

# 1 Pseudocode

## 1.1 The main program

```
Initialize the queues (job queue, memory management queue and Level 0-2 queues);
Fill Job Dispatch queue from job dispatch list file;
Ask the user to enter integer values for parameters;
Initialize the root memory block of size 2048(main memory to use);
While there is a currently running process or either queue is not empty
    Unload any arrived pending processes from the Job Dispatch queue;
    Dequeue process from Job Dispatch queue and Enqueue on arrived job queue;
    If there's a task on the arrived job queue
        dequeue the task;
        try to allocate the task on the main memory to a block fitting its size;
        if a memory block is assigned to the task
            store the pointer of the block in a PCB field;
            enqueue the process in level 0 queue;
        else there's no block available
            queue the task back in the front of the arrived job queue again;
    If a process is currently running
        Decrease process remaining_cpu_time by quantum;
        If times up
            Send SIGINT to the process to terminate it;
            Calculate the turnaround time;
            Update the overall turnaround time;
            Free the memory block the process was allocated to;
            Free up process structure memory;
        Else
            Send SIGTSTP to suspend currently running process;
            If the process is in Level 0 queue
                Enqueue it to the tail of Level 1 queue;
            Else if it is in Level 1 queue
                Increase the iteration counter by 1;
                If the counter is over the iteration limit
                    Enqueue it to the tail of Level 2 queue;
                Else
                    Enqueue it back to the tail of Level 1 queue again;
            Else if it is in Level 2 queue
                Enqueue it back to the head of Level 2 queue;
        Mark the process has been run;
    If no process is currently running
        Dequeue a process from the head of the first available non empty queue;
        Mark the corresponding level the process is at;
        If the process job is a suspended process
            Send SIGCONT to resume it;
        Else
            Start it (fork & exec);
Set the process as the currently running process;
```

```

If there's a process to run
    Update the quantum by the smaller one between the cpu remaining time
    and the corresponding time quantum per the queue the process is at;
Else
    Set the quantum as the minimum value among t_0, t_1 and 1;
Sleep for quantum;
Increase timer by quantum;
Go back to the while loop;
Calculate and show the average turnaround time and the average waiting time;
Terminate the Round Robin dispatcher;

```

## 1.2 MAB subroutines

- allocate a process to a memory block of an appropriate size

```

If the size of the process is less than 8 (MB)
    Pad the size to be 8;
Call the block splitting oracle routine
to find a suitable block on the main memory;
If there's a free block
    Mark the block as being in use;
    Return the pointer of the block;
Else
    Return a null pointer

```

- block split oracle

```

//Base Case:
If the block is in use
    Return a null pointer;
If the block can't be split evenly into 2 blocks to store the process
    If a part of the block is not free
        Return a null pointer;
    Else
        Return the pointer of the block;
//Recursive Step:
If the block has been split already
    Attempt to search down from its left part;
    If there's no space under its left child block
        Attempt to search down from its right part;
Else
    Create two children blocks evenly;
    Attempt to use this oracle to split its left child further;

```

- free the memory block a process occupies

```

Mark the block as being free;
Call the merge oracle routine to handle the merging protocol.

```

- merge oracle

```

//Base Case:
If the block is the root
    Do nothing and return;
//Recursive Step:
If the sibling block is free
(implicitly there's no children for the sibling block)
    Delete the block and its sibling block;
    Merge 2 blocks into 1;
    Use the oracle itself to attempt to merge the combined block;

```

## 2 Testing Boundary Value Analysis

We use the notation  $D_i$  for the memory size requirement for a task  $i$ . Chronological testing has been conducted previously in Stage 1, therefore, processes' time sequences are kept the same. A third field standing for size is appended. Here focused are memory allocation and testing the strictly FCFS blocking mode of the arrived job queue.

### 2.1 Buddy memory allocation API unit testing

Here we use MAB\_MAX (the overall size of the main memory) as 2048 (MB), MAB\_MIN (lower bound size of a block) as 8 (MB).

#### 2.1.1 Alloc

1. allocate a block of size  $7 > D_i \geq 1; \forall j \in [3, \dots, 10] : 2^{j+1} > D_i \geq 2^j; D_i = 2048$  to an empty main memory
2. allocate  $0 < D_i \leq 2048$  when there's a block of 1025 MB in use.
3. allocate  $D_i = 1024; D_i < 1024; D_i > 1024$  when there's a block of  $\leq 1024$  MB in use.
4. allocate 16 blocks of 128 MB to test multiple even blocks
5. allocate 16 blocks of 129 MB to test slightly screwed blocks.

#### 2.1.2 Free

1. free the main memory in use/free
2. free any block not in use(only testable calling a dummy Mab object directly).
3. free the only block in use(free the pointers returned upon success allocation in section 2.1.1)

### 2.1.3 Alloc & Free combination to test Spilt and Merge in a normal workflow

Please check [the online instance](#). You can find a testing C file('test\_mab.c') and replicate the unit testing by 'gcc -o <unittest> mab.c pcb.c test\_mab.c' and run the executable '<unittest>'.

## 2.2 Arrived job queue testing

1. a process of 1025 MB to block the (chronologically) following processes before that process is terminated
2. 2 consecutive processes of 1024 MB to run concurrently.
3. multiple small processes that are not blocked.

Note that user inputs tuning matters to test. For example, 'test\_file\_2' I recommend to use a large number for the iteration cap and 1 for the level 0 quantum to monitor concurrent RRs for multiple processes with compatible size.