



# Recursion in 20''

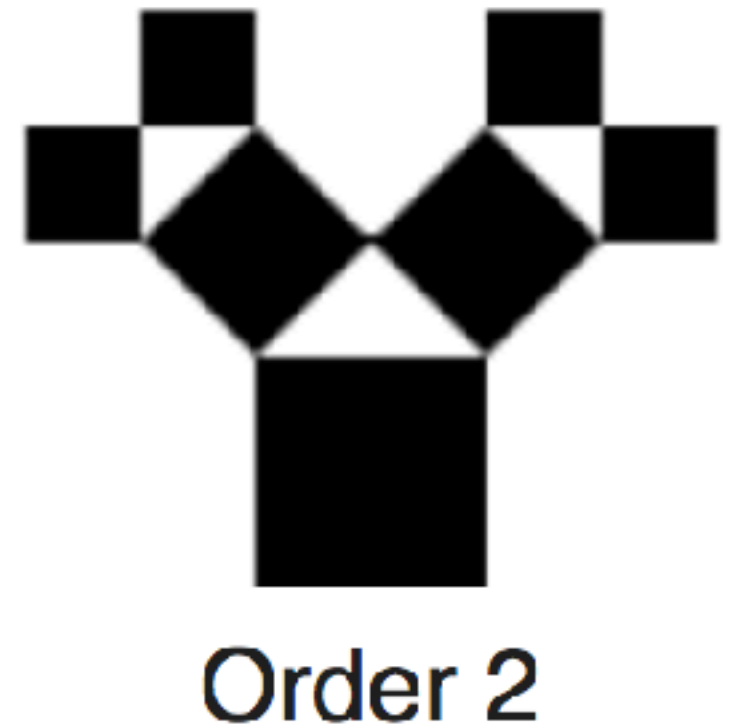
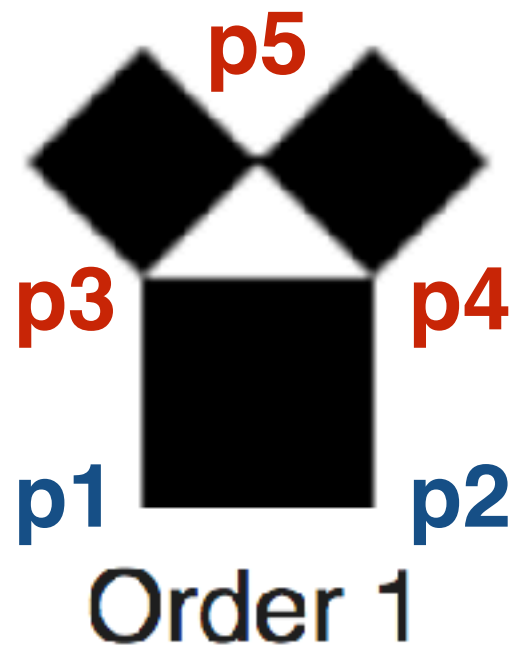
**Mircea F. Lungu**

University of Groningen  
Netherlands

@mircealungu

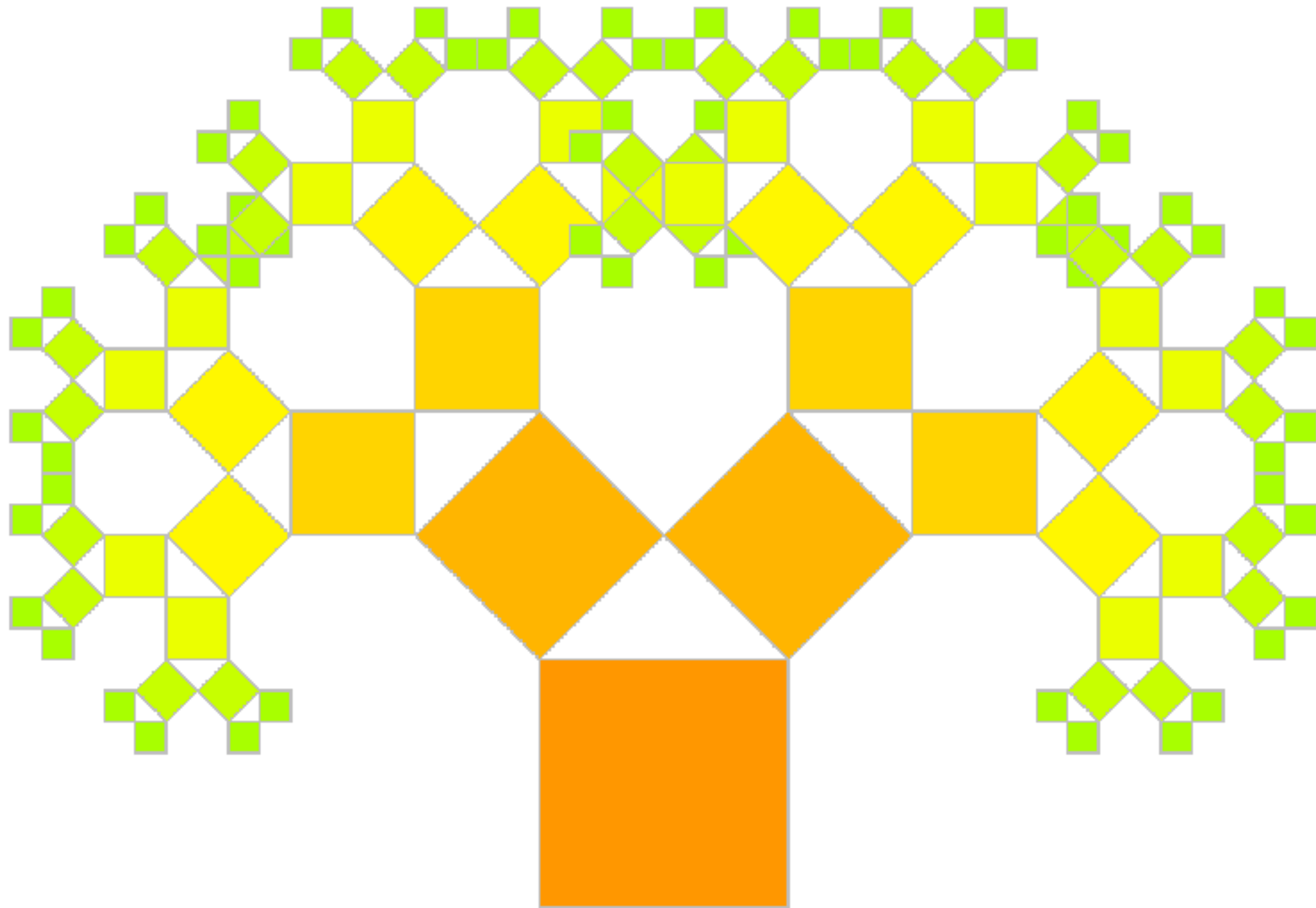
[https://github.com/mircealungu/open\\_lectures/recursion](https://github.com/mircealungu/open_lectures/recursion)

# Pythagoras' Tree



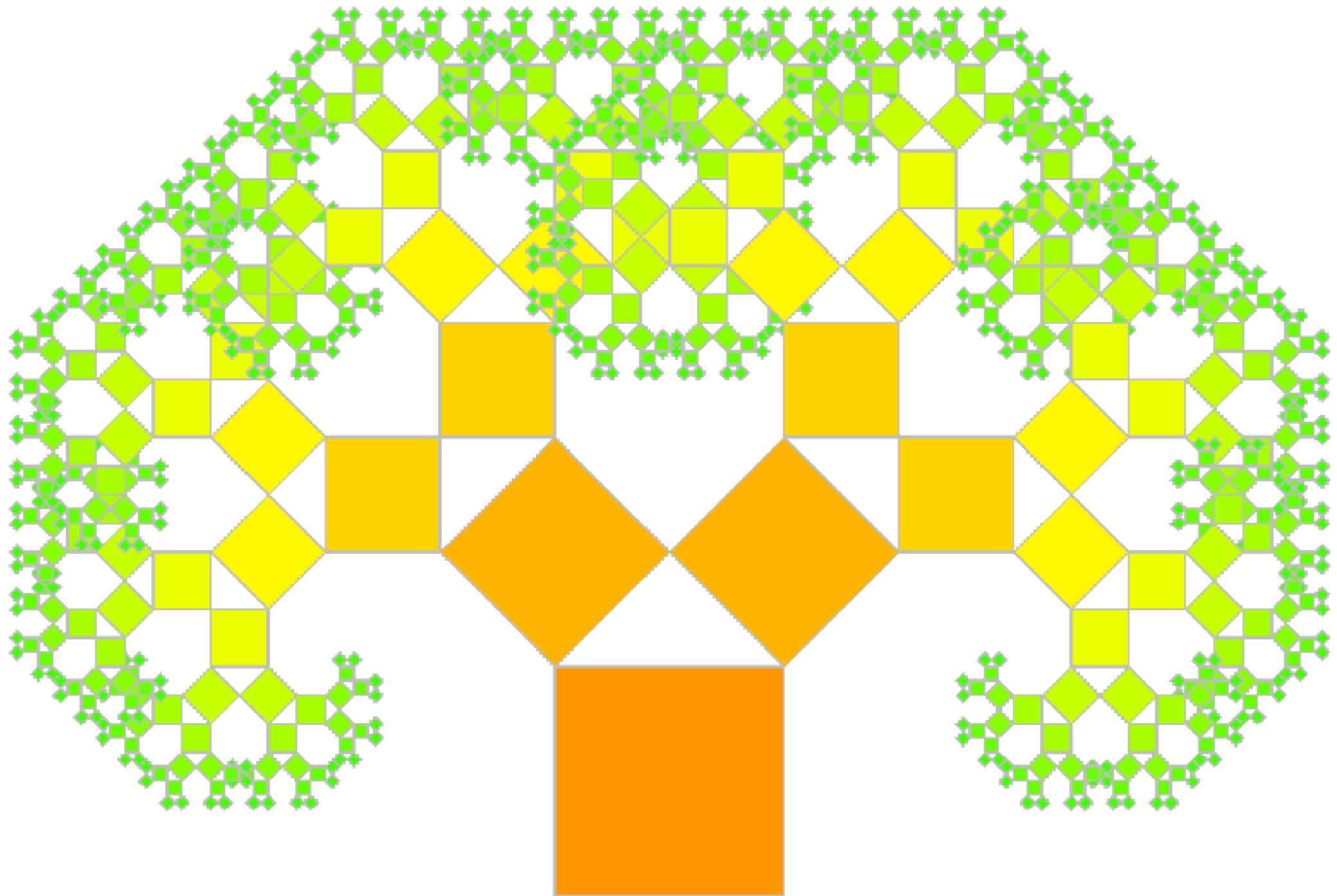
[Code on GitHub](#)

Order 7



[Code on GitHub](#)

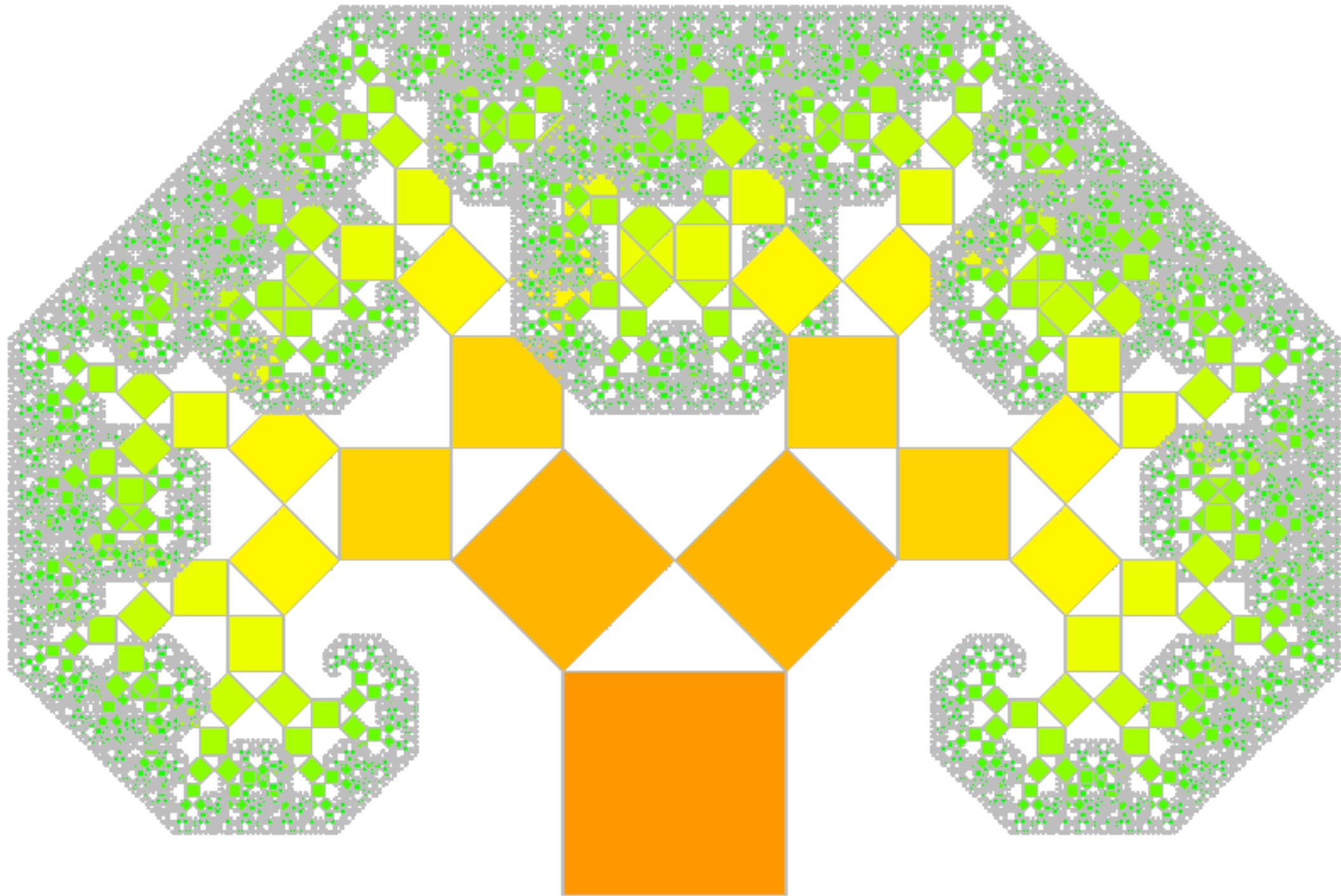
Order 10



[Code on GitHub](#)

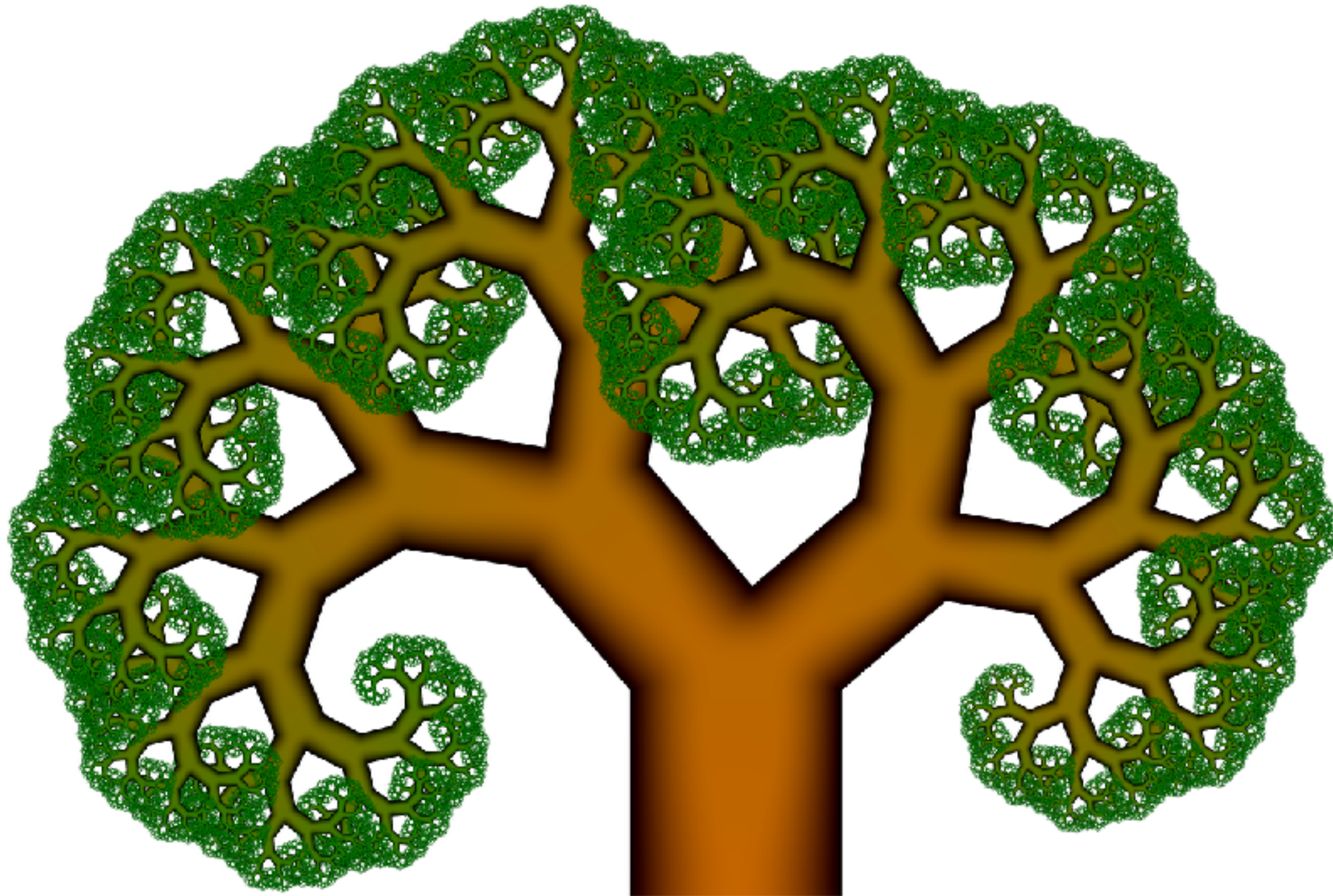


Order 15



[Code on GitHub](#)

Order ??



Warning: not on GitHub anymore!!!

By Atze van der Ploeg Atzecsse - Own work,  
Attribution, <https://commons.wikimedia.org/w/index.php?curid=28183362>



```
private void drawTree(Graphics2D g, Point2D p1, Point2D p2,  
                      int depth) {  
  
    if (depth == depthLimit)  
        return;  
  
    Point2D p3 = new Point2D.Double();  
    Point2D p4 = new Point2D.Double();  
    Point2D p5 = new Point2D.Double();  
  
    computeNewPoints(p1, p2, p3, p4, p5);  
  
    drawSquare(g, p1, p2, p3, p4, depth);  
  
    drawTree(g, p4, p5, depth + 1);  
    drawTree(g, p5, p3, depth + 1);  
}
```

[Code on GitHub](#)

```
private void drawTree(Graphics2D g, Point2D p1, Point2D p2,  
                      int depth) {
```

```
    if (depth == depthLimit)  
        return;
```

**base case**

when reached causes recursion to end

```
    Point2D p3 = new Point2D.Double();  
    Point2D p4 = new Point2D.Double();  
    Point2D p5 = new Point2D.Double();
```

```
    computeNewPoints(p1, p2, p3, p4, p5);
```

```
    drawSquare(g, p1, p2, p3, p4, depth);
```

```
    drawTree(g, p4, p5, depth + 1);  
    drawTree(g, p5, p3, depth + 1);
```

**recursive calls**

```
}
```

[Code on GitHub](#)



# The power of recursion lies in the possibility of defining

- **an infinite set of objects by a finite statement**
- **an infinite number of computations by a finite recursive program**

*Algorithms + Data Structures = Programs,*  
Wirth, Niklaus (1976)

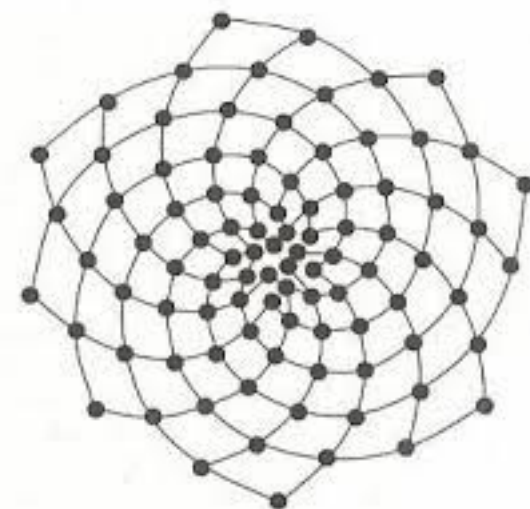
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$$F_{n-2} = F_n - F_{n-1}$$

$F_0$	$F_1$	$F_2$	$F_3$	$F_4$	$F_5$	$F_6$	$F_7$	$F_8$	$F_9$	$F_{10}$	$F_{11}$	$F_{12}$	$F_{13}$	$F_{14}$	$F_{15}$
0	1	1	2	3	5	8	13	21	34	55	89	144	233	377	610





# Fibonacci in Java

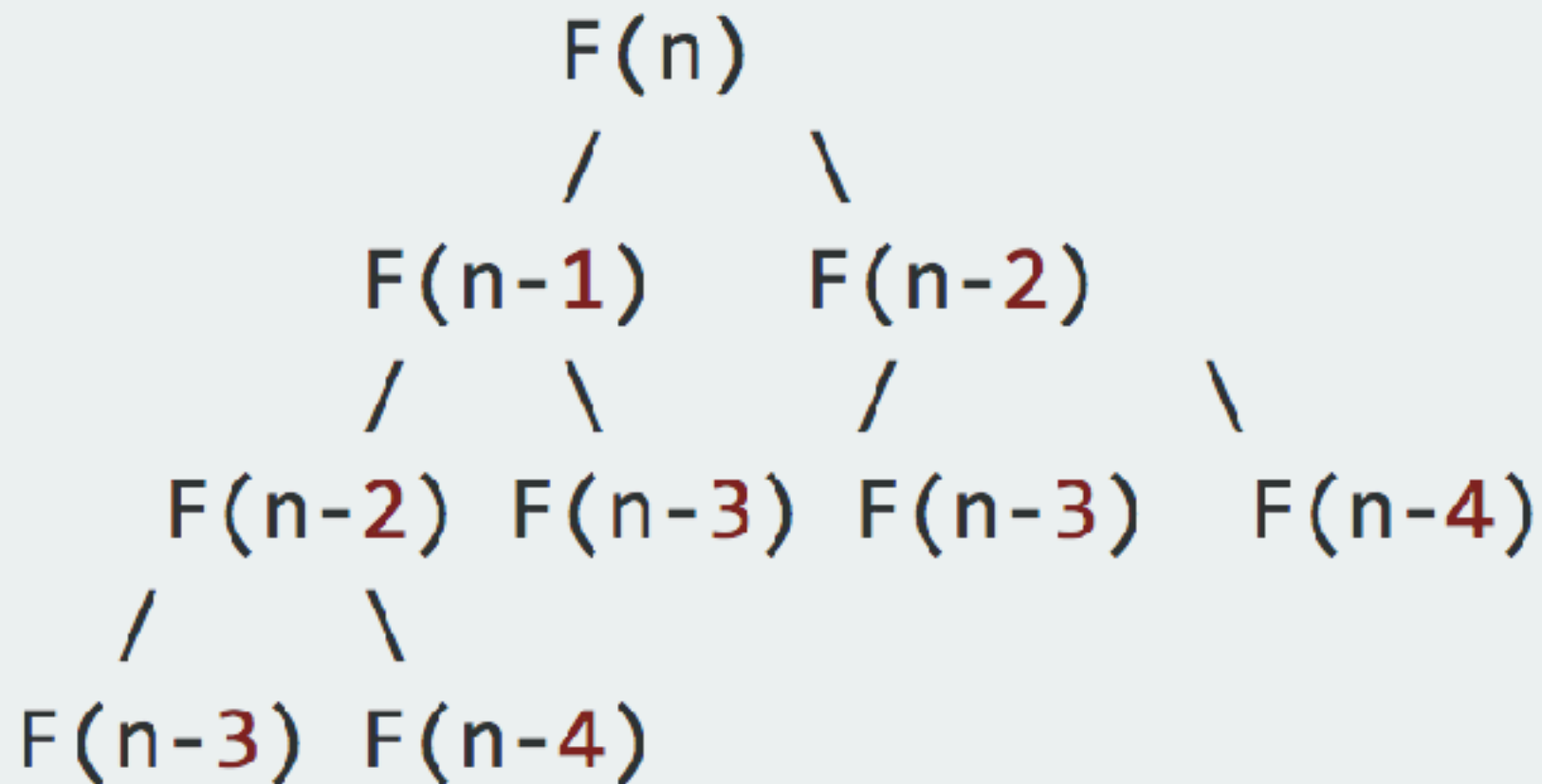
```
public static long fibonacci(final int n)
{
    return (n < 2) ? n : fibonacci(n - 1) + fibonacci(n - 2);
}
```

base case

recursive calls



# Exponential Growth



Dynamic Programming - method for solving a complex problem by breaking it down into a collection of simpler subproblems, solving each of those subproblems just once, and storing their solutions



Demo — scalability of naïve Fibonacci



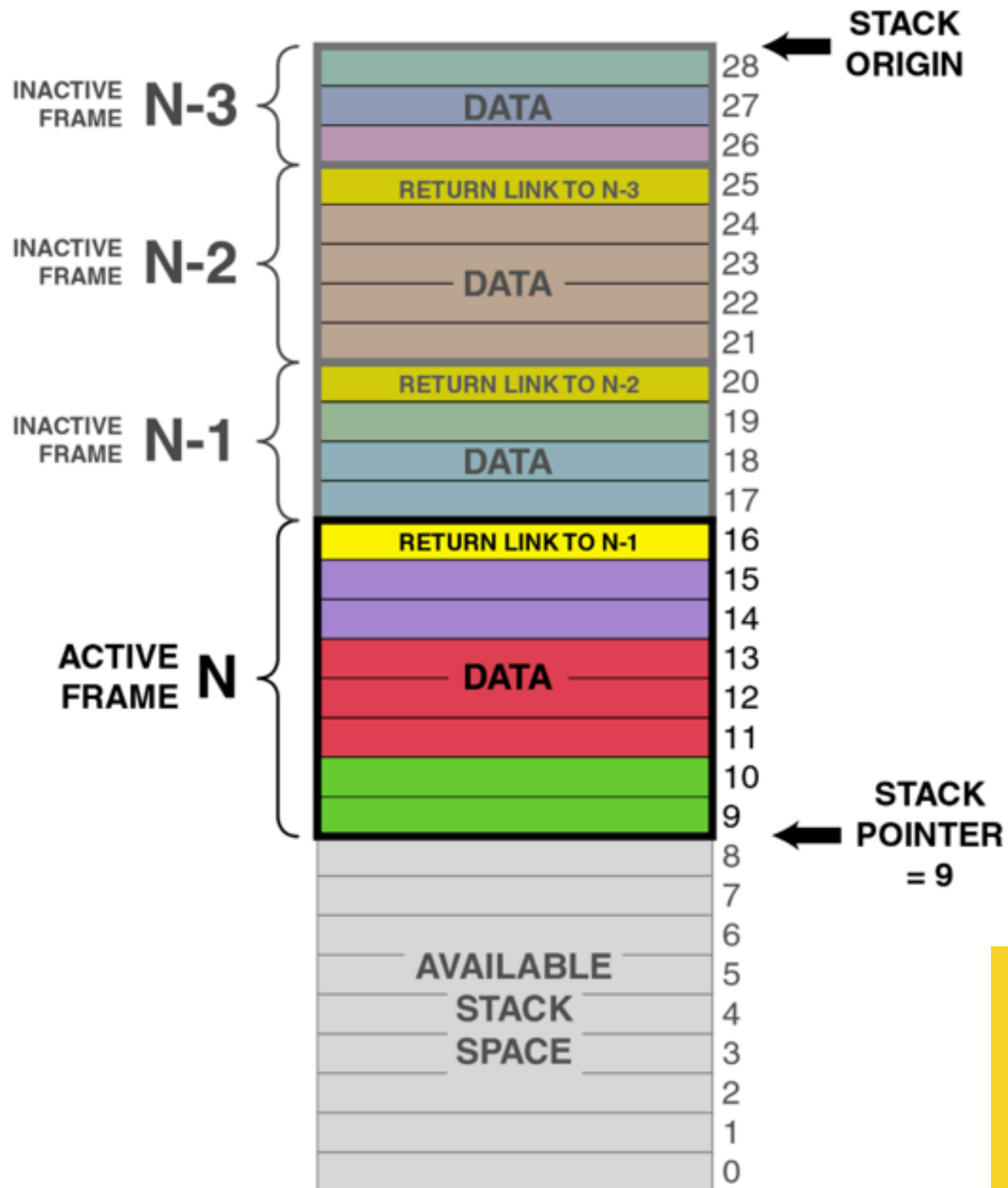
# Accelerated Version

```
public static long fasterFibonacci(final int n)
{
    if ( n <= 1 )
        return n;

    return fibonacciAcc(0, 1, n);
}
```

```
private static long fibonacciAcc(long prev, long curr, int n)
{
    if (n == 1)
        return curr;

    return fibonacciAcc(curr, prev + curr, n - 1);
}
```



# The Call Stack

language implementation  
structure which allows interrupting  
and resuming a function

# Peril: stackoverflow

```
public static long fasterFibonacci(final int n)
{
    if ( n == 1 )
        return n;

    return fibonacciAcc(0, 1, n);
}

public static void main(String[] args) {
    for ( int counter = 0; counter <= 50; counter++ )
        System.out.printf( "Fibo(%d) = %d\n", counter, fasterFibonacci(counter) )
}
```

Fibonacci

/Library/Java/JavaVirtualMachines/jdk1.8.0\_25.jdk/Contents/Home/bin/java ...

Exception in thread "main" java.lang.StackOverflowError

at Fibonacci.fibonacciAcc([Fibonacci.java:11](#))

at Fibonacci.fibonacciAcc([Fibonacci.java:11](#))

at Fibonacci.fibonacciAcc([Fibonacci.java:11](#))



Demo - Testing the depth of Java's Stack

# Tail Recursion

```
private static long fibonacciAcc(long prev, long curr, int n)
{
    if (n == 1)
        return curr;
    return fibonacciAcc(curr, prev + curr, n - 1);
}
```

can be optimized with *tail call optimization*  
easy to transform in iteration

# The power of recursion lies in the possibility of defining

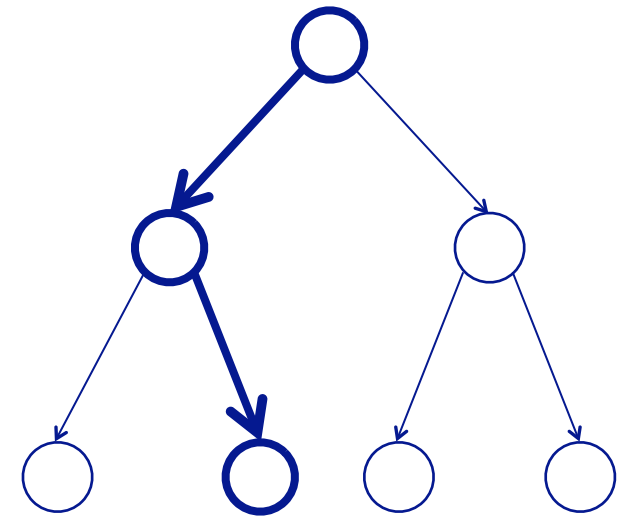
- **an infinite set of objects by a finite statement**
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# Modeling Binary Trees

```
class Node<T> {  
    T value;  
    Node<T> left;  
    Node<T> right;  
  
    Node(T value) {  
        this.value = value;  
    }  
  
    void visit() {  
        System.out.print(this.value + " ");  
    }  
}
```



# Traversing Trees

recursive algos for recursive structures

```
static void traverse(Node<?> node, ORDER order) {  
    if (node == null) {  
        return;  
    }  
    switch (order) {  
        case PREORDER:  
            node.visit();  
            traverse(node.left, order);  
            traverse(node.right, order);  
            break;  
    }  
}
```

base case

recursive calls

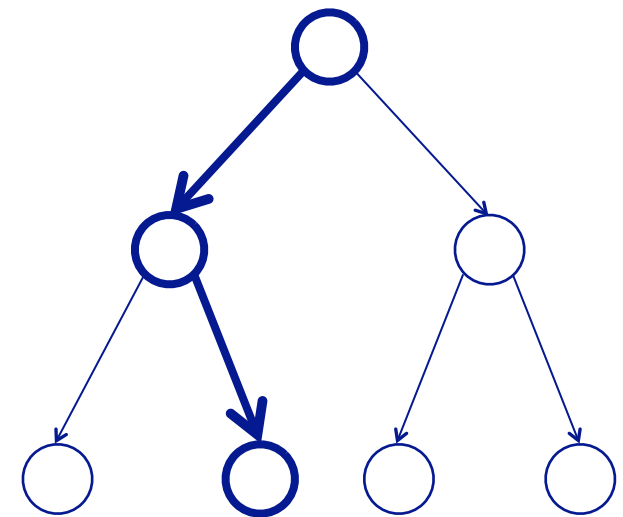
“structural recursion”

# The power of recursion lies in the possibility of defining

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# Binary search





/\*

*a – array with integer elements*

*x – element we want to know whether present*

*left – leftmost position where element could be*

*right – rightmost position where element could be*

\*/

```
private int binarySearch(int[ ] a, int x, int left, int right) {
```

```
    if (left > right) return -1;
```

```
    int mid = (left + right)/2;
```

```
    if (a[mid] == x) return mid;
```

```
    else if (a[mid] < x)
```

```
        return binarySearch(a, x, mid+1, right);
```

```
    else
```

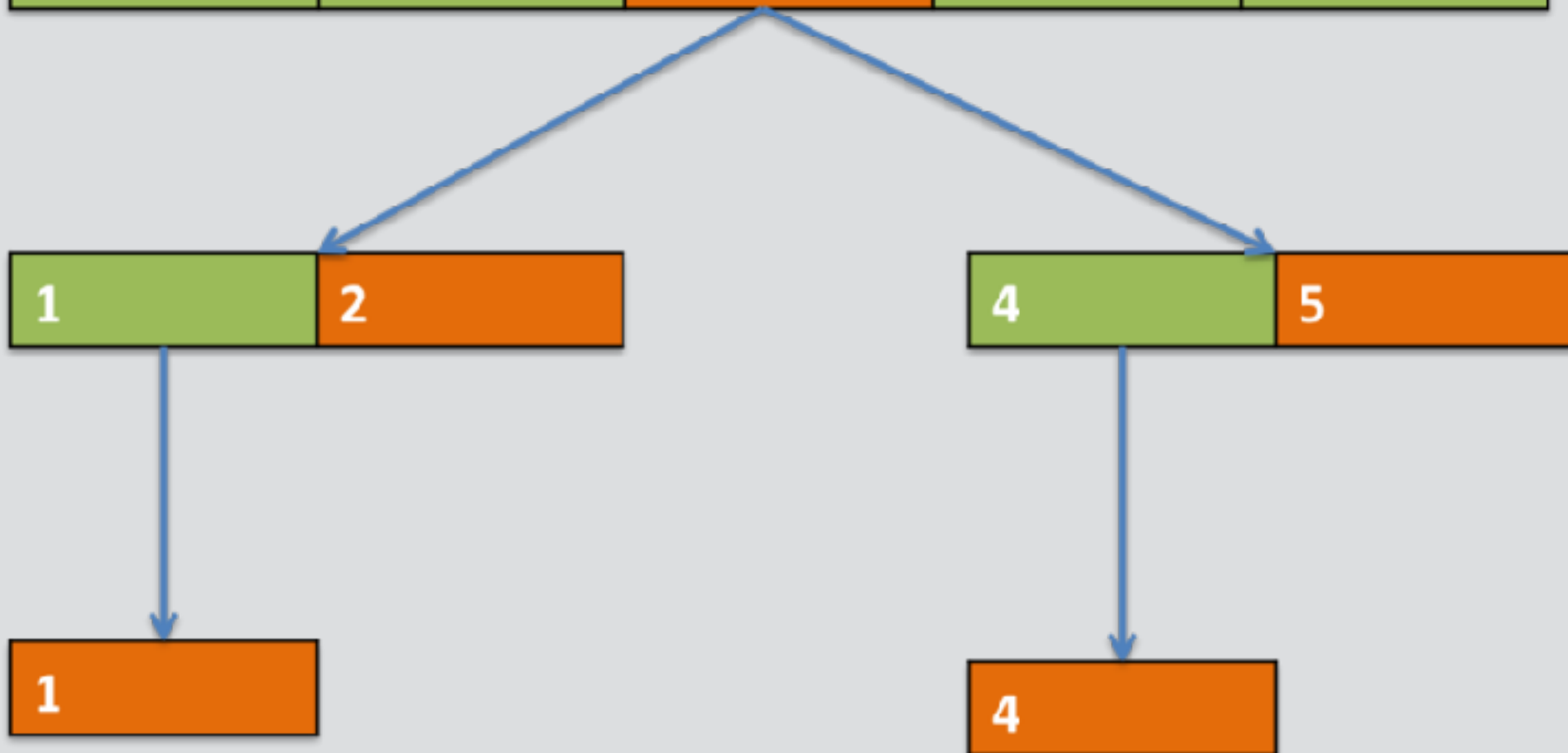
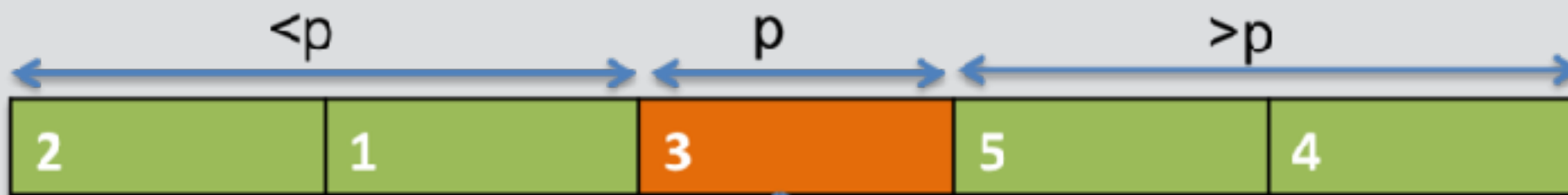
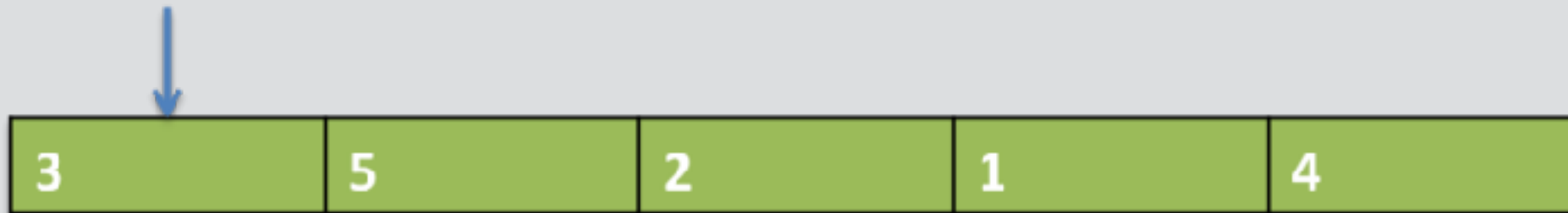
```
        return binarySearch(a, x, left, mid-1);
```

```
}
```

base  
case

recursive  
calls

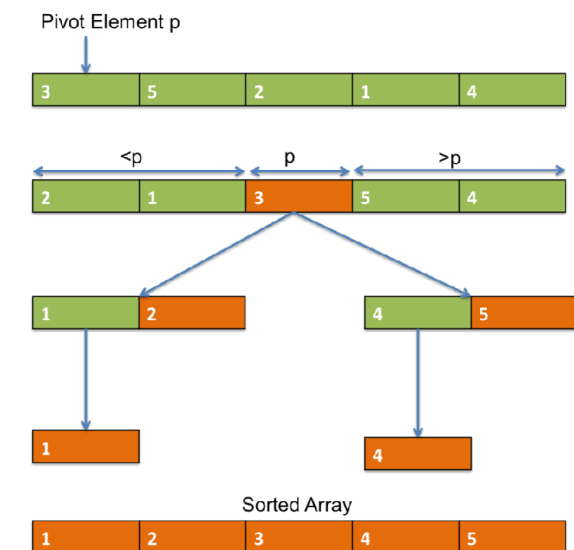
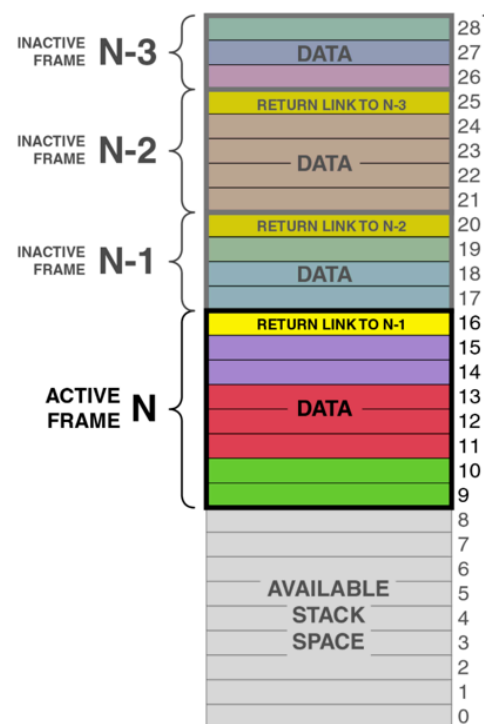
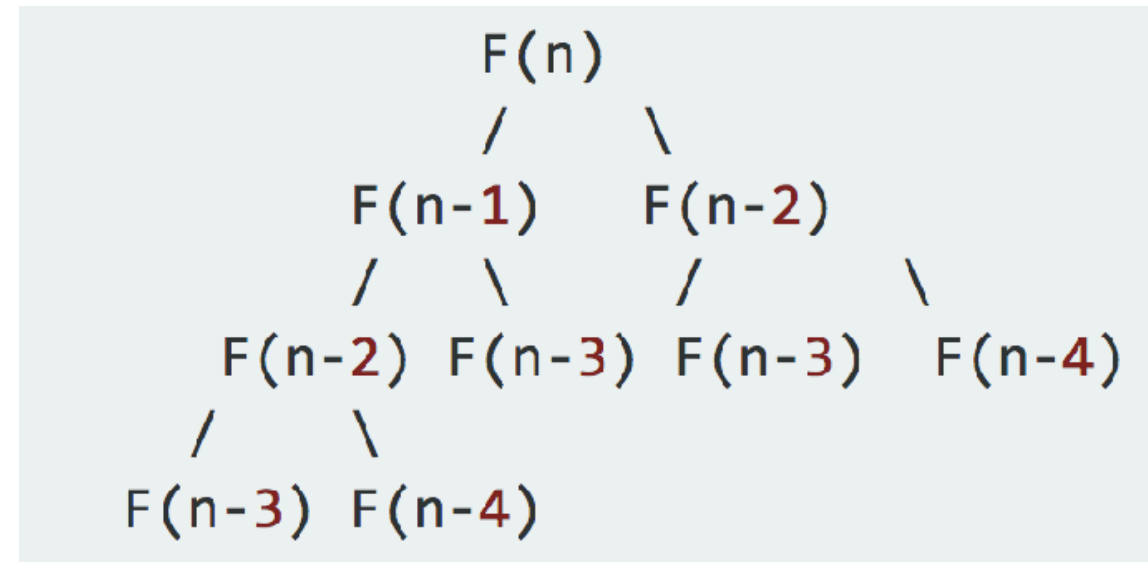
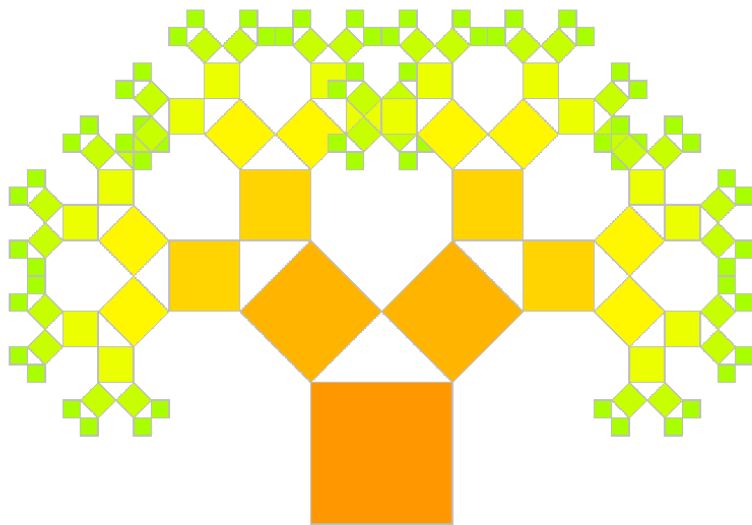
Pivot Element p



Sorted Array



# Quicksort



Code: [https://github.com/mircealungu/open\\_lectures/recursion](https://github.com/mircealungu/open_lectures/recursion)

# Further Reading

- Dynamic Programming (CodeChef)
- Tail Call Optimization on JVM (DrDobbs)
- Replacing Recursion with Iteration (ThoughtWorks)

# Credits

- Cover Art: <http://qcc-art.deviantart.com/art/Tony-Monahan-Turtles-All-the-Way-Down-590318917>



# Backup Materials

# Replacing Recursion with Iteration

```
Stack<Object> stack;  
stack.push(first_object);  
while( !stack.isEmpty() ) {  
    // Do something  
    my_object = stack.pop();  
  
    // Push other objects on the stack.  
  
}
```

```
void quicksort(int *array, int left, int right)
{
    int stack[1024];
    int i=0;

    stack[i++] = left;
    stack[i++] = right;

    while (i > 0)
    {
        right = stack[--i];
        left = stack[--i];

        if (left >= right)
            continue;

        int index = partition(array, left, right);
        stack[i++] = left;
        stack[i++] = index - 1;
        stack[i++] = index + 1;
        stack[i++] = right;
    }
}
```

iterative  
quicksort

# Composite Design Pattern

