BU CS320 Assignment 6: Context Free Grammars

November 6, 2023

1. Given the following grammar where $\langle expr \rangle$ is the starting symbol:

Demonstrate the grammar above is ambiguous.

Suppose we want to derive the statement let z = 3 in z; begin () end

Left-associative:

<expr> ; <expr> let <id> = <expr> in <expr> ; <expr> let z = <expr> in <expr> ; <expr> let z = <dig> in <expr> ; <expr> let z = 3 in <expr> ; <expr> let z = 3 in <id> ; <expr> let z = 3 in z ; <expr> let z = 3 in z ; <expr> let z = 3 in z ; begin <expr> end let z = 3 in z ; begin () end

Right-associative:

```
<expr> ; <expr>
<expr> ; begin <expr> end
<expr> ; begin () end
let <id> = <expr> in <expr> ; begin () end
let <id> = <expr> in <id> ; begin () end
let <id> = <expr> in z ; begin () end
let <id> = <dig> in z ; begin () end
let <id> = 3 in z ; begin () end
let z = 3 in z ; begin () end
```

2. Modify the grammar (reproduced below) to be unambiguous. Hint: There is not just one way.

```
<id> ::= a | b | c | ... | z

<dig> ::= 0 | 1 | 2 | ... | 9

<exprs> ::= () | <dig> | <id> | let <id> = <expr> in <expr> | begin <expr> end

<expr> ::= <exprs> | <exprs> ; <expr>
```

3. Demonstrate your modified grammar fixes the previously shown ambiguity.

Left-associative:

<exprs>; <expr> let <id> = <expr> in <expr>; <expr> let z = <expr> in <expr>; <expr> let z = <exprs> in <expr>; <expr> let z = <dig> in <expr>; <expr> let z = 3 in z; begin <expr> end let z = 3 in z; begin () end

Right-associative:

```
<exprs>; <expr>
<exprs>; <exprs>
<exprs>; begin <expr> end
<exprs>; begin <exprs> end
<exprs>; begin () end
let <id> = <expr> in <expr>; begin () end
let <id> = <expr> in <expr>; begin () end
let <id> = <expr> in <id> ; begin () end
let <id> = <expr> in z; begin () end
let <id> = <expr> in z; begin () end
let <id = <expr> in z; begin () end
let <id = <dig in z; begin () end
let <id = 3 in z; begin () end
let z = 3 in z; begin () end</pre>
```