

Goals

- The goal is to learn:
 - how an MPI program runs on a distributed memory system
 - how to write MPI programs
 - how to use basic point-to-point and collective MPI communication
 - how topologies can be defined
 - some useful examples for further work with MPI



Content

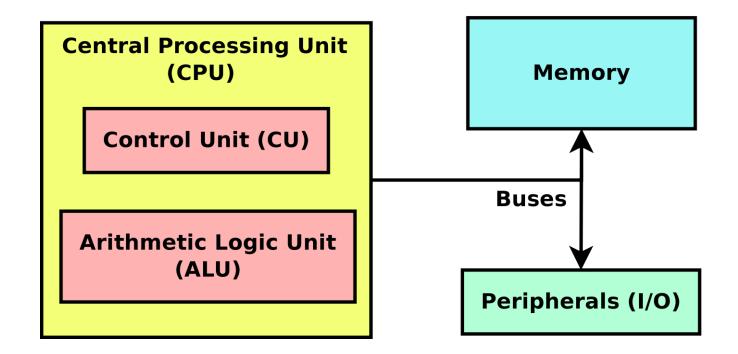
MPI and distributed memory parallelism

Structure of the library functions

Different hands-on exercises



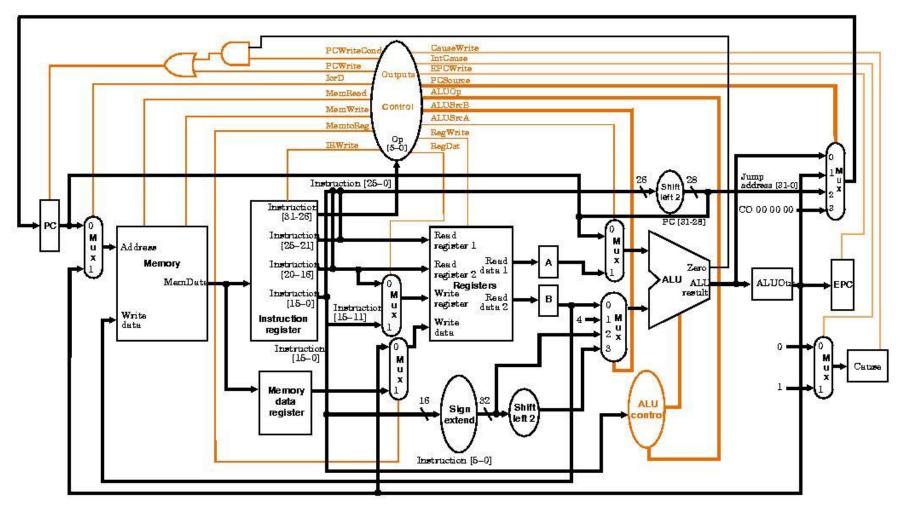
Recap: a computer is...





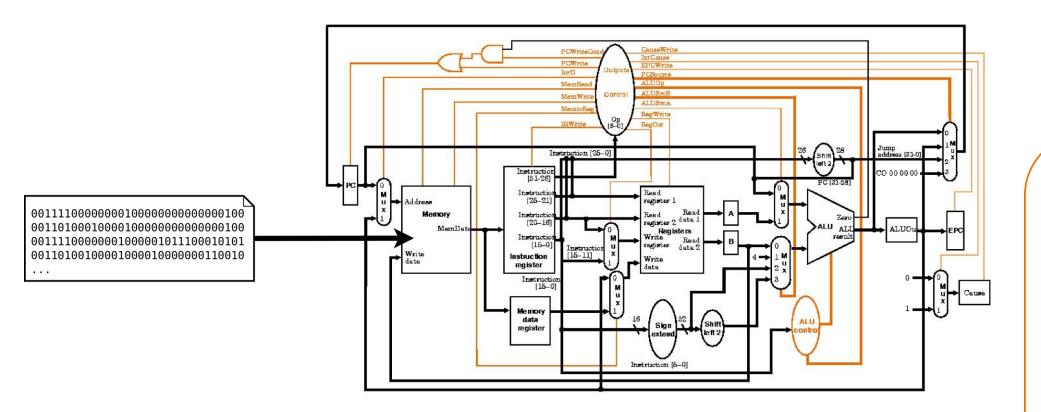
Recap: a CPU is...

Example of the MIPS architecture: control unit lines in orange, data lines in black



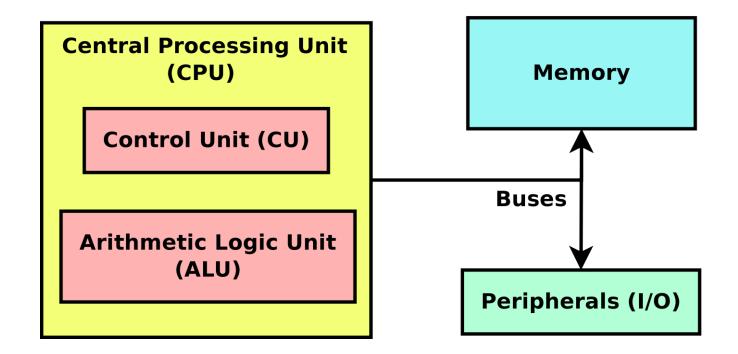


Recap: a program becomes a process...



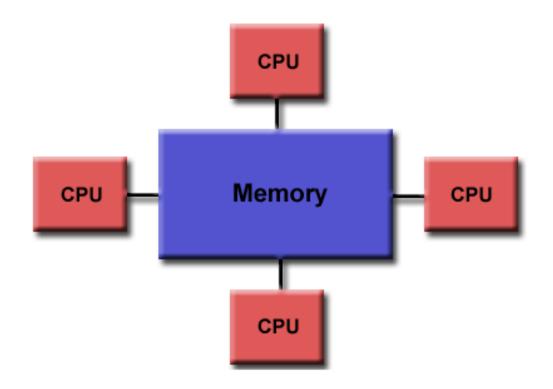


Recap: different types of computer architectures



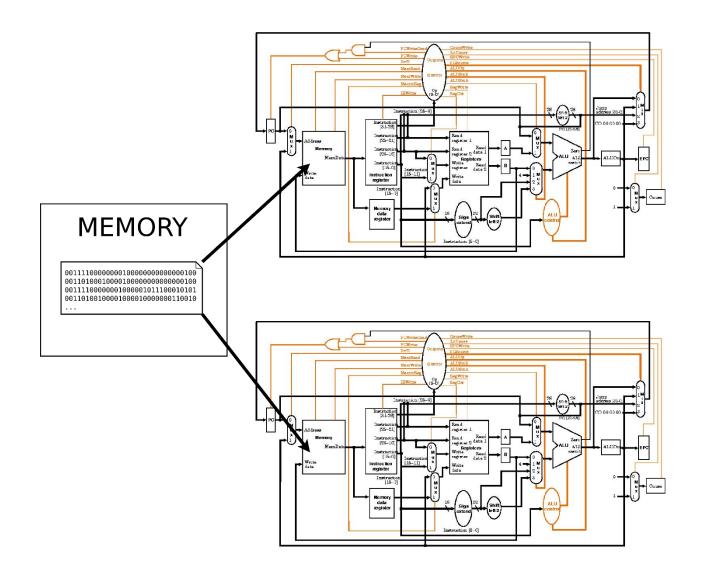


Recap: different types of computer architectures

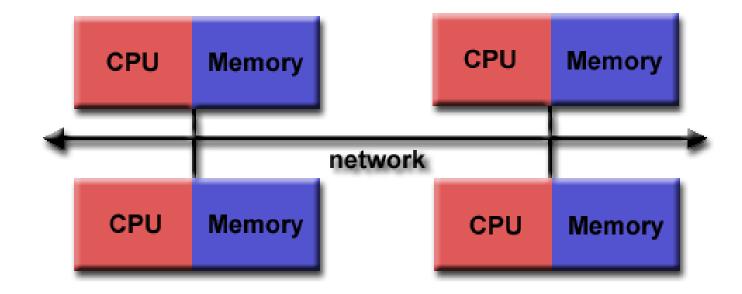




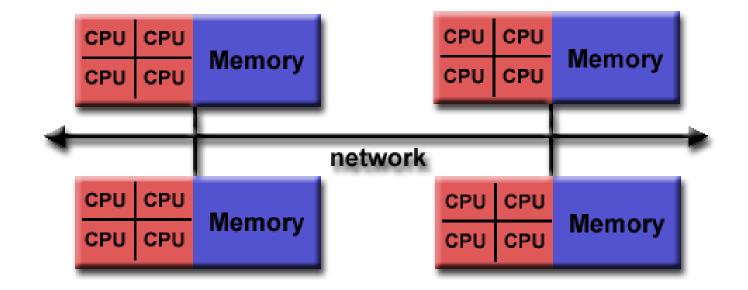
Recap: different types of computer architectures



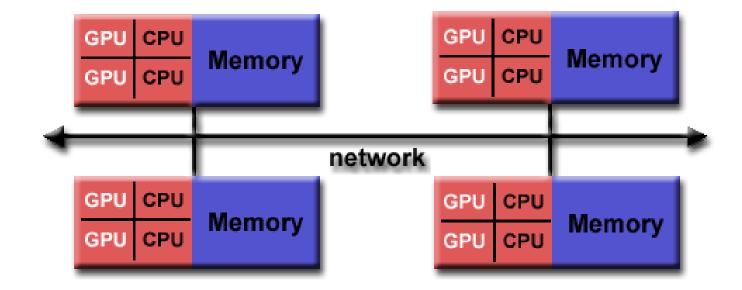




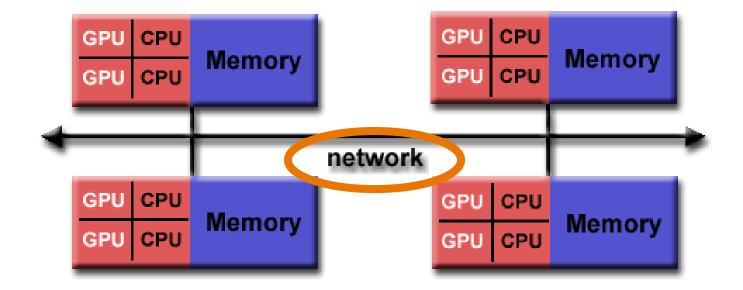






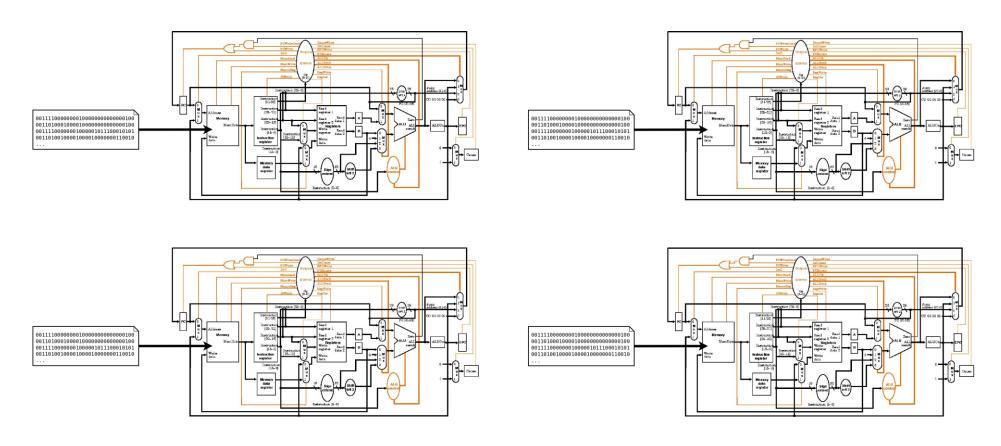






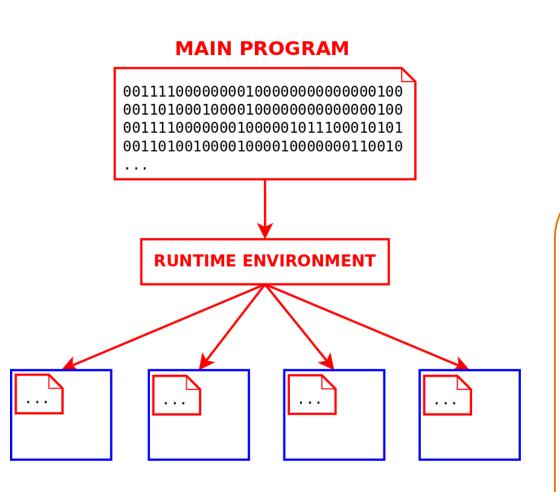


- Basic execution model: SPMD (Single Program, Multiple Data)
- Even MPMD (Multiple Program, Multiple Data) could be possible...



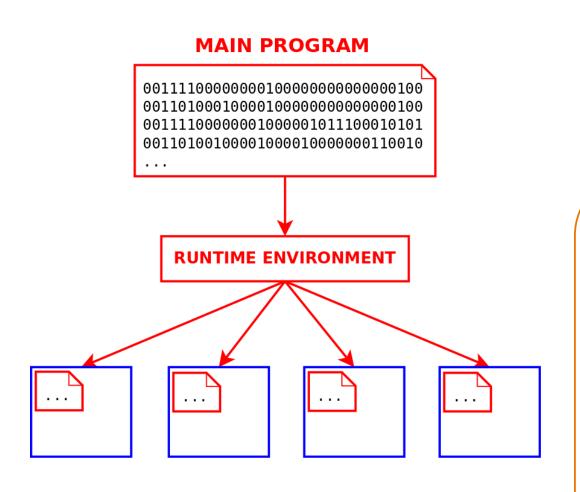


- MPI: Message Passing Interface
- Different processes are created on the one or many machines (nodes)
- Each process executes the same program and has a private copy of all variables
- Inter-process communication by sending/receiving messages andsynchronizations
- Bindings for several languages,
 like C/C++, Fortran, Java or Python



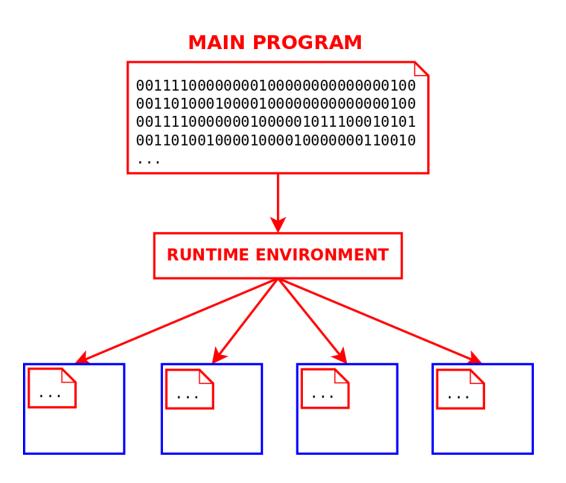


- The interface defines the protocol and the semantic specifications of its functional parts
- Different implementations of the interface can exist
 - IntelMPI
 - OpenMPI
 - MPICH / MVAPICH
- De-facto standard for parallel programming on distributed
 HPC architectures





- In general, MPI requires specific compilation (e.g. mpicc) and specific runtime environment (e.g. mpirun)
- The MPI compiler loads the necessary libraries from the selected MPI implementation
- The MPI runtime is the process that manages communications
 - Internally implemented with daemon processes





Example of multi-process execution with MPI

```
top - 11:05:19 up 41 days, 15:15, 1 user,
                                           load average: 0,52, 0,26, 0,16
Tasks: 473 total, 5 running, 468 sleeping,
                                              0 stopped,
                                                           0 zombie
%Cpu(s): 0,0 us, 0,1 sy, 16,7 ni, 83,2 id, 0,0 wa, 0,0 hi, 0,0 si, 0,0 st
KiB Mem : 49284800 total, 35059112 free,
                                          784368 used, 13441320 buff/cache
KiB Swap: 25165820 total, 25136612 free,
                                           29208 used. 46491000 avail Mem
 PID USER
               PR NI
                         VIRT
                                 RES
                                        SHR S %CPU %MEM
                                                             TIME+ COMMAND
28875 casparl
                   19
                       419216
                                6880
                                       4240 R/ 100,0 \ 0,0
                                                        0:04.27 my mp1 prog
28876 casparl
               39 19
                      419220
                                6876
                                       4236 R 100,0 0,0
                                                          0:04.27 my mpi prog
28877 casparl
                   19
                       419220
                                6868
                                       4240 R 100,0 0,0
                                                          0:04.26 my mpi proq
                                       4236 R 100,0 0,0
28878 casparl
                   19
                       419216
                                6876
                                                          0:04.26 my_mpi_prog
                        15944
                                1820
                                                         53:35.03 pim
3356 root
                    Θ
                                        956 S 1,7 0,0
                                                 .3 0,0 59:26.22 xfsaild/sda3
  477 root
               20
                    Θ
                                          0 S
```

4 processes, each using a single CPU core



Example of multi-process execution with MPI

- Note the difference:
 - OpenMP: 1 process, 4 threads per process

PID USER	PR		VIRT	RES		%CPU %		TIME+ COMMAND
31648 casparl	39	19	228816	1656	1180 R	400,0	0,0	0:22.37 my_omp_prog
84 root	rt	0	Θ	Θ				0:08.41 migration/15
477 root	20	Θ	Θ	Θ	0 S	0,3	Θ,Θ	59:26.57 xfsaild/sda3

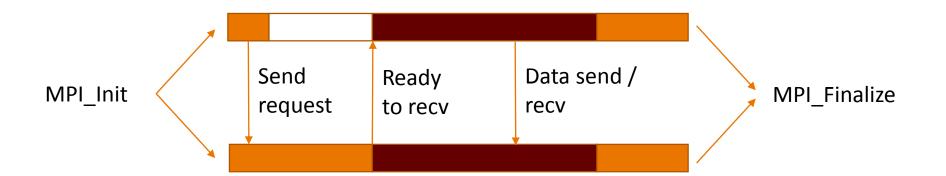
MPI: 4 processes, 1 thread per process

PID U	JSER	PR	NI	VIRT	RES	SHR S	S SCPU	%MEM	TIME+ COMMAND
28875 d	casparl	39	19	419216	6880	4240 F	₹ 100,0	0,0	0:04.27 my_mpi_prog
28876 d	casparl	39	19	419220			100,0		0:04.27 my_mpi_prog
28877 c	casparl	39	19	419220	6868	4240 F	100,0	Θ,Θ	0:04.26 my_mpi_prog
28878 d	casparl	39	19	419216	6876	4236 F	100,0	0,0	0:04.26 my_mpi_prog
3356 r	root	20		15944					53:35.03 pim
477 r	root	20	Θ	Θ	Θ	0 9	0,3	0,0	59:26.22 xfsaild/sda3

... and you may even combine MPI + OpenMP...



- The use of MPI requires some programming efforts to manage all processes coordinately
- In the beginning the MPI environment is initialized, and also finalized at the end of the parallel code
- The basic element is the communicator, which defines
 - communication context
 - set of processes that can communicate between each other inside that context





- The use of MPI requires some programming efforts to manage all processes coordinately
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 - communication context
 - set of processes that can communicate between each other inside that context

	MPI_COMM_WORLD	My_communicator
Process 0	Rank 0	
Process 1	Rank 1	Rank 0
Process 2	Rank 2	Rank 1



General references for MPI

- Specification documents in the web page of the MPI Forum
- Current specification document: <u>version 3.1</u>
- Revisions and discussions for the elaboration of the MPI standard 4.0 are ongoing



- MPI functions are used inside the code and require including a file or module
 - C/C++: #include <mpi.h>
 - Fortran: use mpi (with Fortran 2008, definitely do "use mpi_f08")
- Essential structure for MPI routines
 - C/C++: error = MPI_Routine(arg0, arg1, ...);(the error output may be ignored, but it is very useful for debugging)
 - Fortran: call MPI_ROUTINE(arg1, arg2, ..., ierror)(the error output is always the last argument: VERY IMPORTANT!!!)
- Global communicator: predefined handle
 - C/C++ AND Fortran: MPI_COMM_WORLD (type "MPI_Comm" or "integer")



Every MPI code begins with an initialization of the environment

```
C/C++: int MPI_Init(int *argc, char ***argv)
```

typical call: MPI_Init(&argc, &argv);

argc and argv are taken from the main program, NULL is possible too

Fortran: subroutine MPI_INIT(ierr) integer :: ierr

typical call: call MPI_Init(ierr)

- Every MPI code ends with a finalize statement that stops the environment
 - C/C++: int MPI_Finalize()
 - Fortran: subroutine MPI_FINALIZE(ierr) integer :: ierr



- MPI has different handles that identify specific objects
- Most common handle: predefined constants
 - Defined in mpi.h (C/C++) or in the mpi/mpi_f08 modules (Fortran)
 - E.g. global communicator MPI_COMM_WORLD
- Other handles: type definitions for MPI specific variables or error return codes
 - E.g. communicator types

C/C++:
MPI_Comm my_communicator

Fortran: integer :: my communicator

- with mpi_f08: TYPE(MPI_Comm)



Every MPI process is identified with a number, called "rank"

C/C++: int MPI_Comm_rank(MPI_Comm comm, int *rank)

MPI Communicator

Output argument with size value

ierr argument...

NEVER FORGET

when using the

mpi module!!!



Fortran: subroutine MPI_COMM_RANK(comm, rank, ierr)

"use mpi": integer :: comm, rank, ierr

"use mpi_f08": TYPE(MPI_Comm) :: comm

integer :: rank; integer, optional :: ierr



ierr argument is optional with the module mpi_f08. Try to use this module when possible!



The number of processes in a communicator can be obtained with another routine

C/C++: int MPI_Comm_size(MPI_Comm comm, int *size)

MPI Communicator

Output argument with size value

ierr argument...

NEVER FORGET

when using the

mpi module!!!



Fortran: subroutine MPI_COMM_SIZE(comm, size, ierr)

"use mpi":

integer :: comm, rank, ierr

"use mpi_f08": TYPE(MPI_Comm) :: comm

integer :: rank; integer, optional :: ierr



ierr argument is optional with the module mpi_f08. Try to use this module when possible!



Ready for a first hands-on exercise?

Go to the GitHub web page of the course

https://github.com/sara-nl/PRACE-MPI-OpenMP

- Try to write your first MPI code!
 - C/C++: mpi_hello_world.c
 - Fortran: mpi_hello_world.f

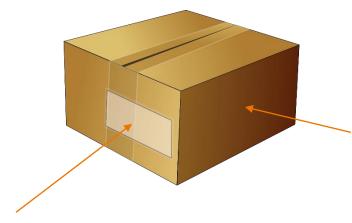


Some ideas about MPI execution model

- The execution environment (e.g. mpirun, srun...) creates the requested processes
- After executing MPI_Init, the processes are located in the same communication environment: MPI_COMM_WORLD
- Before MPI_Init and after MPI_Finalize the processes still exist, but there is no possible interaction or coordination between them
- MPI follows the SPMD (Single Program, Multiple Data) execution model
 - All processes execute exactly the same program
 - The programmer has the responsibility of making processes cooperate
- General way of cooperation between processes
 - Manual distribution of data and/or tasks
 - Message exchanges (send / receive)



- MPI_Send() and MPI_Recv()
 - In order to send and receive it is necessary to create a "parcel"
 - Need for different information: mostly contents and "address label"



- 1. Variable(s) to send/receive
- 2. How many variables to send
- 3. The size of one variable

- 1. Source/Destination address
- 2. Tag
- 3. Communicator (MPI_COMM_WORLD)



- A message is any type of data that is communicated from one MPI process to another (or many) process(es)
- MPI datatypes can be basic or derived
 - Basic types are integer, float, double precision, etc.
 - Derived types are usually sets of basic datatypes and/or other derived datatypes
- The naming of datatypes varies with the programming language



Routine to send a message

```
    C/C++: int MPI_Send(void *buffer, int count, MPI_Datatype datatype,
    int dest, int tag, MPI_Comm comm);
```

Fortran: MPI_SEND(buffer, count, datatype, dest, tag, comm, ierr)
<datatype> buffer(*)
integer :: count, datatype, dest, tag, comm, ierr



Routine to receive a message

```
    C/C++: int MPI_Recv(void *buffer, int count, MPI_Datatype datatype,
    int dest, int tag, MPI_Comm comm, MPI_Status status);
```

Fortran: MPI_RECV(buffer, count, datatype, src, tag, comm, status, ierr)

<datatype> buffer(*)

integer :: count, datatype, src, tag, comm

integer :: status(MPI_STATUS_SIZE), ierr



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

Variable to send



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

#Variables to send



Send / receive: C example

Rank 0 int send_val = 10;

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Datatype of the variable to send



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Datatype of the variable to send



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

Destination rank



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

Tag of the package



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

MPI Communicator



Send / receive: C example

Rank 0 int send_val = 10;

MPI_Send(&send_val, 1, MPI_INT, 1, 10, MPI_COMM_WORLD)

Contains info about message & sender

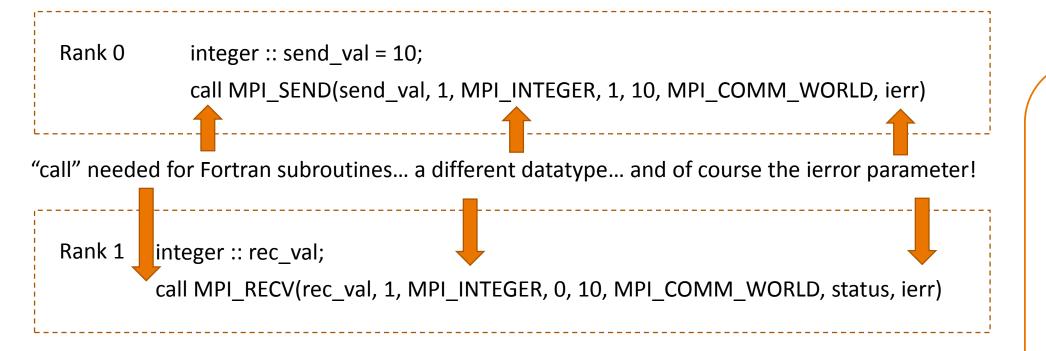
Rank 1 int rec_val;

MPI_Recv(&rec_val, 1, MPI_INT, 0, 10, MPI_COMM_WORLD, &status)

Expects message of size 1 MPI_INT from rank 0 with tag 10



Send / receive: Fortran example (just a couple differences)





C Data Types		Fortran Data Types		
MPI_CHAR	char	MPI_CHARACTER	character(1)	
MPI_WCHAR	wchar_t - wide character			
MPI_SHORT	signed short int			
MPI_INT	signed int	MPI_INTEGER MPI_INTEGER1 MPI_INTEGER2 MPI_INTEGER4	integer integer*1 integer*2 integer*4	
MPI_LONG	signed long int			
MPI_UNSIGNED	unsigned int			
MPI_FLOAT	float	MPI_REAL MPI_REAL2 MPI_REAL4 MPI_REAL8	real real*2 real*4 real*8	
MPI_DOUBLE	double	MPI_DOUBLE_PRECISION	double precision	
MPI_LONG_DOUBLE	long double			
MPI_C_COMPLEX MPI_C_FLOAT_COMPLEX	float _Complex	MPI_COMPLEX	complex	
MPI_C_DOUBLE_COMPLEX	double _Complex	MPI_DOUBLE_COMPLEX	double complex	
MPI_C_BOOL	_Bool	MPI_LOGICAL	logical	
MPI_BYTE	8 binary digits	MPI_BYTE	8 binary digits	
MPI_PACKED	data packed or unpacked with MPI_Pack()/ MPI_Unpack	MPI_PACKED	data packed or unpacked with MPI_Pack()/ MPI_Unpack	



- Matching communication
 - Correct source and destination ranks inside the same communicator
 - Matching tags
 - Matching datatypes
 - ... and a sufficiently large buffer on the receiver's side
- If there are many matching communications, the order is always preserved in time
 - First come, first served
- Some status information is stored in the receiving side
 - It may be ignored by using the constant MPI_STATUS_IGNORE as an argument
 - ... but there is some valuable information inside



Contents of status

C/C++: MPI_Status status

status.MPI_SOURCE

status.MPI_TAG

Fortran: integer status(MPI_STATUS_SIZE)

status(MPI_SOURCE)

status(MPI_TAG)



Contents of status: getting the message count

C/C++: int MPI_Get_count(MPI_Status *status,

MPI_Datatype datatype, int *count)

Fortran: MPI_GET_COUNT(status, datatype, count, ierror)

integer :: status(MPI_STATUS_SIZE)

integer :: datatype, count, ierror

 The "get_count" functions will be very useful in some cases, in order to allocate the receive buffer in advance with enough data for a given message



Hands-on (very quick)

- Try the code on point-to-point communication
 - C/C++: mpi_pnt2pnt.c
 - Fortran: mpi_pnt2pnt.f
- Put the correct source and destination, and the necessary tags



Different modes for sending data

- Synchronous send (MPI_SSEND)
 - Requires synchronization between sender and receiver processes
 - It succeeds when a matching receive call is started
- Buffered send (MPI_BSEND)
 - Asynchronous: the send is always completed, even without a matching receive
 - It requires a buffer declared with MPI_BUFFER_ATTACH
- Generic send (MPI_SEND)
 - It may be either synchronous or buffered send: decision left to the MPI library
- Ready send (MPI_RSEND)
 - Performs an immediate send without checking the receiver side (better avoid...)

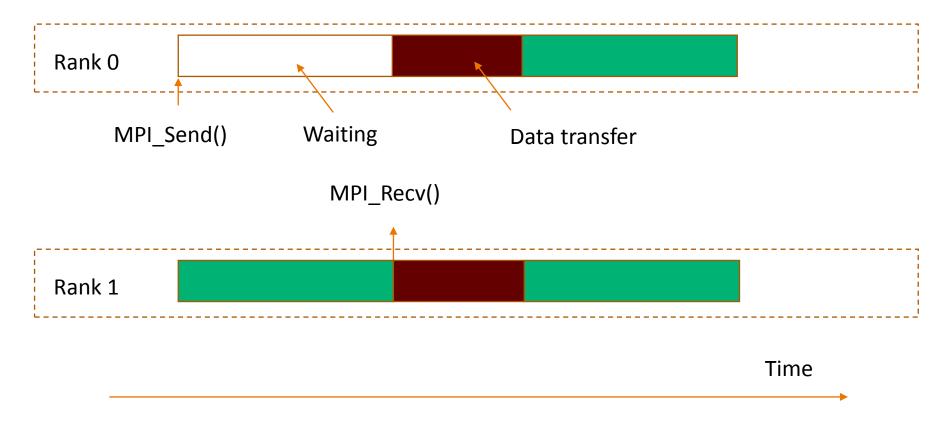


Only one mode for receiving data

- Generic receive (MPI_RECV)
 - Matches all possible send modes
 - It waits until a matching send is posted



MPI_Send (in general) and MPI_Recv are blocking functions



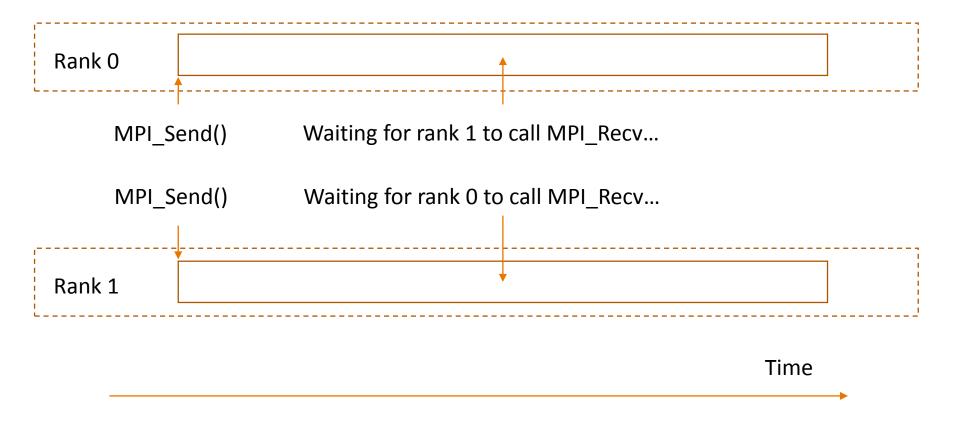


Hands-on

- Try the code of the ping-pong communication
 - C/C++: mpi_pingpong.c
 - Fortran: mpi_pingpong.f
- Have a look at the code and see how it is working
- What can be wrong?

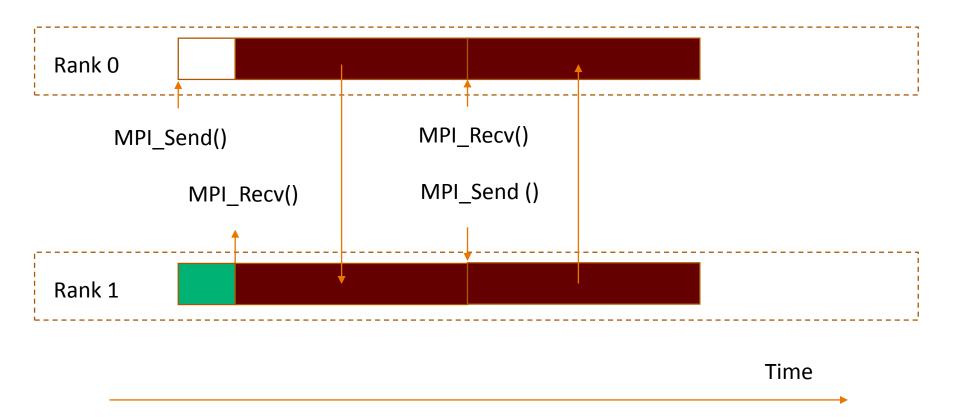


MPI_Send and MPI_Recv are blocking functions: a deadlock may occur!!!





 Solution: reverse the order of MPI_Send and MPI_Recv for one of the communications





General problems with blocking communication

- If the MPI library implementation uses synchronous sends, a programmer may find
 - Deadlocks: complete block of the communication, in which two processes try to send at the same time
 - Serialization: inefficient communication because of a bad scheduling of communications



General problems with blocking communication

- If the MPI library implementation uses synchronous sends, a programmer may find
 - Deadlocks: complete block of the communication, in which two processes try to send at the same time
 - Serialization: inefficient communication because of a bad scheduling of communications
- Solution: non-blocking communications



- It is possible to split the communication process (both send and recv) in two steps
- First step: post a communication request
 - Immediate return of the call and obtain an identifier for the MPI request
 - The effective communication will be done by the MPI library in the background
- Second step: wait for the communication to be completed
 - This will happen after the whole send buffer has been processed and the receive buffer contains all the necessary information



General routine to immediately send a message

```
    C/C++: int MPI_Isend(void *buffer, int count, MPI_Datatype datatype,
    int dest, int tag, MPI_Comm comm, MPI_Request request);
```

Fortran: MPI_ISEND(buffer, count, datatype, dest, tag, comm, request, ierr)<datatype> buffer(*)

integer :: count, datatype, dest, tag, comm, request, ierr



General routine to immediately receive a message

```
    C/C++: int MPI_Irecv(void *buffer, int count, MPI_Datatype datatype,
    int src, int tag, MPI_Comm comm, MPI_Request request);
```

Fortran: MPI_IRECV(buffer, count, datatype, src, tag, comm, request, ierr)
 <datatype> buffer(*)
 integer :: count, datatype, src, tag, comm, request, ierr



General routine wait for a communication to complete

C/C++: int MPI_Wait(MPI_Request request, MPI_Status status);

Fortran: MPI_WAIT(request, status, ierr)

integer :: request, status(MPI_STATUS_SIZE), ierr



- In a practical way, a call to the wait routine right after an immediate send or receive is analogous to a call to the blocking send or receive, respectively
- In any case, the request obtained from the isend or irecv should be matched in wait
- Between the two steps (immediate communication and wait) a process is allowed to perform additional operations that are unrelated to the communication
- This means the following:
 - The sending process is not allowed to access or reuse the sending buffer or any of the request information
 - The completeness of the operation is only ensured after calling the wait function
 - at least for C/C++



Non-blocking communication (Fortran only)

- There is a special issue with Fortran: it is an optimizing language, so a call to the wait routine may not imply that the send or recv buffers are ready for use again
- After MPI 3.0, an additional line of code is required In order to avoid any problem with buffer synchronization in Fortran
- Necessary steps:
 - Use the module mpi_f08 and declare the isend or irecv buffer as asynchronous

```
<datatype>, asynchronous :: buffer
```

Execute the following line of code always immediately after MPI_WAIT

```
IF (.NOT. MPI_ASYNC_PROTECTS_NONBLOCKING)
& CALL MPI_F_SYNC_REG( buf )
```



- It is also possible to perform a wait in a different way using the test routine
- The use of test requires a loop until the request is fulfilled
 - flag == 1 in C or flag == .TRUE. in Fortran
 - Fortran: this is equivalent to wait, so the asynchronous buffer should also be protected using MPI_F_SYNC_REG as indicated
- Test function
 - C/C++: int MPI_Test(MPI_Request request, int flag, MPI_Status status);
 - Fortran: MPI_TEST(request, flag, status, ierr)

integer :: request, status(MPI_STATUS_SIZE), ierr

logical :: flag



- Another features include It is also possible to perform a wait in a different way using the test routine
- The use of test requires a loop until the request is fulfilled
 - flag == 1 in C or flag == .TRUE. in Fortran
 - Fortran: this is equivalent to wait, so the asynchronous buffer should also be protected using MPI_F_SYNC_REG as indicated
- Test function
 - C/C++: int MPI_Test(MPI_Request request, int flag, MPI_Status status);
 - Fortran: MPI_TEST(request, flag, status, ierr)

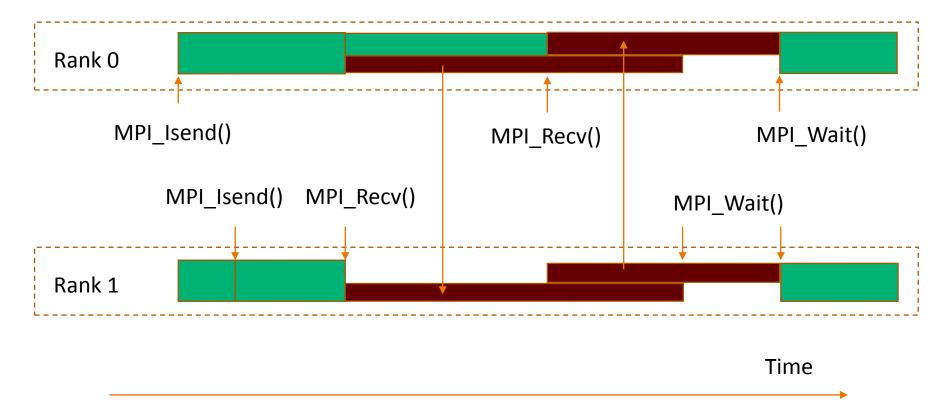
integer :: request, status(MPI_STATUS_SIZE), ierr

logical :: flag



Non-blocking communication example

- The processing of two MPI_Isend() and MPI_Recv() calls may be overlapped
- Deadlocks are gone, and it provides a way to avoid serialization!





Hands-on

- Work with the following code:
 - C/C++: mpi_pingpong_nonblocking.c
 - Fortran: mpi_pingpong_nonblocking.f
- Adapt the code to be really non-blocking
- Use MPI_Isend and/or MPI_Irecv conveniently
- Check the performance of the code compared to the blocking version
 - C/C++: mpi_pingpong_blocking.c
 - Fortran: mpi_pingpong_blocking.f
- Which is the fastest version, and how much fast is it?
- There is another function called MPI_SENDRECV... Can you check the standard specification and see if it would be possible to use it here?



- Sometimes it is necessary to use many point-to-point communications at a time, with specific coordination and synchronization between processes
- Collective functions provide optimized implementations for some common communication patterns
- Different types of collectives
 - Data movement
 - Computation
 - Synchronization

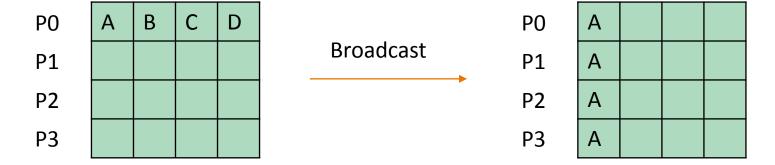


One-to-all

P0	Α	В	U	D
P1				
P2				
Р3				



One-to-all





Broadcast

```
C/C++: int MPI_Bcast(void *buffer, int count, MPI_Datatype datatype, int root, MPI_Comm comm);
```

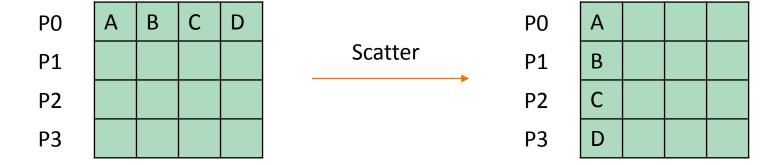
Fortran: MPI_BCAST(buffer, count, datatype, root, comm, ierr)

<datatype> buffer(*)

integer :: count, datatype, root, comm, ierr



One-to-all





```
Scatter (with the option of defining MPI_IN_PLACE as sendbuffer)C/C++: int MPI_Scatter(
```

```
void *sendbuffer, int sendcount, MPI_Datatype senddatatype,
void *recvbuffer, int recvcount, MPI_Datatype recvdatatype,
int root, MPI_Comm comm);
```

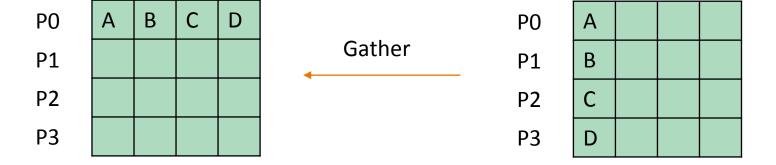
Fortran: MPI_SCATTER(sendbuffer, sendcount, senddatatype, recvbuffer, recvcount, recvdatatype, root, comm, ierr)

<datatype> sendbuffer(*), recvbuffer(*)

integer :: sendcount, senddatatype, recvcount, recvdatatype, root, comm, ierr



All-to-one





Gather (with the option of defining MPI_IN_PLACE as sendbuffer)C/C++: int MPI Gather(

void *sendbuffer, int sendcount, MPI_Datatype senddatatype, void *recvbuffer, int recvcount, MPI_Datatype recvdatatype, int root, MPI_Comm comm);

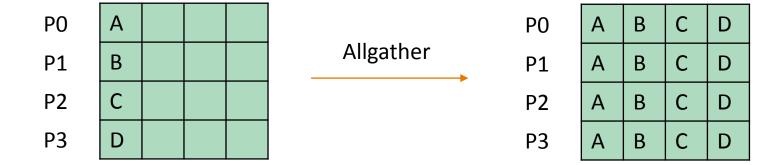
Fortran: MPI_GATHER(sendbuffer, sendcount, senddatatype, recvbuffer, recvcount, recvdatatype, root, comm, ierr)

<datatype> sendbuffer(*), recvbuffer(*)

integer :: sendcount, senddatatype, recvcount, recvdatatype, root, comm, ierr



All-to-all





Allgather int MPI_Allgather(C/C++: void *sendbuffer, int sendcount, MPI_Datatype senddatatype, void *recvbuffer, int recvcount, MPI_Datatype recvdatatype, MPI_Comm comm); MPI_ALLGATHER(sendbuffer, sendcount, senddatatype, Fortran: recvbuffer, recvcount, recvdatatype, comm, ierr) <datatype> sendbuffer(*), recvbuffer(*) integer :: sendcount, senddatatype, recvcount, recvdatatype, comm, ierr



All-to-all

P0	Α	Α	Α	Α		P0	Α	В	C	D
P1	В	В	В	В	Alltoall	P1	Α	В	C	D
P2	С	C	С	С		P2	Α	В	С	D
Р3	D	D	D	D		Р3	Α	В	С	D



```
Alltoall
```

```
int MPI_Alltoall(
C/C++:
              void *sendbuffer, int sendcount, MPI_Datatype senddatatype,
              void *recvbuffer, int recvcount, MPI_Datatype recvdatatype,
              MPI_Comm comm);
            MPI_ALLTOALL(sendbuffer, sendcount, senddatatype,
Fortran:
                           recvbuffer, recvcount, recvdatatype,
                           comm, ierr)
<datatype> sendbuffer(*), recvbuffer(*)
integer :: sendcount, senddatatype, recvcount, recvdatatype, comm, ierr
```



- Computational operations are also possible
 - Most common one: reduce





integer :: count, datatype, op, comm, ierr

Reduce int MPI_Reduce(C/C++: void *sendbuffer, void *recvbuffer, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm); MPI_ALLTOALL(sendbuffer, recvcount, count, datatype Fortran: op, root, comm, ierr) <datatype> sendbuffer(*), recvbuffer(*)



Collective communication: reduction operations

- There are different predefined operations for reduce
 - MPI_MIN, MPI_MAX, MPI_SUM, MPI_PROD
- Another option: custom definition of the operation
 - C/C++: int MPI_Op_create(MPI_User_function *function,int commute, MPI_Op *op);
 - Fortran: MPI_OP_CREATE(function, commute, op, ierr)

external :: function

logical :: commute

integer :: op, ierr



Other collective operation: synchronization barrier

C/C++: int MPI_Barrier(MPI_Comm comm);

Fortran: MPI_BARRIER(comm, ierr)

integer :: comm, ierr



Hands-on (1)

Try to solve the exercise on collective functions proposed in this directory

https://github.com/sara-nl/PRACE-MPI-OpenMP/tree/master/MPI/collectives

The original codes are stored in the PRACE CodeVault

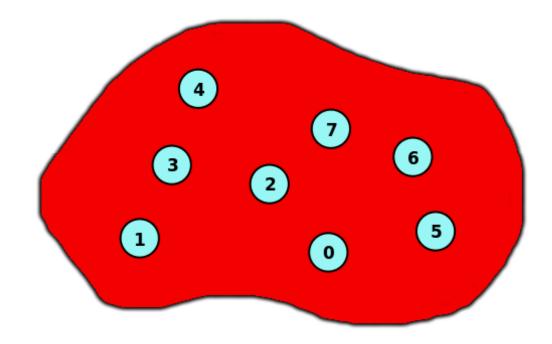


Hands-on (2)

- ... we've already come a long way here. Time for a parallelization from scratch!
- Try to parallelize the computation of pi
 - C/C++: mpi_pi.c
 - Fortran: mpi_pi.f
- Three main tasks need to be performed
 - Distribution of the different iterations between processes
 - Getting the final complete result from all processes
 - Checking that there is scalability with a significant amount of threads



- Well known initial communicator: MPI_COMM_WORLD
- Main representation: pool of processes

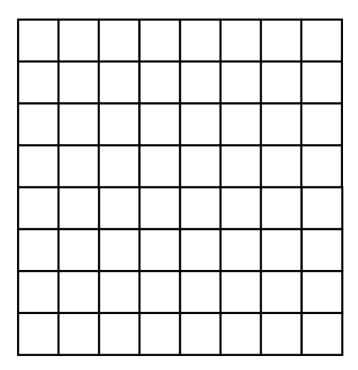




- Possible definition of specific connections between processes to obtain a more structured view of a communicator
- Reasons to obtain this?

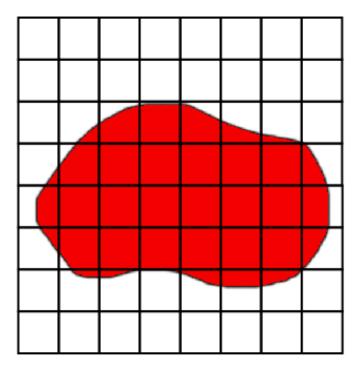


- Possible definition of specific connections between processes to obtain a more structured view of a communicator
- Reasons to obtain this?
- Mimic the structure of the data that we are working on!





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- Mimic the structure of the data that we are working on!





- MPI provides virtual topologies to organize processes
- Easy access to neighbor processes for communication
- Most common case: cartesian topology
 - C/C++: int MPI_Cart_create(MPI_Comm comm_old, int ndims, int dims[], int periods[], int reorder,
 MPI_Comm *comm_cart);
 - Fortran: MPI_CART_CREATE(comm_old, ndims, periods, reorder, comm_cart, ierr)

integer :: comm_old, ndims, dims(*), comm_cart, ierr

logical :: periods(*), reorder



- MPI provides virtual topologies to organize processes
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integer :: comm_old, ndims, dims(*), comm_cart, ierr

logical :: periods(*), reorder



 Connections between processes in cartesian topologies are established by their given coordinates: easy to calculate and to convert to/from rank

C/C++: int MPI_Cart_coords(MPI_Comm comm, int rank, int maxdims, int coords[]);

Fortran: MPI_CART_COORDS(comm, rank, maxdims, coords, ierr)

integer :: comm, rank, maxdims, coords, ierr



 Connections between processes in cartesian topologies are established by their given coordinates: easy to calculate and to convert to/from rank

C/C++: int MPI_Cart_rank(MPI_Comm comm, int coords[], int *rank);

Fortran: MPI_CART_RANK(comm, coords, rank, ierr)

integer :: comm, rank, maxdims, coords, ierr



Cartesian topologies may also be subdivided in their different dimensions

```
C/C++: int MPI_Cart_sub(MPI_Comm comm, int remain_dims[],
MPI_Comm *comm_new);
```

Fortran: MPI_CART_SUB(comm, remain_dims, comm_new, ierr)

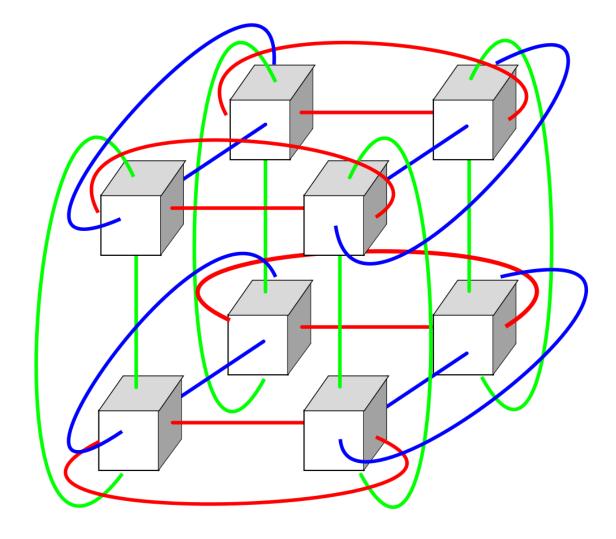
integer :: comm, comm_new, ierr

logical :: remain_dims(*)



Hands-on

- Just do this:
 - C/C++: mpi_torus.c
 - Fortran: mpi_torus.f





- Communicators also have maximum flexibility to be created with a custom amount of threads, and later reorganized
- It is necessary to operate at the level of groups of processes to define a completely custom communicator
 - C/C++: int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI Comm *newcomm);
 - Fortran: MPI_COMM_SPLIT(comm, color, key, newcomm,ierr)

integer :: comm, color, key, newcomm, ierr



- Communicators also have maximum flexibility to be created with a custom amount of threads, and later reorganized
- It is necessary to operate at the level of groups of processes to define a completely custom communicator
 - C/C++: int MPI_Comm_group(MPI_Comm comm, MPI_Group *group);
 - Fortran: MPI_COMM_GROUP(comm, group, ierr)

integer :: comm, group, ierr



- Communicators also have maximum flexibility to be created with a custom amount of threads, and later reorganized
- It is necessary to operate at the level of groups of processes to define a completely custom communicator
 - C/C++: int MPI_Group_incl(MPI_Group group, int nranks, int ranks[],
 MPI Group *newgroup);
 - Fortran: MPI_GROUP_INCL(group,nranks,ranks,newgroup, ierr)
 - integer :: group, nranks, ranks(*), newgroup, ierr



- Communicators also have maximum flexibility to be created with a custom amount of threads, and later reorganized
- It is necessary to operate at the level of groups of processes to define a completely custom communicator
 - C/C++: int MPI_Comm_create(MPI_Comm comm, MPI_Group group,MPI Comm *newcomm);
 - Fortran: MPI_COMM_CREATE(comm, group, newcomm, ierr)
 - integer :: comm, group, newcomm, ierr



Hands-on

- Take the previous example of the 3D torus
- Create communicators for each dimension of it
- Check the values of the ranks for every communicator



General comments (1)

- MPI is based on processes, but also provides threading possibilities
 - MPI_InitThread(int *argc, char ** argv[], int thr_requested, int thr_provided);
 - Varying level of support by the different MPI implementations
- MPI provides a quite complete I/O library with plenty of functionality
 - Recommended: use of high-level libraries that use it with wrappers
 - E.g. PnetCDF, SIONlib
- Derived data types can be defined and may be useful, but relatively cumbersome
 - Try to check if it is possible to codify relatively complex communication using basic data types
 - Use the general type MPI_PACKED and different routines (MPI_Pack and MPI_Unpack to put different data types in the same stream



General comments (& 2)

- MPI also supports one-sided and shared memory communication
 - MPI_Get / MPI_Put / MPI_Win_create ...
 - Efficient way of performing data movements without explicit "handshake"
 - More info in a forthcoming PRACE MOOC on FutureLearn
- MPI has lots of further possibilities
 - Keep checking the <u>standard documentation</u>
 - Keep practising and be challenged!



