

Sequential P Systems with Active Membranes Working on Sets

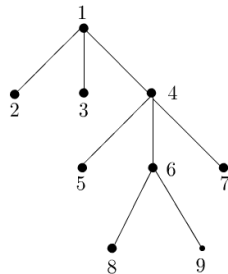
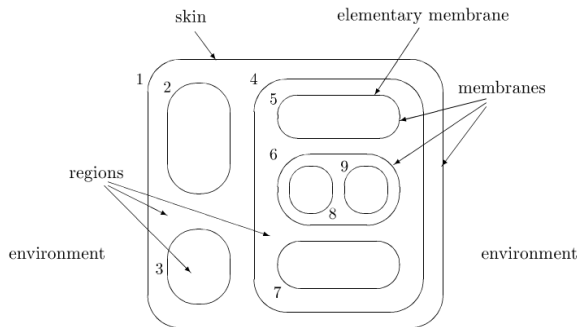
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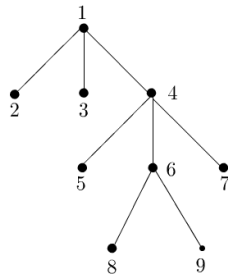
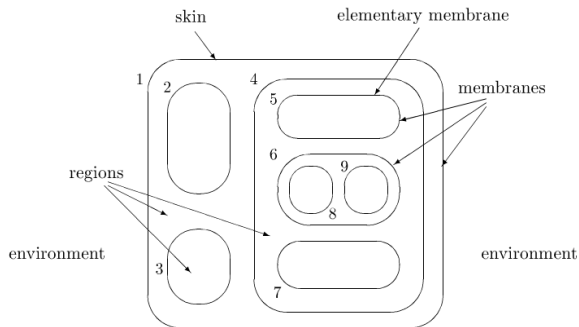
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- 1 Overview of formal models
 - P systems
 - Models with set semantics
- 2 Sequential active set membrane systems
 - Original semantics
 - inject-or-create semantics
 - wrap-or-create semantics

Membrane structure

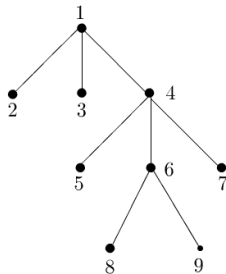
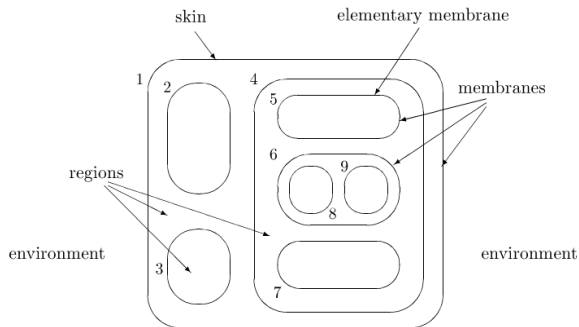


Membrane structure



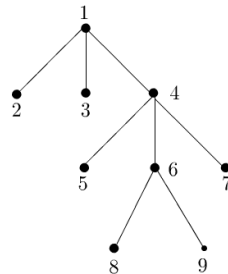
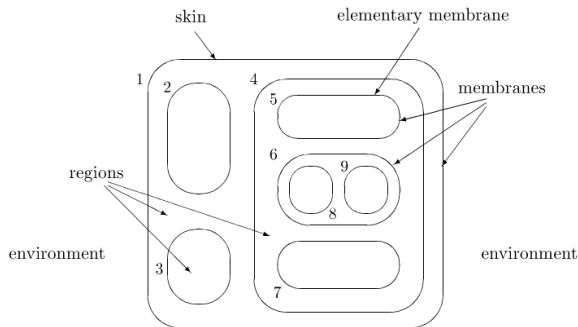
- Multisets

Membrane structure



- Multisets
- Rewriting rules

Membrane structure



- Multisets
- Rewriting rules
- Passive vs. Active

Computation

- Maximal parallel vs. sequential

Computation

- Maximal parallel vs. sequential
- Language
 - generating mode
 - accepting mode

Multiset vs. set semantics

- How realistic is the counting?
- Effectiveness of verification techniques
- No conflict (objects can participate as reactants in as many rules as they want)

Reaction systems

TODO: definition

Set membrane systems

- Alhazov [Alhazov, 2006]: multiplicities of objects are ignored
R, with active membranes universal
- Kleijn, Koutny [Kleijn and Koutny, 2011]:
min-enabled computational step \Rightarrow sequential R
- maximal parallel \Rightarrow deterministic

Sequential active set membrane systems

- $\Pi = (\Sigma, C_0, R_1, \dots, R_m)$

Sequential active set membrane systems

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- $C = (T, l, c)$
 - $l : V(T) \rightarrow \{1, \dots, m\}$
 - $c : V(T) \rightarrow 2^\Sigma$

Sequential active set membrane systems

- $\Pi = (\Sigma, C_0, R_1, \dots, R_m)$
- $C = (T, l, c)$
 - $l : V(T) \rightarrow \{1, \dots, m\}$
 - $c : V(T) \rightarrow 2^\Sigma$
- Rewriting rules
 - $u \rightarrow w$
 - $u \rightarrow w\delta$
 - $u \rightarrow [{}_j v_1]_j v_2,$
where $u \in \Sigma, |u| \geq 1, v_1, v_2 \in \mathbb{N}$ and $w \in (\Sigma \times \{\cdot, \uparrow, \downarrow_j\})$

Proof of universality

We can simulate the register machine.

- Simple simulation: unary encoding of register values by nested membranes. Linear time overhead and linear number of membranes. Alphabet size: 2 times number of instructions + registers.
- Optimized simulation: binary encoding by nested membrane labels. Logarithmic time overhead and logarithmic number of membranes. Alphabet size: 3 times number of instructions + 5.

TODO: add some slides with example simulation

Issues with original semantics

- Explicit membrane creation rule
- Sending an object to a child membrane

inject-or-create

No explicit membrane creation rule TODO: proof of universality

wrap-or-create

Explicit membrane creation rule which is not blocked by presence of another such membrane TODO: proof of universality Additional control objects, but constant time overhead.



Alhazov, A. (2006).

P systems without multiplicities of symbol-objects.
Information Processing Letters, 100(3):124–129.



Kleijn, J. and Koutny, M. (2011).

Membrane systems with qualitative evolution rules.
Fundam. Inf., 110(1-4):217–230.

Thanks for your attention!