# Using Inhibitors to Achieve Universality of Sequential P Systems

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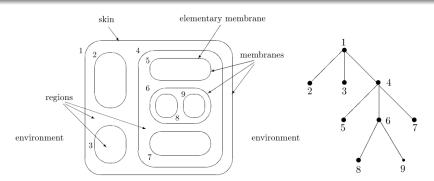
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- Overview of P systems
  - P systems
  - Variants

- Sequential P systems with inhibitors
  - Accepting case
  - Generating case

#### Membrane structure



#### Contents of the membrane

- multiset of objects
  - a | b | b
- rewriting rules
  - $\bullet \ a \mid b \mid b \rightarrow a \mid a_{out} \mid b_{in_6}$
  - $\bullet \ b \to a \mid \delta$

#### P system

We define a P system as

$$\Pi = (V, \mu, w_1, w_2, \dots, w_m, R_1, R_2, \dots, R_m)$$
, where:

- V is an alphabet of objects
- ullet  $\mu$  is a membrane structure
- $w_1, w_2, \dots w_m$  are initial multisets of objects in membranes  $1 \dots m, w_i \subseteq \mathbb{N}^V$
- $R_1, R_2, ..., R_m$  are sets of rewriting rules in membranes 1...m, where

$$R_i \subseteq (\mathbb{N}^V \setminus 0^V) \times \mathbb{N}^{V \times (\{here, out\} \cup \{in_1, ... in_m\})}$$

.

- configuration = membrane structure + contents
- computational step: maximal parallelism

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$$egin{array}{ccccc} a \mid b \mid b 
ightarrow c & & (r_1) \ & b 
ightarrow c \mid c & & (r_2) \end{array}$$

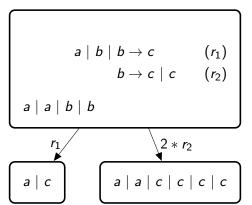
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$$\begin{array}{c|cccc}
a & b & b \rightarrow c & (r_1) \\
& b \rightarrow c & c & (r_2) \\
a & a & b & b
\end{array}$$

$$\begin{array}{c|ccccc}
r_1 \\
& a & c
\end{array}$$

P systems

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- Parikh mapping: PsRE

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- asynchronous (∼ sequential in most cases) [Freund, 2005]
- minimal parallelism (PsRE) [Ciobanu et al., 2007]

## Extensions of sequential P systems

- priorities [Ibarra et al., 2005]
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- inhibitors [Kováč, 2014, submitted]
- further study (rules with emptyness detection, ...)

# Register machine

Minsky register machine is M = (n, P, i, h), where:

- n is the number of registers
- P is a set of labeled instructions of type:
  - (add(r), k, l)
  - (sub(r), k, l)
  - halt
- i is the initial instruction
- h is the final instruction

# Simulation of register machine

- Contents of register j is represented by the multiplicity of the object a<sub>j</sub>
- For an instruction (add(r), k, l) there is a rule  $e \rightarrow a_j | f$
- For an instruction (sub(r), k, l) there are rules
  - $e|a_j \rightarrow f$
  - $e \rightarrow z | a_j$
- Halting rules
  - $h|a_j \to h|\#$  for all  $a \le j \le n$
  - $\bullet$  #  $\rightarrow$  #

# Overview of the simulation for the generating case

- Simulation of a maximal parallel step
- Phases of membranes: RUN and SYNCHRONIZE, represented by objects

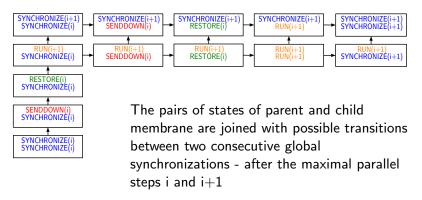
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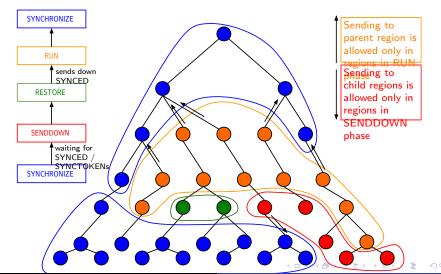
- Preventing the rule application on already rewritten objects in the same maximal parallel step
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  - add RESTORE phase
- Sending objects via membranes
  - add SENDDOWN phase

## Parent and child membrane phases



Obr. : Possible pairs of states of parent and child membrane

#### Snapshot of all membrane states



Thanks for your attention