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Cs477

Hw5

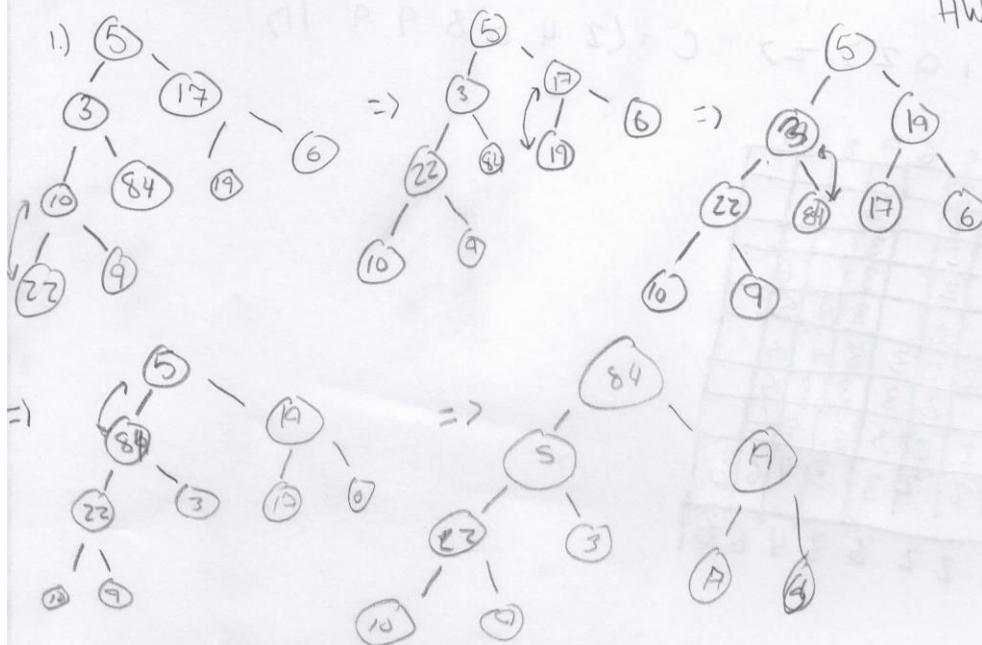
1A.) 6.3-1

Create Max Heap on $\langle 5, 3, 17, 10, 84, 19, 6, 22, 9 \rangle$

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2.) 8.2-3

The algorithm would be correct regardless of the order. However, the algorithm isn't stable. In the final for loop, a repeating element that appears later in the array is placed in the output array. In the original algorithm, a repeating element that appears earlier, would be placed before a repeating element that appears later in the array.

$\hookrightarrow A[1, 5, 5, 6] \Rightarrow B = [2]$ First 5 should be placed first.

\downarrow Note to grader* \uparrow this is 8.2-3, 1B \rightarrow 8.2-1 will be on the next page

(1)

$$C = \langle 2 \ 2 \ 2 \ 2 \ 1 \ 0 \ 2 \rangle \Rightarrow C = \langle 2 \ 4 \ 6 \ 8 \ 9 \ 9 \ 11 \rangle$$

Handwritten multiplication table for base 7. The table is 10x10, with columns and rows labeled 0 through 6. The entries are calculated modulo 7. For example, 3x3=2, 4x4=2, 5x5=6, 6x6=1. The table is drawn on a grid with some cells shaded gray.

	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	0
2	2	3	4	5	6	7	8	9	0	1
3	3	4	5	6	7	8	9	0	1	2
4	4	5	6	7	8	9	0	1	2	3
5	5	6	7	8	9	0	1	2	3	4
6	6	7	8	9	0	1	2	3	4	5
7	7	8	9	0	1	2	3	4	5	6
8	8	9	0	1	2	3	4	5	6	7
9	9	0	1	2	3	4	5	6	7	8

3) On seperate file

40) $6.3 - 2$

b) 6.3-2
we want the loop in index i to decrease from $[A.length/2]$

to 1 because every node would be satisfied. If

we increased from 1, then the root node

would be correct first and not retrocedo. Energy

note would be satisfied except the root node

if we increased by 1 at A. length $\frac{1}{2}$.

②

CA problem3

Array A is not a heap.

Array B is not a heap.

Press any key to continue . . .

50) 8.2-4 (197)

Describe an algorithm that gives n ints in range $0-k$, preprocesses its input, then answers any query about how many of the n ints fall into the range $[a \dots b]$ in $O(1)$.
Algorithm should use $O(n+k)$ preprocessing time

Given: Array $A[1 \dots n]$, range is $0-k$

1.) Two new arrays

↳ $B[0 \dots k]$;

$C[0 \dots k]$;

2.) Every element in B , set to 0

$B[0 \dots k] = \{0\}$;

3.) Make for loop to increment C respective to elements in A array

↳ 3A.) for $(i=1 \rightarrow A.length)$ / $O(n)$ complexity

3B.) $B[A[i]] = B[A[i]] + 1$

3C.) $C[0] = B[0]$

4.) Use for loop to store count to j th element in $B[]$ array

4A.) for $(j=1 \rightarrow k)$ / $O(k)$ time

4B.) $C[j] = B[j] + C[j-1]$

↳ Complexity is $O(n+k)$; $C[b] - C[a] + B[a] = \Rightarrow$ $O(1)$ after pre-processing