

String Searching 1 of 2

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String Searching

 Find the string "vtewfvtxqwfczsrdzcaj" in the following text:

qwerxcvvtewfzxcfasfedrsadfsdacfasdrtvtewqwertcsvte
wfvtxqwfczsrdzfeceeaeszxcvtsafsersdxzcvtedfaevsadv
tewfvtxqwfczsvzxgvtasfvtcasrfvtewqtrwtravtewfxtrac
wrtrdtgfdvxvvsbdgfstqtretydfxvzccadawqeewtertgfvbd
vczfafsvtewfvtxqwfczsgfsdfdxvzvzvtvsdgfsgtfwt6fqwt
qwrcfxtvtewfwtqwfzvwqgtfvtqfwcxetwfazreqresdqxrdqc
fwqdxvgfewcvtwefxvtrfczrqesxqecaqrfzvtqwxvbwyegcbe
bcwtfexvtfwxcrqxeqdcqzrwdfvtwxefvctyvtewfwefxqtfxc
qcdzrqxesrzqxrqcwqtfxtewfcvwerygcvewytxvqewtcxzdcd
qwfxvtewfvtxqwfczsrdzcajwfcsxtqwefdvetwqfvxdtqfwvq

Algorithms for string search

• string: S[0 .. m-1] ananaba

• text: T[0 .. n-1] bannabanabananaban

• Simple:

```
\label{eq:formula} \begin{split} & \text{for } i \leftarrow 0 \quad \text{to } n\text{-}m\text{-}1 \\ & \quad \text{found} \leftarrow \text{true} \\ & \quad \text{for } k \leftarrow 0 \text{ to } m\text{-}1 \\ & \quad \text{if } S[k] \text{!= } T[i\text{+}k] \text{ then } \text{found} \leftarrow \text{false, } \text{continue} \\ & \quad \text{if found } \text{then } \text{return } i \end{split}
```

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String search Simple search abcdmndsjhhhsjgrjgslagfiigirnvk fir- Slide the window by 1 • t = t +1; ananfdfjoijtoiinkjjkjgfjgkjkkhgklhg KMP - Why look at characters in the text multiple times? - Slide the window faster • t = t + s • but sometimes you can not skip s, need go back a little • a table to tell how to back step • t = t + s - M[s]**Knuth Morris Pratt** string: S[0 .. m-1] ananaba • text: T[0 .. n-1] anbananananaba ananaba When there is a mismatch, ⇒ move the string along to the earliest place it could possibly match and keep stepping Need a table to say how far to match: Is there a matching prefix of the match so far. Match so far: 0 1 2 3 4 5 ve string along: 1 1 2 3-1 4-2 5-3 Move string along: 1 1 Next match from: 0 0 0 String search Simple search abcdmndsjhhhsjgrjgslagfligirnvk firabcdefg - Slide the window by 1 • t = t +1; ananfdfjoijtoiinkjjkjgfjgkjkkhgklhg • KMP - Slide the window faster • t = t + s - M[s]- Never recheck the matched characters • If there a "suffix ==prefix"? - No, skip these characters » M[s] = 0 - Yes, reuse, no need to recheck these characters » M[s] is the length of the "reusable" suffix

Knuth Morris Pratt $\textbf{input} : \mathsf{string} \ \mathsf{S[0} \ .. \ \mathsf{m-1]} \ , \quad \mathsf{text} \ \ \mathsf{T[0} \ .. \ \mathsf{n-1]}$ output: the position in T at which S is found, or -1 if not present position of current character in S variables: s ← 0 $t \leftarrow 0$ start of current nM[0 .. m-1] self match table start of current match in T Construct self match table M while t+s < n if S[s] = T[t+s] then // match $s \leftarrow s + 1$ if s = m then return t // found S else if M[s] = -1 then // mismatch, no self overlap $s \leftarrow 0, \quad t \leftarrow t + s + 1,$ // mismatch, with self overlap $t \leftarrow t + s - M[s]$ // match position jumps forward $s \leftarrow M[s]$ return -1 // failed to find S

KMP how far to move along?

string: ananabatext: ...ananx???....

- If mismatch at string position s (and text position t+s)
 - find largest substring ending at s-1 that matches a prefix of string
 - move t to (t + s length of substring)
 - keep matching from $s \leftarrow$ length of substring
- special case:
 - if s = 0, then move t to t + 1 and match from $s \leftarrow 0$

KMP: Building the table.

```
input: S[0 .. m-1] // the string
output: M[0 .. m-1] // match table
initialise: M[0] ← -1
           M[0] \leftarrow 0

j \leftarrow 0
                                                         ananaba
                             // position in prefix
                                                        ananaba
           pos ← 2
                            // position in table
while pos < m
    if S[pos - 1] = S[j]
                               //substrings ...pos-1 and 0..j match
         M[pos] \leftarrow j+1,
         pos++, j++
    else if j > 0
                               // mismatch, restart the prefix
        j \leftarrow M[j]
    else //j = 0
                               // we have run out of candidate prefixes
         M[pos] \leftarrow 0,
         pos++
```

Summary searches forward, never matches a text character twice (and never skips a text character) jumps string forward based on self match within the string: prefix of string matching a later substring. doesn't use the character in the text to determine the jump. Cost?