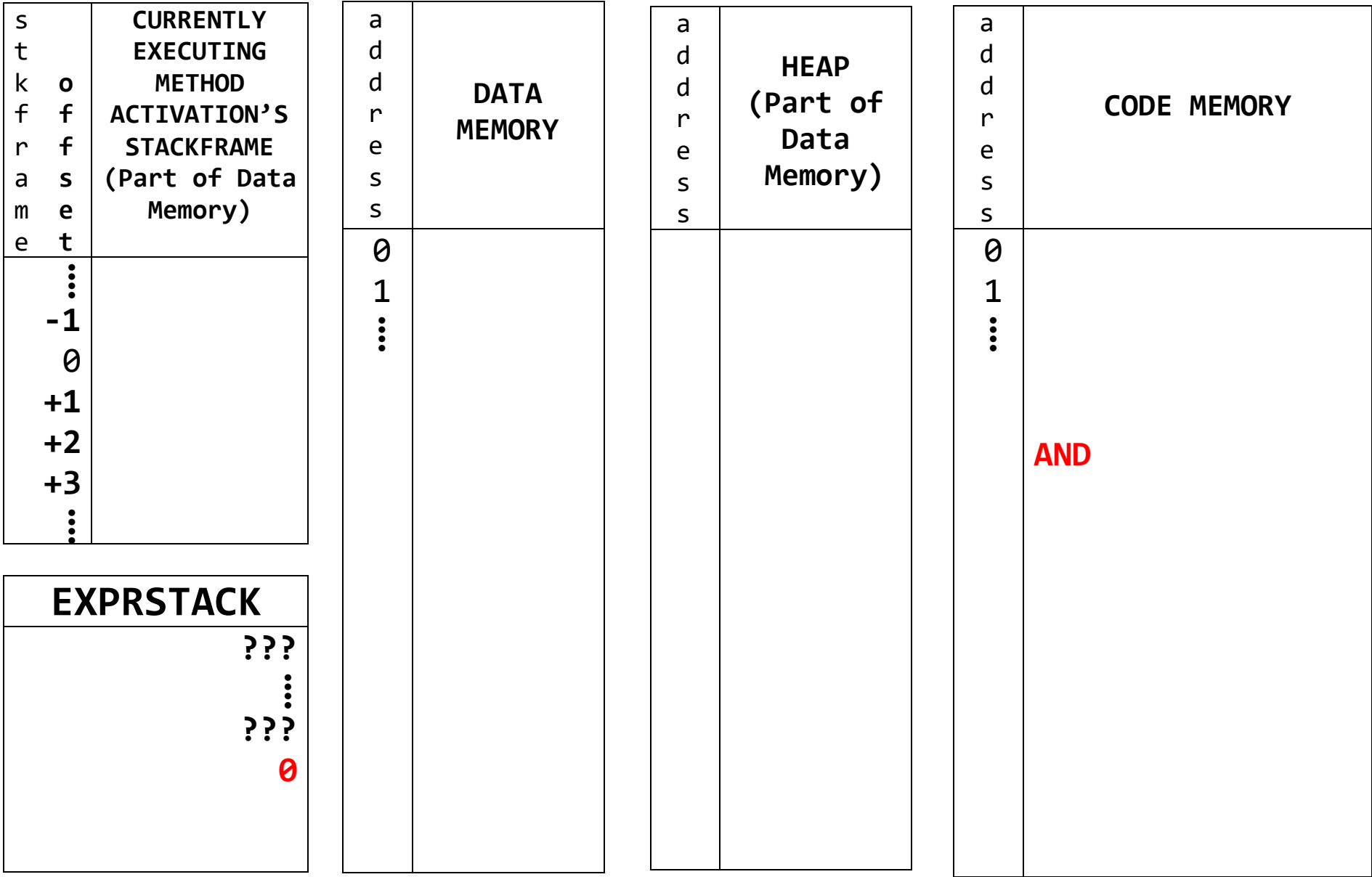
a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0 1 ⋮	SUB

BEFORE execution of: **AND**



AFTER execution of: AND



BEFORE execution of: LE (“Less than or Equal to”)

s t o k e t	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)
⋮ -1 0 +1 +2 +3 ⋮	

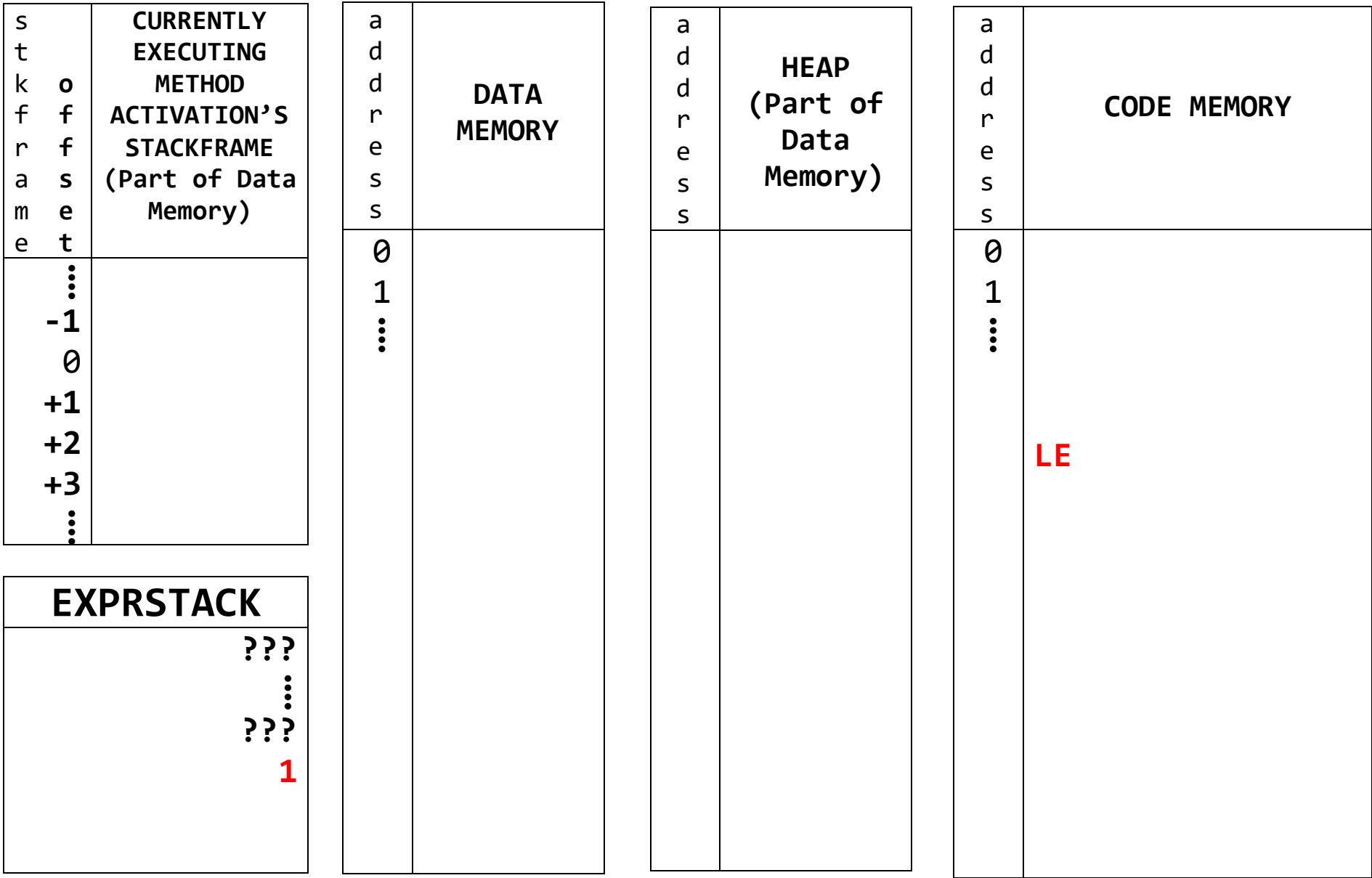
EXPRSTACK
??? ⋮ ??? 7 12

a d d r e s s	DATA MEMORY
0 1 ⋮	

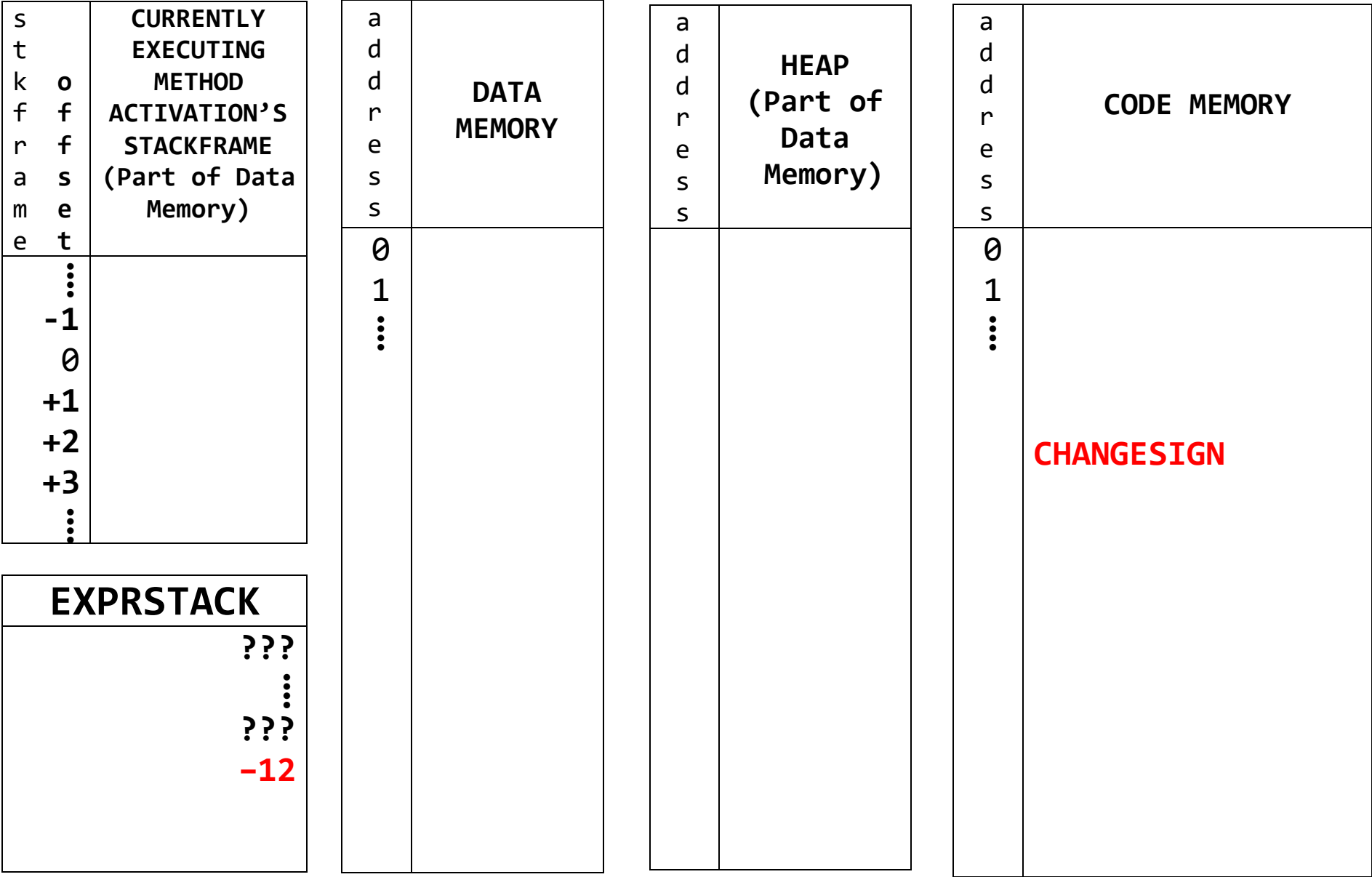
a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0 1 ⋮	LE

AFTER execution of: LE (“Less than or Equal to”)



BEFORE execution of: **CHANGESIGN**



AFTER *execution of:* **CHANGESIGN**

Stackframe Address	Currently Executing Method's Stackframe (Part of Data Memory)
...	
-1	
0	
+1	
+2	
+3	
...	

EXPRSTACK	
???	
⋮	
???	
12	

a d d r e s s	DATA MEMORY
0 1 ⋮	

a d d r e s s	HEAP (Part of Data Memory)

addresses	CODE MEMORY
0 1 ⋮	CHANGESIGN

BEFORE execution of: NOT

s t k o f f r a m e s	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)
⋮ -1 0 +1 +2 +3 ⋮	

EXPRSTACK
??? ⋮ ??? 1

a d d r e s s	DATA MEMORY
0 1 ⋮	

a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0 1 ⋮	NOT

AFTER execution of: NOT

stack offset address method	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)
⋮ -1 0 +1 +2 +3 ⋮	

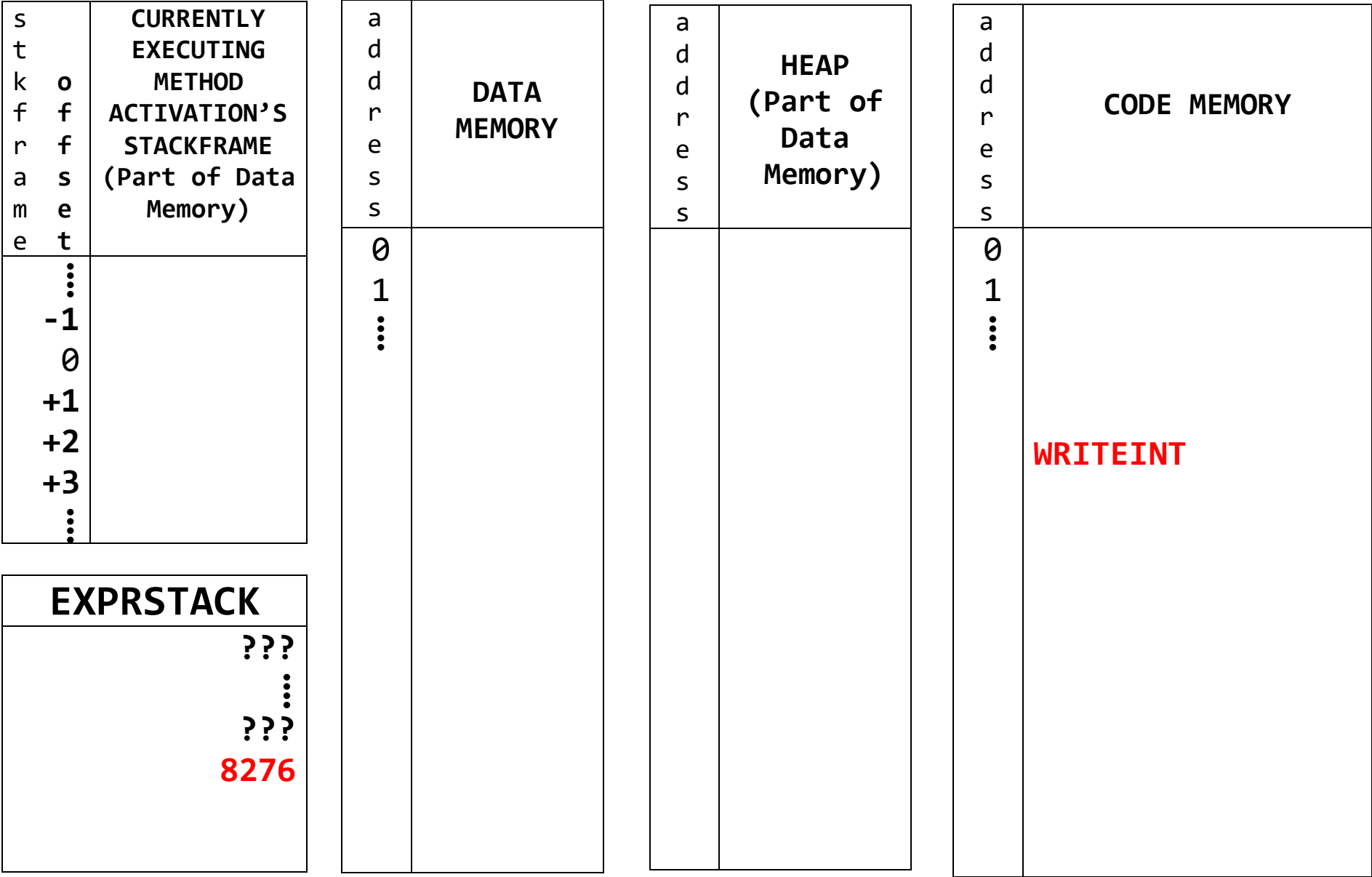
EXPRSTACK
??? ⋮ ??? 0

address	DATA MEMORY
0 1 ⋮	

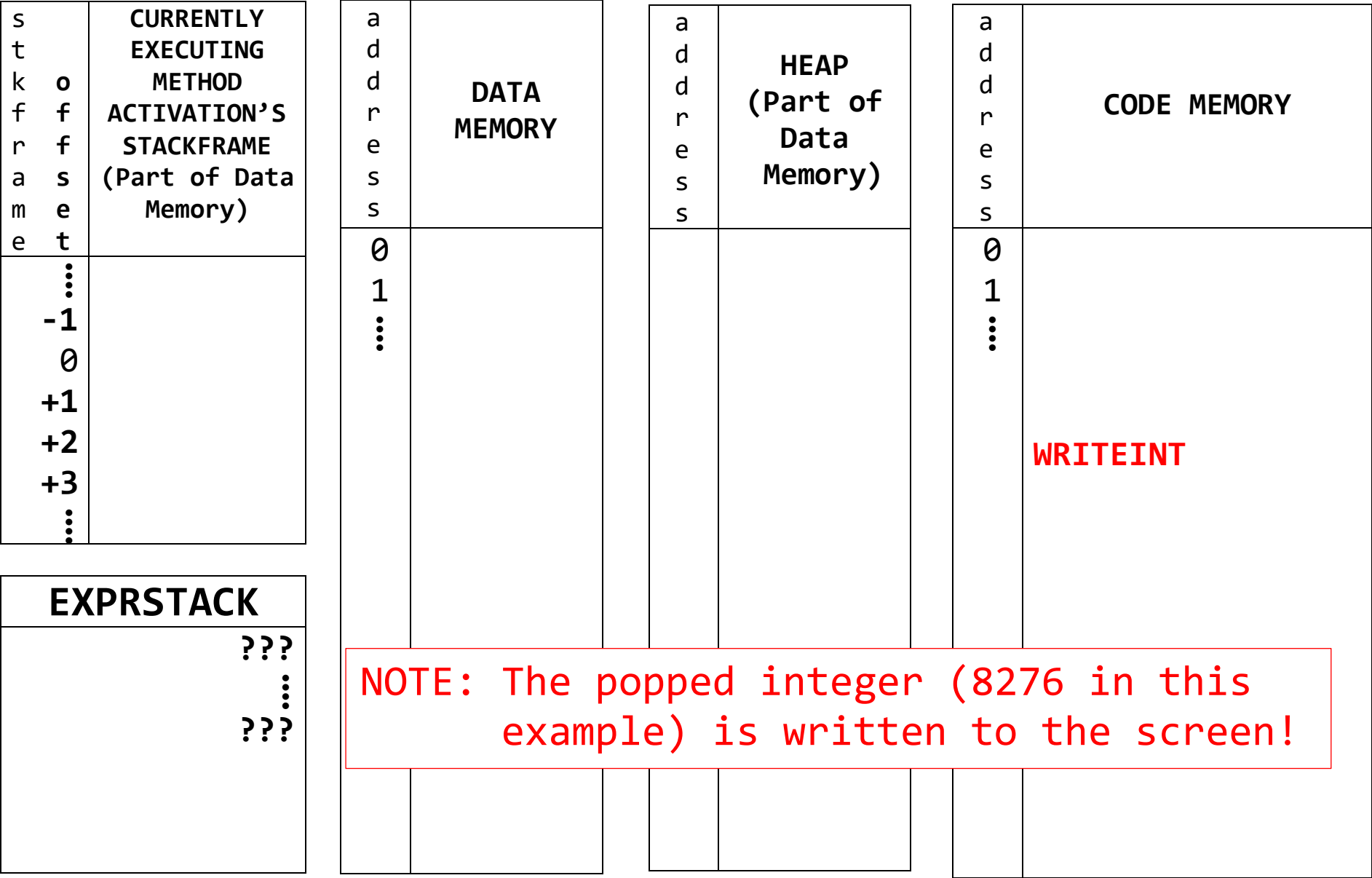
address	HEAP (Part of Data Memory)

address	CODE MEMORY
0 1 ⋮	NOT

BEFORE execution of: **WRITEINT**



AFTER execution of: **WRITEINT**



BEFORE execution of: **DISCARDVALUE**

s t o c k f r a m e s	CURRENTLY EXECUTING METHOD ACTIVATION'S STACKFRAME (Part of Data Memory)
⋮ -1 0 +1 +2 +3 ⋮	

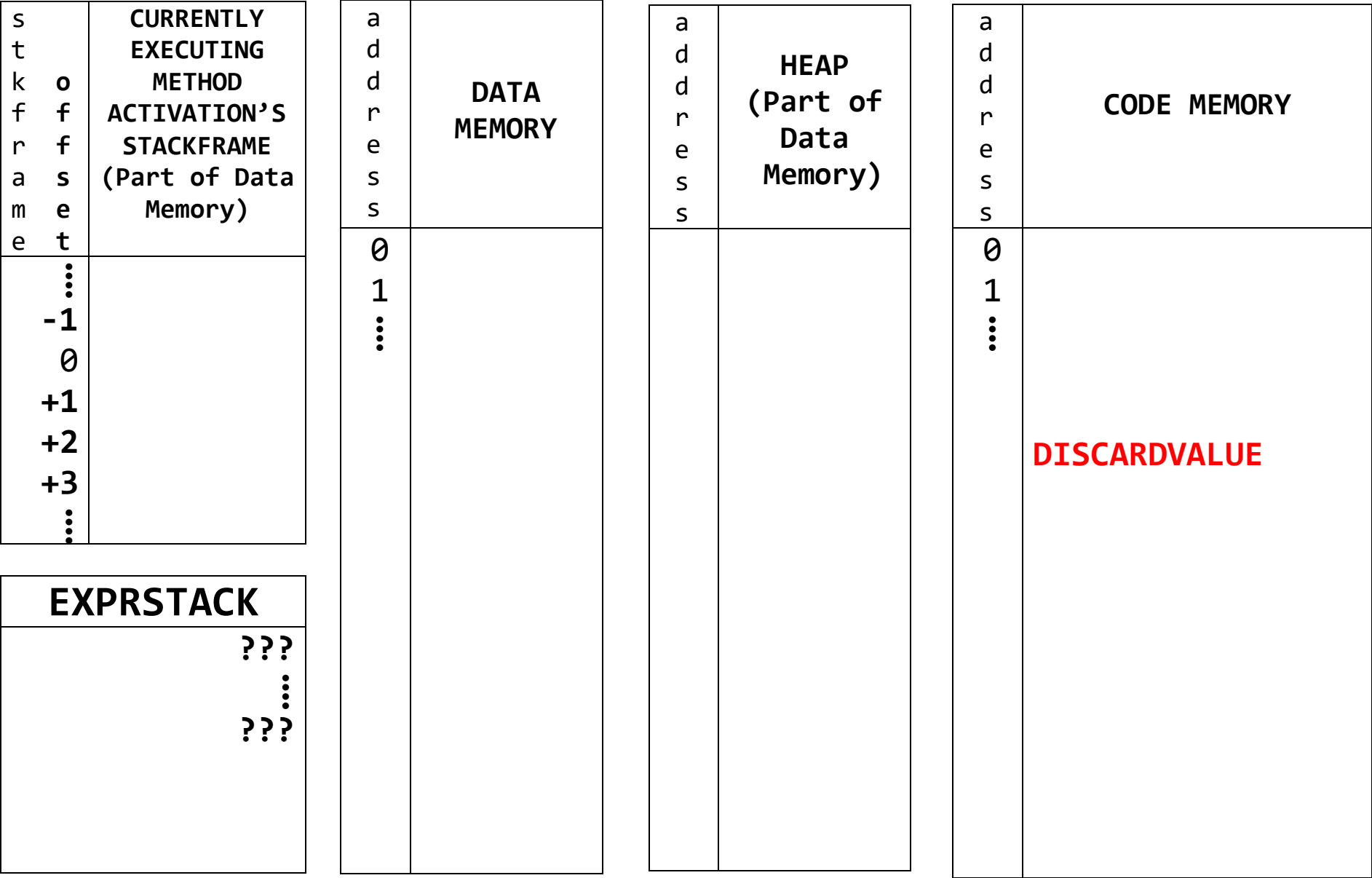
EXPRSTACK
??? ⋮ ??? 7263

a d d r e s s	DATA MEMORY
0 1 ⋮	

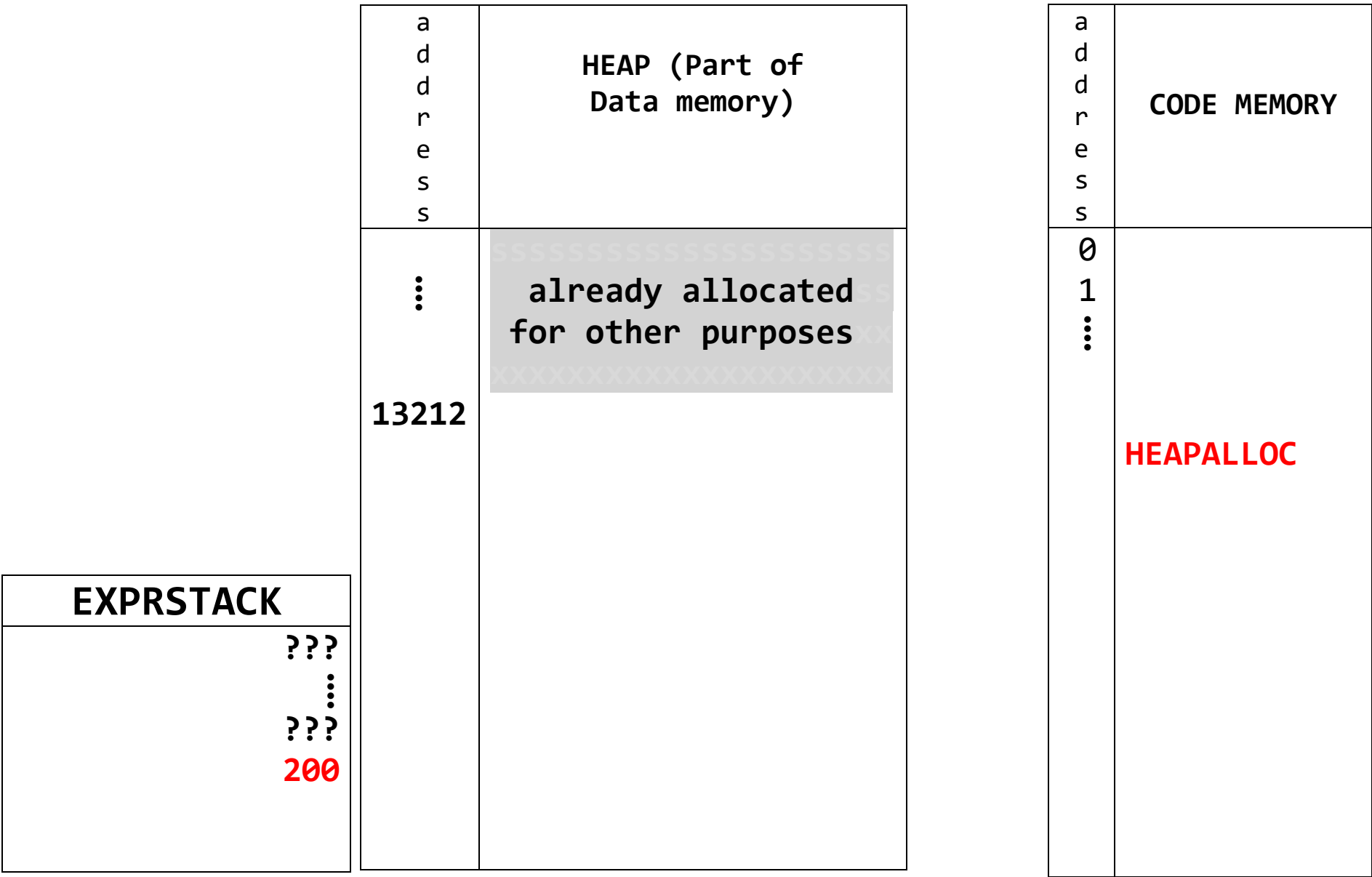
a d d r e s s	HEAP (Part of Data Memory)

a d d r e s s	CODE MEMORY
0 1 ⋮	DISCARDVALUE

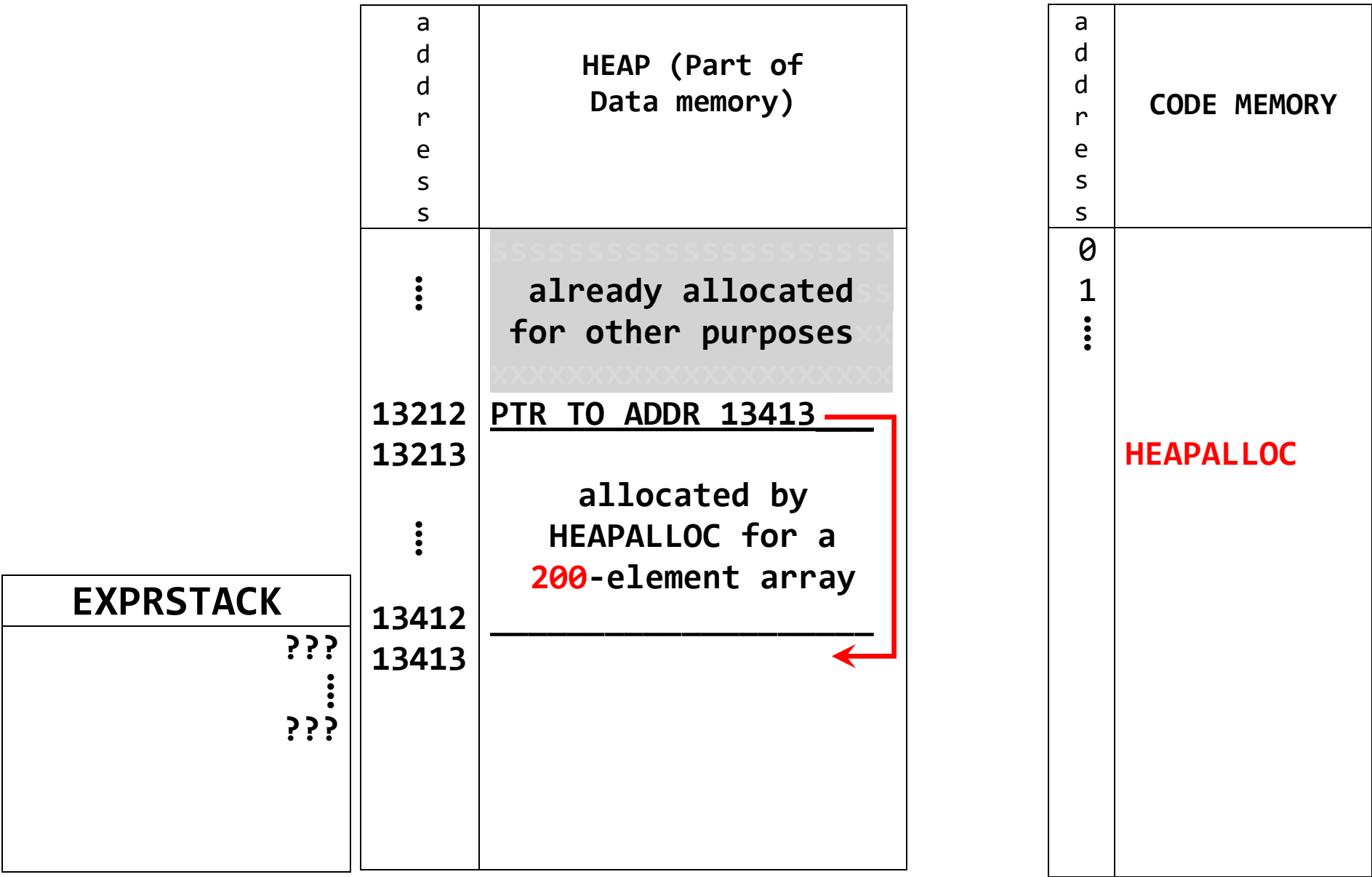
AFTER *execution of:* **DISCARDVALUE**



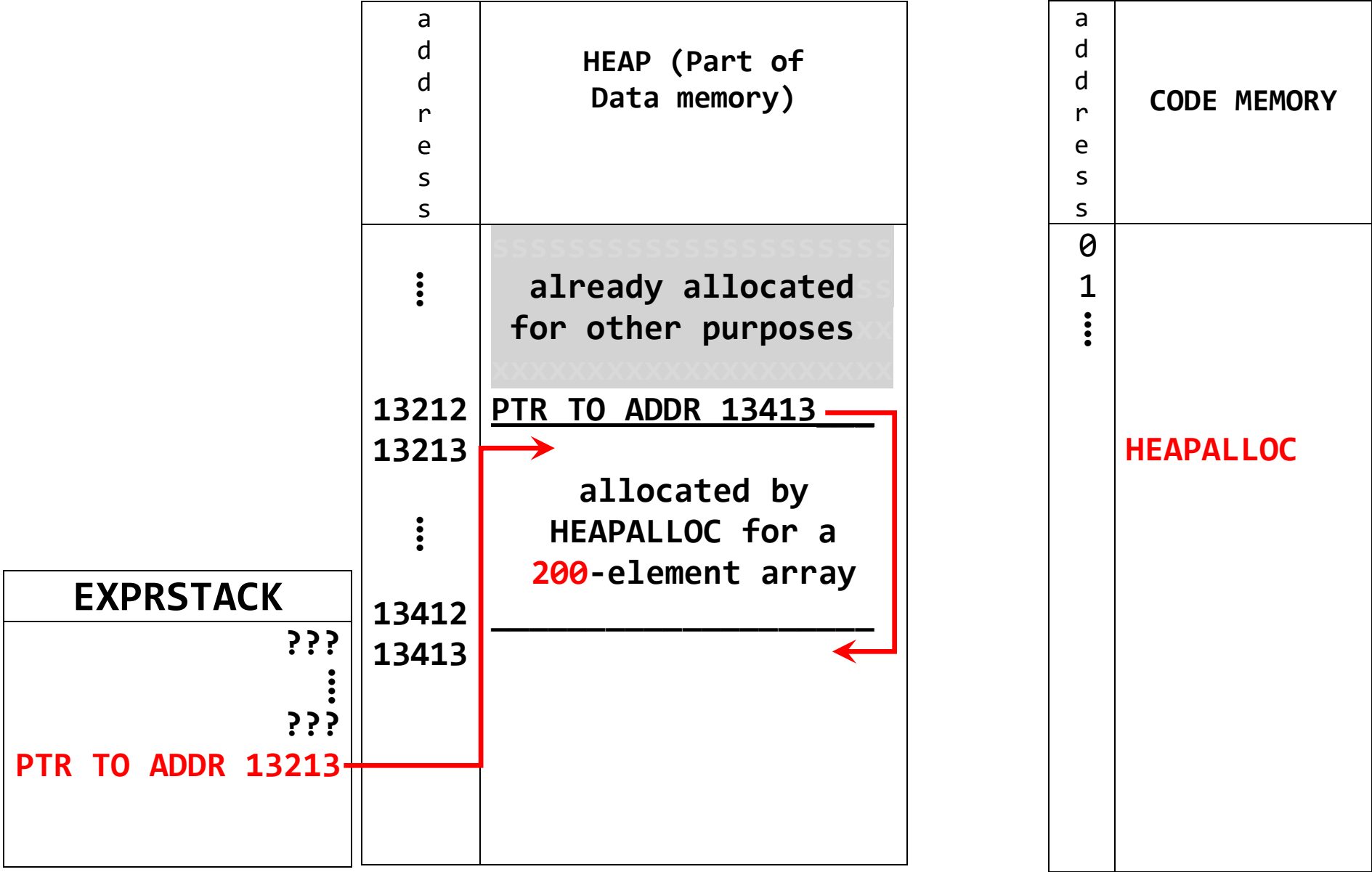
BEFORE execution of: **HEAPALLOC**



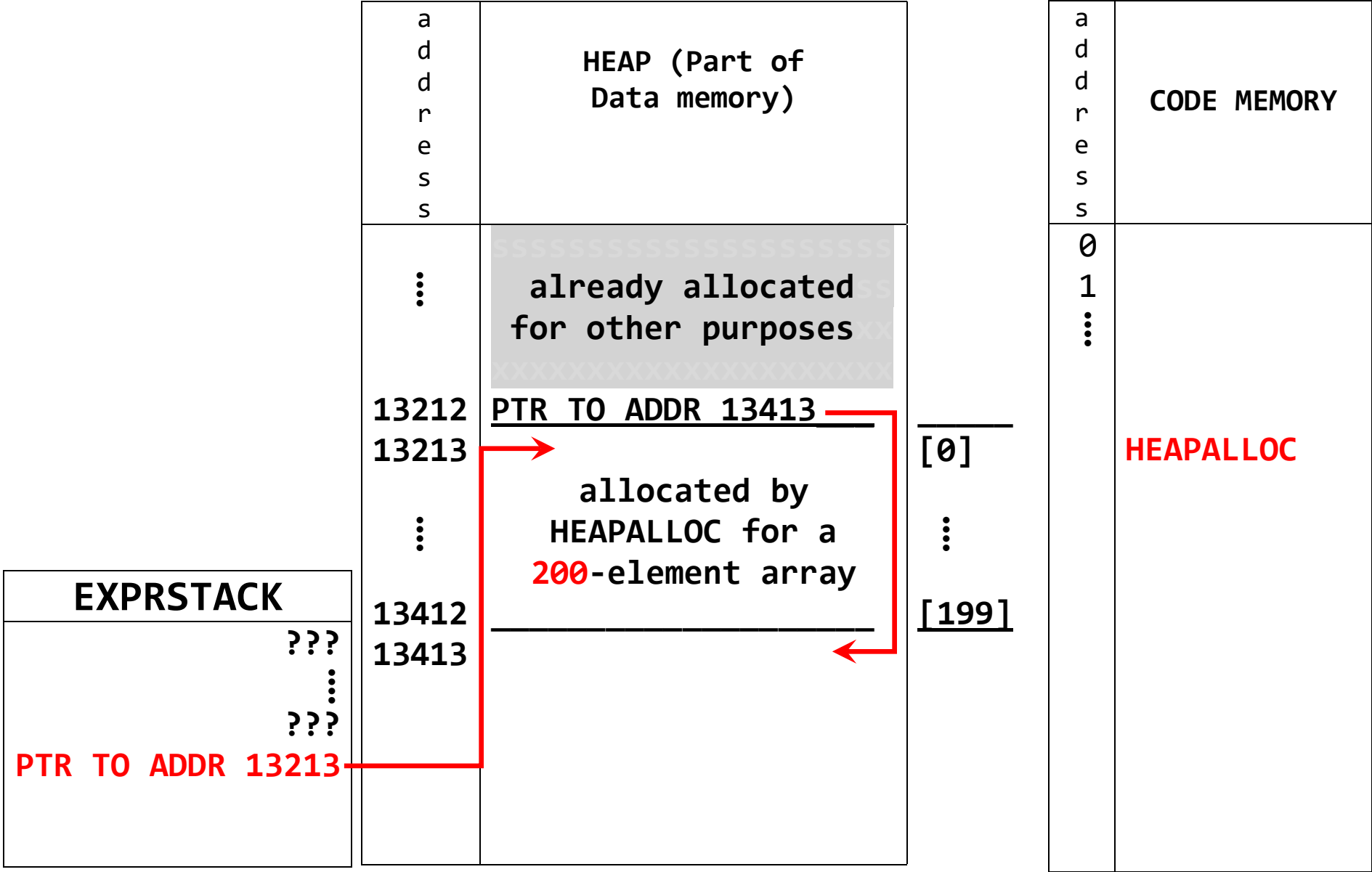
DURING execution of: HEAPALLOC



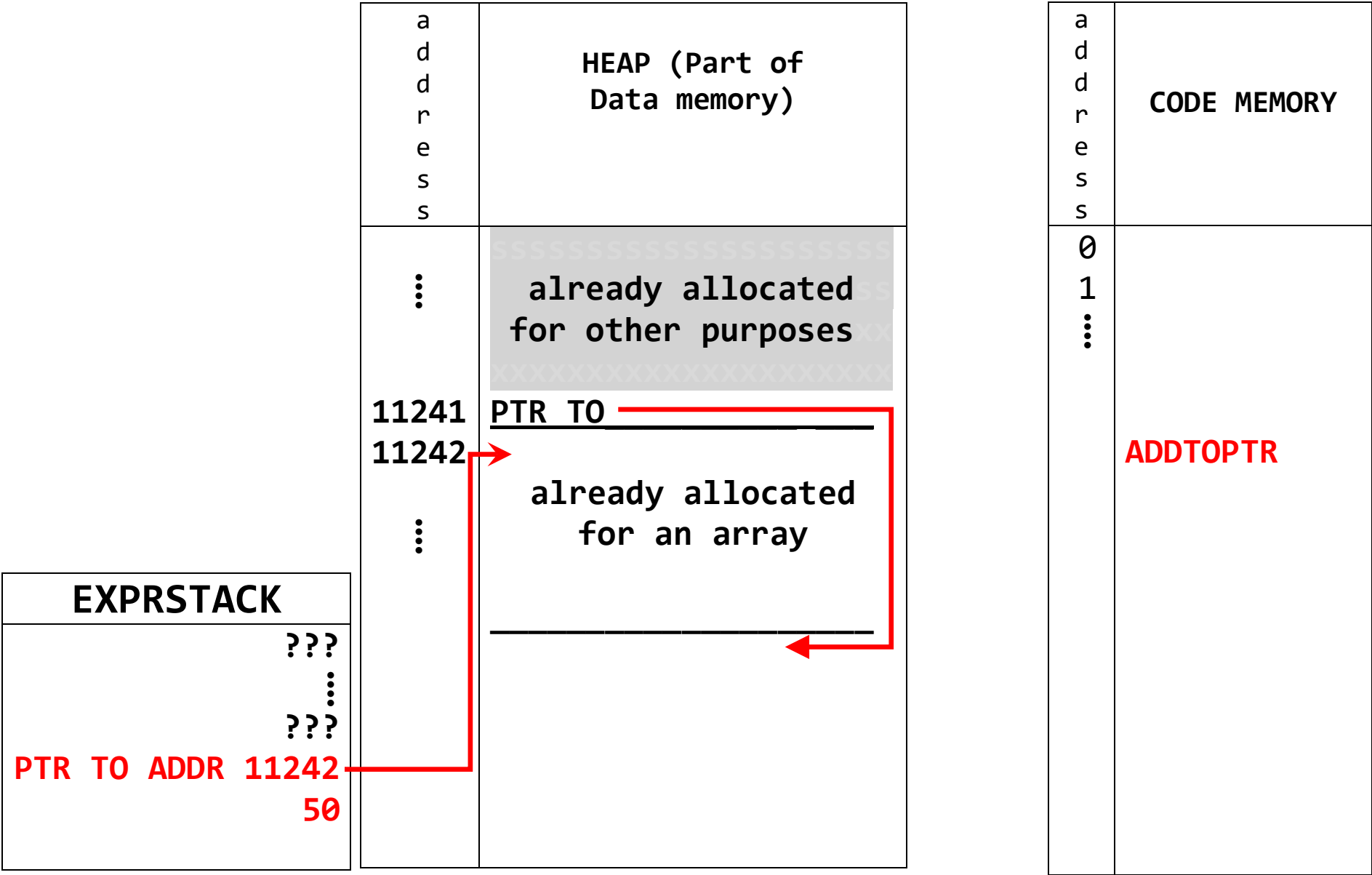
AFTER execution of: HEAPALLOC



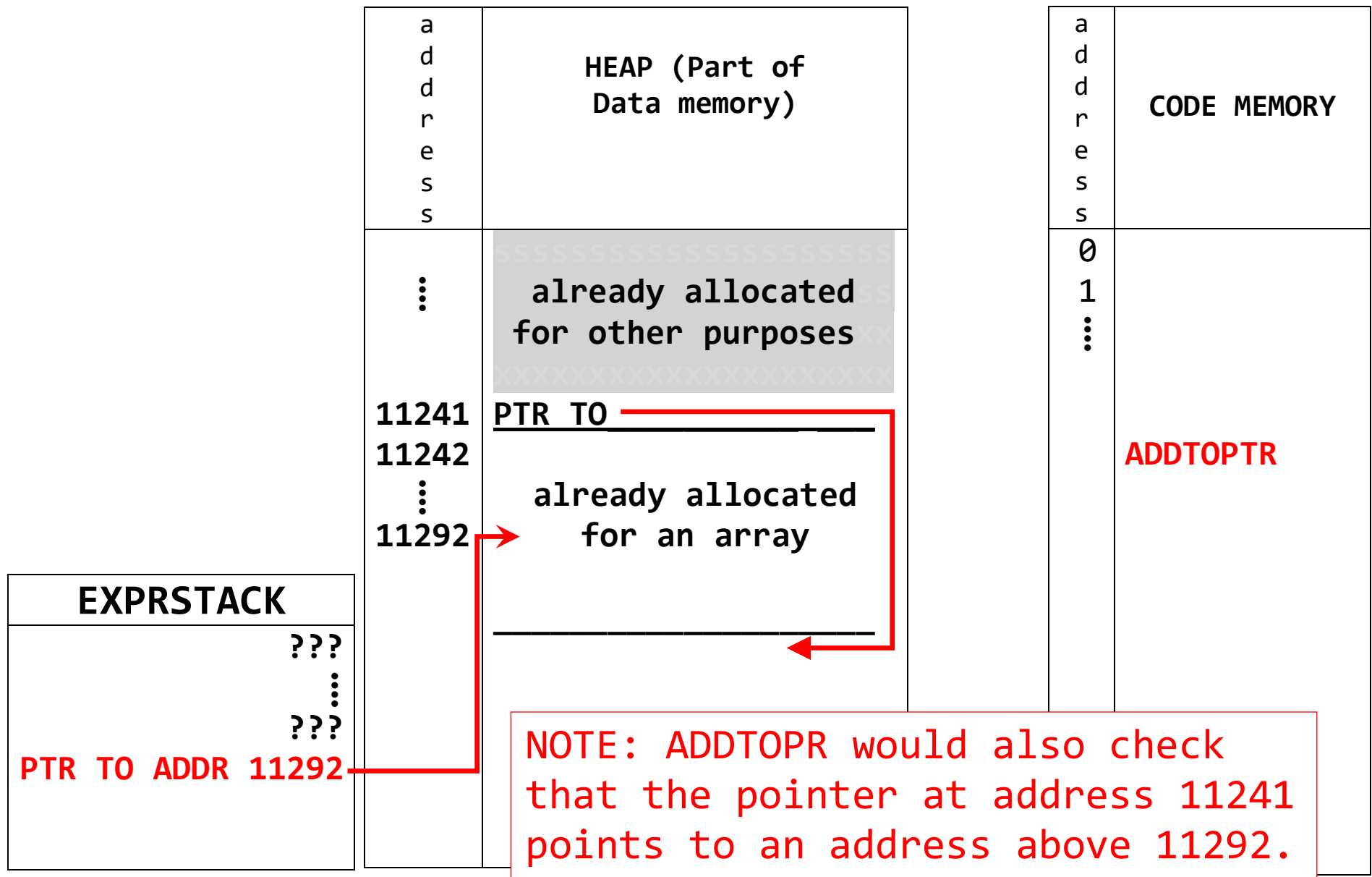
AFTER execution of: HEAPALLOC



BEFORE execution of: ADDTOPTR



AFTER *execution of:* ADDTOPTR



BEFORE execution of: **JUMP 87**

PC 34

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

a d d r e s s	CODE MEMORY
0	
1	
⋮	
33	JUMP 87
⋮	
87	

AFTER execution of: **JUMP 87**

PC

87

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

a d d r e s s	CODE MEMORY
0	
1	
⋮	
33	JUMP 87
⋮	
87	

BEFORE execution of: **JUMPONFALSE 77** (Example 1)

PC 52

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

EXPRSTACK	
	???
	⋮
	???
	1

a d d r e s s	CODE MEMORY
0	JUMPONFALSE 77
1	
⋮	
51	
⋮	
77	

AFTER execution of: **JUMPONFALSE 77** (Example 1)

PC 52

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

EXPRSTACK
???
⋮
???

a d d r e s s	CODE MEMORY
0	JUMPONFALSE 77
1	
⋮	
51	
⋮	
77	

BEFORE execution of: **JUMPONFALSE 77** (Example 2)

PC 52

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

EXPRSTACK	
	???
	⋮
	???
	0

a d d r e s s	CODE MEMORY
0	JUMPONFALSE 77
1	
⋮	
51	
⋮	
77	

AFTER execution of: **JUMPONFALSE 77** (Example 2)

PC

77

NOTE: The PC register is the program counter; it is used to store the code memory address of the next instruction to be executed.

EXPRSTACK

???
⋮
???

a
d
d
r
e
s
s

CODE MEMORY

0
1
⋮
51
⋮
77

JUMPONFALSE 77