# Algorithms

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### Spectral clustering

**Definition 1.** Ncut(A, B) is defined as

$$Ncut(A,B) = cut(A,B)(\frac{1}{d(A)} + \frac{1}{d(B)})$$

where  $d(A) = \sum_{i \in A} d_i$ .

Remark 1. • The intuition is that when A is relatively small, the  $\frac{1}{d(A)}$  will be large, hence discouraging the isolating small groups.

 $\bullet$  Finding minimum Ncut is equivalent to finding vector v that minimizes

$$\frac{v^T L v}{v^t D v} \text{such that } v^t D 1 = 0, v_i \in \{a, b\}$$

where L = D - W.

### When the data is separable linearly

There are several ways. First, linear programming.

#### Algorithm 1 Linear Programming for Classifier

**Objective:** minimize  $\mathbf{u} = (0, \dots, 0)$  (dummy variable)

Subject to: 
$$A\mathbf{w} \geqslant \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

 $A_{i,j} = y_i x_{i,j}$  (where j'th element of the vector  $x_i$ )

Remark 2.  $\bullet$  u is a dummy variable here; we essentially only check if the constraint is satisfied.

- This is only applied for when the data is separable.
- We can formula the regression problem with loss function l(h,(x,y)) = |h(x) = y| using linear programming.

#### Algorithm 2 Batch Perceptron

```
1: function BATCHPERCEPTRON(x_1, y_1, ..., x_m, y_m)
        w(1) \leftarrow (0, \ldots, 0)
 2:
 3:
        for t \leftarrow 1, 2, \dots do
            if there exists i such that y_i\langle w(t), x_i\rangle \leq 0 then
 4:
                 w(t+1) \leftarrow w(t) + y_i x_i
 5:
 6:
            else
                 output w(t)
 7:
 8:
                 break
            end if
 9:
        end for
10:
11: end function
```

Remark 3. • Note that  $y_i(w^{(t+1)}, x_i) = y_i(w^{(t)}, x_i) + ||x_i||^2$ 

- The algorithm must stop after at most  $(RB)^2$  iterations, where  $R = \max_i ||x_i||$  represents a data spread, and  $B = \min\{||w|| : i \in [m], y_i \langle w, x_i \rangle \ge 1\}$  represents margin.
- To prove this, it suffices to show that  $1 \geqslant \frac{\langle w^*, w^{(T+1)} \rangle}{\|w^*\| \|w^{(T+1)}\|} \geqslant \frac{\sqrt{T}}{RB}$ , which we proceed by bounding numerator and denominator separately.
- We can prove that this bound is tight. For some vector  $w^* \in \mathbb{R}^d$ , the algorithm incurs  $m = (BR)^2$  errors (considering m = d).
- Moreover, for d = 3, an algorithm can be designed to commit exactly (m) errors for any given  $m \in \mathbb{N}$ , serving as an upper bound concurrently

#### Logistic Regression

**Definition 2.** Fit the logistic function  $\phi_{sig}(x) = \frac{1}{1 + exp(-\langle w, x \rangle)}$  with minimization scheme  $w = argmin_{w \in \mathbb{R}^d} \frac{1}{m} \sum_{i=1}^m \log(1 + exp(-y_i \langle w, x_i \rangle))$ .

Remark 4. • The explaratory variable is between 0 and 1, making it interpretable as a probability.

- Appropriate for binary classification.
- Logistic loss function is a convex function so it's efficient to minimize.