Be aware that your homework should be your own work. It is a matter of intellectual honesty to write your homework strictly by yourself. Using solutions from any other source is not allowed.

1. Use mathematical induction to show that when n is an exact power of 2,  $T(n) = n \lg n$  is the solution of the recurrence relation

$$T(n) = \begin{cases} 2 & \text{if } n = 2\\ 2T(\frac{n}{2}) + n & \text{if } n = 2^k \text{ for } k > 1. \end{cases}$$

- 2. Suppose we are comparing implementations of Insert-sort and Merge-sort on the same machine. For input of size  $n=2^k$  for  $k \geq 1$ , Insert-sort runs in  $8n^2$  steps, while Merge-sort runs in  $64n \lg n$  steps. For which value of n does Insert-sort beat Merge-sort?
- 3. We can express INSERT-SORT as a recursive procedure as follows. In order to sort A[1...n], we recursively sort A[1...n-1] and then insert A[n] into sorted array A[1...n-1].
  - (a) Write the pseudocode for this recursive version of Insert-sort, name it Insert-sort-recur.
  - (b) Write a recurrence for the running time of of Insert-sort-recur.
  - (c) Find the solution of this recurrence relation.
  - (d) Is Insert-sort-recur more expensive than Insert-sort?
- 4. In this exercise, we consider a Selection-sort algorithm. To sort n numbers stored in array a, we first find the smallest element of S and exchanging it with the element in a[1]. Then find the second smallest element of a, and exchange it with a[2]. Continue in this manner for the first n-1 element of a.
  - (a) Write a pseudocode for the Selection-sort algorithm.
  - (b) Analyze the running times.
- 5. Given an array  $s = \langle s[1], s[2], \dots, s[n] \rangle$ , and  $n = 2^d$  for some  $d \ge 1$ . We want to find the minimum and maximum values in s. We do this by comparing elements of s.
  - (a) The "obvious" algorithm makes 2n-2 comparisons. Explain.
  - (b) Can we do it better? Carefully specify a more efficient divide-and-conquer algorithm.
  - (c) Let T(n) = the number of comparisons your algorithm makes. Write a recurrence relation for T(n).
  - (d) Show that your recurrence relation has as its solution T(n) = 3n/2 2.