

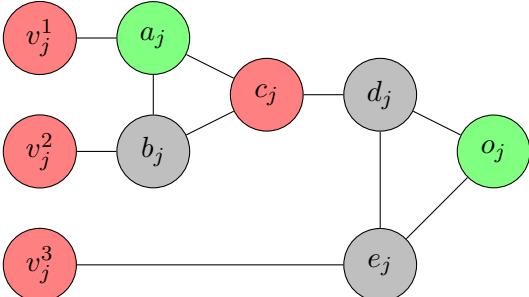
Complexity - Exercise Sheet 3

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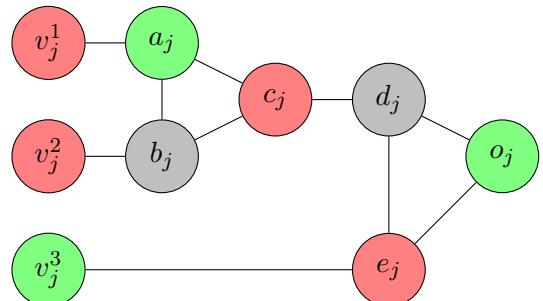
Exercise 3.2. A k -coloring of a graph $G = (V, E)$ is an assignment $c : V \rightarrow \{1, 2, \dots, k\}$ such that if $\{u, v\} \in E$, then $c(u) \neq c(v)$. Let $k\text{-COL} = \{\langle G \rangle \mid G \text{ has a } k\text{-coloring}\}$. Prove that $3\text{SAT} \leq_P 3\text{-COL}$.

Solution. Let us construct a graph $G = \langle V, E \rangle$ from a 3-CNF formula ϕ as follows.

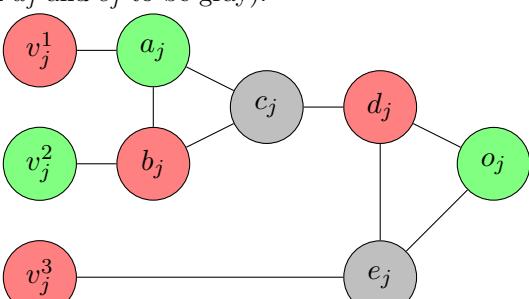
1. Introduce two vertices v_f and v_u in V and connect them by an edge. If a vertex has the same color as v_f , then its corresponding literal or variable is assigned to **False**. If it has the same color as v_u , then it is assigned to **Undefined**.
2. For each clause $C_j = (\ell_j^1 \vee \ell_j^2 \vee \ell_j^3)$ in ϕ , introduce three vertices v_j^1, v_j^2, v_j^3 in V and connect them to v_u . This forces every literal to be assigned to either **True** or **False**.
3. We have to make sure that the assignments of the literals in each clause are consistent. For each variable x in ϕ , introduce two vertices v_x and $v_{\neg x}$ in V and connect them to each other and to v_u . If a vertex v_j^k corresponds to the literal x (resp. $\neg x$), connect v_j^k to $v_{\neg x}$ (resp. v_x).
4. Finally, we have to make sure that for each triple v_j^1, v_j^2, v_j^3 , at least one vertex has the different color from v_f . Let say the color of v_f is red, v_u is gray and the other color is green. We will design a subgraph for each clause. Let o_j be a vertex in this subgraph that is connected to v_f and v_u i.e. it is always green. We want that there is a valid coloring for the subgraph if and only if at least one vertex among v_j^1, v_j^2, v_j^3 is green, which means that at least one literal in the clause is assigned to **True**. The possible cases are shown in Figures (a)-(h).



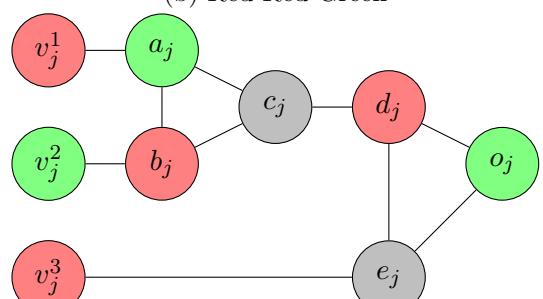
(a) Red-Red-Red: **invalid** (other colorings also force both d_j and e_j to be gray).



(b) Red-Red-Green

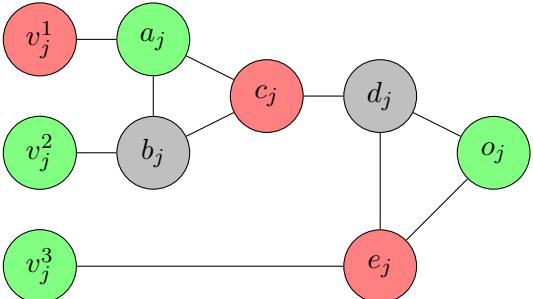


(c) Red-Green-Red

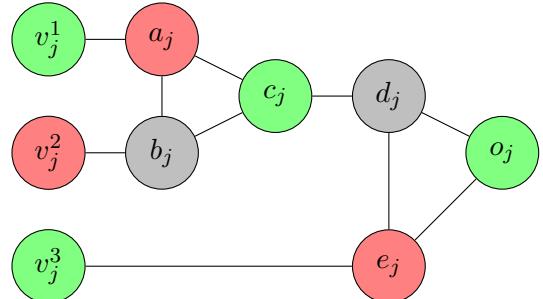


(d) Green-Red-Red

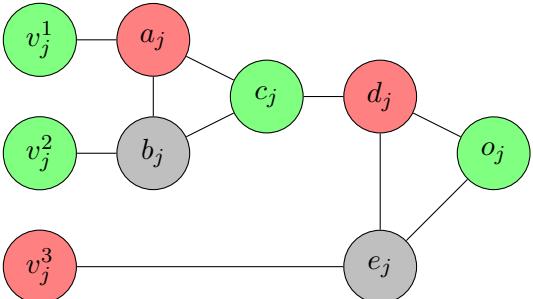
It is clear that if G has a valid 3-coloring, then we can construct a satisfying assignment for ϕ following the colors of v_x 's. Conversely, if ϕ has a satisfying assignment the we color v_f by red, v_u by gray. For each variable x , if x is assigned to **True**, color v_x by green and $v_{\neg x}$ by red; otherwise, color v_x by red and $v_{\neg x}$ by green. Use this information to color v_j^k 's. Since each clause has at least one true literal, we follow Figures (b)-(h) to color



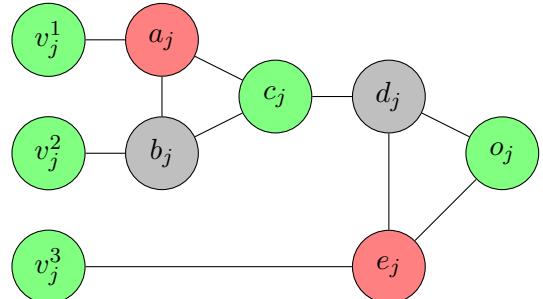
(e) Red-Green-Green



(f) Red-Green-Green



(g) Green-Green-Red



(h) Green-Green-Green

the subgraph of each clause. This gives a valid 3-coloring for G . Therefore, ϕ is satisfiable if and only if G has a valid 3-coloring.

The transformation from ϕ to G can be done in polynomial time of the length of ϕ . Indeed, the creation of v_f and v_u also takes constant time. The number of variables is at most the length of ϕ and the number of clauses is at most one third of the length of ϕ . Each clause introduces a constant number of vertices and edges, and so does each variable. The edges are created at most the square of the number of vertices, which is still polynomial in the length of ϕ .

Thus, $3\text{SAT} \leq_P 3\text{-COL}$.

Exercise 3.3. NAE-3SAT is the problem of deciding for a given 3-CNF formula whether there is a truth assignment to the variables so that in every clause, there is at least one true literal and at least one false literal. Show that NAE-3SAT is NP-hard.

Solution. We will show that $3\text{SAT} \leq_P \text{NAE-3SAT}$. Let ϕ be a 3-CNF formula. Introduce a new variable p that will be used when dealing with all clauses. For each clause $C_j = (\ell_j^1 \vee \ell_j^2 \vee \ell_j^3)$ in ϕ , we will create two clauses $(\ell_j^1 \vee \ell_j^2 \vee w_j)$ and $(\neg w_j \vee \ell_j^3 \vee p)$. The conjunction of all these clauses makes a new formula ϕ' .

We show that if ϕ is satisfiable, then so is ϕ' . In particular, if $\ell_j^1 \vee \ell_j^2 \vee \ell_j^3$ is satisfiable, then $(\ell_j^1 \vee \ell_j^2 \vee w_j) \wedge (\neg w_j \vee \ell_j^3 \vee p)$ is NAE-satisfiable. We set p to false and consider the following cases.

1. If ℓ_j^1 (resp. ℓ_j^2) is true, then we can assign w_j to false. This makes the first clause NAE-satisfied because ℓ_j^1 (resp. ℓ_j^2) is true and w_j is false. The second clause is also NAE-satisfied because $\neg w_j$ is true and p is false.
2. If ℓ_j^3 is true, then we can assign w_j to $\neg(\ell_j^1 \wedge \ell_j^2)$. This makes the first clause NAE-satisfied because we have excluded the cases that ℓ_j^1, ℓ_j^2 and w_j are all true and ℓ_j^1, ℓ_j^2 and w_j are all false. The second clause is also NAE-satisfied because ℓ_j^3 is true and p is false.

Conversely, if ϕ' is satisfiable, then so is ϕ . Let \mathcal{M} be a NAE-satisfying assignment for the variables in ϕ' . Note that since the literals in each clause of ϕ' are not all the same cannot be all true, the assignment $\overline{\mathcal{M}}$ such that for each variable in ϕ' , $\overline{\mathcal{M}}(x) = \neg \mathcal{M}(x)$ is also a NAE-satisfying assignment. Therefore, we can choose \mathcal{M} such that p is false. We will show that the restriction of \mathcal{M} to the variables in ϕ is a satisfying assignment for ϕ . In particular, if $(\ell_j^1 \vee \ell_j^2 \vee w_j)$ and $(\neg w_j \vee \ell_j^3 \vee p)$ is NAE-satisfiable, then $\ell_j^1 \vee \ell_j^2 \vee \ell_j^3$ is satisfiable. Suppose that this is not the case i.e. $\ell_j^1, \ell_j^2, \ell_j^3$ are all false. Then w_j is true to make the first clause NAE-satisfied. Plugging this into the second clause, we have $(\neg w_j \vee \ell_j^3 \vee p) = (\text{false} \vee \text{false} \vee \text{false})$, which is not NAE-satisfiable. This is a contradiction.

Therefore, ϕ is satisfiable if and only if ϕ' is NAE-satisfiable.

The transformation from ϕ to ϕ' can be done in polynomial time of the length of ϕ . Indeed, the number of clauses in ϕ is at most one third of the length of ϕ . Each clause introduces a constant number of clauses and variables.

Thus, $\text{3SAT} \leq_P \text{NAE-3SAT}$ and so NAE-3SAT is NP-hard.