

Shell sort 와 Merge sort(1,2) 를 이용한
space-time trade-off 분석

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Lorem Ipsum is simply dummy text of the printing and typesetting industry.

space-time tradeoff ?

컴퓨터 분야에서 space-time tradeoff란

- 큰 메모리로 적은 시간의 문제를 품
 - 작은 메모리로 큰 시간의 문제를 품
- 의 상관 관계를 의미함

데이터 저장
(압축)

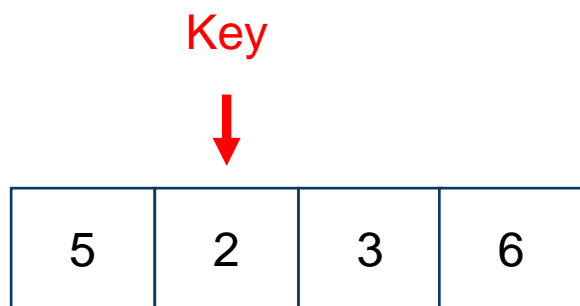
Look up table
(cache)

Merge
sort

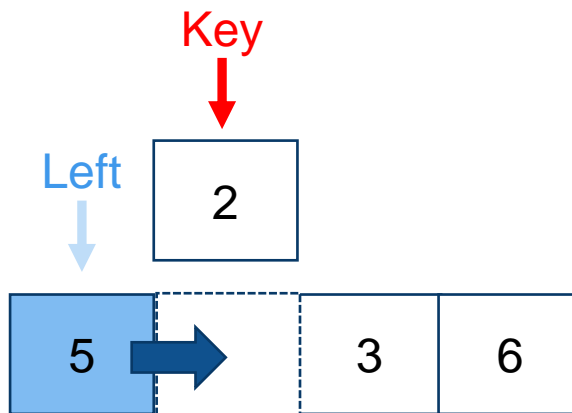


Shell sort

- selction sort
- shell sort
- code

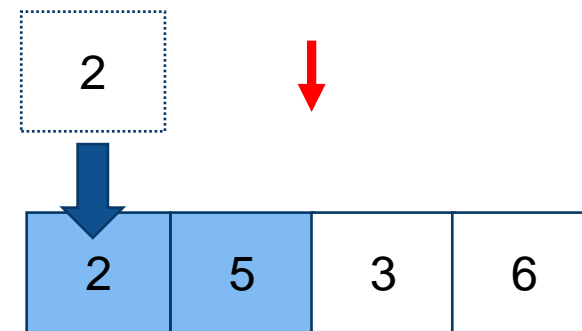


1.Key 값 선정

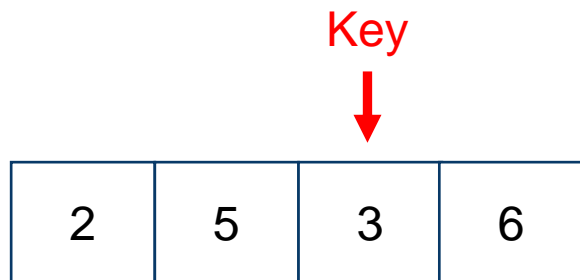


If Left > Key ? move :stop

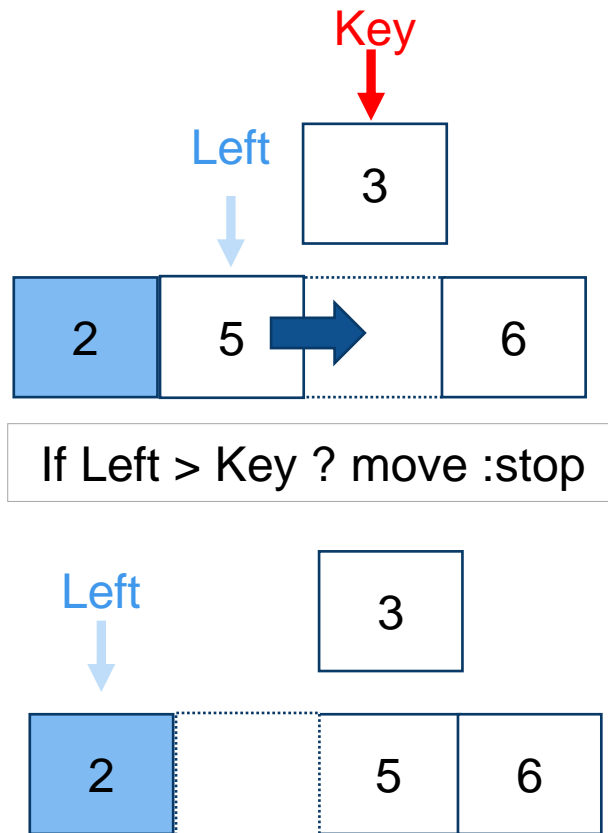
2.왼쪽 요소와 비교(반복)



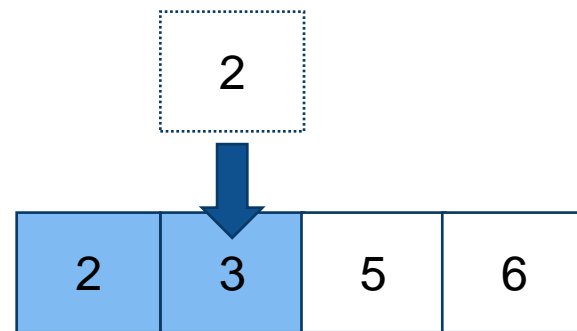
3.Stop 시 삽입



1. Key 값 선정



2. 5를 옮긴 후, 2보다 크니
Stop



3. Stop후 Key 값 삽입

Code

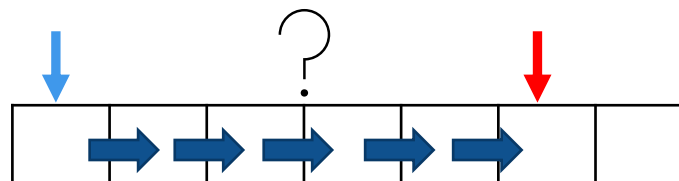
```
def selection_sort1(arr):  
    i=j=key =0  
    length=len(arr)  
  
    for i in range(1,length):  
        key=arr[i] ## key 값 선택 → 1.Key값 선정  
        j=i-1 ## key 왼쪽과 비교  
        while j>=0 and arr[j] > key: → 2. 왼쪽 요소와 비교(반복) 후 move  
            arr[j+1] = arr[j]  
            j=j-1  
        arr[j+1]=key → 3. stop 후 삽입  
  
    return arr
```

1 Donald L. Shell

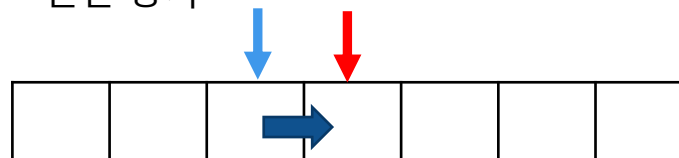


-1995년 Donald L. Shell이 제안

2 삽입 정렬의 장,단점

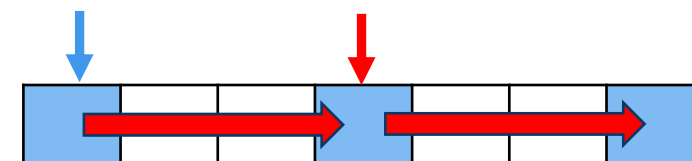


-비교 값들의 간의 거리가 멀면 이동 연산 증가

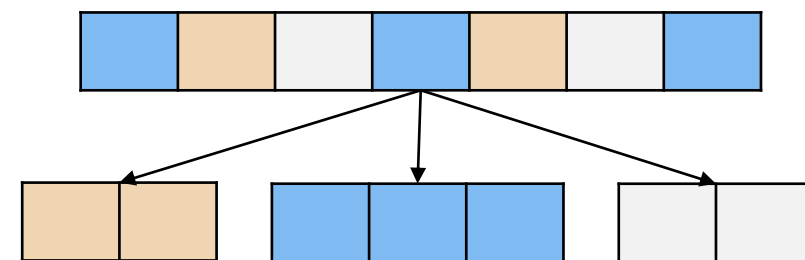


-어느 정도 정렬된 배열 내에서는 속도 빠름

3 삽입 정렬 응용



-한칸 씩 이동하던 거리를 gap만큼 이동



-Gap만큼 떨어진 부분 리스트들끼리 삽입 정렬시킨다

Code

```
def shell_sort(arr):
    size=len(arr)
    gap=size//2
    while gap > 0:
        for i in range(0,gap):
            selection_sort(arr,i,size,gap)
        gap= gap//2
    return arr
```

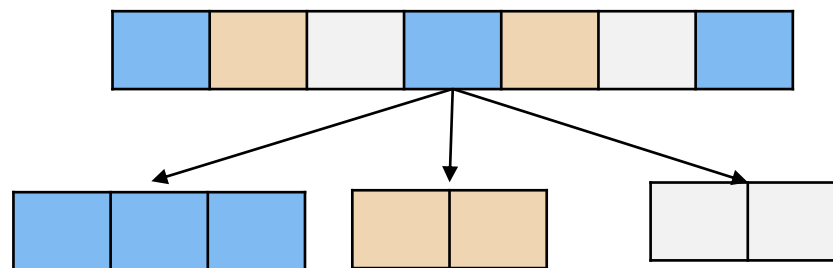
```
def selection_sort(arr,first,size,gap):
    for i in range(first+gap,size):
        key=arr[i]
        j=i-gap
        while j>=first and arr[j] > key:
            arr[j+gap] = arr[j]
            j=j-gap
        arr[j+gap]=key
```

→ gap은 1/2로 줄어나감

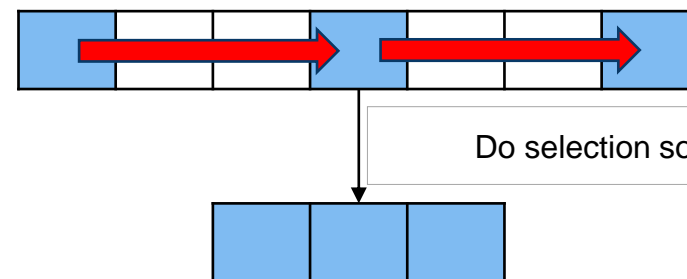
→ 부분 리스트 생성

→ 2. (시작점,size,gap)을
이용해 삽입정렬

[부분 리스트 생성]



[삽입정렬]

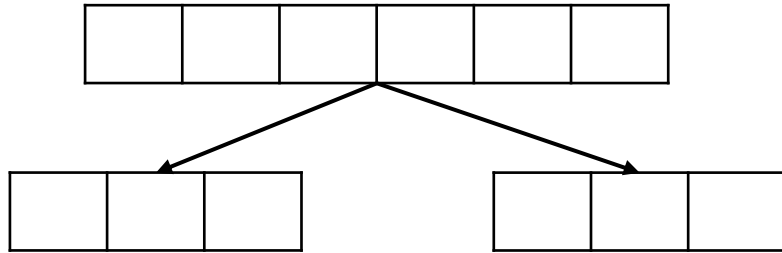


Merge sort

- Merge sort 1,2

1 Merge sort 1(out-of-place)

[두개의 배열 생성]



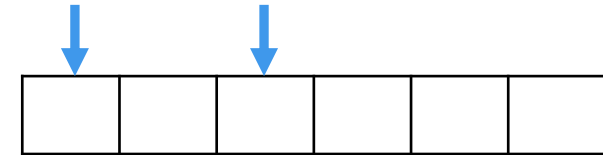
```
void mergesort(int n, keytype S[]) {
    const int h = n/2, m = n - h;
    keytype U[1..h], V[1..m];    // 추가 메모리(array)
```

- 두개의 하위 배열 생성
- 추가 메모리 필요

VS

1 Merge sort 2(in-place)

[인덱스 전달]



```
void merge(h, m, U[], V[], S[]) {
    index i, j, k;    // 추가 메모리(no array)
```

- 전역 배열의 인덱스를 넘겨줌
- merge1에 비해 약간 느림

2

Merge sort 1(extra array)

Merge sort1 code

Code

```
def merge_sort1(arr):
    half=len(arr)//2
    if half==0:
        return arr
    arr_left=arr[:half]
    arr_right=arr[half:]
    left_result=merge_sort1(arr_left)
    right_result=merge_sort1(arr_right)
    return merge(left_result,right_result)
```

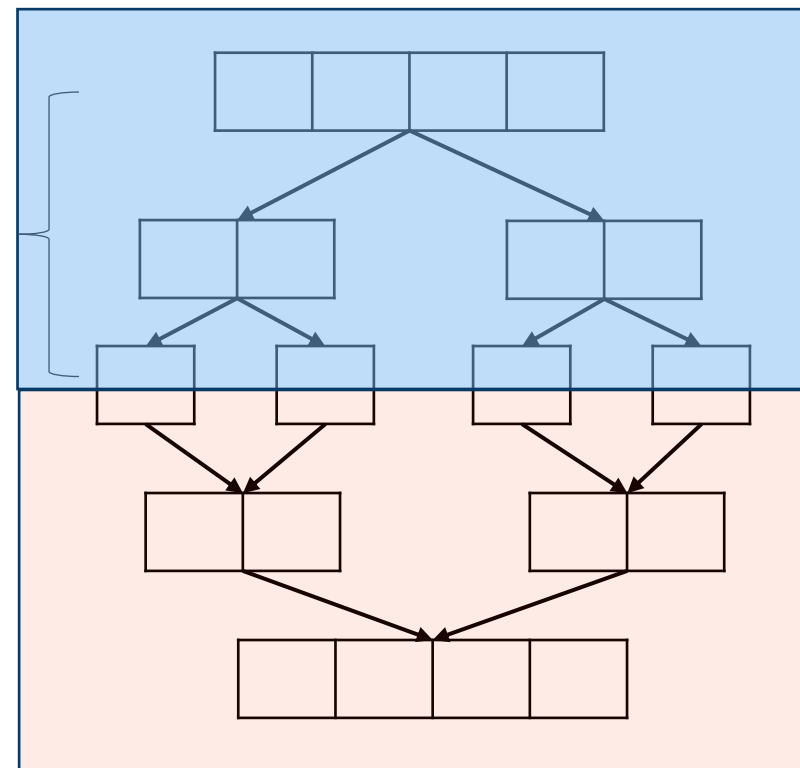
1. 배열 생성 후, 분할

```
def merge(left_arr,right_arr):
    i=0;j=0
    sum_arr=[]

    while(i<len(left_arr) and j<len(right_arr)):
        if (left_arr[i] > right_arr[j]):
            sum_arr.append(right_arr[j])
            j+=1
        else:
            sum_arr.append(left_arr[i])
            i+=1
    if (i==len(left_arr)):
        sum_arr.extend(right_arr[j:])
    if(j==len(right_arr)):
        sum_arr.extend(left_arr[i:])
    return sum_arr
```

2. 좌우 배열 비교

3. 남은 배열 삽입



2

Merge sort 2(in-place)

Merge sort 2 (in-place)

Code

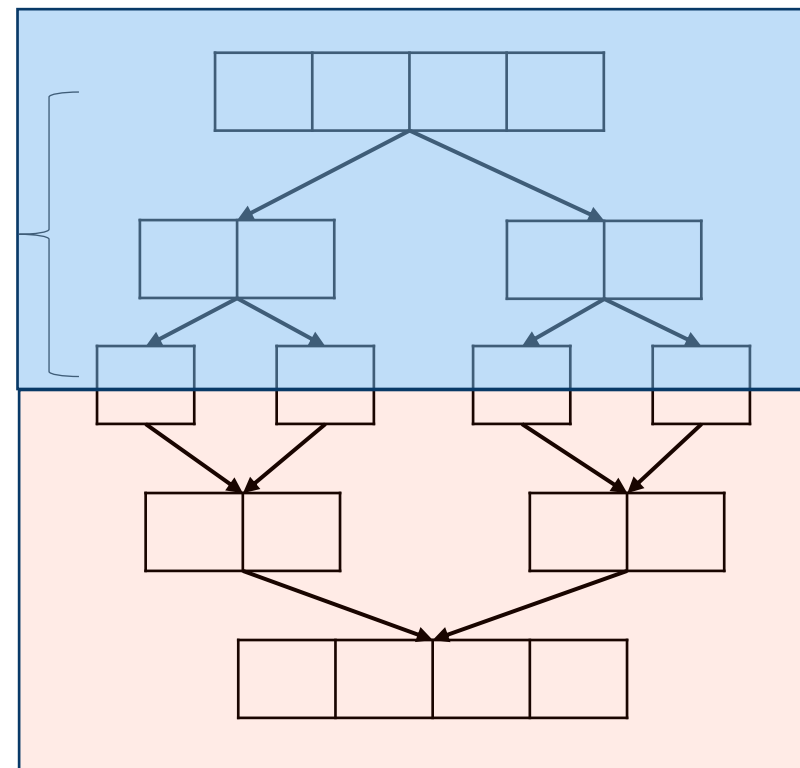
```
def mergesort2(arr, left, right):
    if(right-left>=1):
        mid=(right+left)//2
        mergesort2(arr, left, mid)
        mergesort2(arr, mid+1, right)
        return merge_def(arr, left, mid, right)
    else:
        return
```

```
def merge_def(arr, left, mid, right):
    sum_arr = []
    i = left
    j = mid+1
    while (i <= mid and j <= right):
        if (arr[i] < arr[j]):
            sum_arr.append(arr[i])
            i+=1
        else:
            sum_arr.append(arr[j])
            j+=1
    if (i>mid):
        sum_arr.extend(arr[j:right+1])
    if (j>right):
        sum_arr.extend(arr[i:mid+1])
    arr[left:right+1] = sum_arr
    return arr
```

1. 분할 과정에서
인덱스 전달

2. 좌우 배열 비교

3. 남은 배열 삽입



성능 비교

- 시간 복잡도
- 공간 복잡도

Lorem Ipsum is simply dummy text of the printing and typesetting industry.

개발환경,언어

PyCharm Community Edition 2020.1.2
Python 3.6.4

CPU

Intel(R) Core(TM) i7-4790 CPU @ 3.60GHz 3.60 GHz

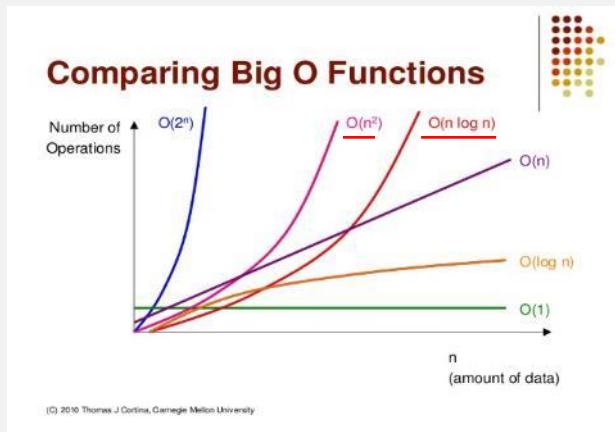
RAM

8.00GB

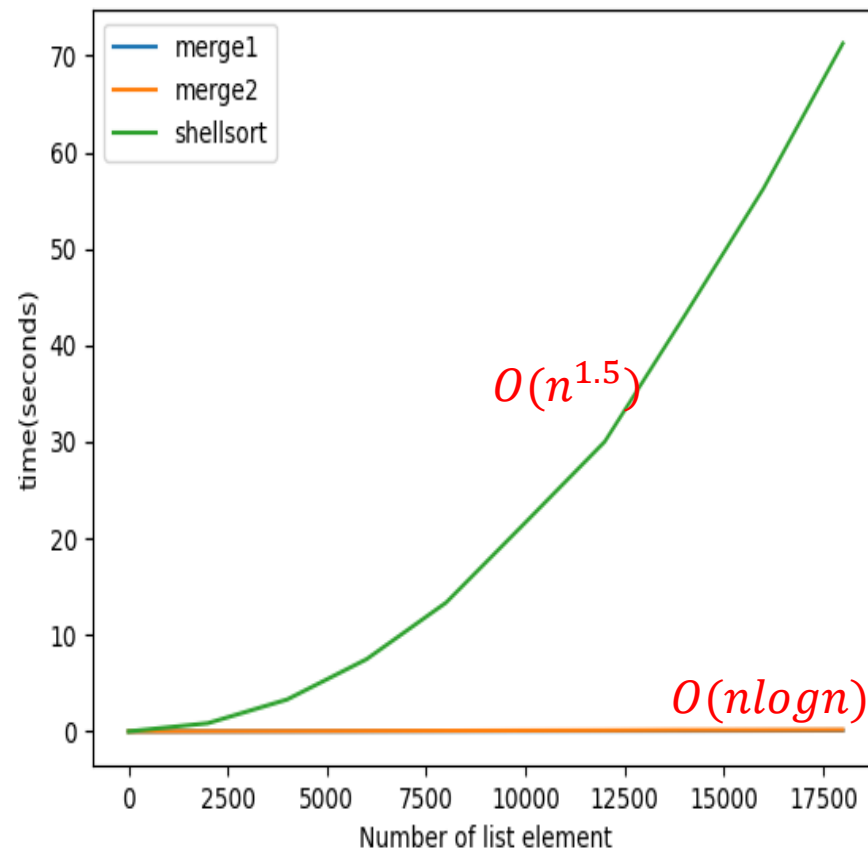
	Best	Avg	Worst	Space complexity
Shell sort	$O(n)$	$O(n^{1.5})$	$O(n^2)$	$O(1)$
Merge sort(out-of-place)	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$
Merge sort(in-place)	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$

Lorem Ipsum is simply dummy text of the printing and typesetting industry.

- $O(n^{1.5})$ 인 shell sort가 가장 느린 모습을 보인다
- 같은 시간 복잡도인 merge sort는 비슷한 시간 소요를 보인다



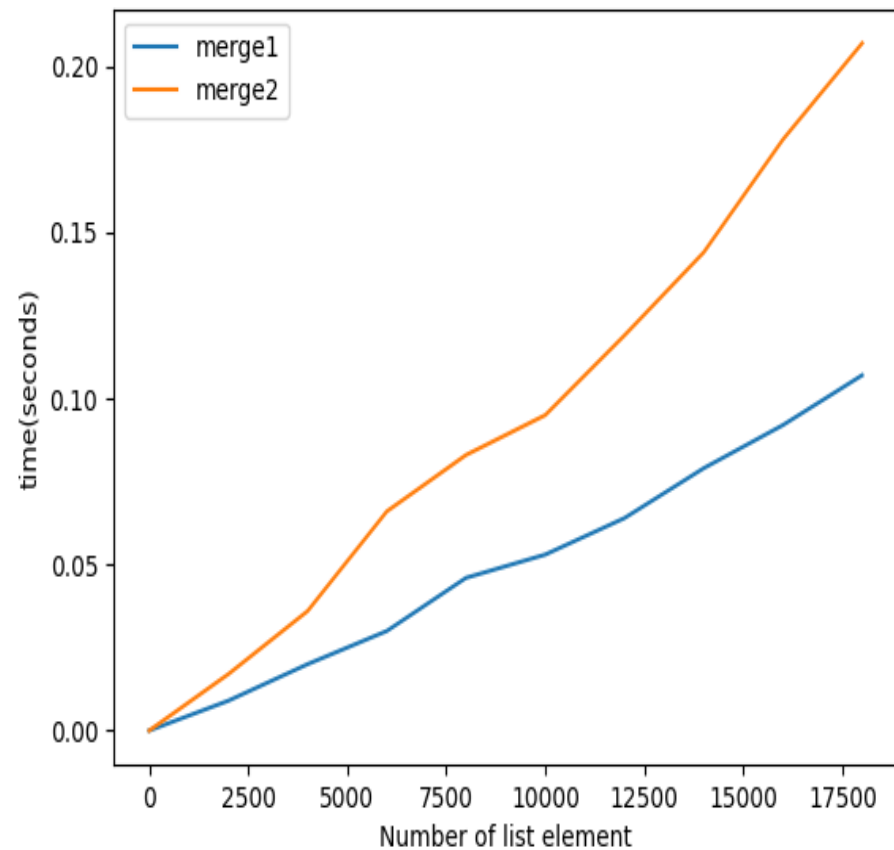
	Best	Avg	Worst	Space complexity
Shell sort	$O(n)$	$O(n^{1.5})$	$O(n^2)$	$O(1)$
Merge sort1	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$
Merge sort2	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$



Lorem Ipsum is simply dummy text of the printing and typesetting industry.

- 같은 시간 복잡도인 $O(n \log n)$
- In-place 인 merge sort 2 가 조금 느린 모습 (마지막 배열 복사 과정)

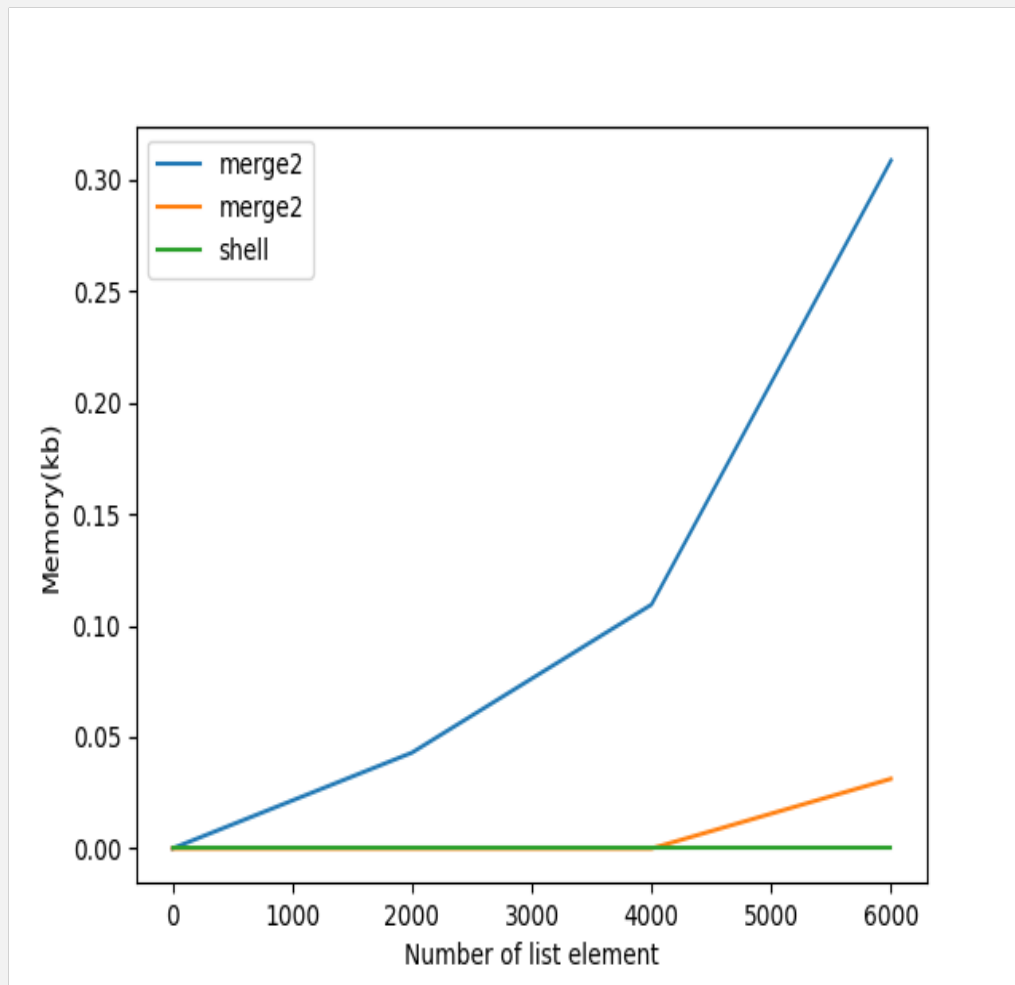
	Best	Avg	Worst	Space complexity
Shell sort	$O(n)$	$O(n^{15})$	$O(n^2)$	$O(1)$
Merge sort1	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$
Merge sort2	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$



Lorem Ipsum is simply dummy text of the printing and typesetting industry.

- 공간복잡도는 shell sort – $O(1)$ / merge sort – $O(n)$
- 실제로는 추가 배열이 필요한 merge1은 더 많은 메모리 사용량을 보임

	Best	Avg	Worst	Space complexity
Shell sort	$O(n)$	$O(n^{15})$	$O(n^2)$	$O(1)$
Merge sort1	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$
Merge sort2	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$



Lorem Ipsum is simply dummy text of the printing and typesetting industry.

1 결론

-시간 복잡도:

shell>merge2(in-place)>=merge1

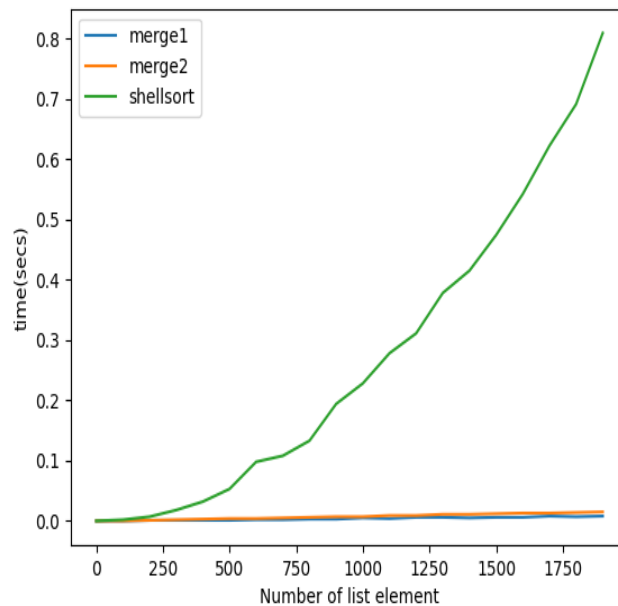
-공간복잡도:

merge1>=merge2(in-place)>shell

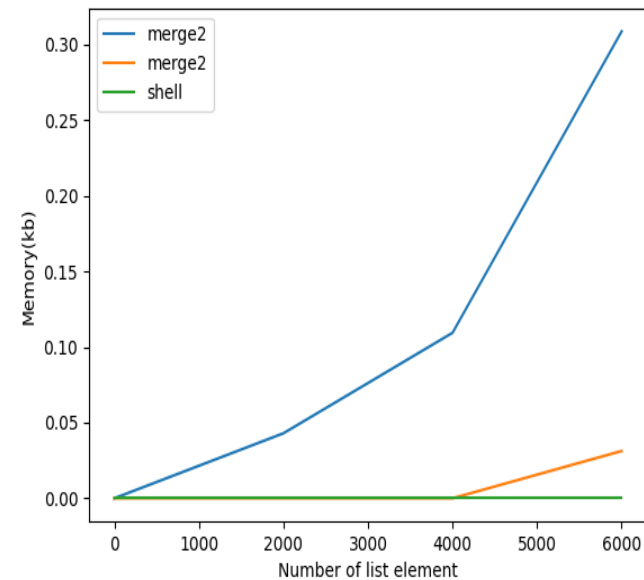
→적당히 작은 n에서 merge,shell sort
는 **space-time trade-off** 가 존재함을
보여준다.

.

2 소요 시간 관찰 결과



3 메모리 사용량 관찰 결과



THANK YOU