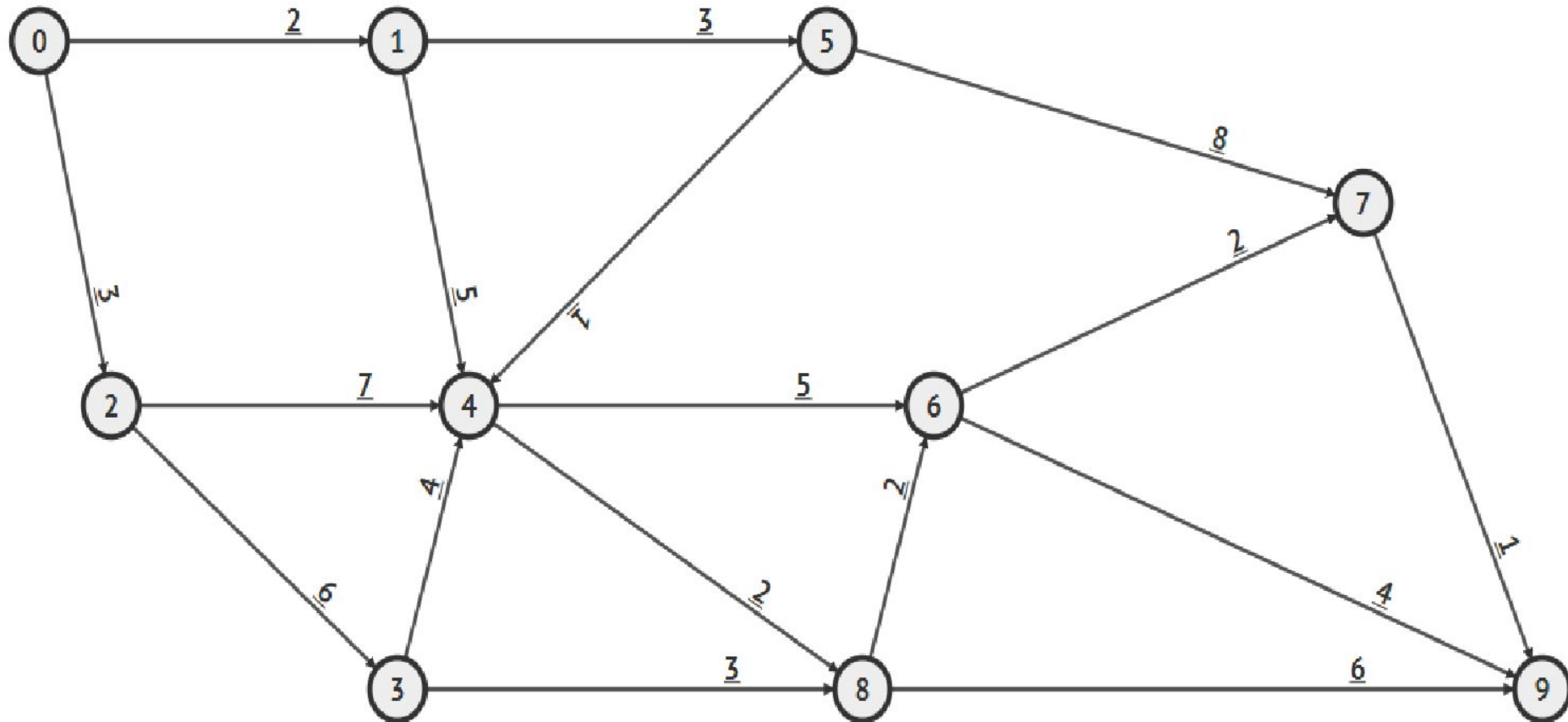


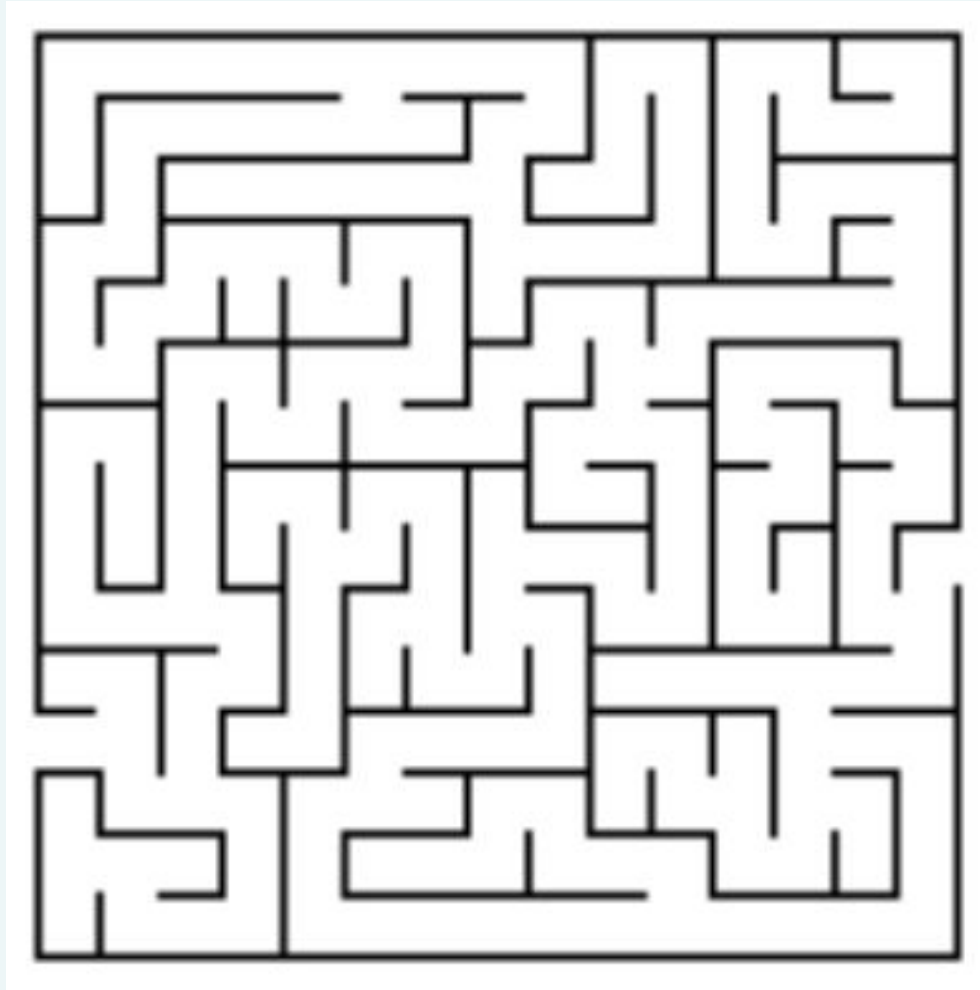
GRAPH TRAVERSAL

Example

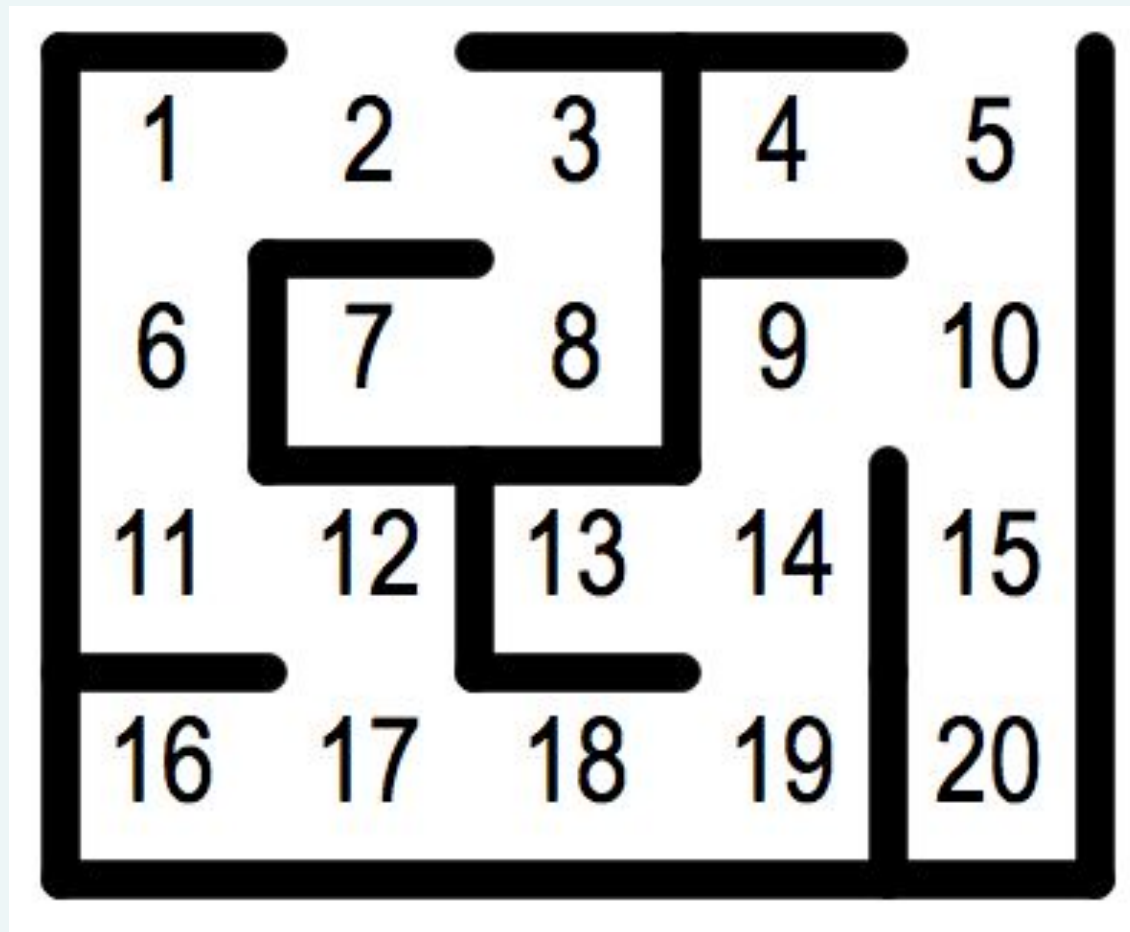
Is there a path from 0 to 9? If there is, list one of them.



Maze Traversal Puzzle



Maze Traversal



Graph Traversal

Graph traversal (also known as graph search) refers to the process of visiting (checking and/or updating) each vertex in a graph. Such traversals are classified by the order in which the vertices are visited

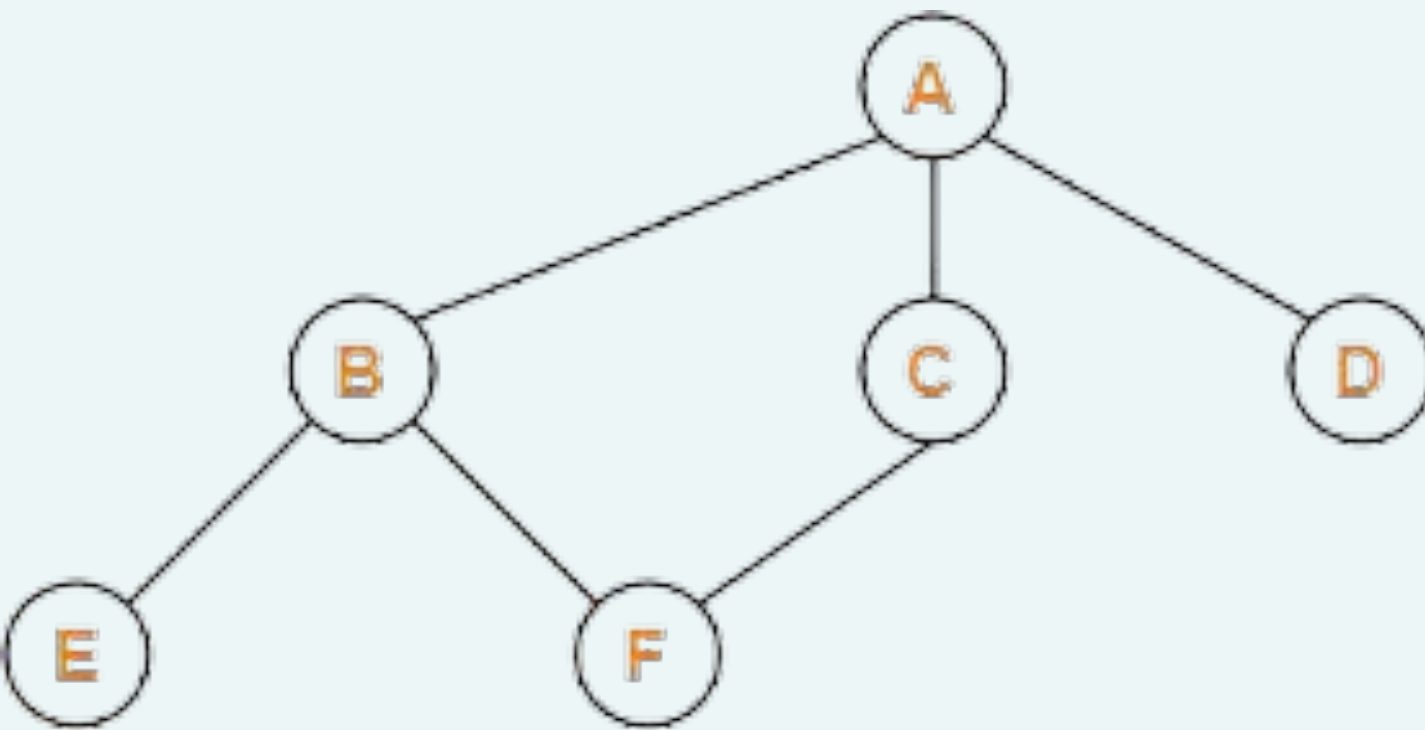


Depth First Search - DFS

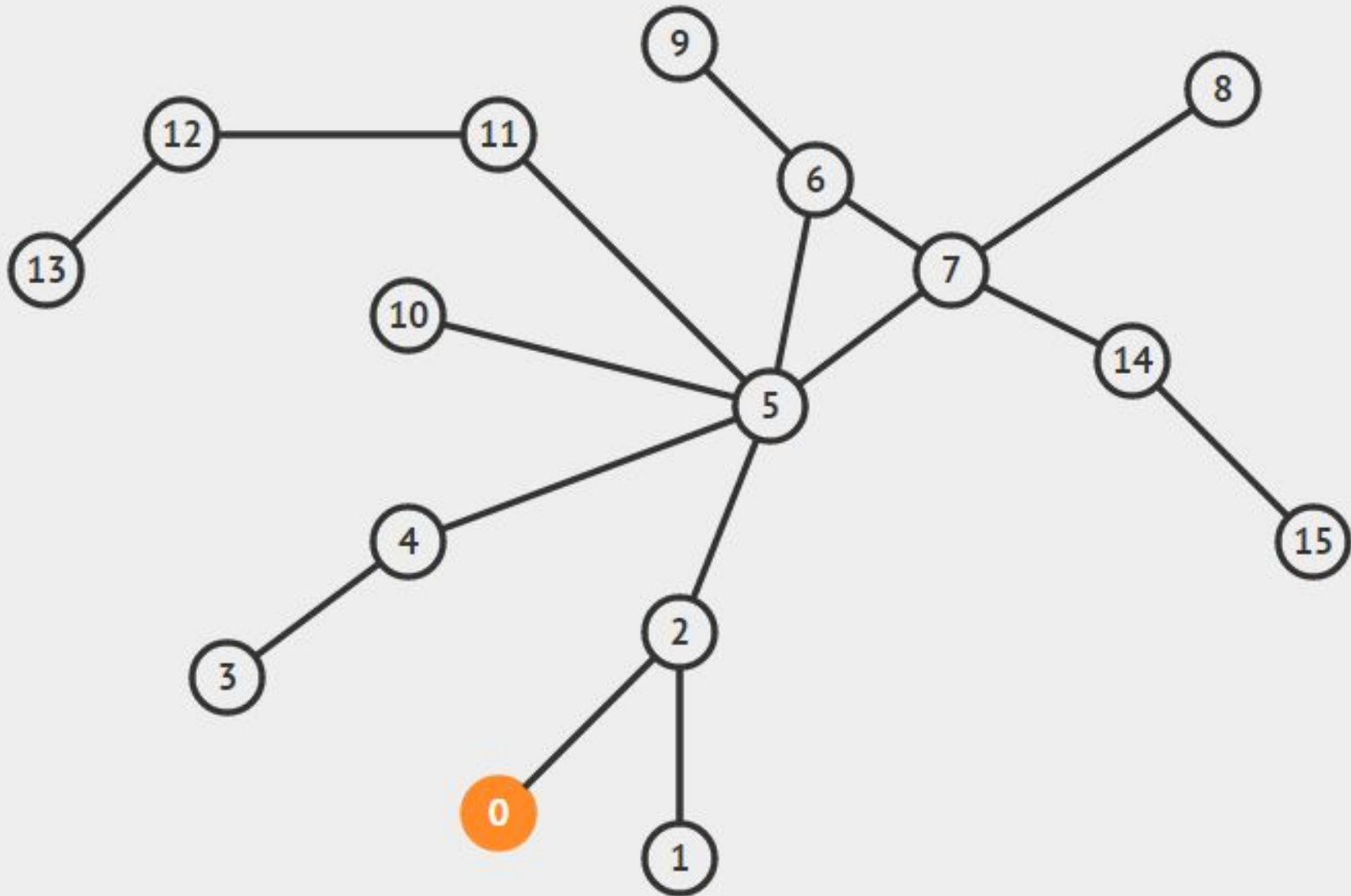
A depth-first search (DFS) is an algorithm for traversing a finite graph. DFS visits the child vertices before visiting the sibling vertices; that is, it traverses the depth of any particular path before exploring its breadth. A stack (often the program's call stack via recursion) is generally used when implementing the algorithm.

```
procedure DFS( $G, v$ ):  
    label  $v$  as discovered  
    for all edges from  $v$  to  $w$  in  $G.adjacentEdges(v)$  do  
        if vertex  $w$  is not labeled as discovered then  
            recursively call DFS( $G, w$ )
```

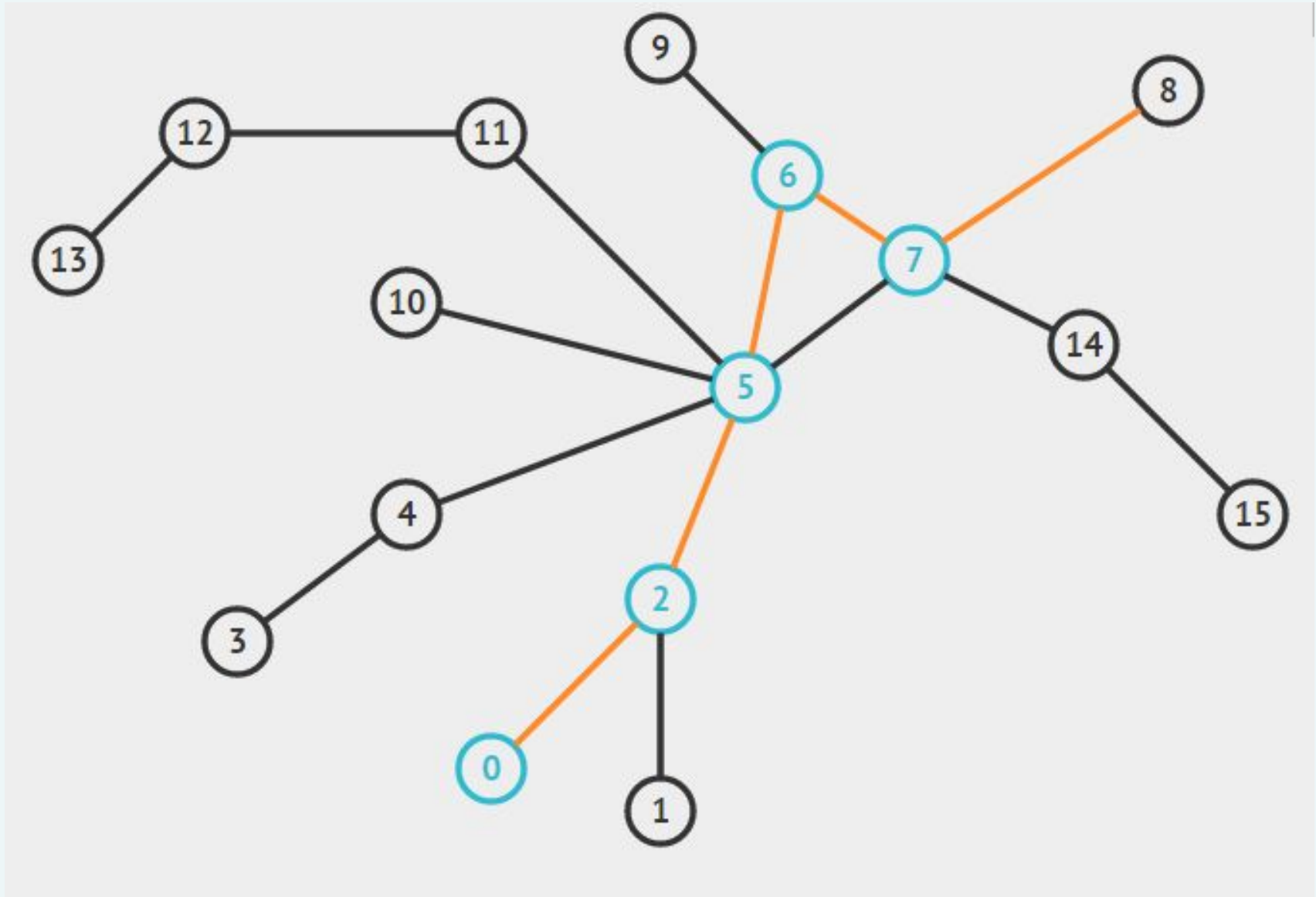
```
procedure DFS( $G, v$ ):  
  label  $v$  as discovered  
  for all edges from  $v$  to  $w$  in  $G.\text{adjacentEdges}(v)$  do  
    if vertex  $w$  is not labeled as discovered then  
      recursively call DFS( $G, w$ )
```



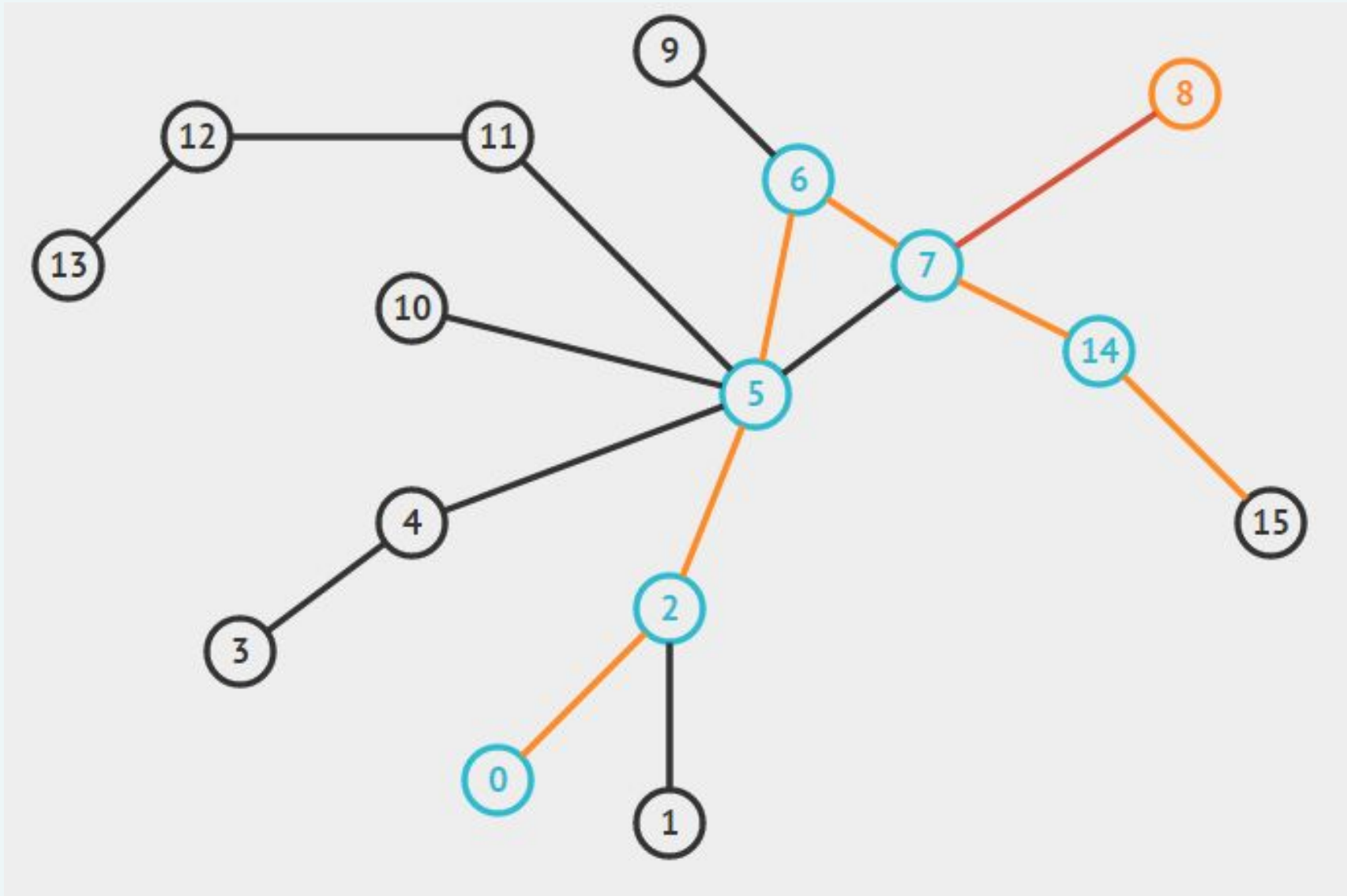
Depth First Search - DFS



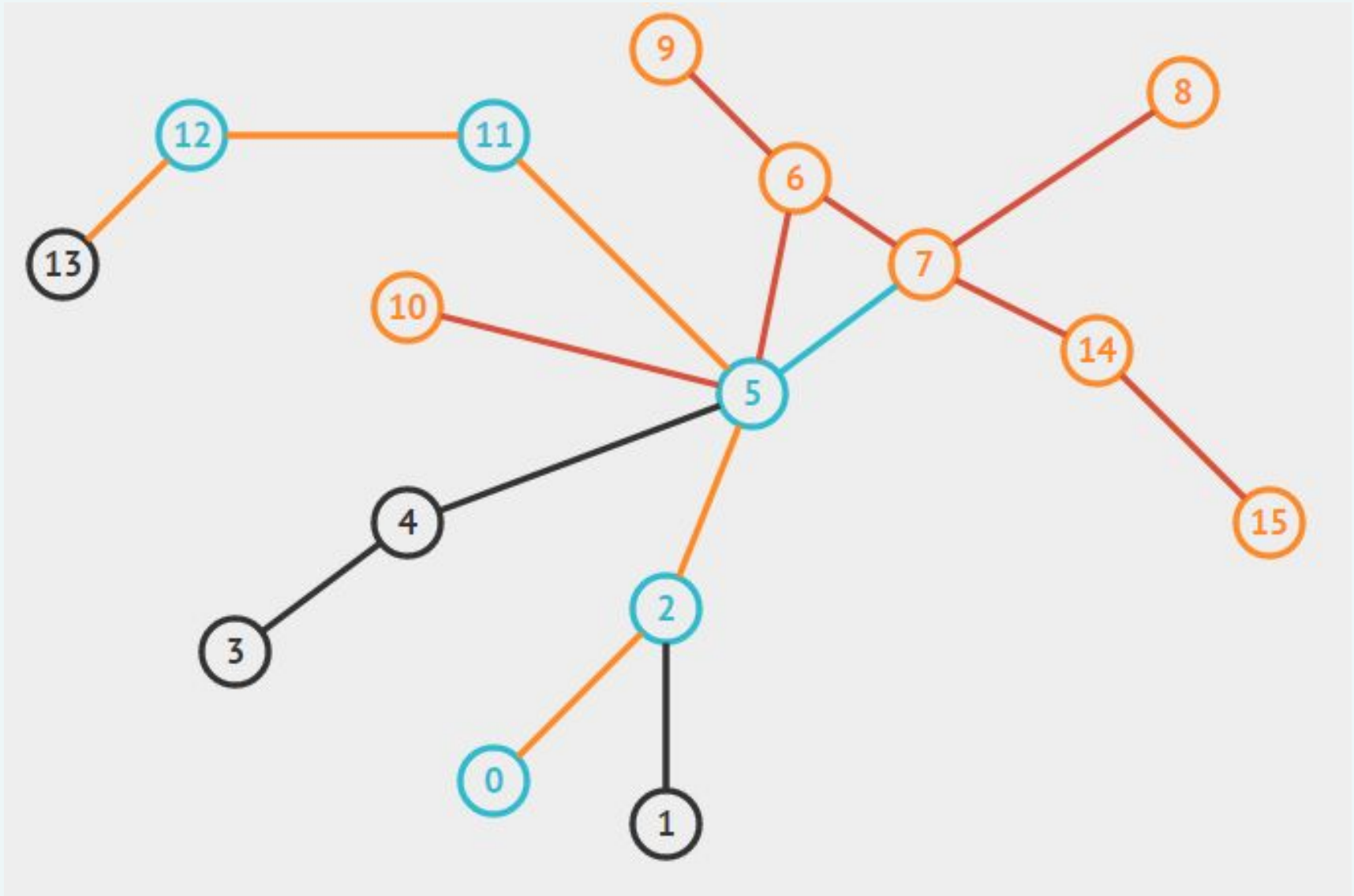
Depth First Search - DFS



Depth First Search - DFS

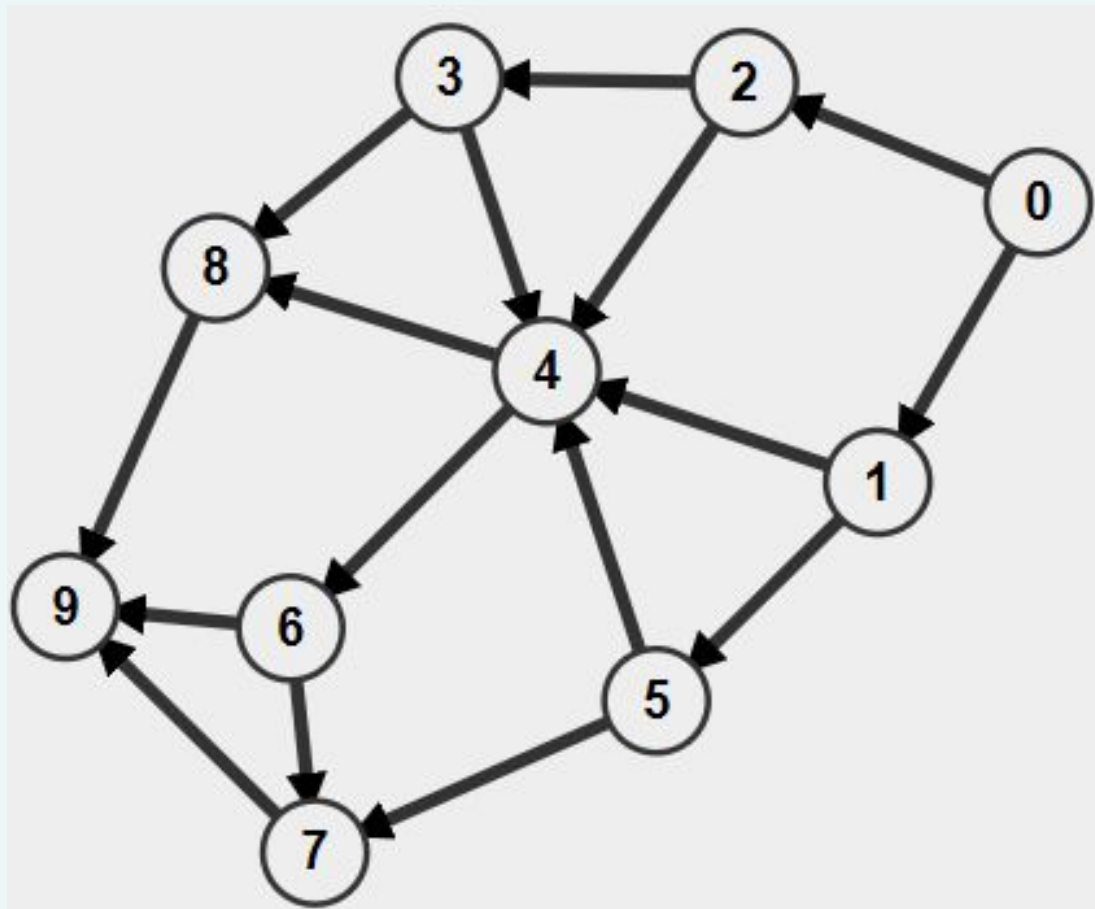


Depth First Search - DFS



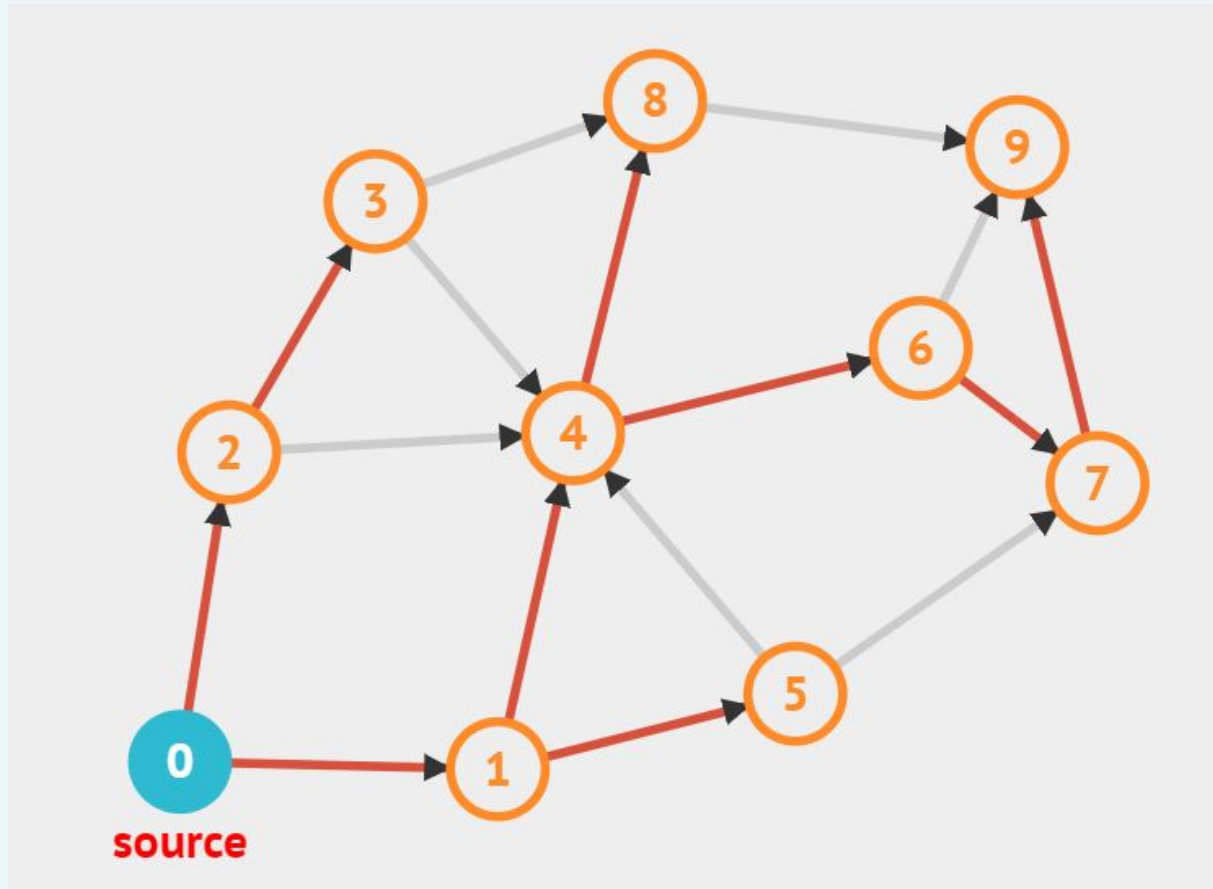
DFS

- Print the lexicographically largest DFS of the following graph starting from 0



DFS

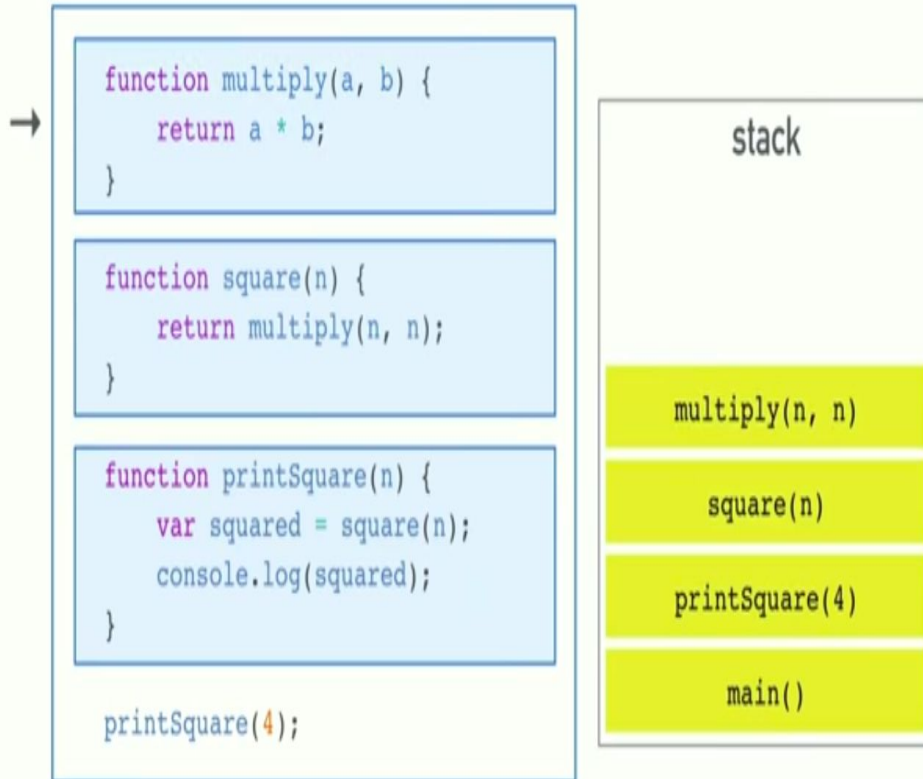
- Print the lexicographically smallest DFS of the following graph



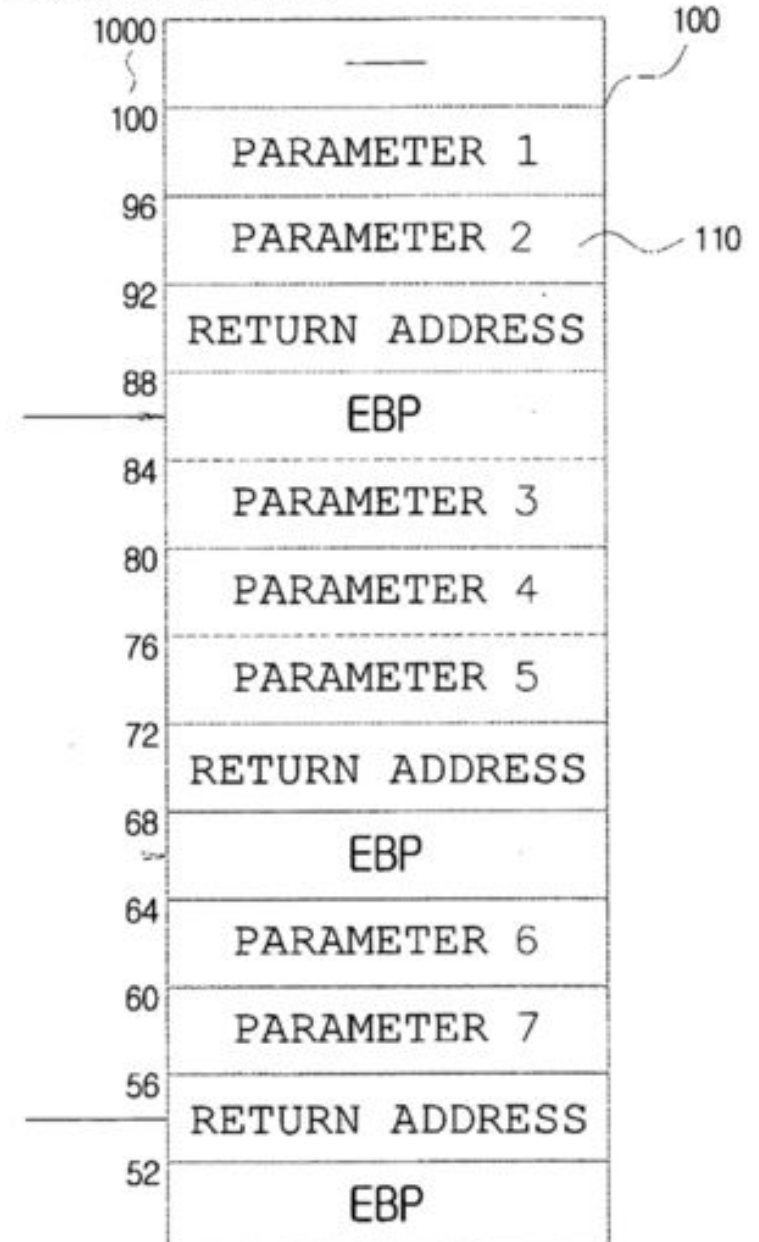
0 1 4 6 7 9 8 5 2 3

Call Stack

JS Call Stack



STACK ADDRESS



Call Stack



```
function multiply(a, b) {  
    return a * b;  
}
```

```
function square(n) {  
    return multiply(n, n);  
}
```

```
function printSquare(n) {  
    var squared = square(n);  
    console.log(squared);  
}
```

```
printSquare(4);
```

stack

multiply(n, n)

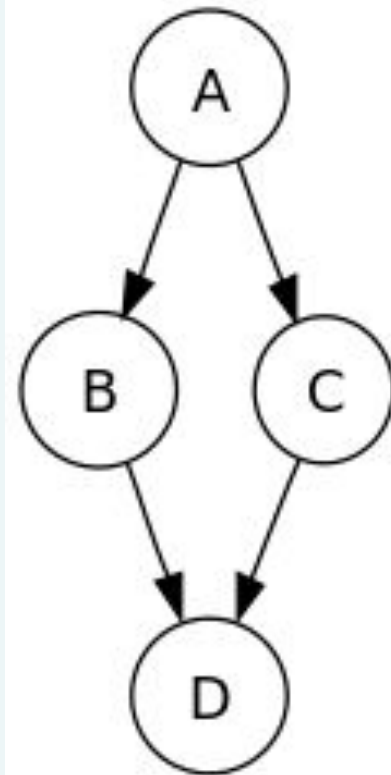
square(n)

printSquare(4)

main()

Traversal Order

- A **preordering** is a list of the vertices in the order that they were first visited by the depth-first search algorithm.
 - E.g. A B D C / A C D B
- A **postordering** is a list of the vertices in the order that they were *last* visited by the algorithm.
 - E.g. D B C A / D C B A
- A **reverse postordering** is the reverse of a postordering, i.e. a list of the vertices in the opposite order of their last visit. Reverse postordering is not the same as preordering.
 - E.g. A C B D / A B C D

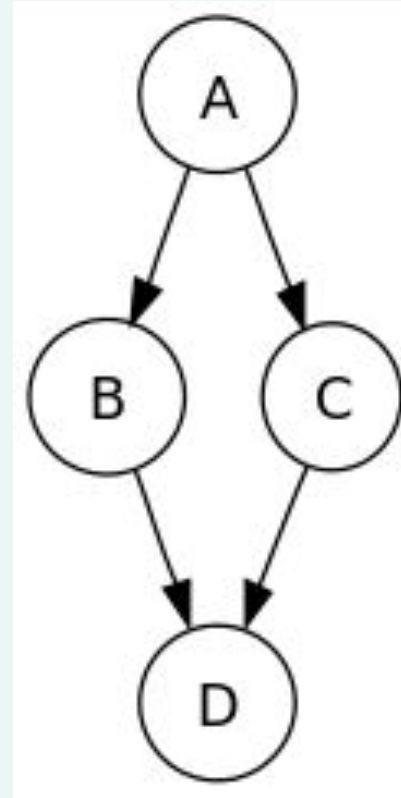


Depth First Search - DFS

```
1  procedure DFS-iterative( $G, v$ ):  
2      let  $S$  be a stack  
3       $S.push(v)$   
4      while  $S$  is not empty  
5           $v = S.pop()$   
6          if  $v$  is not labeled as discovered:  
7              label  $v$  as discovered  
8              for all edges from  $v$  to  $w$  in  $G.adjacentEdges(v)$  do  
9                   $S.push(w)$ 
```

DFS Iterative

- Visited order?
 - A B D C
 - A C D B



Breath First Search

A breadth-first search (BFS) is another technique for traversing a finite graph. BFS visits the sibling vertices before visiting the child vertices

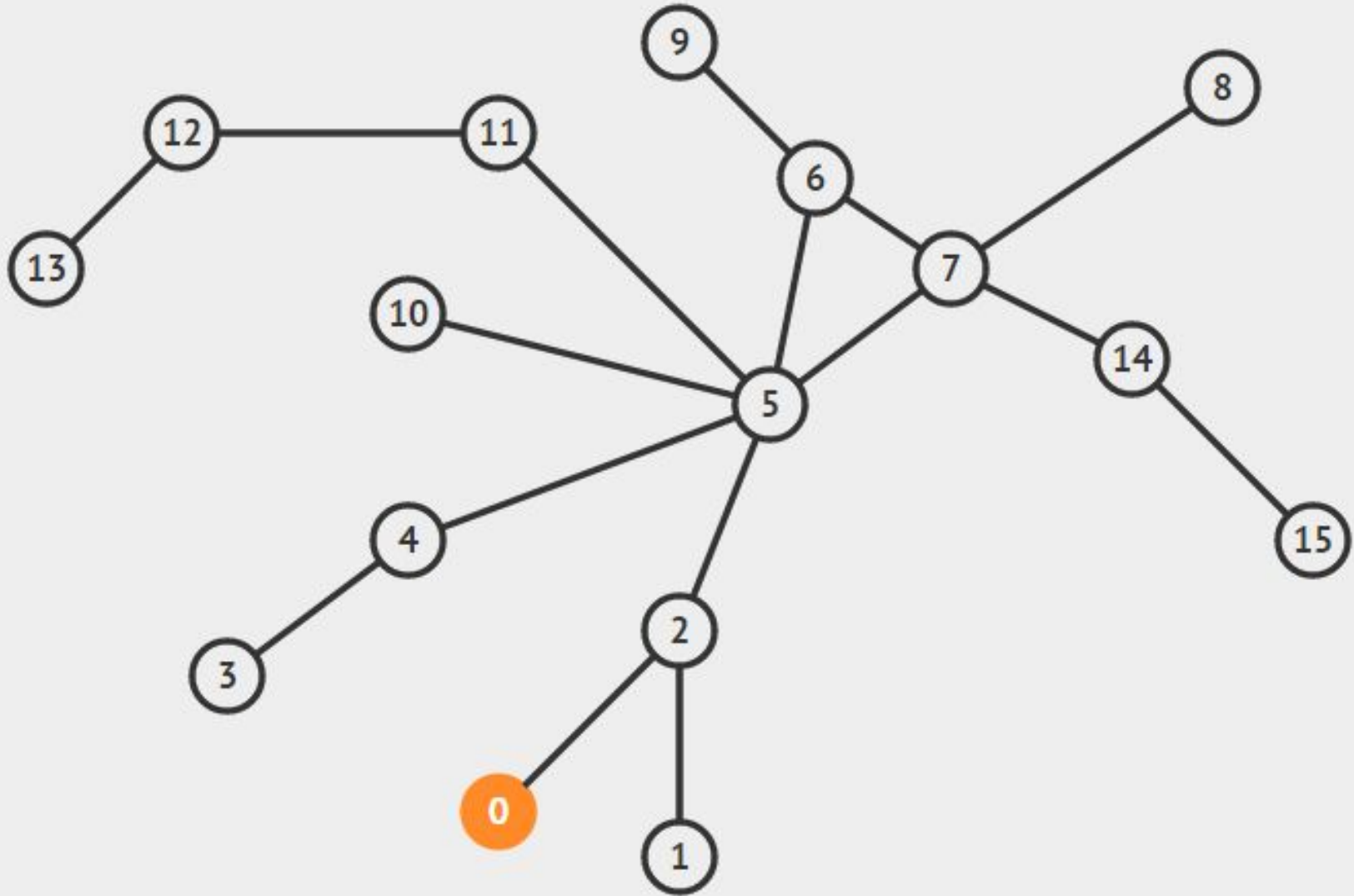
A queue is used in the search process. This algorithm is often used to find the shortest path from one vertex to another.

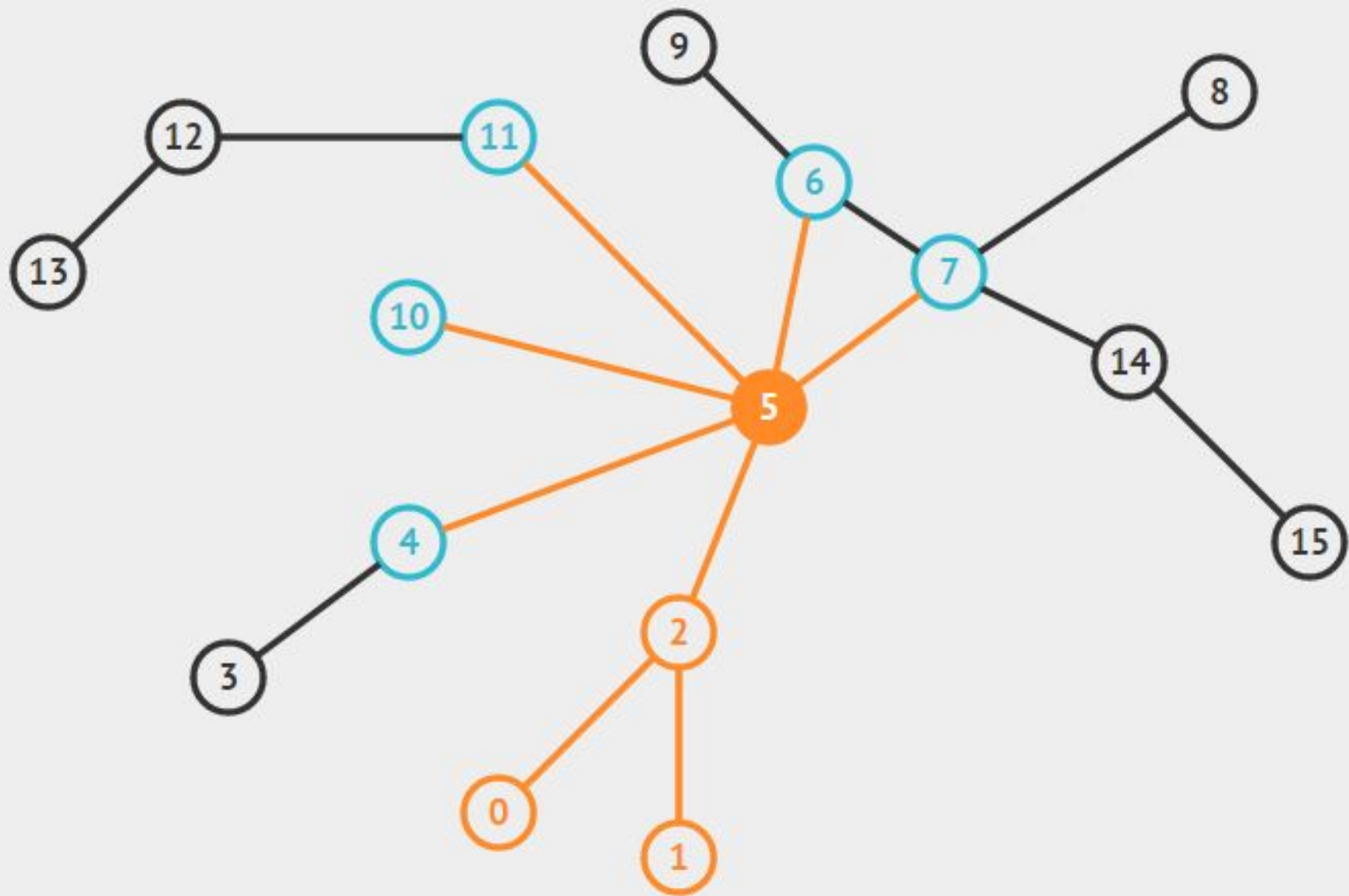
```
procedure BFS( $G, v$ ) is
    create a queue  $Q$ 
    enqueue  $v$  onto  $Q$ 
    mark  $v$ 
    while  $Q$  is not empty do
         $w \leftarrow Q.dequeue()$ 

        for all edges  $e$  in  $G.adjacentEdges(w)$  do
             $x \leftarrow G.adjacentVertex(w, e)$ 
            if  $x$  is not marked then
                mark  $x$ 
                enqueue  $x$  onto  $Q$ 
```

```
procedure BFS( $G, v$ ) is  
    create a queue  $Q$   
    enqueue  $v$  onto  $Q$   
    mark  $v$   
    while  $Q$  is not empty do  
         $w \leftarrow Q.dequeue()$   
  
        for all edges  $e$  in  $G.adjacentEdges(w)$  do  
             $x \leftarrow G.adjacentVertex(w, e)$   
            if  $x$  is not marked then  
                mark  $x$   
                enqueue  $x$  onto  $Q$ 
```

Breadth First Search





Breath First Search

```
1 Breadth-First-Search(Graph, root):
2
3     for each node n in Graph:
4         n.distance = INFINITY
5         n.parent = NIL
6
7     create empty queue Q
8
9     root.distance = 0
10    Q.enqueue(root)
11
12    while Q is not empty:
13
14        current = Q.dequeue()
15
16        for each node n that is adjacent to current:
17            if n.distance == INFINITY:
18                n.distance = current.distance + 1
19                n.parent = current
20                Q.enqueue(n)
```


Applications

- https://www.spoj.com/EIUPROGR/problems/dis2_02/
 - DFS Directed Graph
 - DFS Undirected Graph
 - BFS Directed Graph
 - BFS Undirected Graph
 - Thành phần liên thông
 - Thành phần liên thông 2

Complexity

- $O(V+E)$
- Each vertex is only visited once due to the fact that DFS will only recursively explore a vertex u if $\text{status}[u] = \text{unvisited}$ — $O(V)$
- Every time a vertex is visited, all its k neighbors are explored and therefore after all vertices are visited, we have examined all E edges — ($O(E)$ as the total number of neighbors of each vertex equals to E).

Backtracking

Backtracking is a general algorithm for finding all (or some) solutions to some computational problems, notably constraint satisfaction problems, that incrementally builds candidates to the solutions, and abandons each partial candidate c ("backtracks") as soon as it determines that c cannot possibly be completed to a valid solution.

- Ref:

- https://en.wikipedia.org/wiki/Depth-first_search
- https://en.wikipedia.org/wiki/Breadth-first_search
- <https://en.wikipedia.org/wiki/Backtracking>
- <http://visualgo.net/dfsbfbs.html>