

# yacoub: a Python package for Simulating Generalized Fading Channels

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**Abstract**—We present a well tested Python-based library for simulating and computing generalized fading channels, named *yacoub*. We describe the applicability of *yacoub* using examples in recent communications systems challenges, namely: spectrum sensing, bit error rate computation, and parameter estimation in generalized fading channels. For the latter, we develop an iterative algorithm based on the Majorization-Minimization framework. The development of *yacoub* is open source and its code is available at <http://github.com/mirca/yacoub>.

## I. INTRODUCTION

### Note on notation

Scalars and random variables are denoted as *italic* small-case letters *e.g.*  $x$ ; vectors and random vectors are denoted as *italic*, boldface, small-case letters *e.g.*  $\mathbf{x}$ . The  $n$ -th component of a vector  $\mathbf{x}$  is denoted as  $x_n$ . A complex vector of length  $n$  is defined as  $\mathbf{x} \in \mathbb{C}^{n \times 1}$ . All vectors are column vectors. Matrices are denoted as *italic*, boldface, capital letters as in  $\mathbf{X}$ ; the identity matrix of order  $n$  is denoted as  $\mathbf{I}_n$ . We define a discrete-time circularly symmetric Gaussian process  $\mathbf{z}$  as any collection of random variables  $\mathbf{z} = \mathbf{x} + j\mathbf{y}$ ,  $j \triangleq \sqrt{-1}$ , such that  $\mathbf{x}$  and  $\mathbf{y}$  are iid jointly Gaussian with zero mean vector and covariance matrix given by  $\mathbb{E}[\mathbf{z}\mathbf{z}^\dagger]$ , in which  $\mathbf{z}^\dagger$  means the conjugate transpose of  $\mathbf{z}$ . The expectation value wrt a random variable  $x$  is denoted as  $\mathbb{E}_x$ . The probability of an event  $A$  is denoted as  $\mathbb{P}(A)$ . The indicator function is denoted as  $\mathbb{I}(\cdot)$  and it evaluates to one if its argument is true and zero otherwise.

## II. THE ACCEPTANCE-REJECTION SAMPLER IN LOG-SPACE

## III. MAJORIZATION-MINIMIZATION ALGORITHMS

## IV. EXAMPLES

### A. Spectrum Sensing in Complex Generalized Fading Channels

The spectrum sensing problem consists in deciding whether or not a given channel frequency band is being occupied by a licensed (primary) user and, in case that such frequency band is available, how to opportunistically allocate secondary users such that the interference on the primary user is negligible.

From a probabilistic point of view, the spectrum sensing problem may be framed as a decision theory problem, as

follows

$$H_0 : \mathbf{y} = \mathbf{w}, \quad (1)$$

$$H_1 : \mathbf{y} = h\mathbf{s} + \mathbf{w}, \quad (2)$$

in which  $\mathbf{y} \in \mathbb{C}^{n \times 1}$  is the decoded received vector signal,  $\mathbf{w} \in \mathbb{C}^{n \times 1}$  is complex Gaussian noise process with zero mean vector and covariance matrix given as  $\sigma^2 \mathbf{I}_n$ , and  $h$  is the channel gain.

In [?], the authors have shown that the probability distribution of the energy statistic  $\tilde{y} \triangleq \mathbf{y}^\dagger \mathbf{y}$  conditioned on the knowledge of  $h$ , in case that  $\mathbf{s}$  is an  $M$ -PSK signal such that every symbol has the same probability of occurrence,  $\mathbb{P}(s_n = s) = \frac{1}{M}$ , is given as

$$p(\tilde{y}|h, H_1) = 1 - Q_n \left( \sqrt{\frac{2n|h|^2 E_s}{\sigma^2}}, \sqrt{\frac{2\tilde{y}}{\sigma^2}} \right), \quad (3)$$

in which  $Q_n$  is the Marcum- $Q$  function and  $E_s$  is the energy per symbol.

The pdf of  $\tilde{y}$  can be written using the Law of Total Expectation

$$p(\tilde{y}|H_1) = \mathbb{E}_h [p(\tilde{y}|h, H_1)] = \int_{-\infty}^{+\infty} p(\tilde{y}|h, H_1) p(h) dh. \quad (4)$$

Recall that the energy detection rule can be expressed as

$$d_\delta(\tilde{y}) = \mathbb{I}(\tilde{y} > \delta) \quad (5)$$

in which  $\delta$  is a strictly positive real number known as energy threshold, and  $d_\delta(\tilde{y}) = j$ ,  $j \in \{0, 1\}$ , means that the detector has decided in favor of the hypothesis  $H_j$ .

As a result, the probabilities of false alarm and miss detection can be written as

$$p_f \triangleq \mathbb{P}(d_\delta(\tilde{y}) = 1|H_0) = 1 - p(\delta|H_0), \quad (6)$$

$$p_d \triangleq \mathbb{P}(d_\delta(\tilde{y}) = 0|H_1) = \mathbb{E}_h [p(\delta, h|H_1)], \quad (7)$$

### B. Parameter Estimation in Nakagami- $m$ fading

$$p(\mathbf{h}) = \prod_{i=1}^n \frac{2m^m}{\Gamma(m)\Omega^m} h_i^{2m-1} \exp\left(-\frac{mh_i^2}{\Omega}\right) \quad (8)$$

$$= \left(\frac{2m^m}{\Gamma(m)\Omega^m}\right)^n \exp\left(-\frac{m \sum_{i=1}^n h_i^2}{\Omega}\right) \prod_{i=1}^n h_i^{2m-1} \quad (9)$$

$$\log p(\mathbf{h}) \cong n(m(\log m - \log \Omega) - \log \Gamma(m)) - m \left( \frac{\sum_{i=1}^n h_i^2}{\Omega} - 2 \sum_{i=1}^n \log h_i \right) \quad (10)$$

A direct maximum likelihood estimator for (10) has been investigated to be infeasible [?].

Therefore, we use a Majorization-Minimization algorithm to find smooth and easy to optimize upper bounds for  $\log p(\mathbf{h})$ .

*A parte problemática aqui é encontrar um limitante inferior para a função Gamma. Talvez serie de Taylor resolva. Isto é, precisamos encontrar inequacoes da forma*

$$m \log m \geq \quad (11)$$

$$- \log \Gamma(m) \geq \quad (12)$$

### C. BER in Complex $\alpha - \mu$ Fading

Consider the system

$$\mathbf{y} = h\mathbf{s} + \mathbf{w} \quad (13)$$

in which  $\mathbf{s} \in \mathbb{C}^{n \times 1}$  is a complex On-Off Keying (OOK) signal,  $h$  is a complex  $\alpha - \mu$  random variable and  $\mathbf{w}$  is a complex Gaussian process with zero mean vector and covariance matrix equals  $\sigma^2 \mathbf{I}_n$ , and  $\mathbf{y}$  is the received complex vector signal.

Assume that the OOK symbols are equiprobable and that there exist no interference between the in-phase and quadrature components, then the probability of one bit error is given as

$$p_e = \frac{1}{2} (\mathbb{P}(\hat{y}_i = 0 | s_i = 1) + \mathbb{P}(\hat{y}_i = 1 | s_i = 0)). \quad (14)$$

Assume that the decoded vector  $\hat{\mathbf{y}}$  is estimated using the minimum distance decoding rule, i.e.,

## V. CONCLUSIONS

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