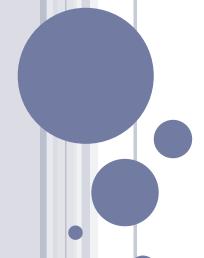
ON CLUSTERING USING RANDOM WALKS

DAVID HAREL AND YEHUDA KOREN

PRESENTED AT FOUNDATIONS OF SOFTWARE TECHNOLOGY AND THEORETICAL COMPUTER SCIENCE 2001



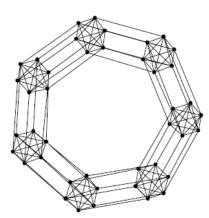
Talk by Kota Hara

OVERVIEW

- Important properties
- Method
 - Flow of the method
 - Creating structured graph
 - Random walk
 - Separating operator
 - Integration with agglomerative clustering
- Results
- Summary

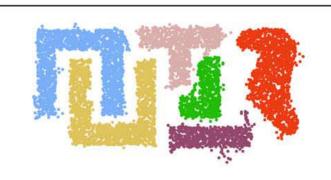
PROPERTIES OF METHOD

- Exploit structure of graph
- Based on deterministic analysis of random walks on weighted graphs (no random simulation)
- Doesn't require a prespecified number of cluster
- No eigenvalue problem involved
- Work well on clustering of 2d spatial points









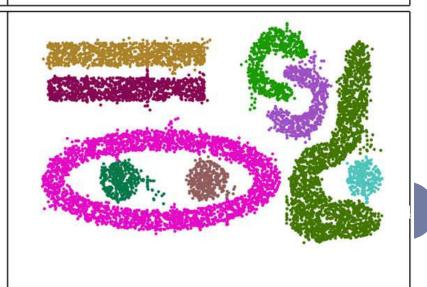






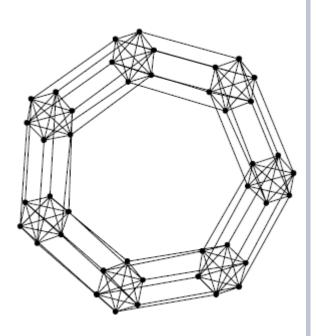


DS6: 10000 points



HOW IT WORKS? (SIMPLE EXAMPLE)

- Want to divide points into 7 clusters
- Create structured graph
- 'Separating operator' decreases weights of 'external edge' while increases those of 'internal edge'
- Remove edges with small weights
- Or apply agglomerative clustering



CREATING STRUCTURED GRAPH G(V, E, w)

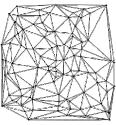
Given data points....

Weight of $exp(-\frac{d(a,b)^2}{C})$ when clustering spatial points

• Uniform in some examples

Edge

- k-mutual neighborhood
- Remove edges whose weights are below threshold
- Delaunay triangulation



RANDOM WALK ON GRAPH

- 'Random walker' walks on graph according to probability of transition (Markov process)
- Probability of transition from node *i* to *j*

$$p_{ij} = \frac{w(i,j)}{d_i}$$
 where $d_i = \sum_{\langle i,k \rangle} w(i,k)$

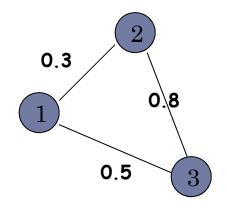
Transition matrix

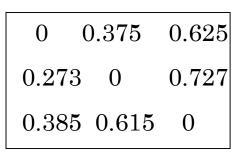
$$M_{ij} = \begin{cases} p_{ij} & \langle i, j \rangle \in E \\ 0 & otherwise \end{cases}$$

TRANSITION MATRIX

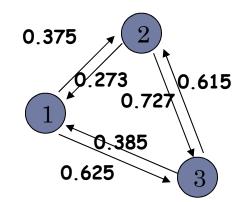
| 0 | 0.3 | 0.5 |
|-----|-----|-----|
| 0.3 | 0 | 0.8 |
| 0.5 | 0.8 | 0 |

Similarity matrix A





Transition matrix M

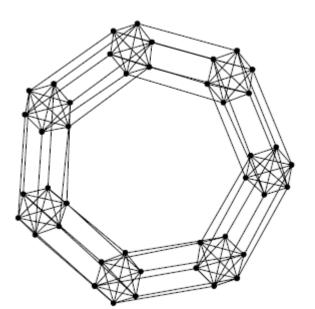


TRANSITION PROBABILITY

- Vector whose j-th component is probability that a random walk originating at i will visit node j in its k-th step $P_{visit}^k(i) \in \mathbb{R}^n$
- $P_{visit}^k(i)$ is the *i*-th row in M^k

SEPARATING OPERATOR

- Based on deterministic analysis of random walks
- Applied to the graph several times to increase weights of intra-cluster edges and decrease inter-cluster edges



SEPARATING OPERATOR 1 (NS)

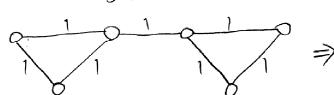
• Vector whose *j*-th component is probability that a random walk originating at *i* will visit node *j* within *k* steps (*k*=3)

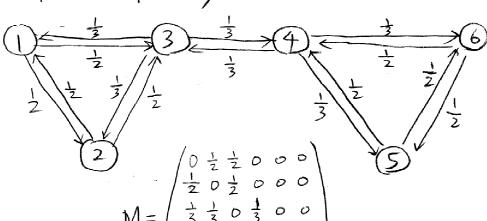
$$P_{visit}^{\leq k}(i) = \sum_{s=1}^{k} P_{visit}^{s}(i)$$

• Use similarity between $P_{visit}^{\leq k}(v)$ and $P_{visit}^{\leq k}(u)$ as a similarity between v and u

• Replace $w(u, v) \quad \forall \langle u, v \rangle \in E$ by

$$\exp(2k - \|P_{visit}^{\leq k}(v) - P_{visit}^{\leq k}(u)\|_{L_1}) - 1$$





$$\begin{bmatrix}
P_{visit}^{1}(4) = \begin{pmatrix} 0 & 0 & \frac{1}{3} & 0 & \frac{1}{3} & \frac{1}{3} \\
V_{visit}^{2}(4) = \begin{pmatrix} \frac{1}{9} & \frac{1}{9} & \frac{1}{9} & \frac{1}{9} & \frac{1}{9} & \frac{1}{9} & \frac{1}{9} \\
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$$\begin{array}{l}
P_{\text{visit}}^{1}(4) = (\overline{q} \ \overline{q} \ \overline{b} \ \overline{b})
\end{array}$$

$$\begin{array}{l}
P_{\text{visit}}^{1}(5) = (600 \frac{1}{2} 0 \frac{1}{2}) \\
P_{\text{visit}}^{2}(5) = (000 \frac{1}{4} \frac{3}{5} \frac{2}{3})
\end{array}$$

$$\begin{array}{l}
P_{\text{visit}}^{2}(5) = (000 \frac{1}{4} \frac{3}{5} \frac{2}{3})
\end{array}$$

$$\|P_{\text{vist}}^{k\leq 2}(4) - P_{\text{visit}}^{k\leq 2}(5)\|_{L_1} = \frac{17}{18} \implies W(4,5) = 20,233$$

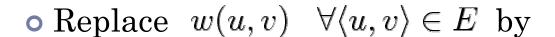
$$\|P_{visit}^{k\leq 2}(4) - P_{visit}^{k\leq 2}(3)\|_{L^{1}} = \frac{32}{18} \implies W(3,4) = 6.2218$$

SEPARATING OPERATOR 2 (CE)

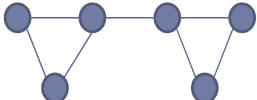
• Probability that random walker starting from *v* visits *u* before returning to *u* (escape probability)

$$P_{escape}(v,u)$$

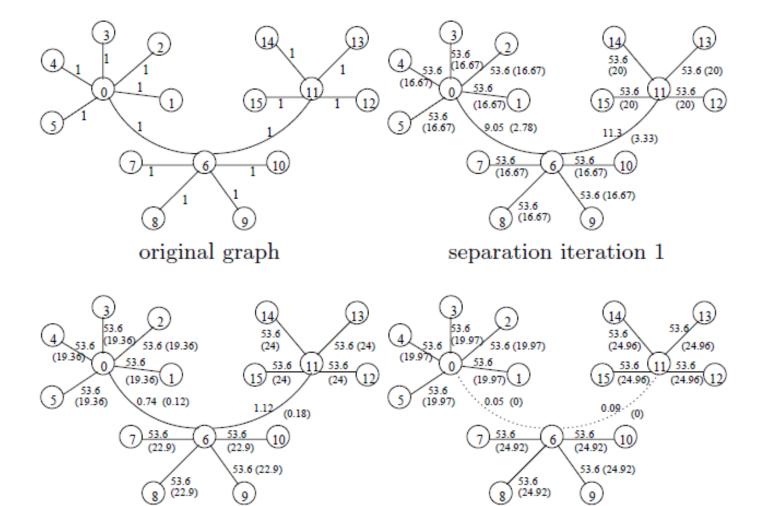
• If v and u are in different cluster, $P_{escape}(v, u)$ will be low



$$P_{escape}(v,u) \cdot P_{escape}(u,v)$$



EXAMPLE

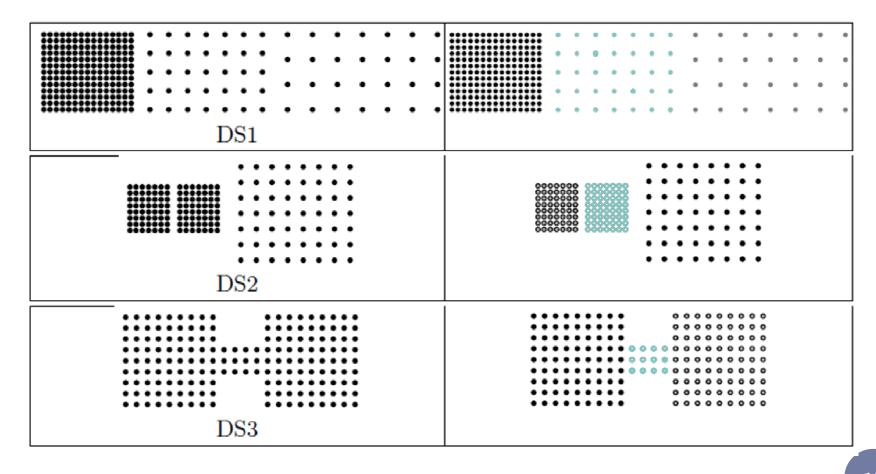


separation iteration 3

separation iteration 2

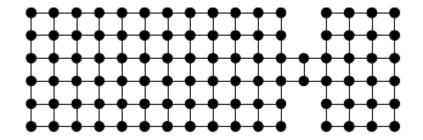
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CLUSTERING SPATIAL POINT-SETS



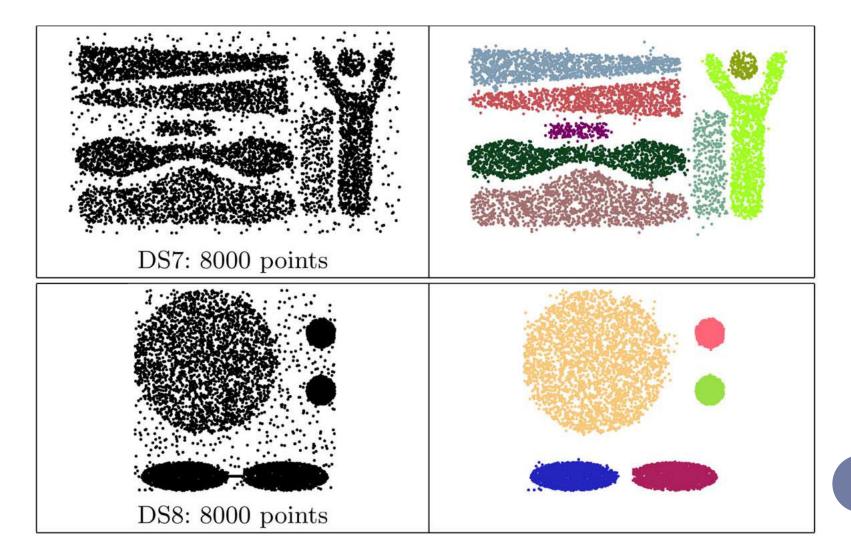
INTEGRATION WITH AGGLOMERATIVE CLUSTERING

- Use separation operation as a preprocessing stage before agglomerative clustering
- Prevent bad local merging that works against the graph structure

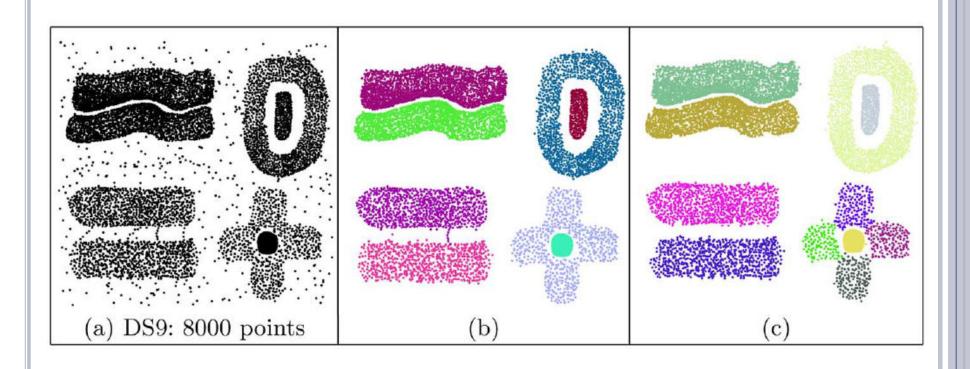


• (Variant) Repeatedly apply separation operation during agglomerative clustering

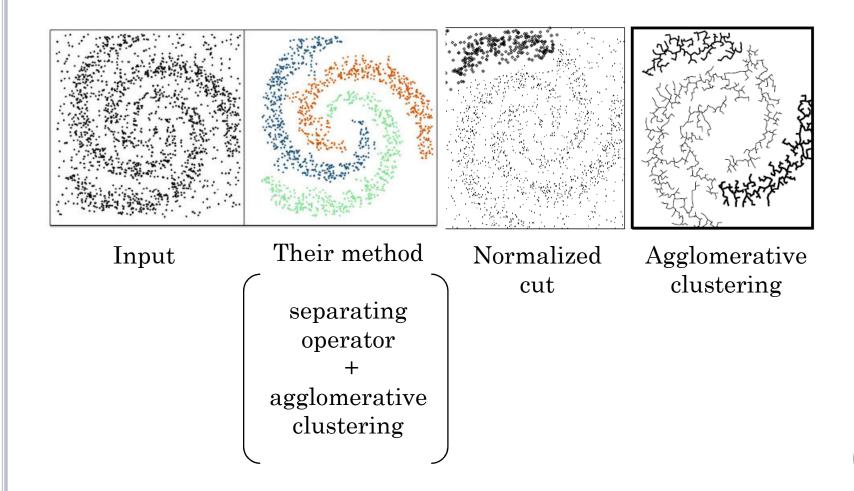
RESULTS



RESULTS (TWO LEVELS IN HIERARCHY)



COMPARATIVE RESULTS



SUMMARY

- Deterministic analysis of random walks on weighted graph (does not use random algorithms to simulate random walks)
- Exploit structure of graph in addition to weights
- Separating operators 'sharpen' the distinction between weights of inter-cluster edge and intracluster edge
- More powerful when incorporated with agglomerative clustering

REFERENCE

- On Clustering Using Random Walks, David Harel and Yehuda Koren, FSTTCS 2011
- Stochastic Image Segmentation by Typical Cuts,
 Fort Collins, Colorado, CVPR 1999
- CHAMELEON: A Hierarchical Clustering Algorithm Using Dynamic Modeling, George Karypis, Eui-Hong (Sam) Han, Vipin Kumar, IEEE Computer 32(1999), 68-75
- http://stat-www.berkeley.edu/users/aldous/RWG/book.html

QUESTION?

COMPUTATION TIME

Table 1. Running time (in seconds; non-optimized) of the various components of the clustering algorithm

| Data Set | Size | Graph construction | Separation | Agglomeration | Overall | Ratio $\frac{Points}{Sec}$ |
|----------|-------|--------------------|------------|---------------|---------|----------------------------|
| DS4 | 8000 | 0.4 | 0.88 | 0.19 | 1.47 | 5434 |
| DS5 | 8000 | 0.41 | 0.83 | 0.19 | 1.43 | 5587 |
| DS6 | 10000 | 0.5 | 1.12 | 0.26 | 1.88 | 5311 |
| DS7 | 8000 | 0.4 | 0.89 | 0.2 | 1.49 | 5358 |
| DS8 | 8000 | 0.39 | 0.93 | 0.2 | 1.52 | 5256 |
| DS9 | 8000 | 0.33 | 0.66 | 0.21 | 1.2 | 6656 |
| DS10 | 3374 | 0.14 | 0.26 | 0.07 | 0.47 | 7178 |

700MHz Pentium III PC

BASIC NOTATION

- Weighted graph G(V, E, w)
- Set of nodes connected to S by a path with at most k edges $V^k(S)$
- Maxima number of edges incident to some single node of G deg(G)
- ullet Subgraph of G induced by S

G(S)