

## More Old Exam Questions

1.[5 pts.] The language TeenyJ is defined like TinyJ except that the syntax of <expr1> is given by:

<expr1> ::= UNSIGNEDINT | new int '[' <expr3> ']' { '[' ']' }

Suppose the Parser class you completed for TinyJ Assignment 1 is to be modified so that it will parse TeenyJ programs (instead of TinyJ programs). Show how you would complete the following parsing method for <expr1>. (No code generation is expected.) [3 solutions are given on pp. 13-14.](#)

```
private static void expr1() throws SourceFileErrorException
{
    TJ.output.printSymbol(NTexpr1);
    TJ.output.incTreeDepth();
}
```

### Solution to Problem 2:

0:	PUSHSTATADDR	0	
1:	PUSHNUM	315	
2:	HEAPALLOC		
3:	SAVETOADDR		
4:	INITSTKFRM	1	
5:	PUSHLOCADDR	1	
6:	PUSHNUM	19	
7:	PASSPARAM		
8:	CALLSTATMETHOD	26	
9:	SAVETOADDR		
10:	PUSHSTATADDR	0	
11:	LOADFROMADDR		
12:	PUSHNUM	271	
13:	ADDTOPTR		
14:	PUSHLOCADDR	1	
15:	LOADFROMADDR		
16:	SAVETOADDR		
17:	WRITESTRING	1	9
18:	PUSHSTATADDR	0	
19:	LOADFROMADDR		
20:	PUSHNUM	271	
21:	ADDTOPTR		
22:	LOADFROMADDR		
23:	WRITEINT		
24:	WRITELNOP		
25:	STOP		
26:	INITSTKFRM	0	
27:	PUSHLOCADDR	-2	
28:	LOADFROMADDR		
29:	PUSHNUM	3	
30:	LT		
31:	JUMPONFALSE	35	
32:	PUSHNUM	0	
33:	RETURN	1	
34:	JUMP	45	
35:	PUSHLOCADDR	-2	
36:	LOADFROMADDR		
37:	PUSHLOCADDR	-2	
38:	LOADFROMADDR		
39:	PUSHNUM	1	
40:	SUB		
41:	PASSPARAM		
42:	CALLSTATMETHOD	26	
43:	SUB		
44:	RETURN	1	

```
TJ.output.decTreeDepth();
}
```

2.[10 pts.] Complete the table below the following program to show the TinyJ virtual machine instructions that should be generated by TJasn.TJ (after completion of TinyJ Assignment 2) for this TinyJ program.

```
class ExamQ {
    static int b[] = new int[315];

    public static void main (String args[])
    {
        int i = g(19);
        b[271] = i;
        System.out.print("g(19) is ");
        System.out.println(b[271]);
    }

    static int g(int m)
    {
        if (m < 3) return 0;
        else return m-g(m-1);
    }
}
```

[The solution is given on p. 1.](#)

0: PUSHSTATADDR 0	16: _____	32: _____
1: PUSHNUM 315	17: _____	33: _____
2: _____	18: _____	34: JUMP 45
3: SAVETOADDR	19: _____	35: _____
4: _____	20: _____	36: _____
5: _____	21: _____	37: _____
6: _____	22: _____	38: _____
7: _____	23: _____	39: _____
8: _____	24: _____	40: _____
9: _____	25: _____	41: _____
10: _____	26: _____	42: _____
11: _____	27: _____	43: _____
12: _____	28: _____	44: RETURN 1
13: _____	29: _____	
14: _____	30: _____	
15: _____	31: _____	

**Hint:** Among the 40 instructions you are asked to write, there are 7 LOADFROMADDR instructions, 6 PUSHNUM instructions, 5 PUSHLOCADDR instructions, 2 each of the ADDTOPTR, CALLSTATMETHOD, INITSTKFRM, PASSPARAM, PUSHSTATADDR, SAVETOADDR, and SUB instructions, and 1 each of the HEAPALLOC, JUMPONFALSE, LT, RETURN, STOP, WRITEINT, WRITELNOP, and WRITESTRING instructions.

While reading this page and the next, you should refer back when necessary to the pages of <https://euclid.cs.gc.cuny.edu/316/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf> that specify the effects of executing each VM instruction.

#### Comments on Problem 2 Regarding the Translation of the Statements

`b[271] = i;    and    System.out.println(b[271]);`

**Note:** The **EXPRSTACK** column on the right shows the items on the expression evaluation stack immediately *after* each VM instruction has been executed. The stack grows downwards—when more than one item is on the stack the first line below the word **EXPRSTACK** refers to the *bottom* item on the stack.

`b[271] = i;` is translated into the seven VM instructions that are shown on the left below. These instructions are put into code memory at addresses 10 – 16, as indicated on p. 1.

PUSHSTATADDR 0	Pushes <b>pointer to b</b> .	<b>EXPRSTACK</b> ptr to b
LOADFROMADDR	Pops <b>pointer to b</b> . Pushes the <b>pointer to b[0]</b> that is stored in b's location.	<b>EXPRSTACK</b> ptr to b[0]
PUSHNUM 271	Pushes the <b>integer 271</b> .	<b>EXPRSTACK</b> ptr to b[0] 271
ADDTOPTR	Pops <b>271</b> and <b>pointer to b[0]</b> . Pushes <b>(pointer to b[0]) + 271</b> (i.e., <b>pointer to b[271]</b> ).	<b>EXPRSTACK</b> ptr to b[271]
PUSHLOCADDR 1	Pushes <b>pointer to i</b> .	<b>EXPRSTACK</b> ptr to b[271] ptr to i
LOADFROMADDR	Pops <b>pointer to i</b> . Pushes the <b>value stored in i's location</b> (i.e., the <b>value of i</b> ).	<b>EXPRSTACK</b> ptr to b[271] value of i
SAVETOADDR	Pops <b>value of i</b> and <b>pointer to b[271]</b> . Saves <b>value of i</b> into the location of <b>b[271]</b> .	<b>EXPRSTACK</b> <i>is empty</i>

`System.out.println(b[271]);` is translated into the seven VM instructions that are shown on the left below. These instructions are put into code memory at addresses 18 – 24, as indicated on p. 1.

PUSHSTATADDR 0	Pushes <b>pointer to b</b> .	<b>EXPRSTACK</b> ptr to b
LOADFROMADDR	Pops <b>pointer to b</b> . Pushes the <b>pointer to b[0]</b> that is stored in b's location.	<b>EXPRSTACK</b> ptr to b[0]
PUSHNUM 271	Pushes the <b>integer 271</b> .	<b>EXPRSTACK</b> ptr to b[0] 271
ADDTOPTR	Pops <b>271</b> and <b>pointer to b[0]</b> . Pushes <b>(pointer to b[0]) + 271</b> (i.e., <b>pointer to b[271]</b> ).	<b>EXPRSTACK</b> ptr to b[271]
LOADFROMADDR	Pops <b>pointer to b[271]</b> . Pushes the <b>value stored in b[271]'s location</b> (i.e., the <b>value of b[271]</b> ).	<b>EXPRSTACK</b> value of b[271]
WRITEINT	Pops <b>value of b[271]</b> . Writes the value on the screen.	<b>EXPRSTACK</b> <i>is empty</i>
WritelnOP	Writes a newline to the screen.	<b>EXPRSTACK</b> <i>is empty</i>

Further problems to test your understanding:

3. Suppose we delete the line `static int b[] = new int[315];` from the TinyJ program of problem 2 but insert a line `int b[] = new int[536];` at the beginning of the body of `main`. (Thus `b` would become the first local variable of `main`, and `i` would become the second local variable of `main` rather than the first local variable.) How would the 14 instructions shown on the previous page change?

Answer: `PUSHLOCADDR 1` would be changed to `PUSHLOCADDR 2`.

Each occurrence of `PUSHSTATADDR 0` would be changed to `PUSHLOCADDR 1`.

4. Suppose that the first variable declaration in a certain TinyJ program is `static int b[][][];` Suppose also that this variable `b` is used in the following statement later in the program: `System.out.print(b[7][29][5]);` What TinyJ VM instructions would the TinyJ compiler translate the latter statement into? [Note: Although Exam 2 may have questions relating to arrays, Exam 2 will not have any question such as this one that involves an indexed variable with more than one actual index. However, there may be a question on the Final Exam that involves an indexed variable with more than one index.]

**Answer to problem 4, and explanation of the generated instructions:**

<code>PUSHSTATADDR 0</code>	Pushes <b>pointer to b</b> .	<b>EXPRSTACK</b> <b>ptr to b</b>
<code>LOADFROMADDR</code>	Pops <b>pointer to b</b> . Pushes the <b>pointer to b[0]</b> that is stored in <b>b's location</b> .	<b>EXPRSTACK</b> <b>ptr to b[0]</b>
<code>PUSHNUM 7</code>	Pushes the <b>integer 7</b> .	<b>EXPRSTACK</b> <b>ptr to b[0]</b> <b>7</b>
<code>ADDTOPTR</code>	Pops <b>7</b> and <b>pointer to b[0]</b> . Pushes <b>(pointer to b[0]) + 7</b> (i.e., <b>pointer to b[7]</b> ).	<b>EXPRSTACK</b> <b>ptr to b[7]</b>
<code>LOADFROMADDR</code>	Pops <b>pointer to b[7]</b> . Pushes the <b>pointer to b[7][0]</b> that is stored in <b>b[7]'s location</b> .	<b>EXPRSTACK</b> <b>ptr to b[7][0]</b>
<code>PUSHNUM 29</code>	Pushes the <b>integer 29</b> .	<b>EXPRSTACK</b> <b>ptr to b[7][0]</b> <b>29</b>
<code>ADDTOPTR</code>	Pops <b>29</b> and <b>pointer to b[7][0]</b> . Pushes <b>(pointer to b[7][0]) + 29</b> (i.e., <b>pointer to b[7][29]</b> ).	<b>EXPRSTACK</b> <b>ptr to b[7][29]</b>
<code>LOADFROMADDR</code>	Pops <b>pointer to b[7][29]</b> . Pushes the <b>pointer to b[7][29][0]</b> that is stored in <b>b[7][29]'s location</b> .	<b>EXPRSTACK</b> <b>ptr to b[7][29][0]</b>
<code>PUSHNUM 5</code>	Pushes the <b>integer 5</b> .	<b>EXPRSTACK</b> <b>ptr to b[7][29][0]</b> <b>5</b>
<code>ADDTOPTR</code>	Pops <b>5</b> and <b>pointer to b[7][29][0]</b> . Pushes <b>(pointer to b[7][29][0]) + 5</b> (i.e., <b>pointer to b[7][29][5]</b> ).	<b>EXPRSTACK</b> <b>ptr to b[7][29][5]</b>
<code>LOADFROMADDR</code>	Pops <b>pointer to b[7][29][5]</b> . Pushes <b>value stored in b[7][29][5]'s location</b> (i.e., <b>value of b[7][29][5]</b> ).	<b>EXPRSTACK</b> <b>value of b[7][29][5]</b>
<code>WRITEINT</code>	Pops <b>value of b[7][29][5]</b> . Writes this value to the screen.	<b>EXPRSTACK</b> <i>is empty</i>

## Some Properties of the Code That is Generated by the TinyJ Compiler

When answering certain exam questions, it may be useful to remember that the code generated by the TinyJ compiler when it translates a TinyJ program has the following properties:

1. The code generated for any assignment statement consists of code whose execution will push a pointer to the target variable's location, followed by code whose execution will push what we want to store, followed by a SAVETOADDR instruction.
2. If the target variable of an assignment statement is not an indexed variable, then the assignment statement is translated into code that begins with a PUSHSTATADDR or PUSHLOCADDR instruction that is not immediately followed by a LOADFROMADDR instruction.
3. With the exceptions of the PUSHSTATADDR and PUSHLOCADDR instructions referred to in item 2, every PUSHSTATADDR or PUSHLOCADDR instruction is immediately followed by a LOADFROMADDR instruction.
4. The code generated to push the value of an indexed variable onto EXPRSTACK consists of code whose execution will push a pointer to the indexed variable's location and an additional LOADFROMADDR instruction at the end.
5. The code generated for any method call includes a PASSPARAM instruction for each argument of the call, and includes a CALLSTATMETHOD instruction; if the call has arguments, then the CALLSTATMETHOD instruction is the next instruction after the last PASSPARAM.

Much further information about the generated code is provided in the document

<https://euclid.cs.gc.cuny.edu/316/Slides/Memory-allocation-VM-instruction-set-and-hints-for-asn-2.pdf>

that has been discussed in class—especially in the section **Code Generation Rules Used by the TinyJ Compiler** and the sections relating to `whileStmt()` and `ifStmt()`.

## More Hand-Translation Examples

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[Note: You can also make up your own hand-translation examples: If X.java is any valid TinyJ program, then the correct solution to the problem of translating X.java can be obtained by running my solution to TinyJ Assignment 2 with X.java as the input file.]

- (a) Suppose `Instruction.getNextCodeAddress() == 35` when a correct solution to TinyJ Assignment 2 begins to translate the following two methods. What code is generated?

```
static void m()
{
    int x = 12, y = 9;
    System.out.print(p(17, y, x+5));
}

static int p (int a, int b, int c)
{
    int u = a - b;
    return c - u;
}
```

### SOLUTION:

35:	INITSTKFRM	2
36:	PUSHLOCADDR	1
37:	PUSHNUM	12
38:	SAVETOADDR	
39:	PUSHLOCADDR	2
40:	PUSHNUM	9
41:	SAVETOADDR	
42:	PUSHNUM	17
43:	PASSPARAM	
44:	PUSHLOCADDR	2
45:	LOADFROMADDR	
46:	PASSPARAM	
47:	PUSHLOCADDR	1
48:	LOADFROMADDR	
49:	PUSHNUM	5
50:	ADD	
51:	PASSPARAM	
52:	CALLSTATMETHOD	55
53:	WRITEINT	
54:	RETURN	0
55:	INITSTKFRM	1
56:	PUSHLOCADDR	1
57:	PUSHLOCADDR	-4
58:	LOADFROMADDR	
59:	PUSHLOCADDR	-3
60:	LOADFROMADDR	
61:	SUB	
62:	SAVETOADDR	
63:	PUSHLOCADDR	-2
64:	LOADFROMADDR	
65:	PUSHLOCADDR	1
66:	LOADFROMADDR	
67:	SUB	
68:	RETURN	3

(b) An example involving arrays:

```
class ArrayTest {
    static int b[] = new int[10];

    public static void main (String args[])
    {
        int a = 1;
        b[3] = a;
        System.out.println(b[3]+a);

        b = new int[5];

        int c[][] = new int [7][];
        c[4] = b;
    }
}
```

What would a correct solution to TinyJ Assignment 2 translate this into?

SOLUTION:

0:	PUSHSTATADDR	0
1:	PUSHNUM	10
2:	HEAPALLOC	
3:	SAVETOADDR	
4:	INITSTKFRM	2
5:	PUSHLOCADDR	1
6:	PUSHNUM	1
7:	SAVETOADDR	
8:	PUSHSTATADDR	0
9:	LOADFROMADDR	
10:	PUSHNUM	3
11:	ADDTOPTR	
12:	PUSHLOCADDR	1
13:	LOADFROMADDR	
14:	SAVETOADDR	
15:	PUSHSTATADDR	0
16:	LOADFROMADDR	
17:	PUSHNUM	3
18:	ADDTOPTR	
19:	LOADFROMADDR	
20:	PUSHLOCADDR	1
21:	LOADFROMADDR	
22:	ADD	
23:	WRITEINT	
24:	WRITELNOP	
25:	PUSHSTATADDR	0
26:	PUSHNUM	5
27:	HEAPALLOC	
28:	SAVETOADDR	
29:	PUSHLOCADDR	2
30:	PUSHNUM	7
31:	HEAPALLOC	
32:	SAVETOADDR	
33:	PUSHLOCADDR	2
34:	LOADFROMADDR	
35:	PUSHNUM	4
36:	ADDTOPTR	
37:	PUSHSTATADDR	0
38:	LOADFROMADDR	
39:	SAVETOADDR	
40:	STOP	

(c) Example involving a while loop:

```
class Fall02a {
    static int a[] = new int[10];

    public static void main (String args[])
    {
        int x = 100;

        while (x > 10)
            x = f(x, 2);
    }

    static int f(int m, int n)
    {
        a[3] = m / n;
        return a[3];
    }
}
```

What would a correct solution to TinyJ Assignment 2 translate this into?

SOLUTION:

0:	PUSHSTATADDR	0
1:	PUSHNUM	10
2:	HEAPALLOC	
3:	SAVETOADDR	
4:	INITSTKFRM	1
5:	PUSHLOCADDR	1
6:	PUSHNUM	100
7:	SAVETOADDR	
8:	PUSHLOCADDR	1
9:	LOADFROMADDR	
10:	PUSHNUM	10
11:	GT	
12:	JUMPONFALSE	22
13:	PUSHLOCADDR	1
14:	PUSHLOCADDR	1
15:	LOADFROMADDR	
16:	PASSPARAM	
17:	PUSHNUM	2
18:	PASSPARAM	
19:	CALLSTATMETHOD	23
20:	SAVETOADDR	
21:	JUMP	8
22:	STOP	
23:	INITSTKFRM	0
24:	PUSHSTATADDR	0
25:	LOADFROMADDR	
26:	PUSHNUM	3
27:	ADDTOPTR	
28:	PUSHLOCADDR	-3
29:	LOADFROMADDR	
30:	PUSHLOCADDR	-2
31:	LOADFROMADDR	
32:	DIV	
33:	SAVETOADDR	
34:	PUSHSTATADDR	0
35:	LOADFROMADDR	
36:	PUSHNUM	3
37:	ADDTOPTR	
38:	LOADFROMADDR	
39:	RETURN	2



(d) Another example involving a while loop:

```
class Fall02b {  
    static int a[] = new int [25];  
  
    public static void main (String args[])  
    {  
        a[5] = 900;  
        System.out.print(g(7));  
    }  
  
    static int g(int m)  
    {  
        int i = a[5];  
  
        while (i > 30)  
            i = i / m;  
  
        return i;  
    }  
}
```

What would a correct solution to TinyJ Assignment 2 translate this into?

SOLUTION:

0:	PUSHSTATADDR	0
1:	PUSHNUM	25
2:	HEAPALLOC	
3:	SAVETOADDR	
4:	INITSTKFRM	0
5:	PUSHSTATADDR	0
6:	LOADFROMADDR	
7:	PUSHNUM	5
8:	ADDTOPTR	
9:	PUSHNUM	900
10:	SAVETOADDR	
11:	PUSHNUM	7
12:	PASSPARAM	
13:	CALLSTATMETHOD	16
14:	WRITEINT	
15:	STOP	
16:	INITSTKFRM	1
17:	PUSHLOCADDR	1
18:	PUSHSTATADDR	0
19:	LOADFROMADDR	
20:	PUSHNUM	5
21:	ADDTOPTR	
22:	LOADFROMADDR	
23:	SAVETOADDR	
24:	PUSHLOCADDR	1
25:	LOADFROMADDR	
26:	PUSHNUM	30
27:	GT	
28:	JUMPONFALSE	37
29:	PUSHLOCADDR	1
30:	PUSHLOCADDR	1
31:	LOADFROMADDR	
32:	PUSHLOCADDR	-2
33:	LOADFROMADDR	
34:	DIV	
35:	SAVETOADDR	
36:	JUMP	24

```

37:    PUSHLOCADDR      1
38:    LOADFROMADDR
39:    RETURN           1

```

(e) The next example involves if as well as while:

```

import java.util.Scanner;

class Spring99 {

    static Scanner input = new Scanner(System.in);
    static int x;

    public static void main (String args[])
    {
        int a;

        x = input.nextInt();
        if (x > 1 & x < 20000) {
            while (x <= 20000) {
                int b = 2;
                x = x * 3;
            }

            int c = x;
            System.out.print(c);
        }
    }
}

```

What would a correct solution to TinyJ Assignment 2 translate this into?

SOLUTION:

Note that the local variables b and c both have a stackframe offset of 2. At the point where c is declared, b no longer exists--b's scope is confined to the body of the while loop. Thus stackframe offset 2 can be reallocated to c.

```

0:    INITSTKFRM      2
1:    PUSHSTATADDR    0
2:    READINT
3:    SAVETOADDR
4:    PUSHSTATADDR    0
5:    LOADFROMADDR
6:    PUSHNUM         1
7:    GT
8:    PUSHSTATADDR    0
9:    LOADFROMADDR
10:   PUSHNUM         20000
11:   LT
12:   AND
13:   JUMPONFALSE     36
14:   PUSHSTATADDR    0
15:   LOADFROMADDR
16:   PUSHNUM         20000
17:   LE
18:   JUMPONFALSE     29
19:   PUSHLOCADDR     2
20:   PUSHNUM         2
21:   SAVETOADDR
22:   PUSHSTATADDR    0
23:   PUSHSTATADDR    0

```

24:	LOADFROMADDR	
25:	PUSHNUM	3
26:	MUL	
27:	SAVETOADDR	
28:	JUMP	14
29:	PUSHLOCADDR	2
30:	PUSHSTATADDR	0
31:	LOADFROMADDR	
32:	SAVETOADDR	
33:	PUSHLOCADDR	2
34:	LOADFROMADDR	
35:	WRITEINT	
36:	STOP	

- (f) Suppose `Instruction.getNextCodeAddress() == 45` when a correct solution to TinyJ Assignment 2 begins to translate the following method, and suppose `z` is a static variable with address 2. What TinyJ VM instructions would this method be translated into?

```
static int p(int x)
{
    int y = 3, w;

    x = z + y;

    if (x < 10) z = x;
    else z = y;

    while (z <= 100) {
        System.out.println(z);
        z = z + y;
    }

    return z - x;
}
```

SOLUTION:

45:	INITSTKFRM	2
46:	PUSHLOCADDR	1
47:	PUSHNUM	3
48:	SAVETOADDR	
49:	PUSHLOCADDR	-2
50:	PUSHSTATADDR	2
51:	LOADFROMADDR	
52:	PUSHLOCADDR	1
53:	LOADFROMADDR	
54:	ADD	
55:	SAVETOADDR	
56:	PUSHLOCADDR	-2
57:	LOADFROMADDR	
58:	PUSHNUM	10
59:	LT	
60:	JUMPONFALSE	66
61:	PUSHSTATADDR	2
62:	PUSHLOCADDR	-2
63:	LOADFROMADDR	
64:	SAVETOADDR	
65:	JUMP	70
66:	PUSHSTATADDR	2
67:	PUSHLOCADDR	1
68:	LOADFROMADDR	
69:	SAVETOADDR	
70:	PUSHSTATADDR	2

71:	LOADFROMADDR	
72:	PUSHNUM	100
73:	LE	
74:	JUMPONFALSE	87
75:	PUSHSTATADDR	2
76:	LOADFROMADDR	
77:	WRITEINT	
78:	WRITELNOP	
79:	PUSHSTATADDR	2
80:	PUSHSTATADDR	2
81:	LOADFROMADDR	
82:	PUSHLOCADDR	1
83:	LOADFROMADDR	
84:	ADD	
85:	SAVETOADDR	
86:	JUMP	70
87:	PUSHSTATADDR	2
88:	LOADFROMADDR	
89:	PUSHLOCADDR	-2
90:	LOADFROMADDR	
91:	SUB	
92:	RETURN	1

First Solution:

```
private void expr1() throws SourceFileErrorException
{
    TJ.output.printSymbol(NTErr1);
    TJ.output.incTreeDepth();

    switch (getCurrentToken()) {
        case UNSIGNEDINT:
            nextToken();
            break;

        case NEW:
            nextToken();
            accept(INT);
            accept(LBRACKET);
            expr3();
            accept(RBRACKET);
            while (getCurrentToken() == LBRACKET) {
                nextToken();
                accept(RBRACKET);
            }
            break;

        default:
            throw new SourceFileErrorException("Expected UNSIGNEDINT or new");
    }

    TJ.output.decTreeDepth();
}
```

Second Solution:

```
private void expr1() throws SourceFileErrorException
{
    TJ.output.printSymbol(NTErr1);
    TJ.output.incTreeDepth();

    if (getCurrentToken() == UNSIGNEDINT) {
        nextToken();
    }
    else if (getCurrentToken() == NEW) {
        nextToken();
        accept(INT);
        accept(LBRACKET);
        expr3();
        accept(RBRACKET);
        while (getCurrentToken() == LBRACKET) {
            nextToken();
            accept(RBRACKET);
        }
    }
    else throw new SourceFileErrorException("Expected UNSIGNEDINT or new");

    TJ.output.decTreeDepth();
}
```

Third Solution:

```
private void expr1() throws SourceFileErrorException
{
    TJ.output.printSymbol(NTErr1);
    TJ.output.incTreeDepth();

    if (getCurrentToken() == UNSIGNEDINT) {
        nextToken();
    }
    else {
        accept(NEW);
        accept(INT);
        accept(LBRACKET);
        expr3();
        accept(RBRACKET);
        while (getCurrentToken() == LBRACKET) {
            nextToken();
            accept(RBRACKET);
        }
    }

    TJ.output.decTreeDepth();
}
```

COMMENT: The third solution is slightly more concise than the first two solutions, but it gives a slightly less informative error message if expr1 is called when currentToken is neither UNSIGNEDINT nor NEW.

## A Mistake to Avoid When Writing Recursive Descent Parsing Code

A common mistake in writing recursive descent parsing code is to write

```
getCurrentToken() == X
```

or `accept(X)` [which performs a `getCurrentToken() == X` test]

using a `Symbols` constant `X` that represents a nonterminal. This is wrong, as `getCurrentToken()` returns a `Symbols` constant that represents a token. Here are two examples of this kind of mistake.

1. In TinyJ Assignment 1, the method `argumentList()` should be based on the following EBNF rule:

```
<argumentList> ::= '(' [<expr3>{,<expr3>} ] ')'
```

When writing this method it would be wrong to write:

```
accept(LPAREN);
if (getCurrentToken() == NExpr3) /* INCORRECT! */ {
    expr3();
    ... // a while loop that deals with {,<expr3>}
}
accept(RPAREN);
```

Here it would be correct to write code of the following form:

```
accept(LPAREN);
if (getCurrentToken() != RPAREN) /* CORRECT */ {
    expr3();
    ... // a while loop that deals with {,<expr3>}
}
accept(RPAREN);
```

2. When writing the method `expr1()` for TinyJ Assignment 1, one case that needs to be dealt with relates to the following part of the TinyJ EBNF rule that defines `<expr1>`:

```
IDENTIFIER ( . nextInt '(' ')' | [<argumentList>] { '[' <expr3> ']' } )
```

Here it would be wrong to write something like:

```
case IDENT:
    nextToken();
    if (getCurrentToken() != DOT) {
        if (getCurrentToken() == NArgumentList /* INCORRECT! */ ) argumentList();
        ... // a while loop that deals with { '[' <expr3> ']' }
    }
    else {
        ... // code to deal with . nextInt '(' ')'
    }
    break;
```

Instead, you can write something like:

```
case IDENT:
    nextToken();
    if (getCurrentToken() != DOT) {
        if (getCurrentToken() == LPAREN /* CORRECT */ ) argumentList();
        ... // a while loop that deals with { '[' <expr3> ']' }
    }
    else {
        ... // code to deal with . nextInt '(' ')'
    }
    break;
```

The use of `LPAREN` in the above code is correct because the first token of any instance of `<argumentList>` must be a left parenthesis, as we see from the EBNF rule

```
<argumentList> ::= '(' [<expr3>{,<expr3>} ] ')'
```