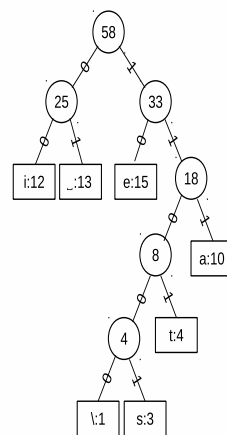


1. (a) The (global) optimal solution to the problem *contains within it* an (local) optimal solutions to subproblems. For example, LCS.
Note: this is one of the hallmarks of the applicability of dynamic programming.
- (b) $O(n \cdot 2^m)$
 $O(m \cdot n)$
- (c) $(A_1 A_2)((A_3 A_4)((A_5(A_6 A_7)))$
2. (a) **False.** It's only true for a connected graph. First, it's easy to see $|E| = O(|V|^2)$. If the graph is connected, $|E| \geq |V| - 1 = \Omega(|V|)$. Therefore, $\Theta(\lg E) = \Theta(\lg V)$.
- (b) **False.** In DFS, the vertices are processed in a LIFO order (stack).
- (c) **False.** It's true for a DAG (Directed Acyclic Graph).
3. Run DFS, if find a back edge, there is a cycle. The run time in this case is $O(|V|)$, not $O(|V| + |E|)!$ If ever see $|V|$ distinct edges, must have seen a back edge because in acyclic undirected forest, $|E| \leq |V| - 1$.

4.

char	frequency	code	# bits
a	10	111	3
e	15	10	2
i	12	00	2
s	3	11001	5
t	4	1101	4
□	13	01	2
\	1	11000	5

Show your work. When you're finished, fill in the column of the table labeled "code" with the binary codeword for each character.



5. (a) For example, for the change of $N = 55$ cents,

$$\begin{array}{l|l|l} \text{greedy:} & 50 + 5 \cdot 1 & 6 \text{ coins} \\ \text{optimal:} & 25 + 3 \cdot 10 & 4 \text{ coins} \end{array}$$

(b) $\text{CHANGE}(N) = \min\{\text{COINS}(N), 1 + \text{COINS}(N - 50), 2 + \text{COINS}(N - 100)\}$

(c) The optimal solution has to either have 0, 1 or 2 half-dollars in it. The solution finds the best of options. Running time $= 3 \cdot O(N) = O(N)$.