

How does server use session key after tls handshake?

Asked 2 years, 11 months ago Active 2 years, 11 months ago Viewed 1k times

3 After TLS handshake both server and client agreed about `session key` to use as symmetric key. So server must store several `session keys` which each one belongs to a specific client.

How server determines for a received request which `session key` must be used (i.e. `session key` of which client)? (Is this true: Client appends its `session id` in each request (somewhere like header?) and server has a map between `session id` and `session key`.)

How long a server uses a `session key`? Is it bad practice to use a session key for long time? Why they change `session key` while it transferred in a secure (i.e. encrypted with public key) manner?

4 Normally, once the TLS handshake is complete, the client and server can exchange encrypted traffic on the same connection without having to communicate the `session id`. This is because the server associates the SSL context information (including the `session key`) with the TCP socket on its end. You can see this in this simplified example of a TCP server:

✓ https://wiki.openssl.org/index.php/Simple_TLS_Server

However, there are two scenarios where the session id is communicated after the initial handshake - session `resumption` and session `renegotiation`. **예외!**

Session resumption allows the same TLS session (containing all the parameters agreed in the handshake) on a *new* TCP connection after closing the original one. This means the connection could time out, and your browser would detect this case and resume the TLS session on a new connection by sending the session id in a new Client Hello. More details on session resumption here:

<http://vincent.bernat.im/en/blog/2011-ssl-session-reuse-rfc5077.html>

Session renegotiation is used to change the parameters to be different from the initial handshake, but on the *same* TCP connection.

As you mention, the storage of a session cache on the server compromises security. An attacker could gain access to the session cache and then decrypt traffic. More details on this here - <https://blog.compass-security.com/2017/06/about-tls-perfect-forward-secrecy-and-session-resumption/>

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answered Sep 29 '18 at 21:28

 Ben Rowland
156 ● 1

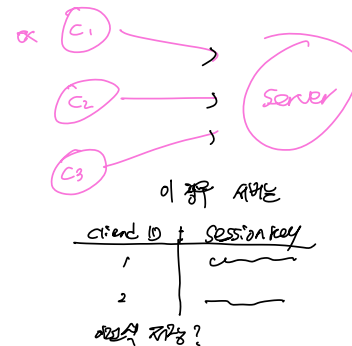
Thanks for your detailed answer. So for each request (post, get, put, delete,...) it must do a handshake? I previously assumed since handshake is a cpu intensive operation it should do a handshake and use its result for multiple requests... - Bonje Fir Sep 30 '18 at 5:30

1 @BonjeFir: you are mixing several things together. TLS is a protection for the TCP connection. get, put ... are actions at the application protocol level. Depending on the application protocol it might be possible to have several actions within a single TCP connection. Specifically with HTTP/1 (when using keep-alive) or with HTTP/2 a TCP connection will be created, upgraded to TLS and then several HTTP requests and responses will be done using the same TCP connection (and the same TLS session) - i.e. one TLS handshake and multiple HTTP requests. - Steffen Ullrich Sep 30 '18 at 6:02

@SteffenUllrich Thanks, very helpful clarification. - Bonje Fir Sep 30 '18 at 6:28

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Q TLS handshake 이후 PMS가 교환된 대칭키 (session key)는 어디에 저장되나요?



Ans)

TLS handshake 이후 서버는 SSL context information과 TCP socket은 연결시키기 때문에, 쿼리스트림은 TCP 연결을 재사용하고 있다. (keep-alive)

그래서 session key은 재사용 필요없다.

HTTP 1. / HTTP 2

↓
connection이 유지되면, 여러 http request
아래와 같은 tcp connection을 사용한다.