LAB 10

Midterm 2 Review

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Apr 3, 2023

LOGISTICS

- Homework 07 due Thu 04/06
- Midterm ••
 - Lectures this week are in scope
 - The midterm 2 study guide will be out soon make sure to check it beforehand! (check Ed 2069 for updates)
 - Check Ed for more resources!!

TODAY'S OUTLINE 🔆

- Topical Review
 - Efficiency + Lab Q11
 - Iterators and generators + Lab Q4
 - Recursion and tree recursion + Lab Q1-2
 - Lab Q10 (trees) walkthrough, if time allows
- If you want individual help in the meantime, just walk to one of our Als!

ORDER OF GROWTH

ORDER OF GROWTH

- Order of growth / efficiency runtime as a **function** of the input size
- Input size
 - numeric magnitude of the number
 - Python lists length of the list
 - linked list/trees number of nodes
- Runtime
 - the number of operations
- Theta notation for approximate runtime
 - ullet $\Theta(f(n))$ approximately f(n) by a constant factor
 - Only consider the term that grows fastest
- constant < logarithmic < linear < quadratic < exponential

ORDER OF GROWTH - CONSTANT

- Constant $\leftrightarrow \Theta(1)$
- Runtime does not change as the input size changes
- For example:

```
def square (x):

return x * x
```

input	function call	return value	operations
1	square(1)	1*1	1
2	square(2)	2*2	1
100	square(100)	100*100	1
n	square(n)	n*n	1

ORDER OF GROWTH - LOGARITHMIC

- Logarithmic $\leftrightarrow \Theta(\log n)$
- Often when we keep dividing the input by a constant

```
def foo (x):
     while x > 0:
         print('hey')
         x //= 2
```

Let's say the while loop runs n times before x reaches 0.

divide x by 2 for n times before it gets to 1

$$\implies \frac{x}{2^n} = 1 \implies 2^n = x \implies n = \log_2 x$$

ORDER OF GROWTH - LINEAR

- Linear $\leftrightarrow \Theta(n)$
- ullet Often when a loop runs n times, each time doing work in constant time

```
def factorial(x):
    prod = 1
    for i in range(1, x + 1):
        prod *= i # execute x times in total
    return prod
```

input	function call	return value	operations
1	factorial(1)	1*1	1
2	factorial(2)	2*1*1	2
100	factorial(100)	100*99**1*1	100
n	factorial(n)	n*(n-1)**1*1	n

ORDER OF GROWTH - QUADRATIC

- Quadratic $\leftrightarrow \Theta(n^2)$
- ullet Often when a nested loop runs n^2 times, each time doing work in constant time

```
def bar(n):
    for a in range(n):
        for b in range(n):
            print(a, b) # print n * n times in total
```

input	function call	operations (prints)
1	bar(1)	1
2	bar(2)	4
100	bar(100)	10000
n	bar(n)	n^2

ORDER OF GROWTH - EXPONENTIAL

- Exponential $\leftrightarrow \Theta(c^n)$, where c is a constant
- Often in tree recursion

```
def rec(n):
    if n == 0:
        return 1
    else:
        return rec(n - 1) + rec(n - 1)
```

input	function call	return value	operations
1	rec(1)	2	1
2	rec(2)	4	3
10	rec(10)	1024	1023
n	rec(n)	2^n	2^n

ORDER OF GROWTH - TREE RECURSION

```
def rec(n):
    if n == 0:
        return 1
    else:
        return rec(n - 1) + rec(n - 1)
```

- Try drawing out the recursion tree diagram
- n + 1 levels
- The i^{th} level has 2^i nodes (root at level 0)
- Each node does constant work (addition)
- Total = $(1 + 2 + ... + 2^{n+1}) \cdot constant = \Theta(2^n)$

ORDER OF GROWTH - TREE RECURSION

What about this? 🤔

```
def rec(n):
    if n == 0:
        return 1
    else:
        return rec(n // 2) + rec(n // 2)
```

- Draw out the recursion tree diagram
- $\log_2 n$ levels
- The i^{th} level has 2^i nodes (root at level 0)
- Each node does constant work (addition)
- Total = $(1 + 2 + ... + 2^{\log_2 n}) \cdot constant = \Theta(n)$

ORDER OF GROWTH - LOOPS

In general, (# times the loops run) * (work done each time)

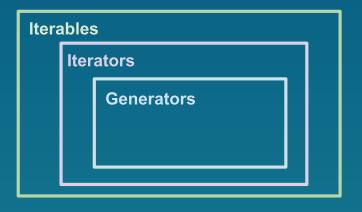
```
def factorial (n):
     # returns n! in linear time
def foo (n):
     # nested loop runs n^2 times
     for i in range(n):
          for j in range(n):
               # each time this takes theta(n) time
               print(factorial(n))
```

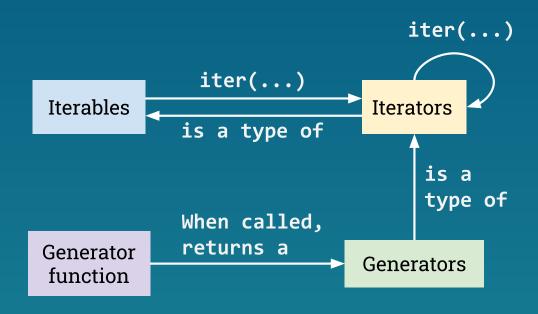
• foo(n) runs in $\Theta(n^3)$ time

LAB Q11

ITERATORS AND GENERATORS

OVERVIEW





ITERABLES AND ITERATORS

- Iterables
 - Can be iterated through using a for loop:

```
for elem in iterable:
# do something
```

- E.g., sequences (lists, tuples, dictionaries, ranges, etc.)
- iter(sequence) returns a new iterator
- Iterators
 - "Tracker" for an iterable
 - next(iterator) gives the next element
 - cannot go back
 - iter(iterator) returns the iterator itself, without resetting
 the position

MORE ON ITERATORS

- Iterators can be "depleted" implicitly with:
 - for loop
 - list comprehension
 - list(iterator)
- To iterate through an iterator:
 - use a for loop

```
for elem in iterator:
    # do sth with elem
```

use a while loop

```
while ...:
   elem = next(iterator)
   # do sth with elem
```

GENERATOR FUNCTIONS AND GENERATORS

- Generator functions
 - Have by at least one yield or yield from statement
 - When called, return a <u>new</u> generator, without evaluating the function body yet
 - Define customized iterators
- Generators
 - next(generator) "lazy evaluation"
 - Keep evaluating until a value is yielded or the function teriminates

YIELD VS YIELD FROM

- Both defines elements in the generator
- yield only one element at a time
- yield from any iterables (sequences, iterators, generators, etc.)

```
for elem in iterable:
    yield elem
# is equivalent to
yield from iterable
```

- for loop and yield can modify the element before yielding
- yield from can only yield the exact elements from the iterable

RECURSIVE GENERATORS

- Base case
 - No more elements to yield often just return
- Recursive case
 - Often iterate through the recursive call (generator) with a for loop
- Combination step
 - What does the recursive generator contain? recursive leap of faith!
 - Given that, how to yield for the current call?

```
def yield_paths (t, value):
    """Yields all possible paths from the root of t
    to a node with the label value as a list."""
    if t.label == value:
        yield [value]
    for b in t.branches:
        for path in yield_paths(b, value):
            yield [t.label] + path
```

LAB Q4

RECURSION, TREERECURSION 6

RECURSION

- Recursive function a function calls itself
 - Base case
 - Recursive case
 - Combination step
- Tree recursive function a function makes more than one call to itself
 - Often times one recusive call represents one choice
 - What choices do we have and how do we combine them together?

BASE CASE

- Simplest input give an answer without further computation
- Often hinted in the doctest
- Sometimes may need an "overshot" case
 - Could it overshoot in the first place?
 - E.g., input n is positive may need to check if n <= 0
- Some common base cases *
 - positive integer n <= 0 or n == 1</p>
 - sequence len(seq) == 0
 - Trees t.is_leaf()
 - Linked lists lnk is Link.empty

^{*} The goal of these is to help you develop some intuition. Please do NOT take these for granted on an exam - it really depends on the specific problem!

RECURSIVE CASE

- Break down the problem into a smaller one
 - How should each argument change in the recursive call, if they can change at all?
- Make a recursive call for the smaller problem
- Some common recursive cases *
 - positive integer n some_number or n // 10
 - sequence seq[1:]
 - Trees often one recursive call on each branch
 - Linked lists lnk.rest

^{*} The goal of these is to help you develop some intuition. Please do NOT take these for granted on an exam - it really depends on the specific problem!

COMBINATION STEP

- What does the recursive call return?
 - based on the problem description, NOT your implementation
 - recursive leap of faith assume that the recursive call always gives the correct result
- Given the solution to a smaller problem, how to solve the bigger problem?

HELPER FUNCTIONS FOR RECURSION

- When to use:
 - Need to keep track of more variables than the given parameters of the outer function
- Where
 - usually nested within the original function
- How
 - define the helper function, and return a call to it with appropriate initial arguments

LAB Q1-2

AND THAT'S IT...



GOOD LUCK with the midterm!

We got this >:)



