

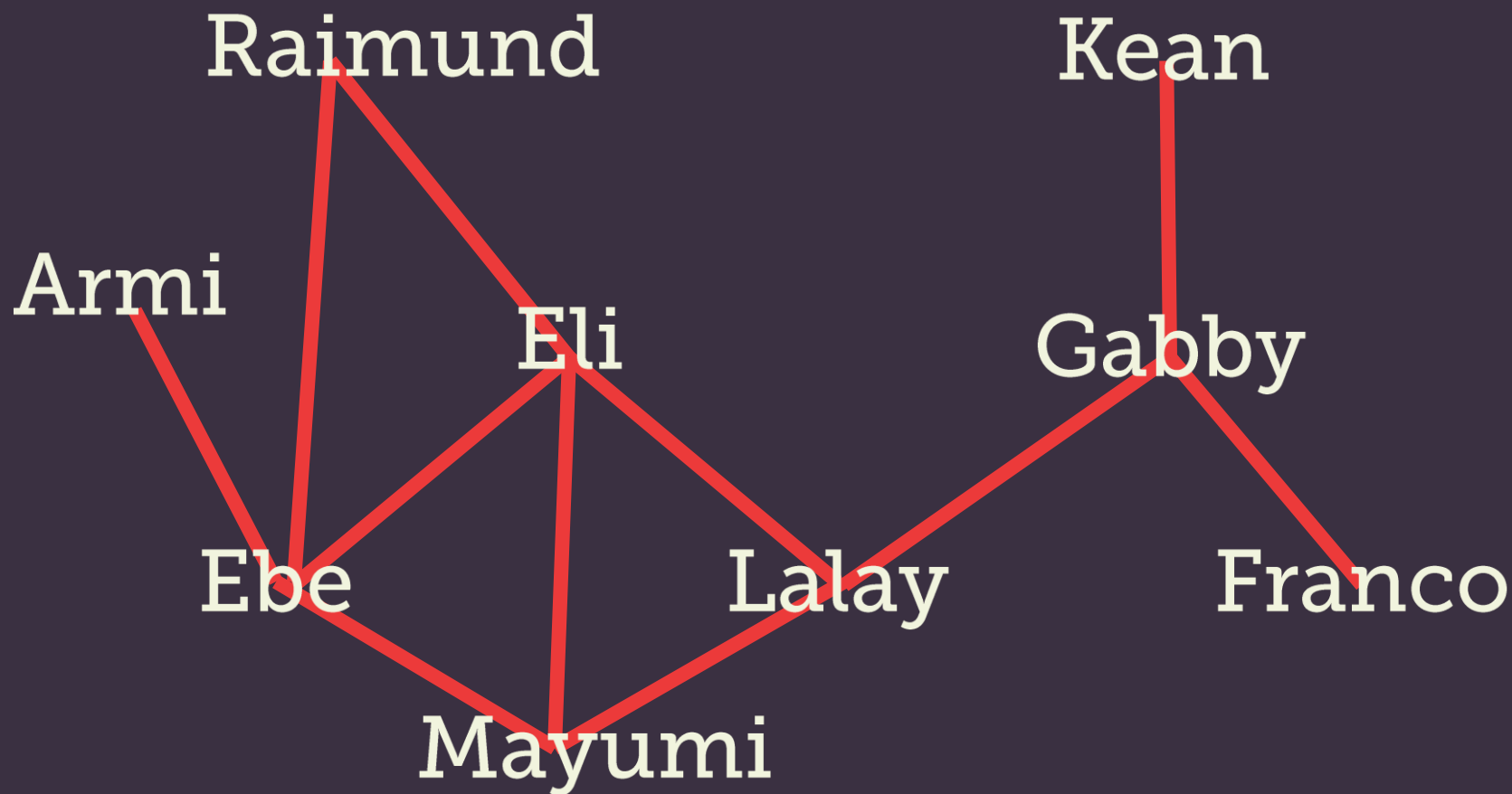
# GRAPH ALGORITHMS

+ DATA STRUCTURES

# GRAPH

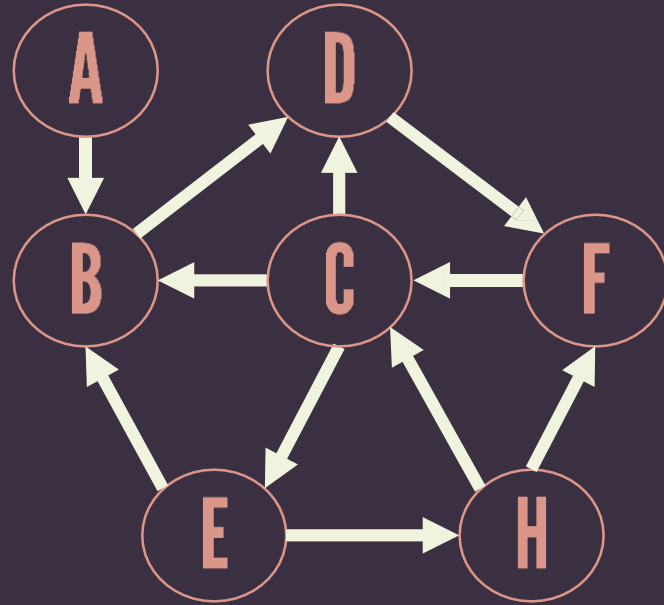
Graph consists of a set of vertices and a set of edges.

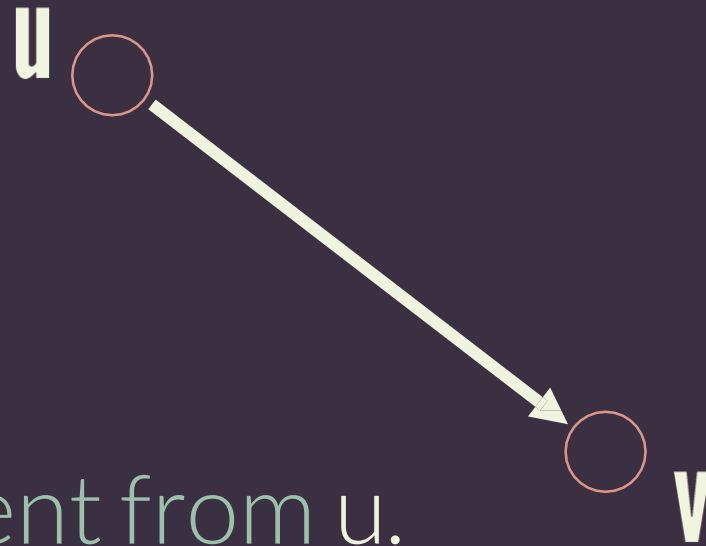
$$G = (V, E)$$



types of  
**GRAPHS**

directed  
**GRAPH**



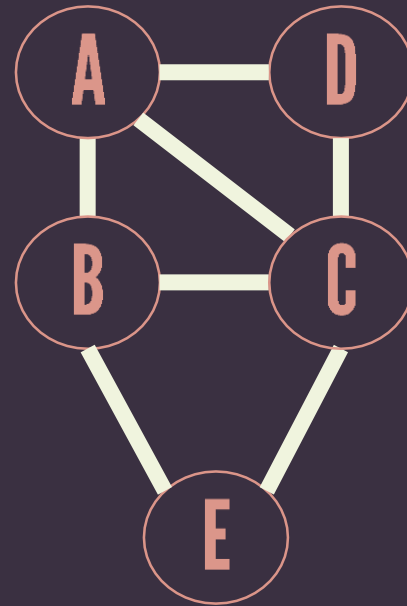


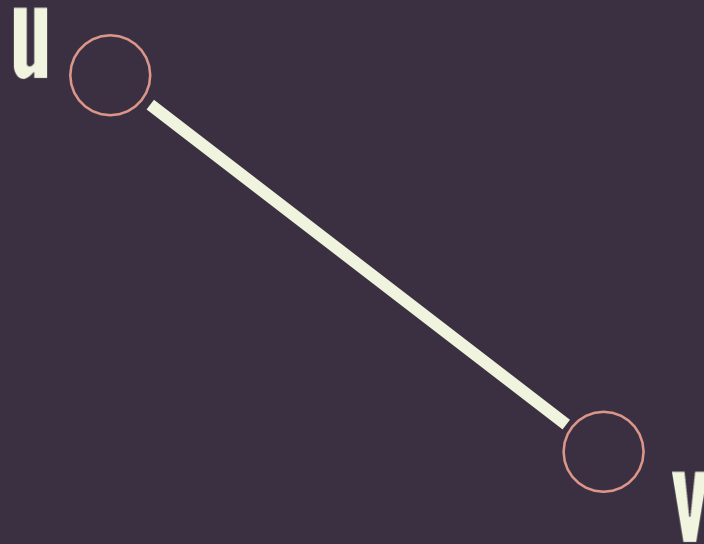
Edge  $(u,v)$  is incident from  $u$ .

Edge  $(u,v)$  is incident to  $v$ .

Vertex  $v$  is adjacent to vertex  $u$ .

undirected  
**GRAPH**

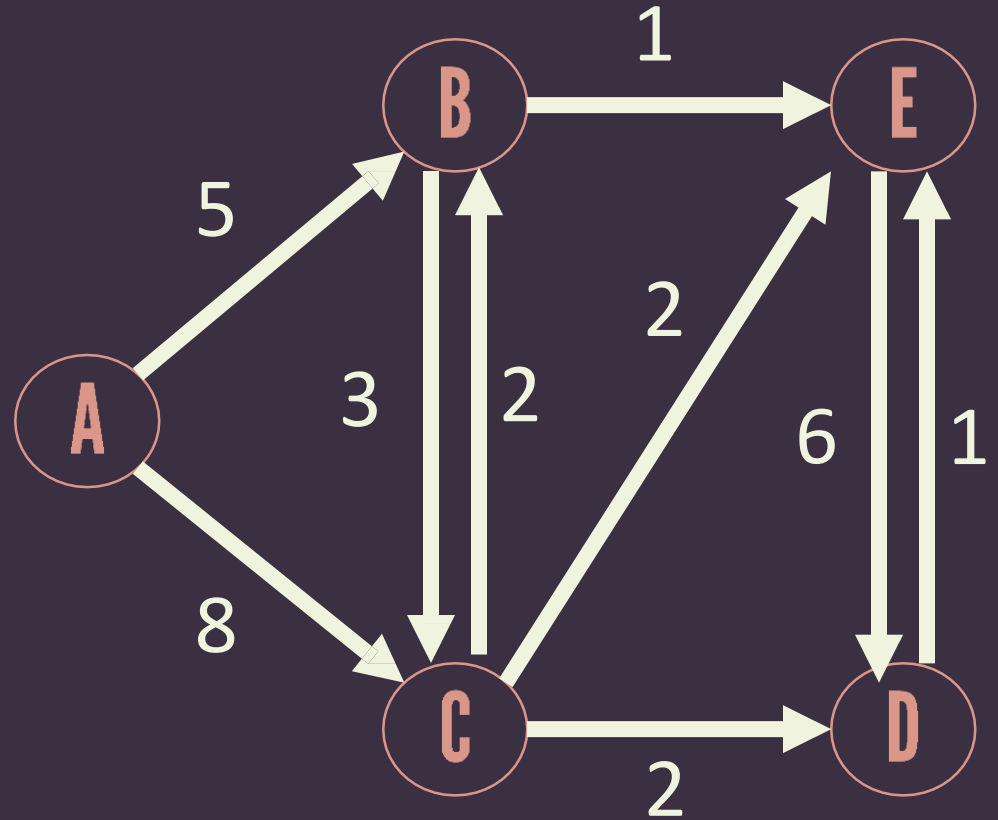




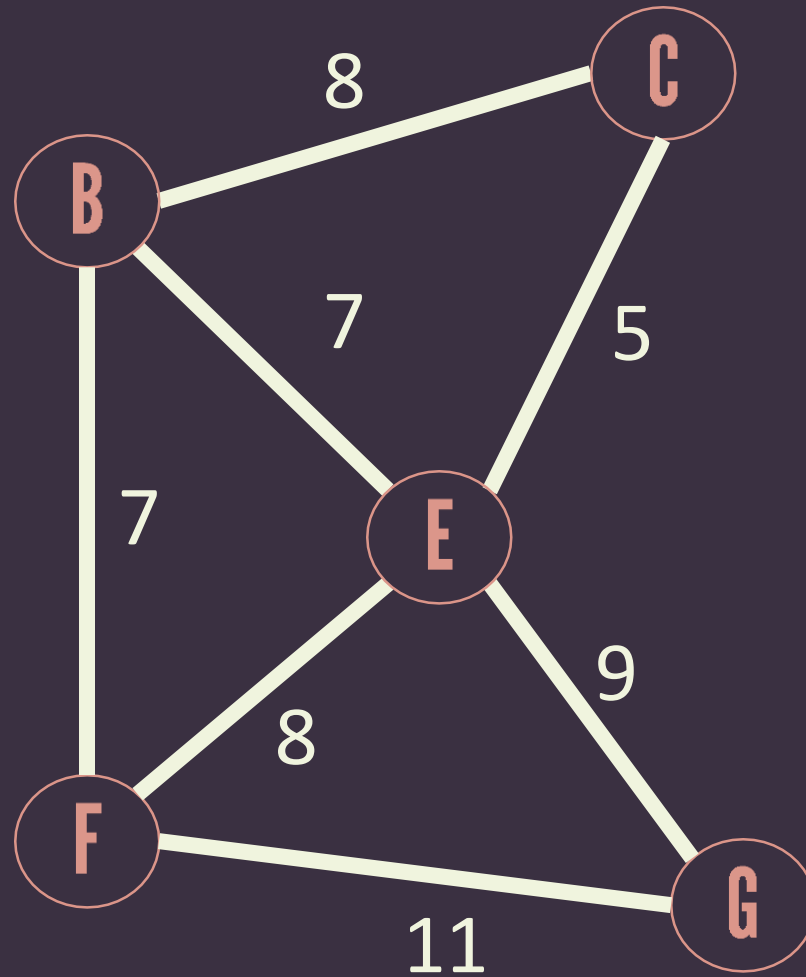
Edge  $(u,v)$  or  $(v,u)$  is incident on  $u$  and  $v$ .  
Vertex  $v$  is adjacent to vertex  $u$ .  
Vertex  $u$  is adjacent to vertex  $v$ .



weighted  
**GRAPH**



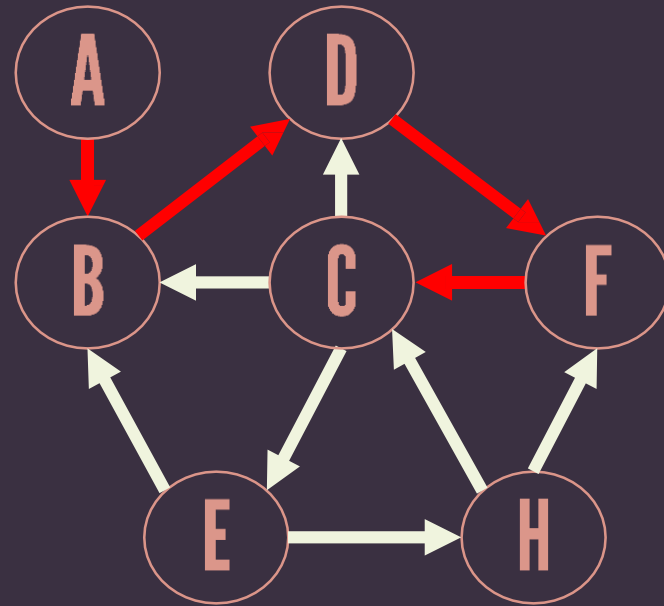
weighted  
**GRAPH**



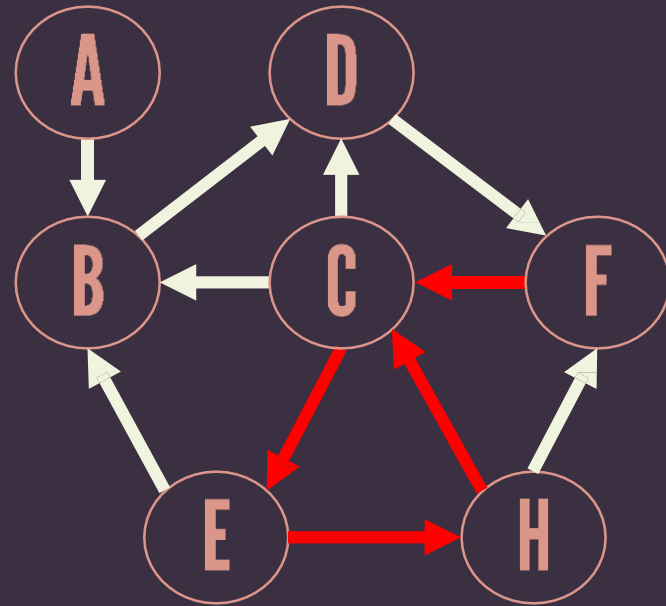
# path in a **GRAPH**

Path represents a sequence of edges between the two vertices.

path in a  
**GRAPH**



# path in a GRAPH



in-degree,  $\rho^+(v)$

# of edges incident to  $v$

out-degree,  $\rho^-(v)$

# of edges incident from  $v$

degree of vertex  $v$ ,  $\rho(v)$

# of edges incident on  $v$

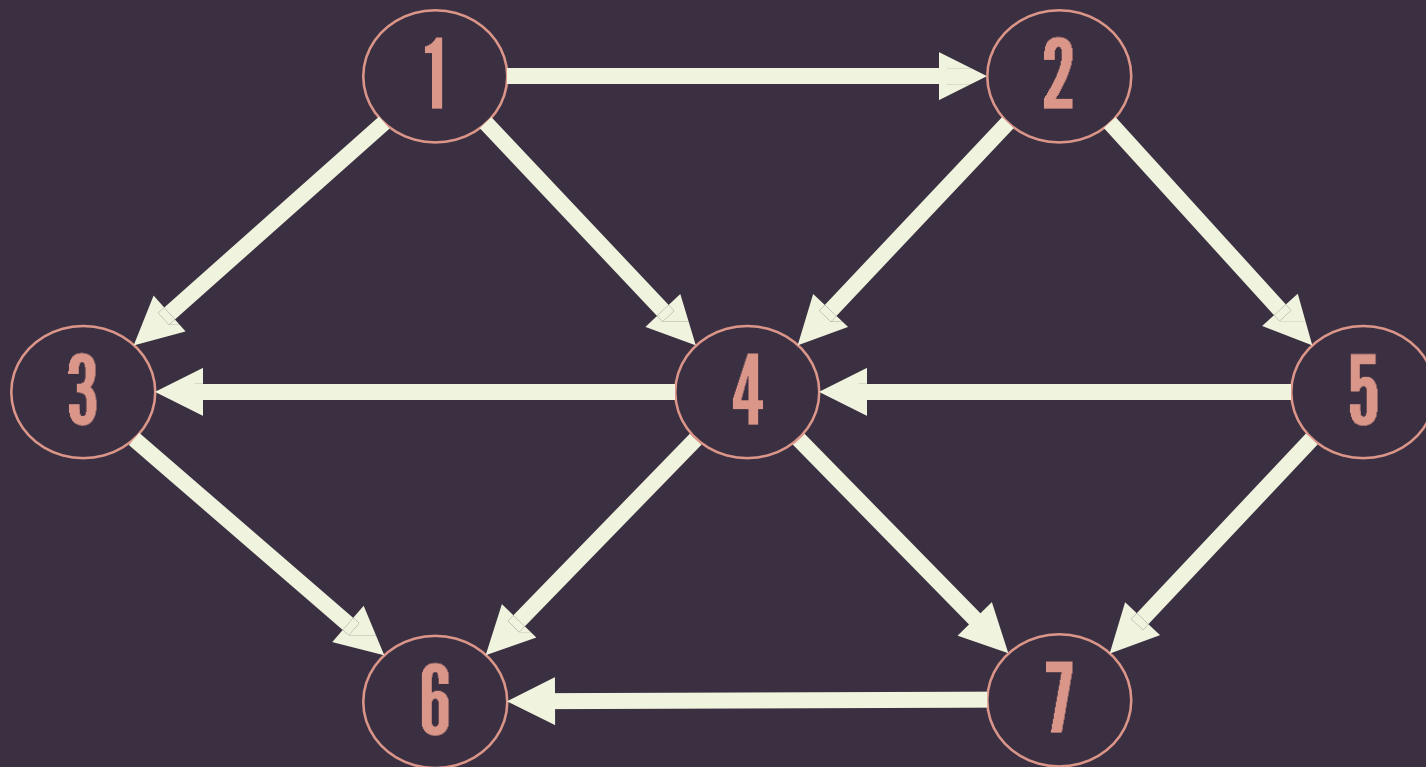
$$= \rho^-(v) + \rho^+(v)$$

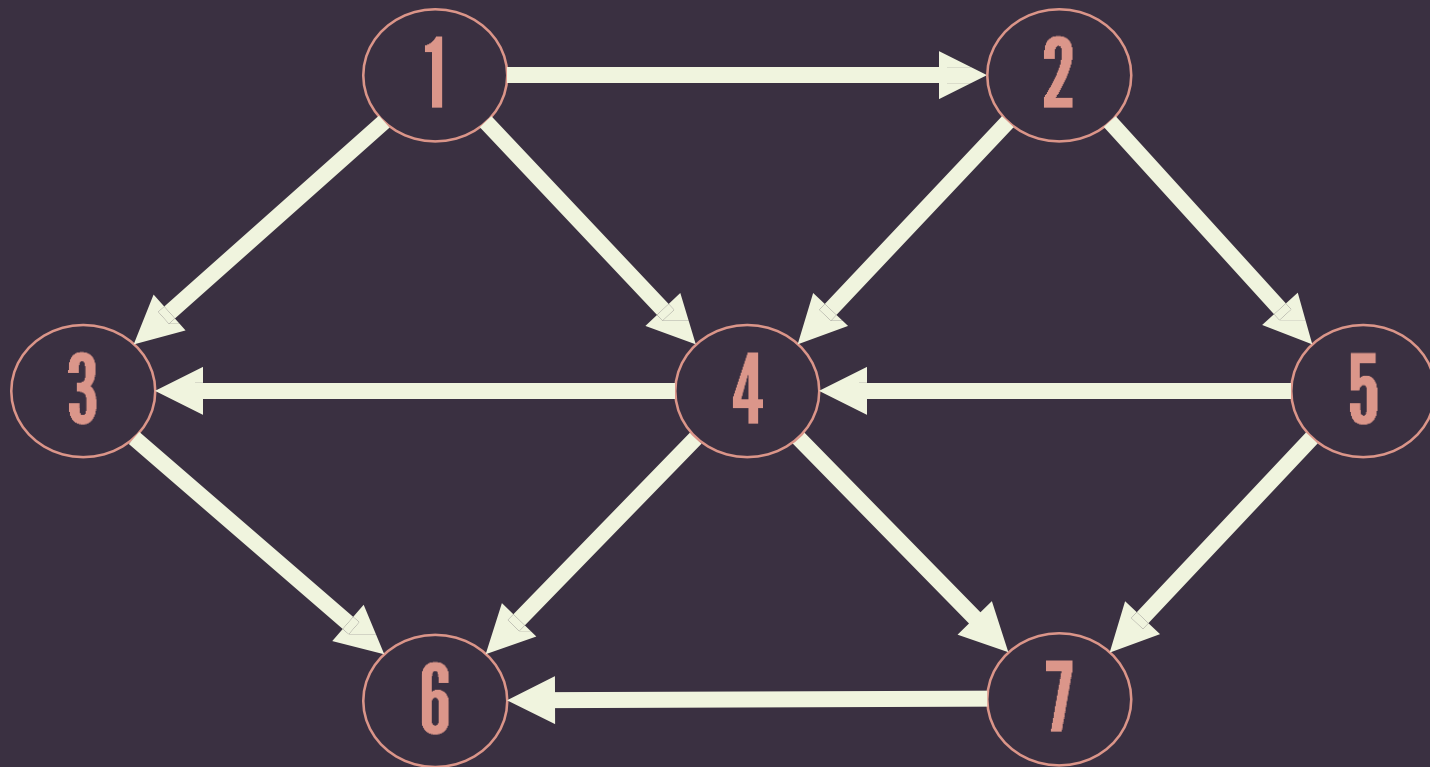
degree of vertex  $v$ ,  $\rho(v)$

# of edges incident on  $v$ .

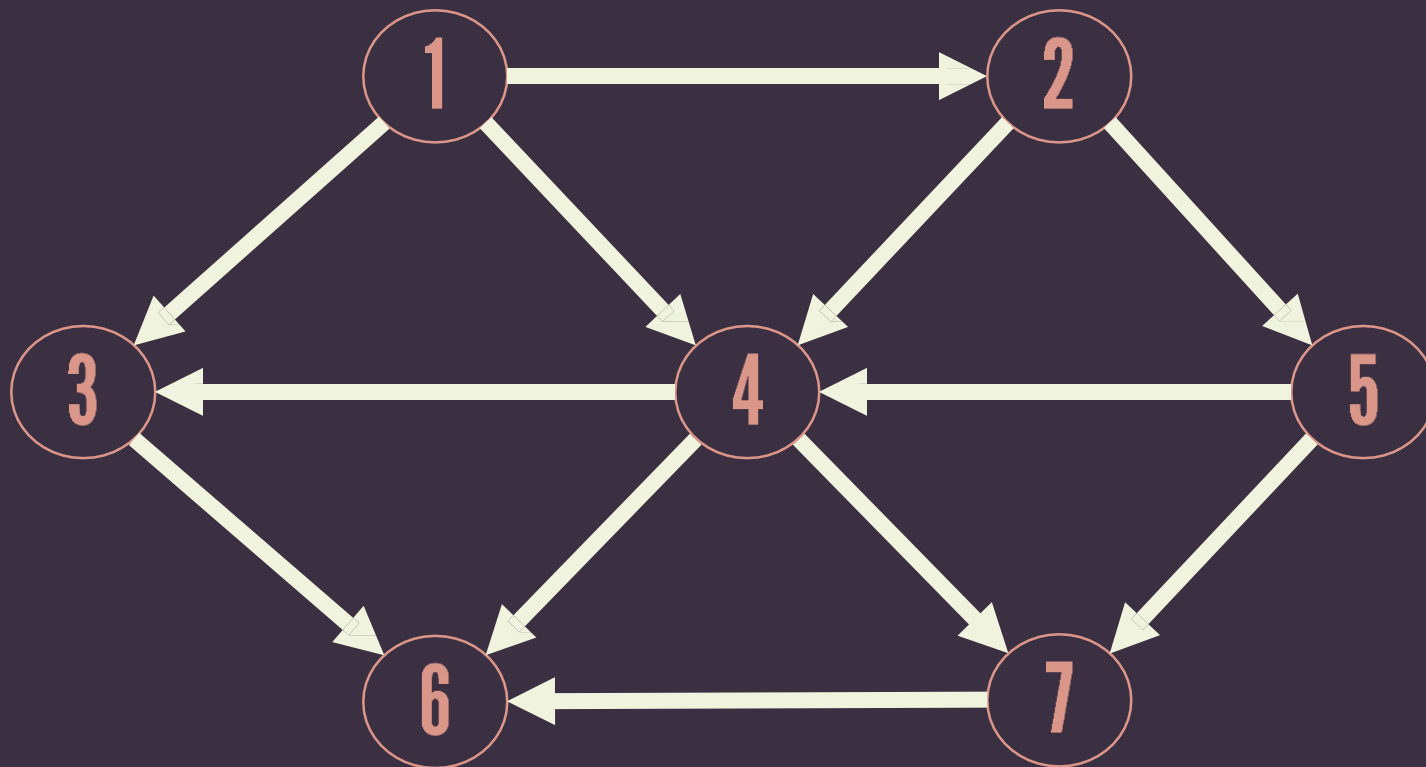
(also applicable for undirected graphs)



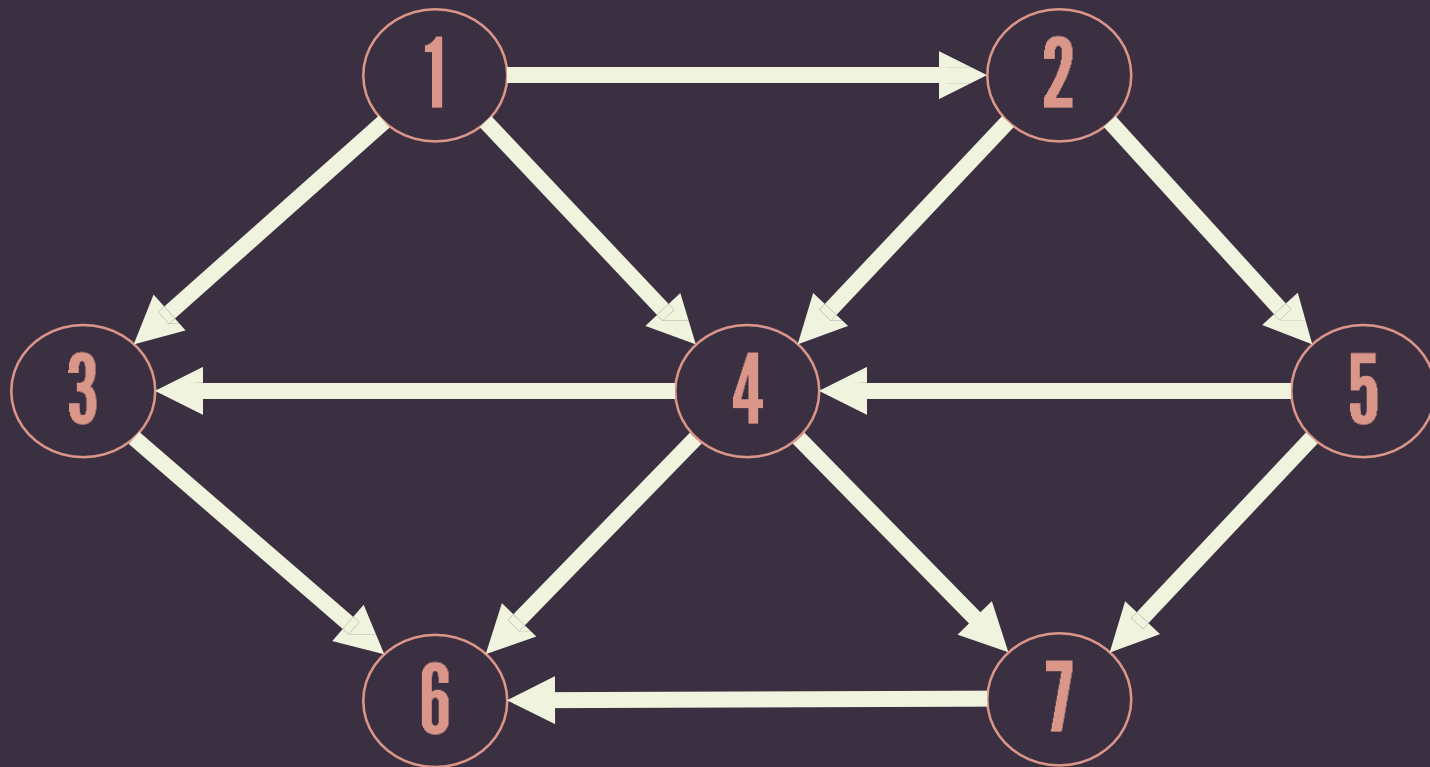




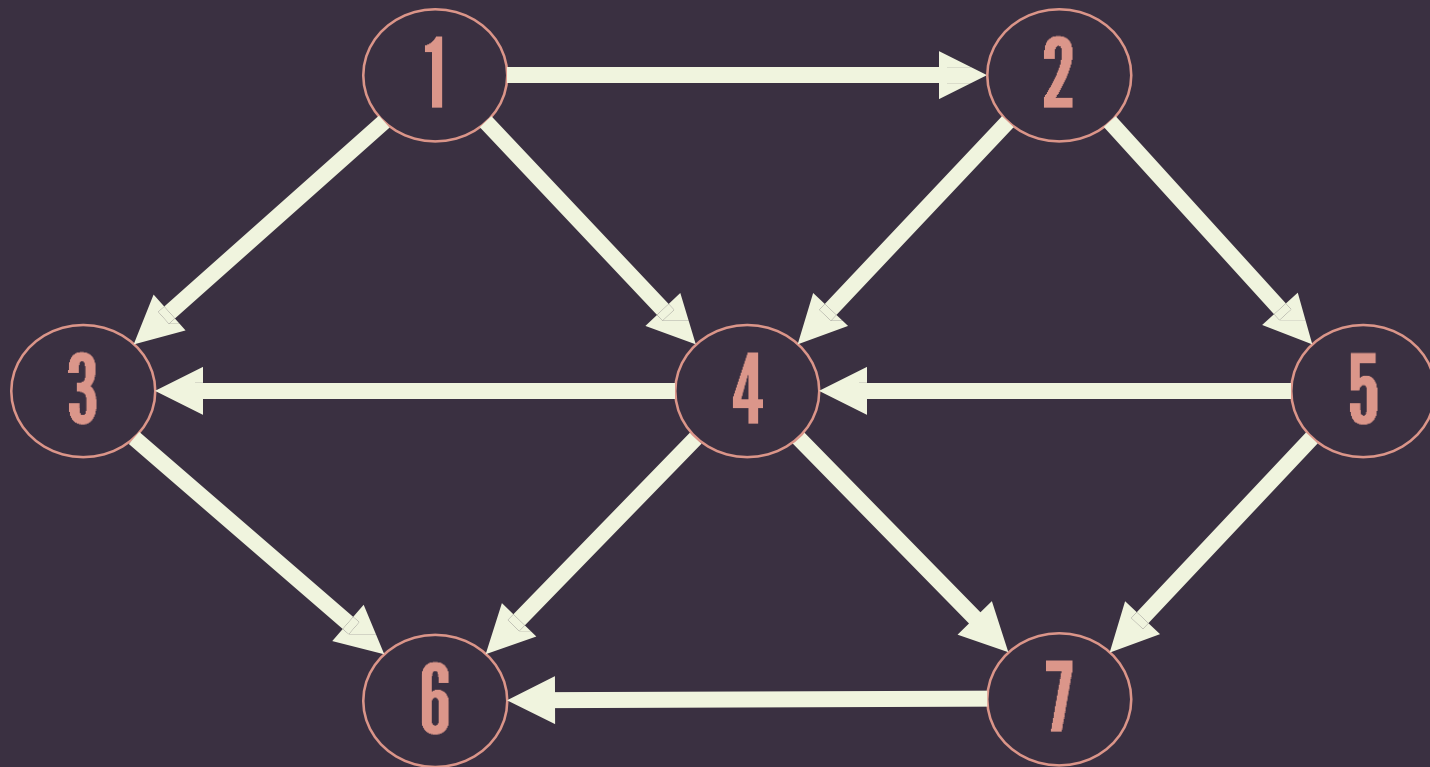
$$\rho^+(4) : 3$$



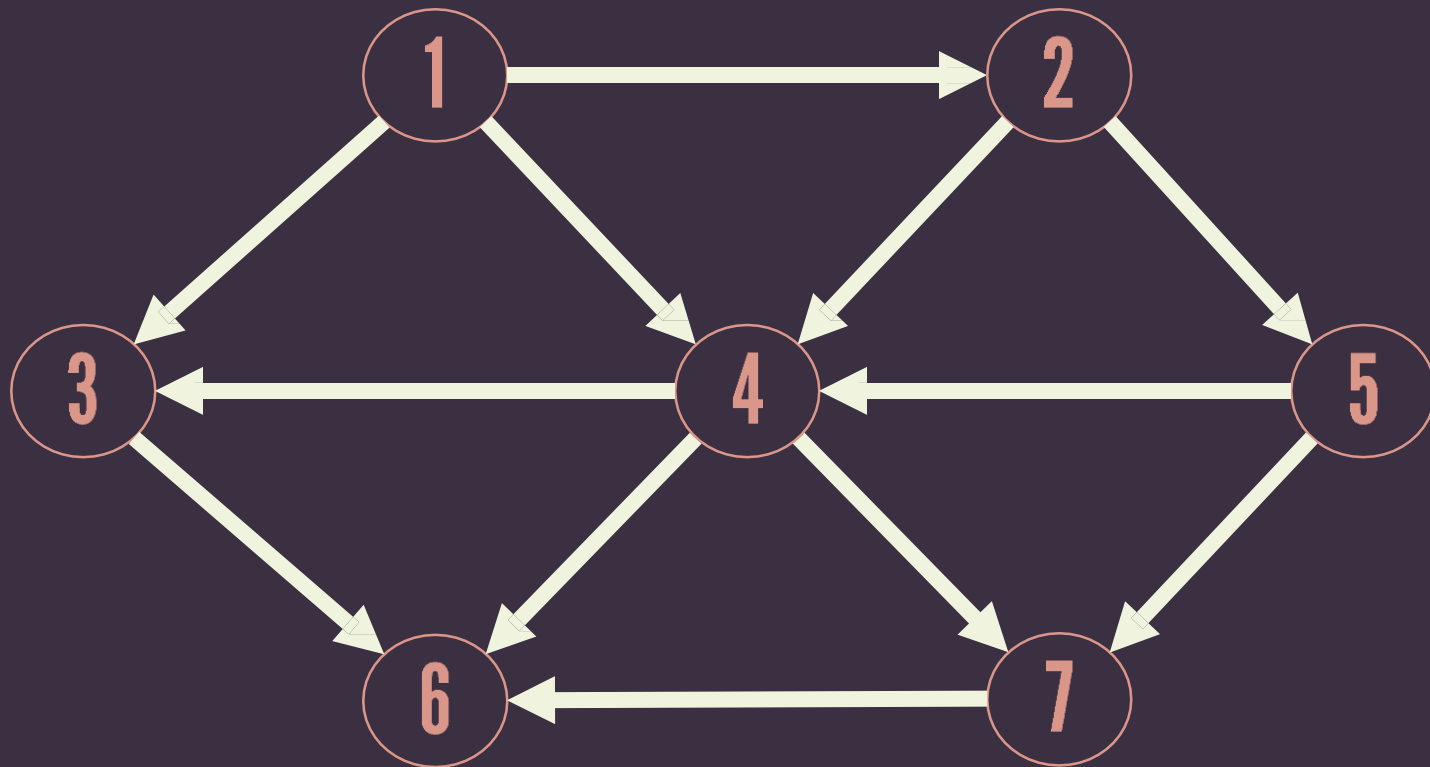
$$\rho^+(7) : 2$$



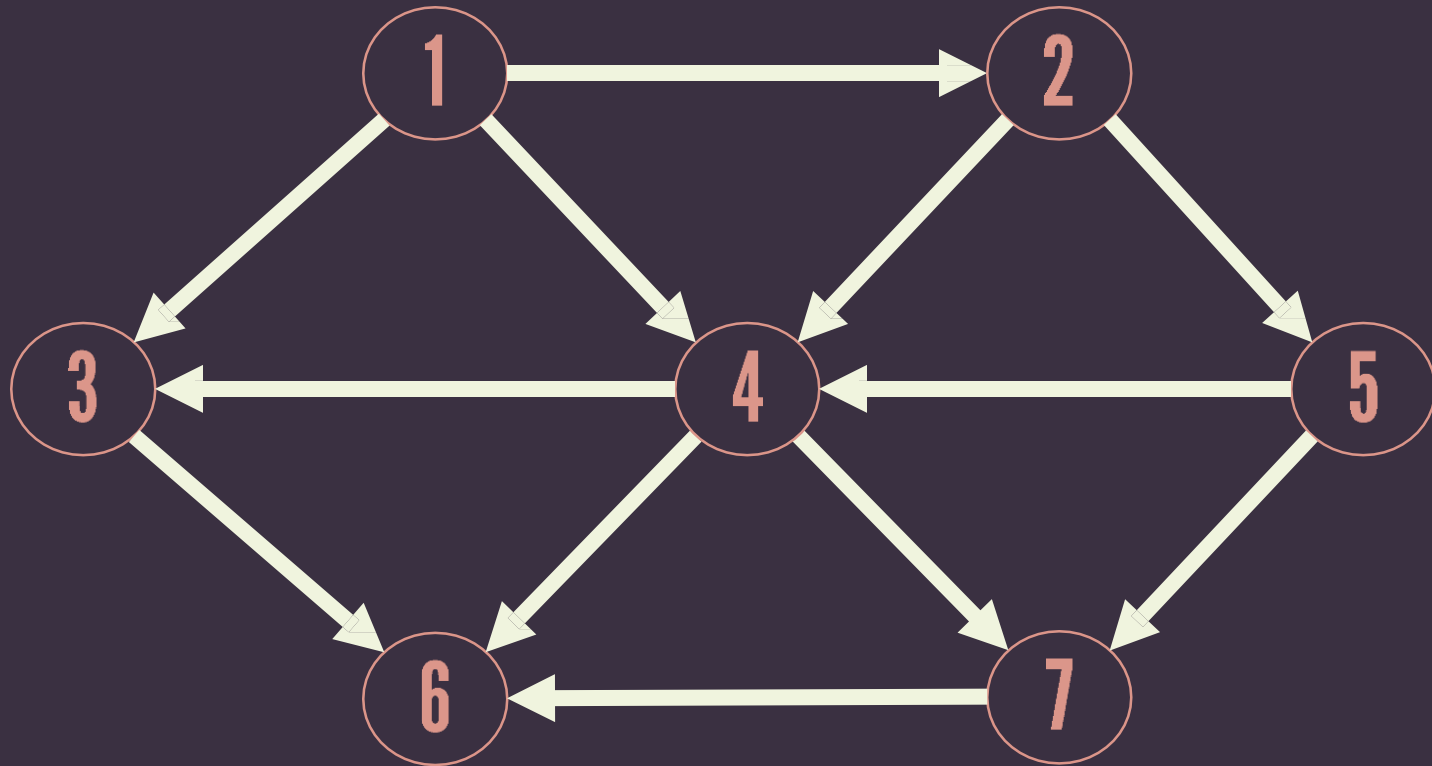
$\rho^-(7) : 1$



$\rho^-(5) : 2$



$\rho(5) : 3$



$\rho(4) : 6$

# GRAPH

## Representations



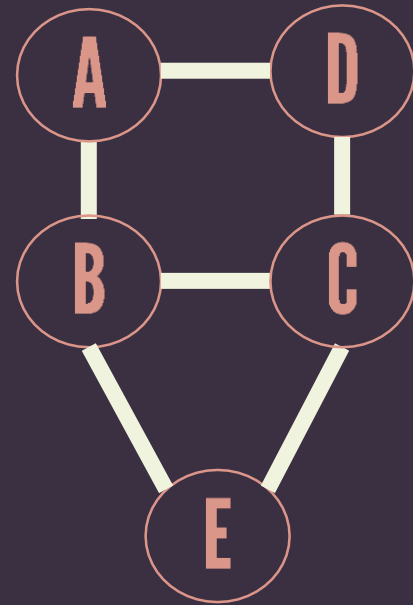
adjacency  
**MATRIX**

2D  
**ARRAY**

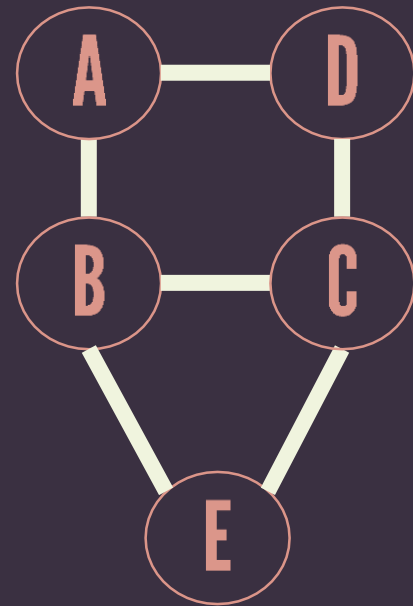
adjacency  
**LIST**

list of  
**LISTS**

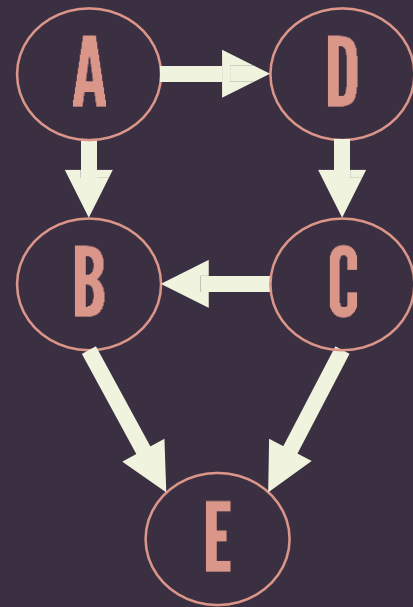
	A	B	C	D	E
A	0	1	0	1	0
B	1	0	1	0	1
C	0	1	0	1	1
D	1	0	1	0	0
E	0	1	1	0	0



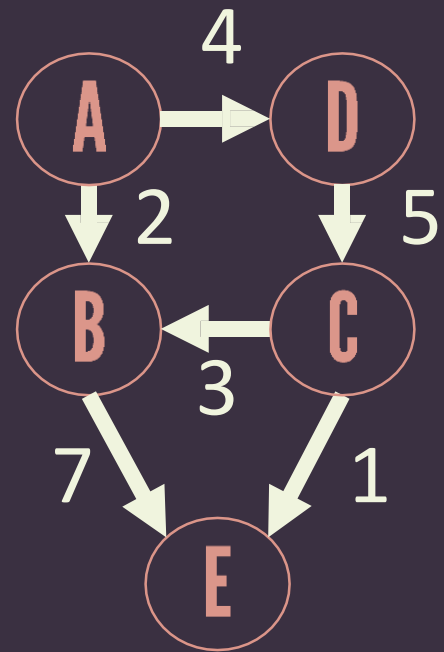
	A	B	C	D	E
A	$\infty$	1	$\infty$	1	$\infty$
B	1	$\infty$	1	$\infty$	1
C	$\infty$	1	$\infty$	1	1
D	1	$\infty$	1	$\infty$	$\infty$
E	$\infty$	1	1	$\infty$	$\infty$

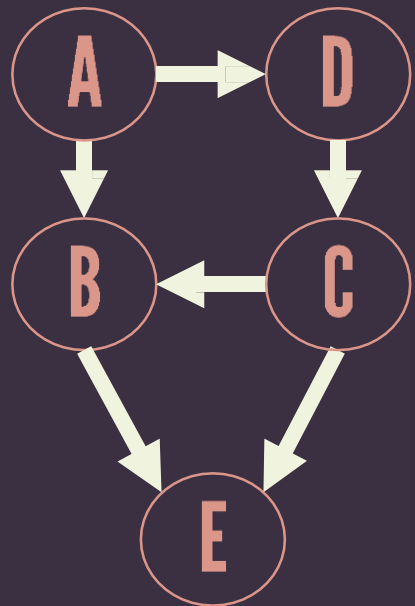
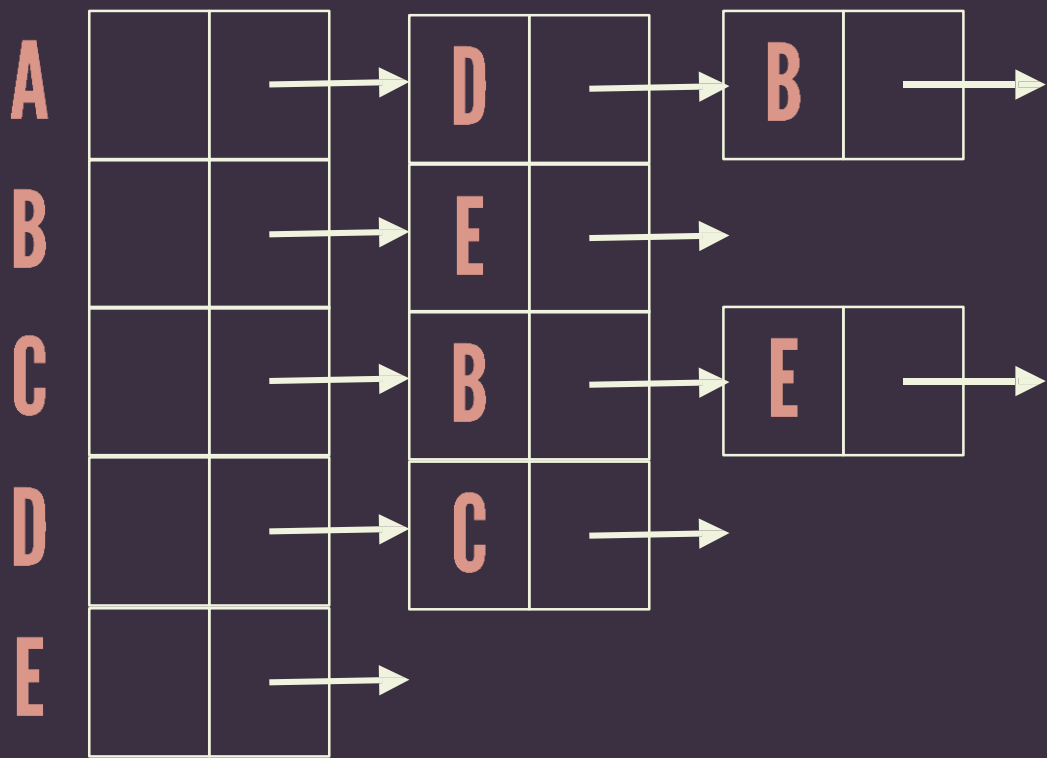


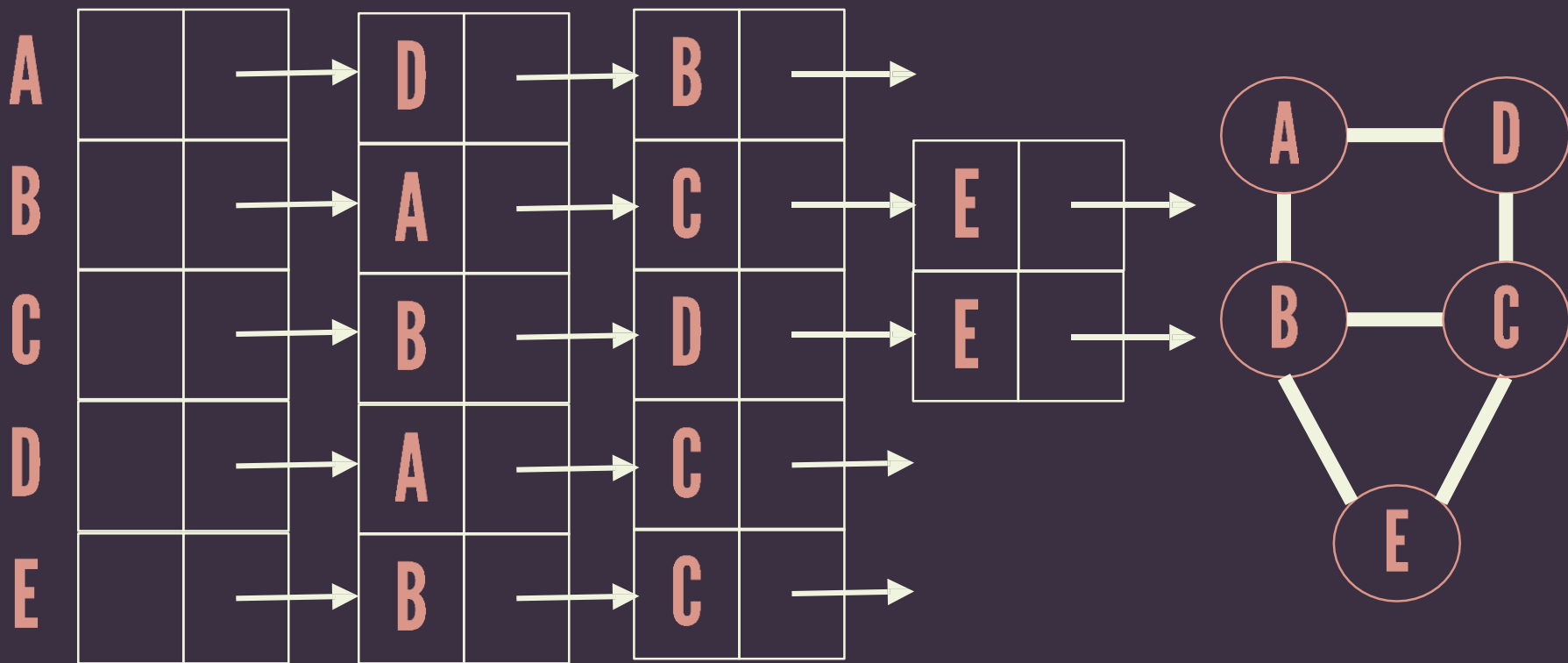
	A	B	C	D	E
A	0	1	0	1	0
B	0	0	0	0	1
C	0	1	0	0	1
D	0	0	1	0	0
E	0	0	0	0	0



	A	B	C	D	E
A	0	2	0	4	0
B	0	0	0	0	7
C	0	3	0	0	1
D	0	0	5	0	0
E	0	0	0	0	0





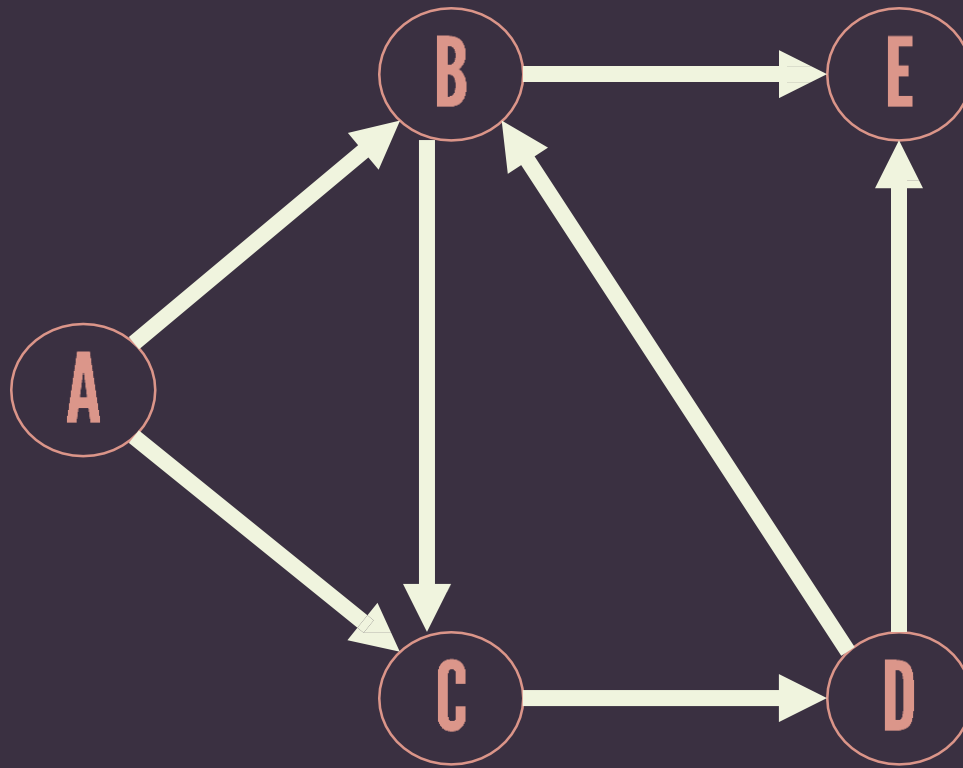




# GRAPH ALGORITHMS

# topological **SORT**

An ordering of vertices in a directed acyclic graph such that if there is a path from  $v_i$  to  $v_j$ , then  $v_j$  appears after  $v_i$  in the ordering.



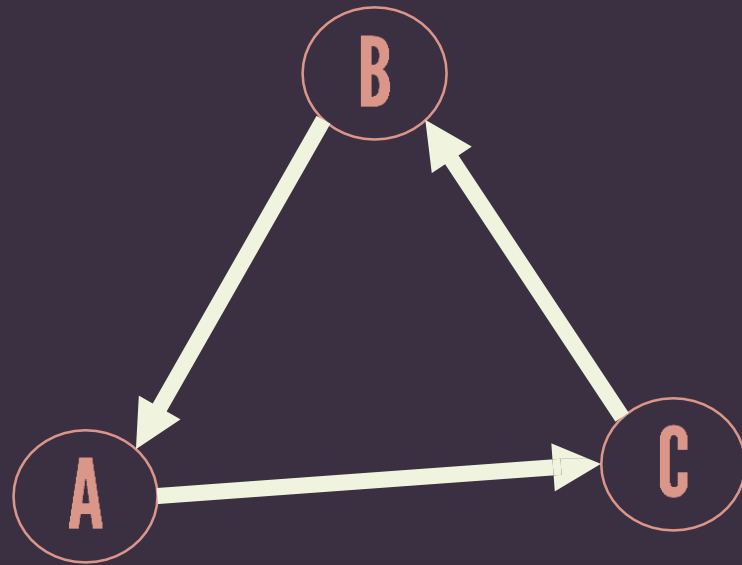
A, B, C, D, E

A, C, D, B, E

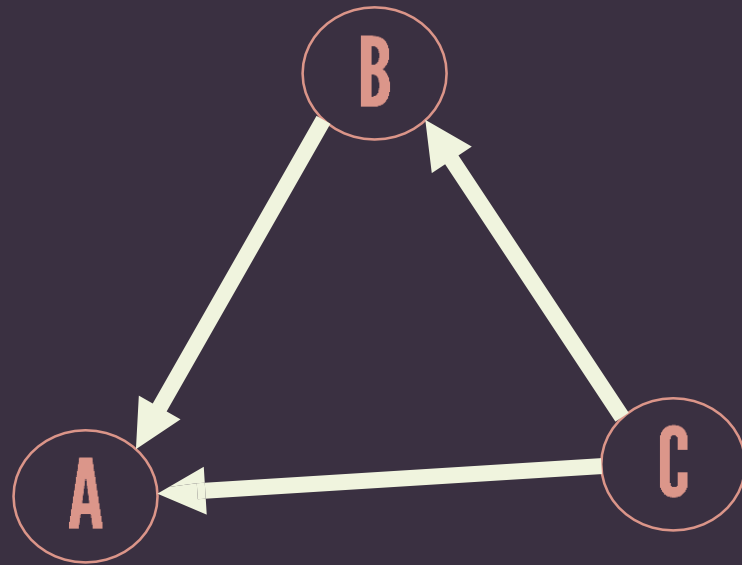
directed  
acyclic graph  
**DAG**

A directed graph with no  
directed cycles.

directed  
acyclic graph  
**DAG**

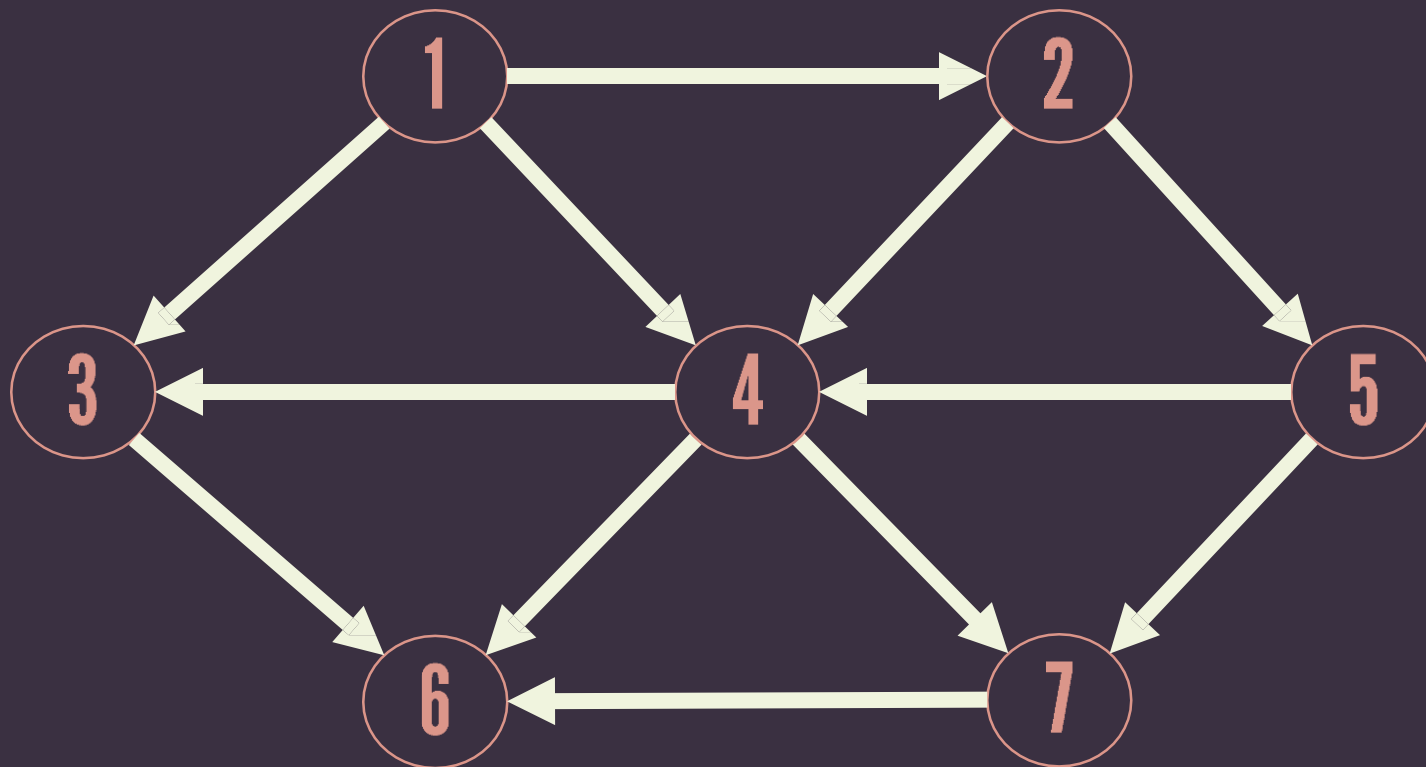


directed  
acyclic graph  
**DAG**

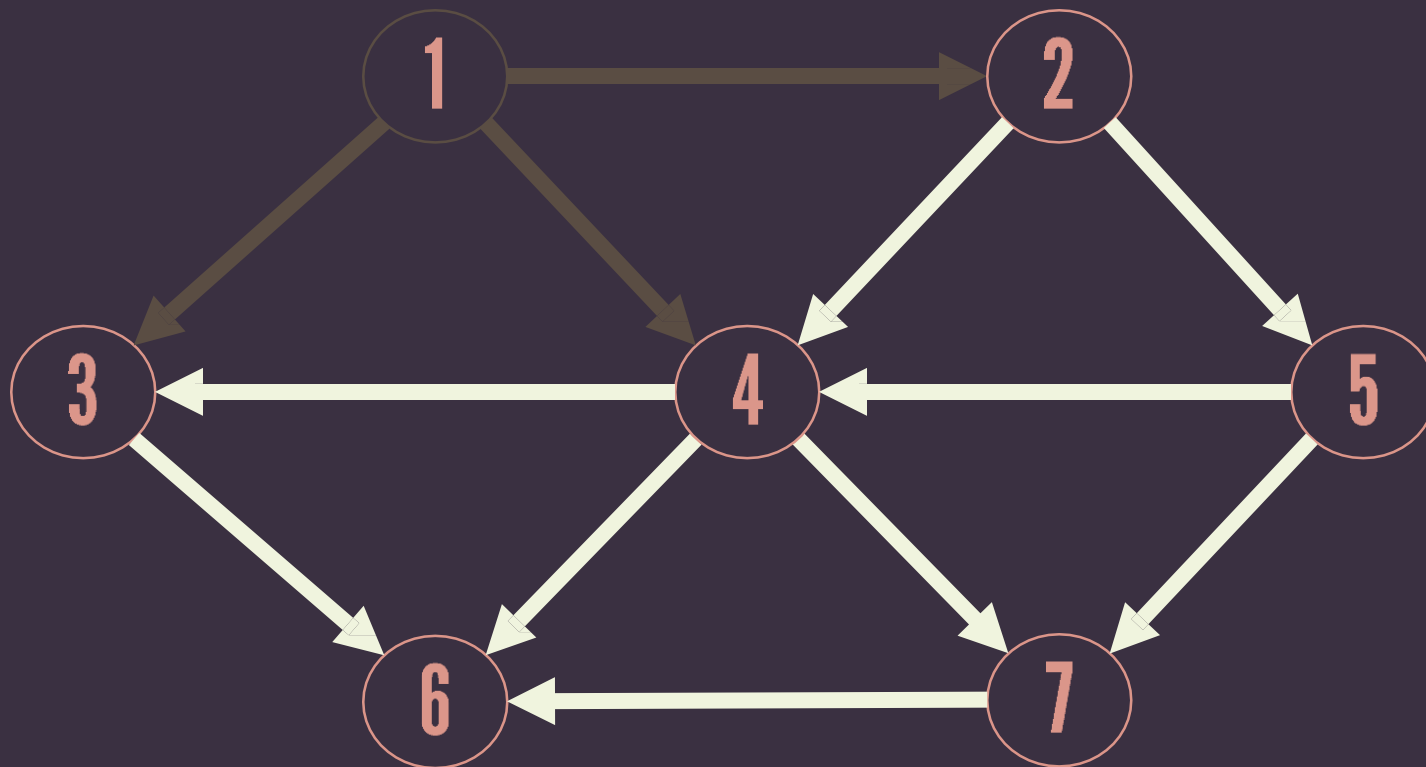


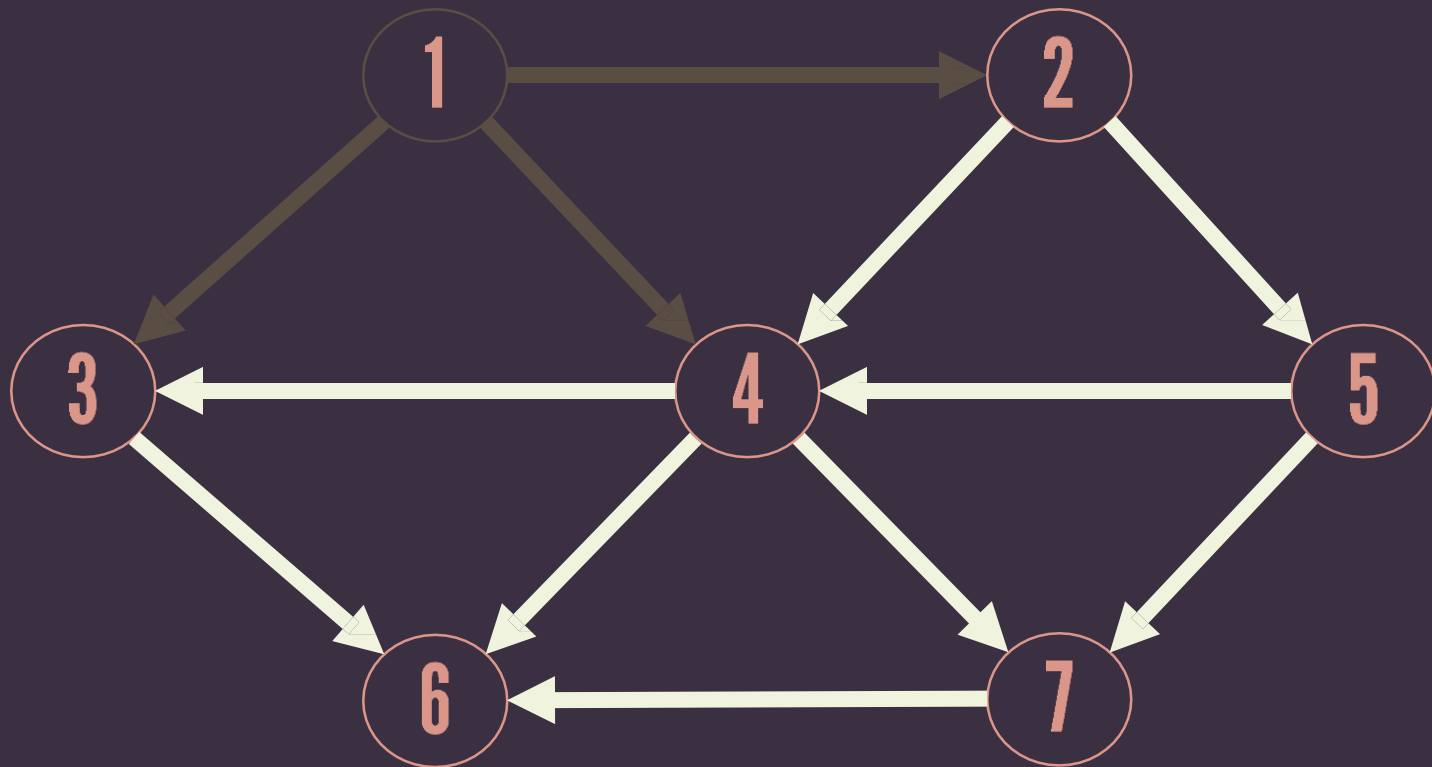
## Algorithm:

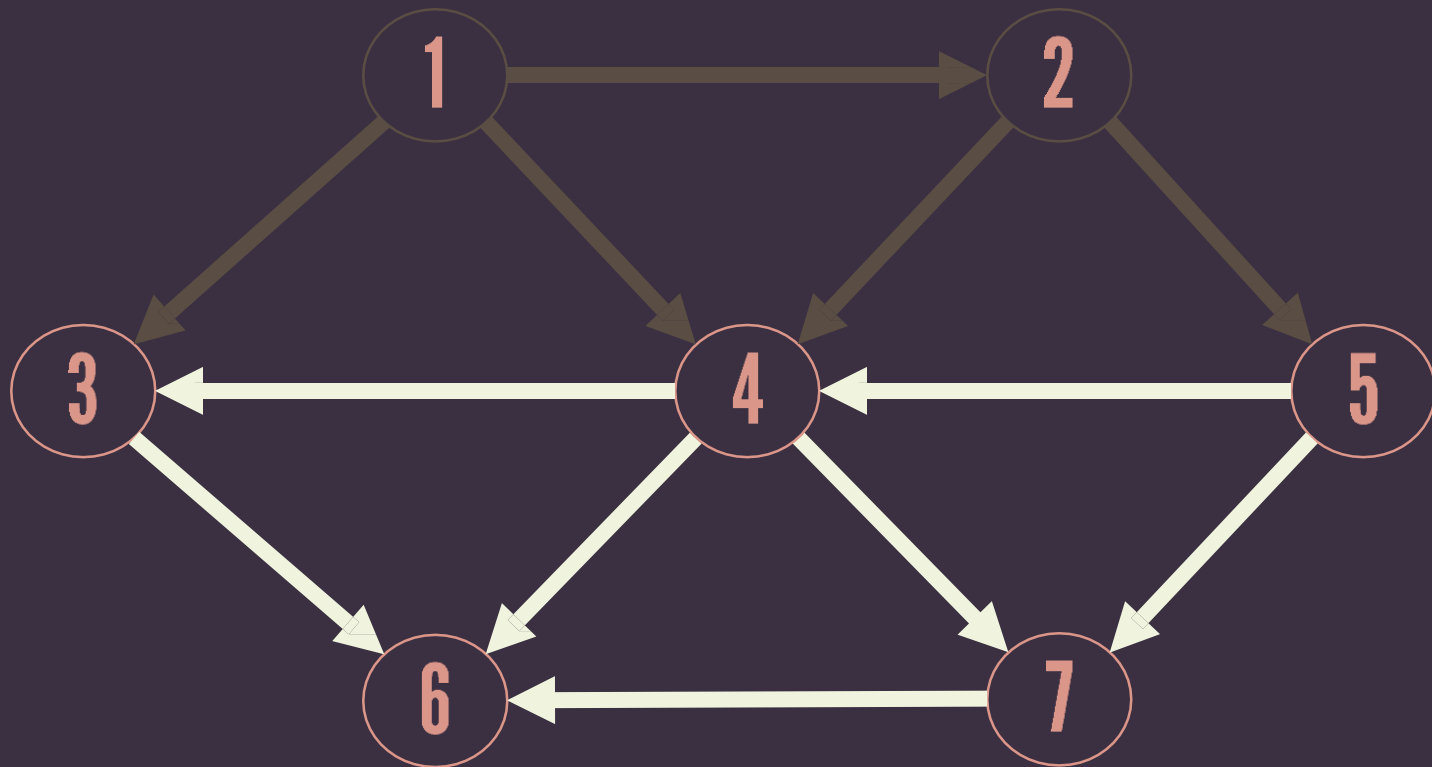
- Find any vertex with no incoming edges (in-degree is 0).
- Print this vertex and remove it along with its edges from the graph.
- Repeat steps above.

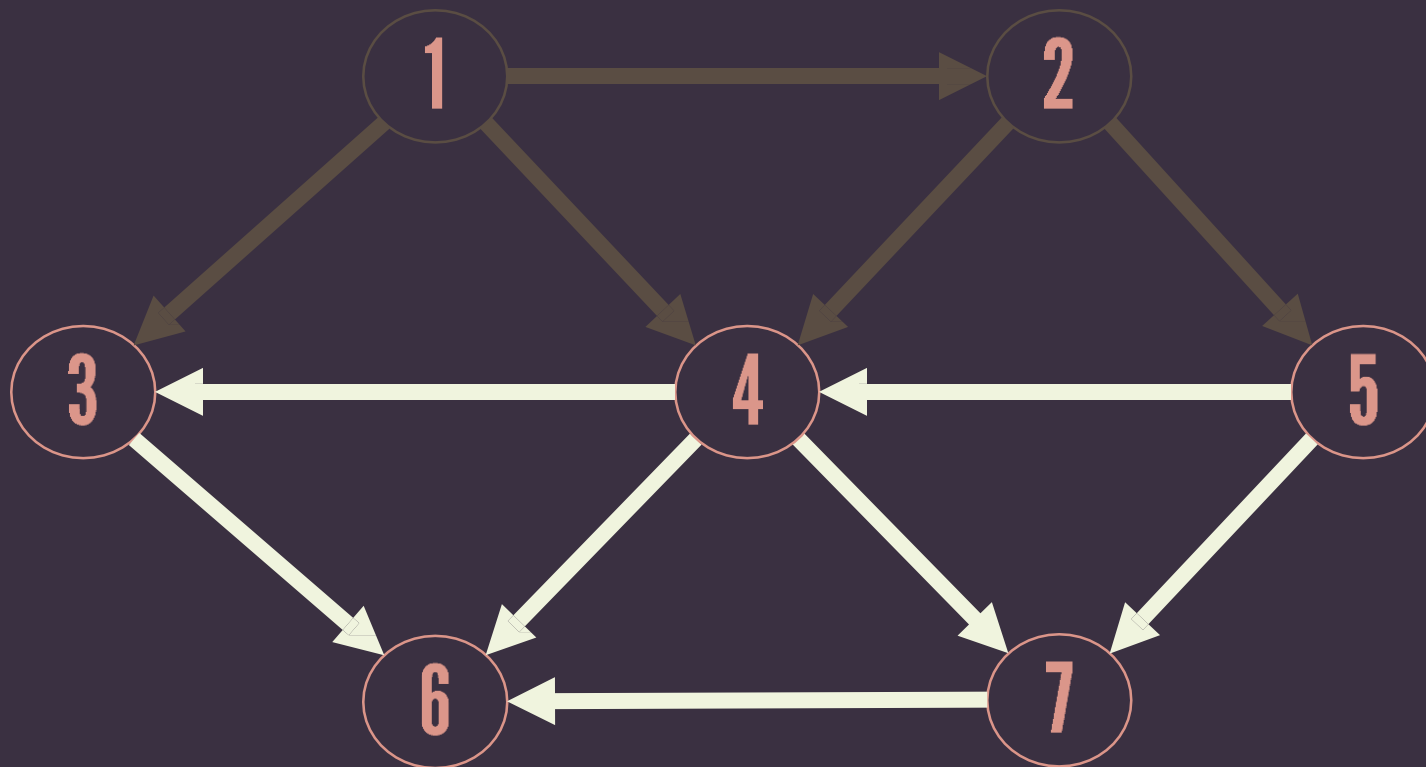




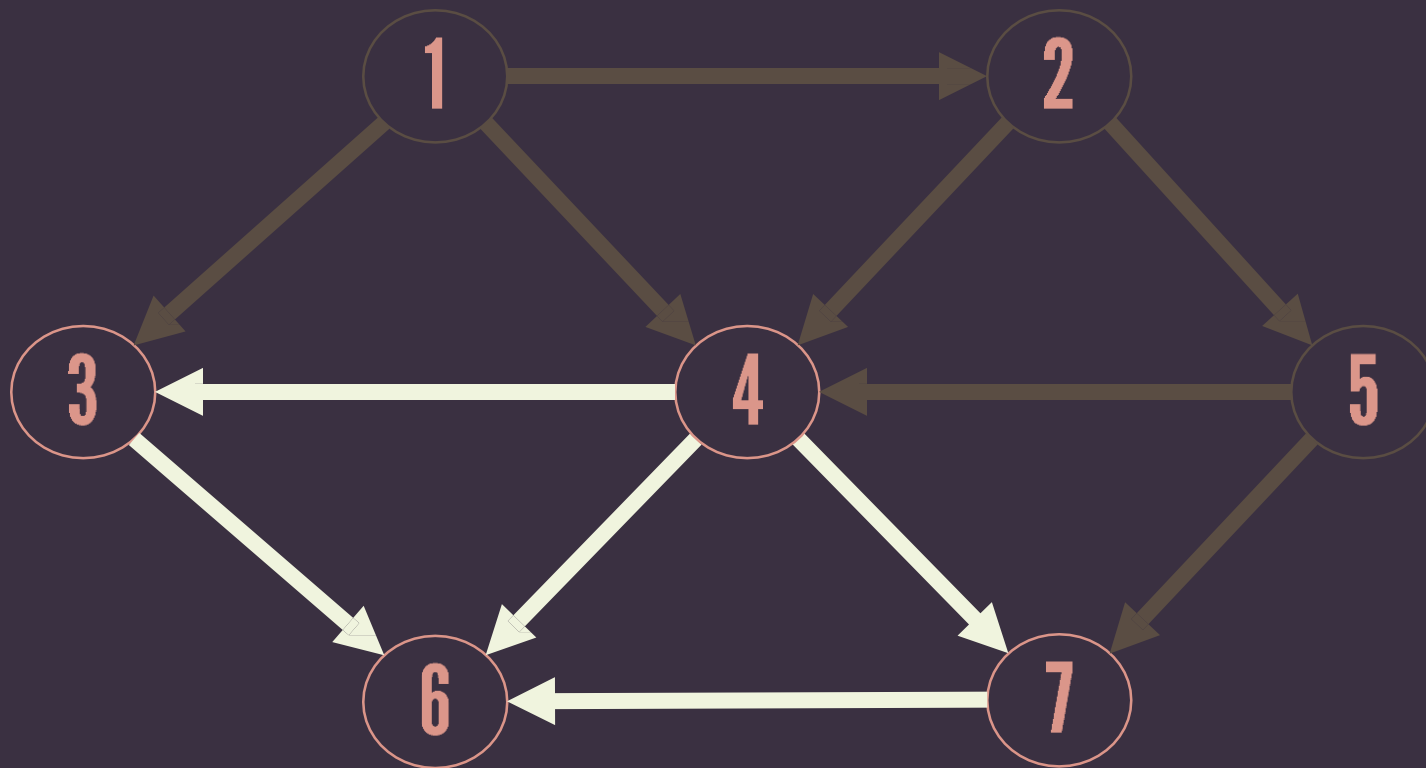




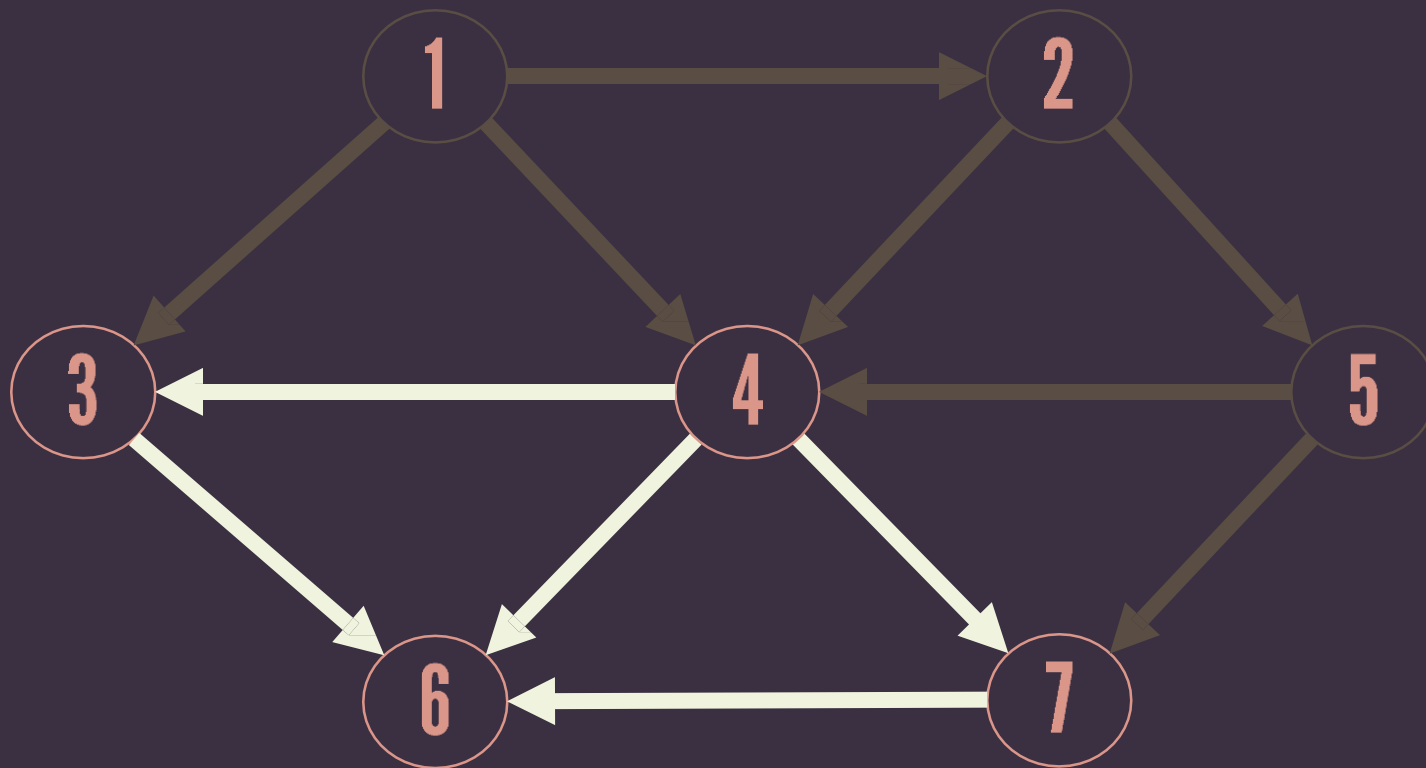




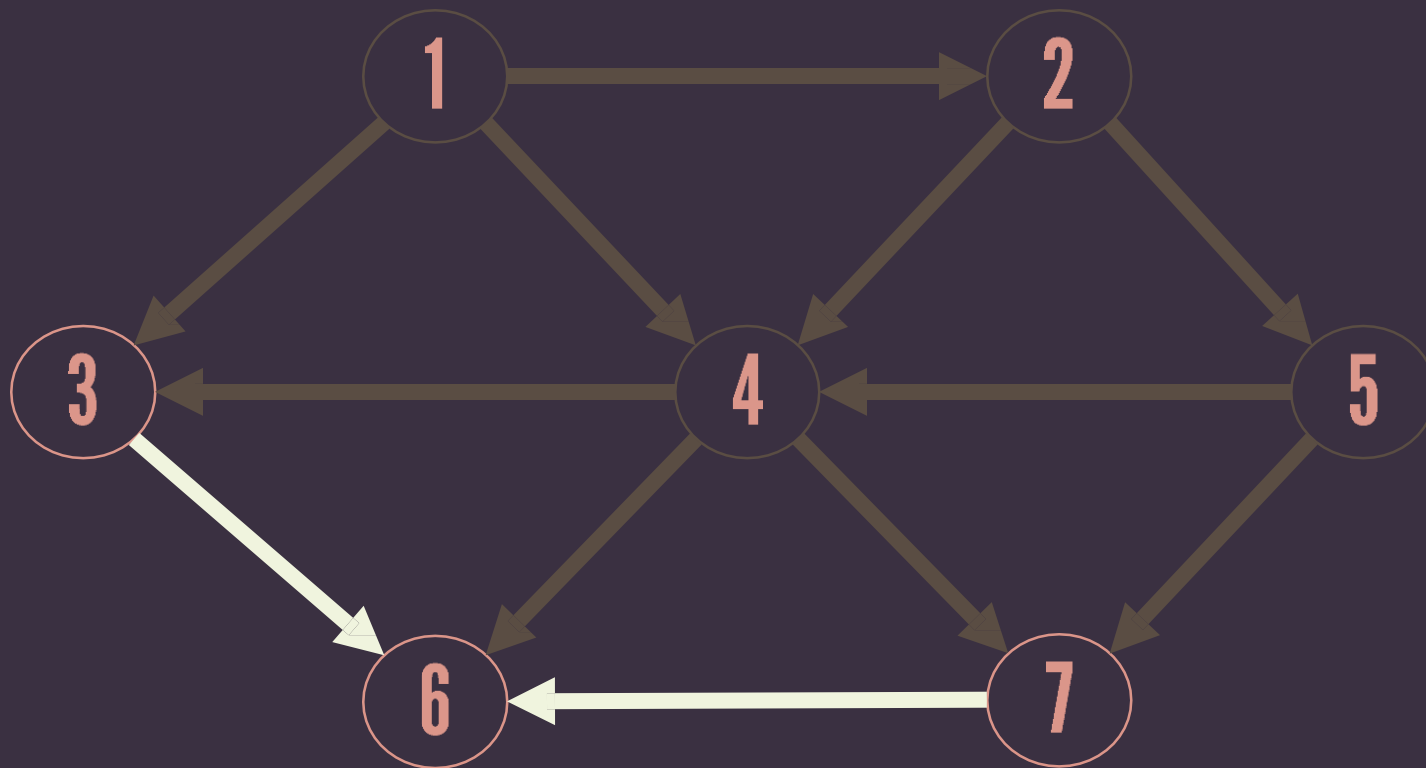
1 2 5



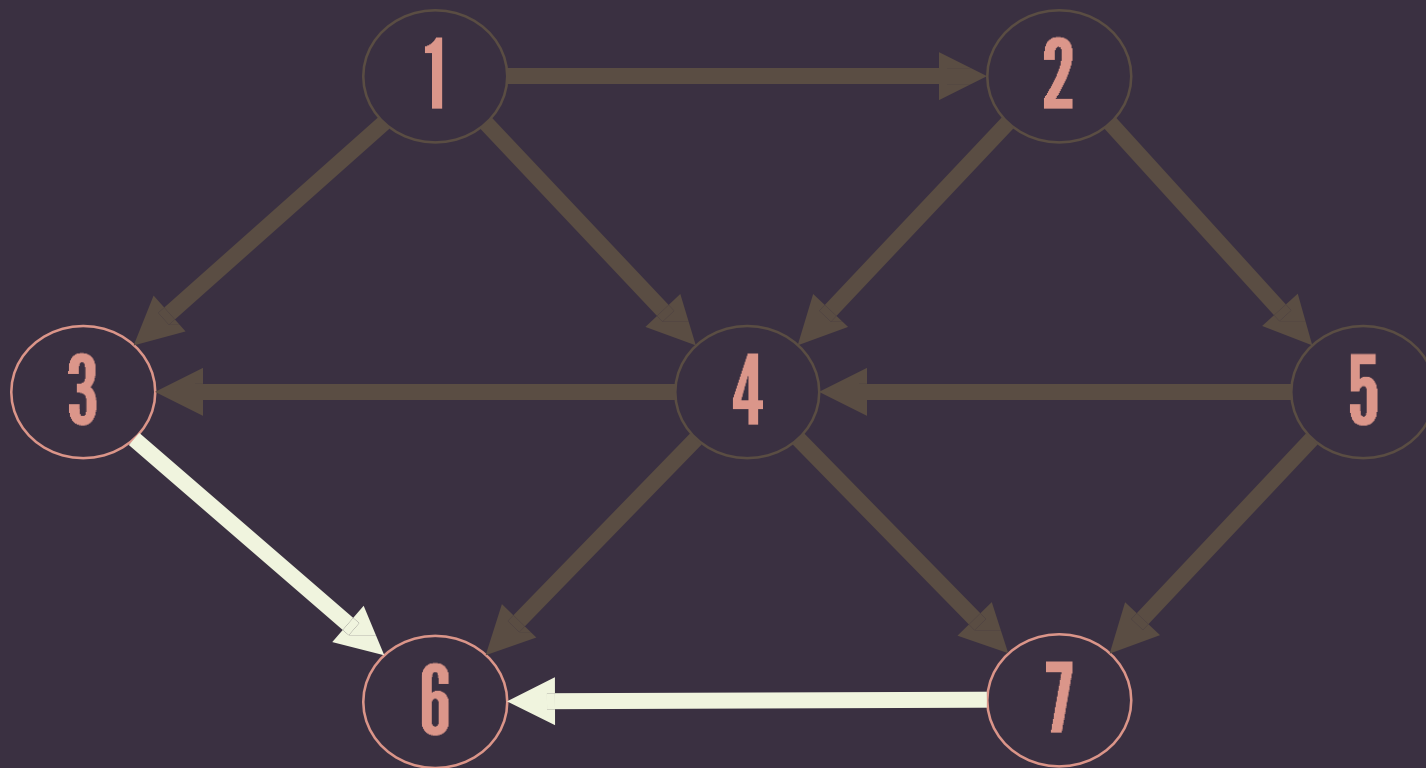
1 2 5



1 2 5 4

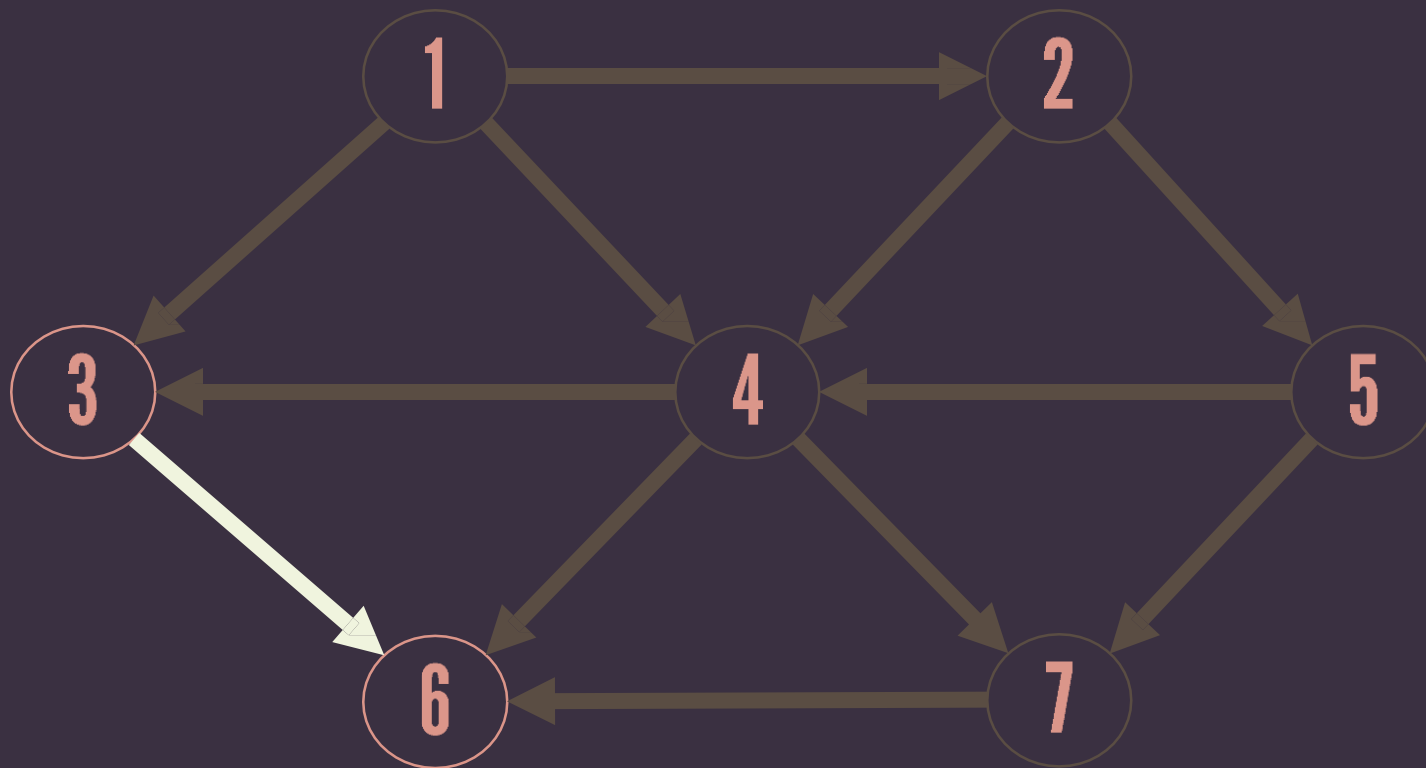


1 2 5 4

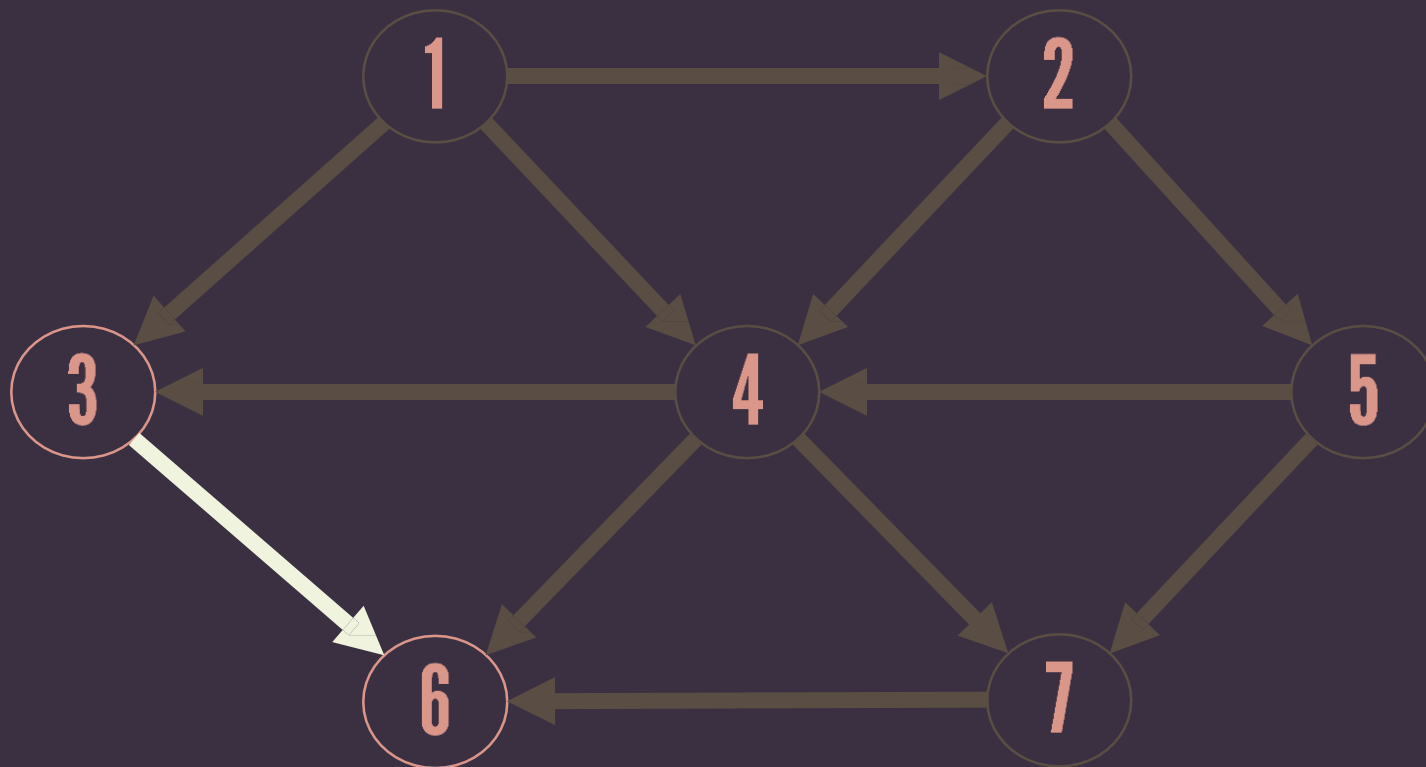


1 2 5 4 7

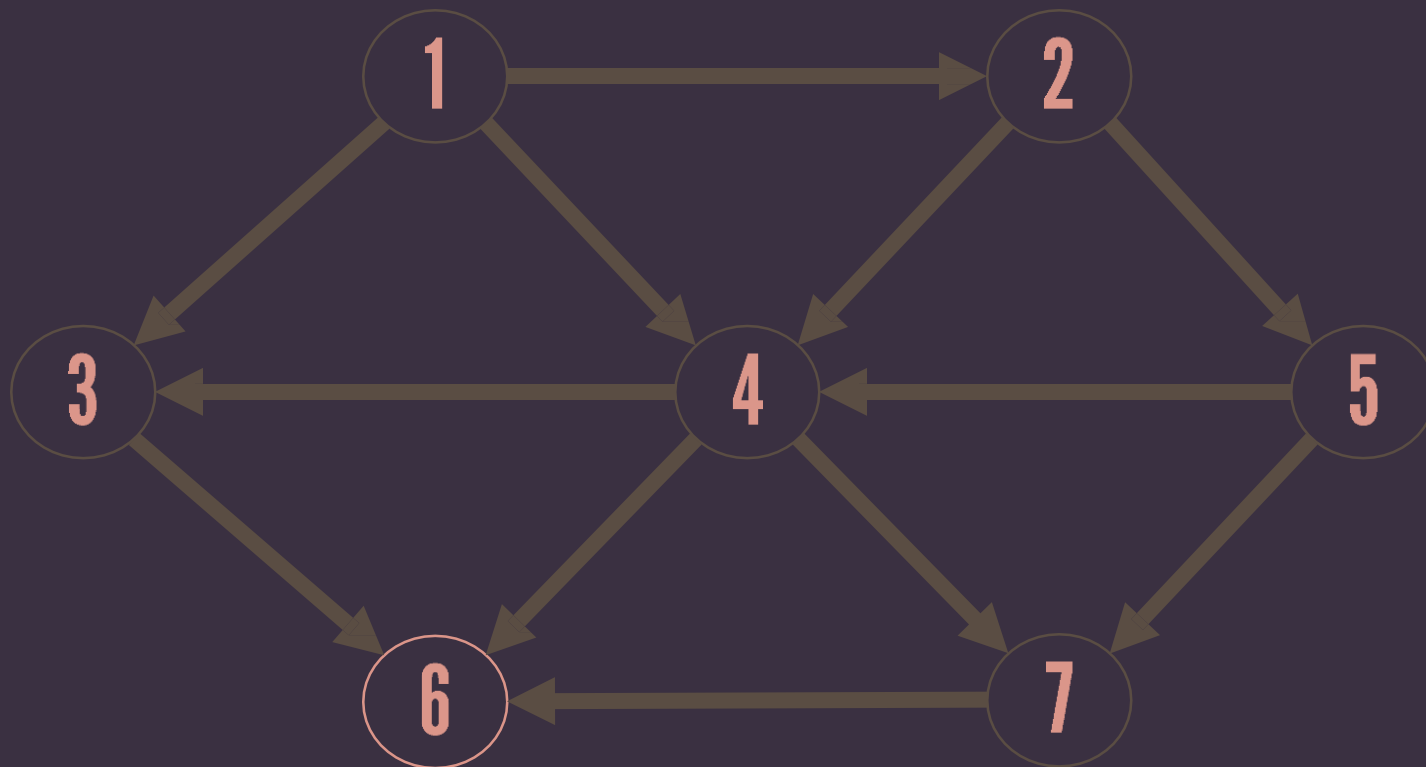




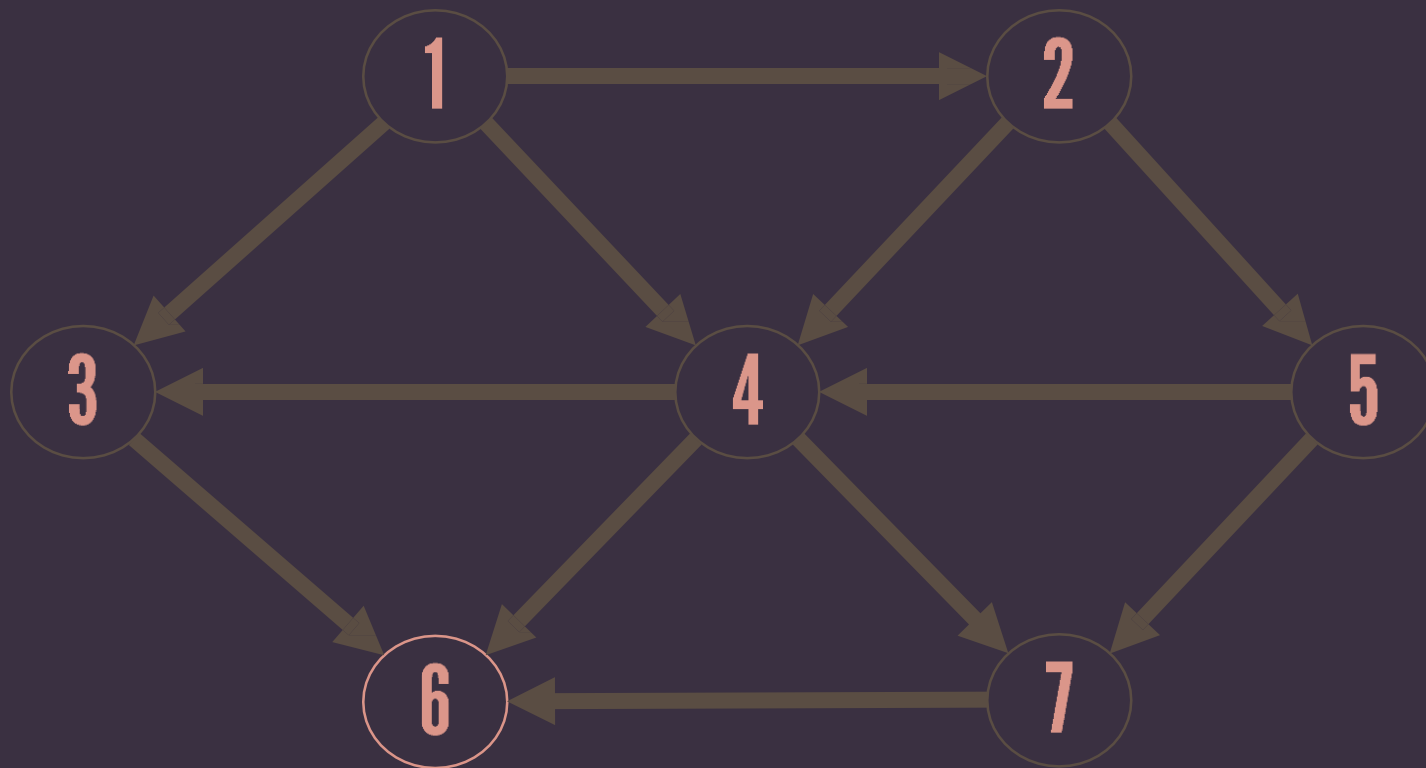
1 2 5 4 7



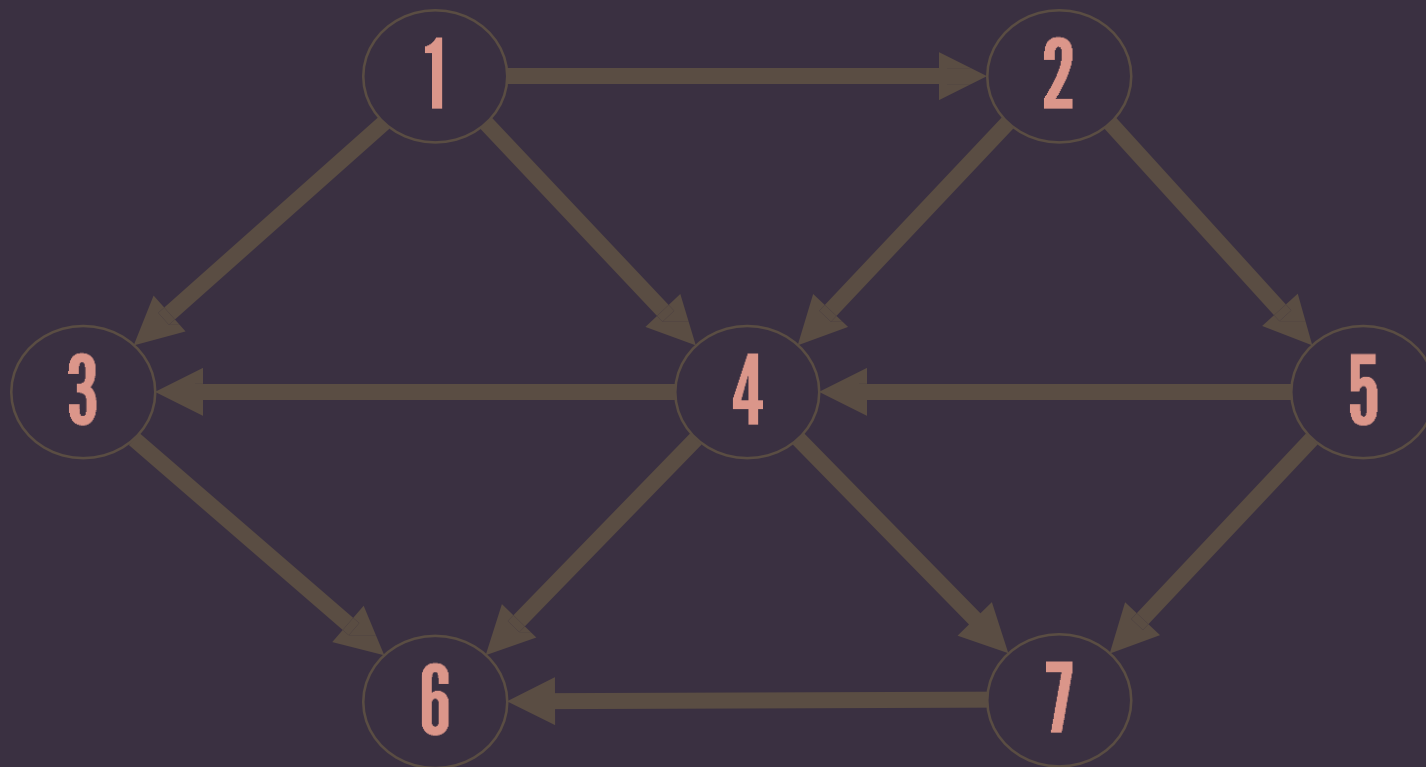
1 2 5 4 7 3



1 2 5 4 7 3



1 2 5 4 7 3 6



1 2 5 4 7 3 6

graph

# TRAVERSALS

**DEPTH**

**FIRST**

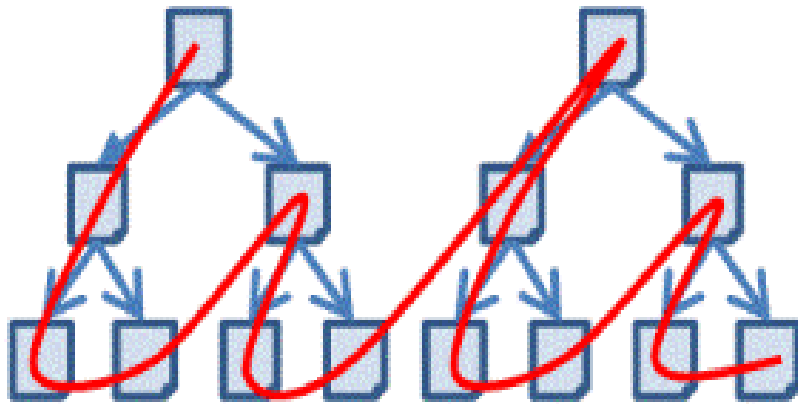
search

**BREADTH**

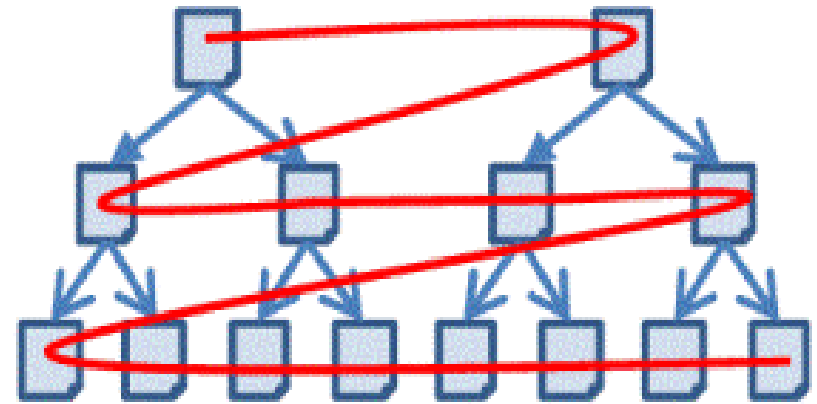
**FIRST**

search





a)



b)

# DEPTH FIRST search

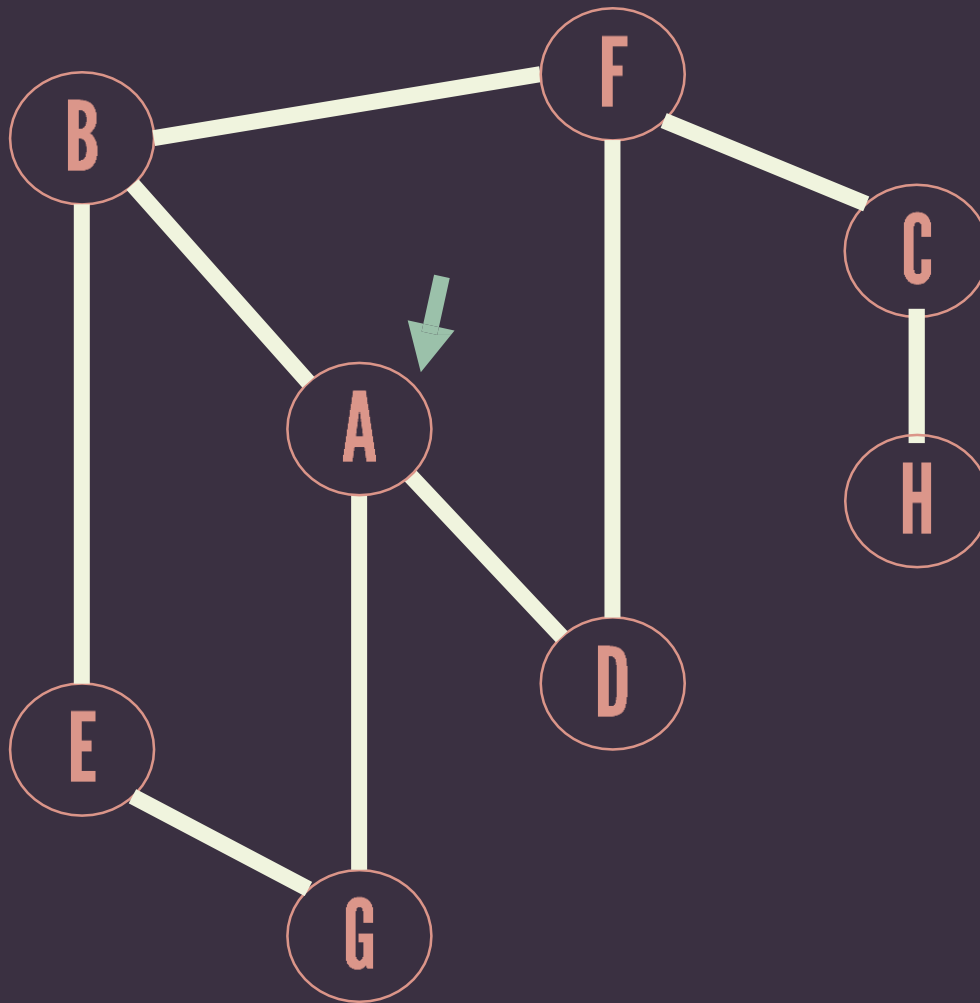
Generalization of preorder traversal.

# DEPTH FIRST search

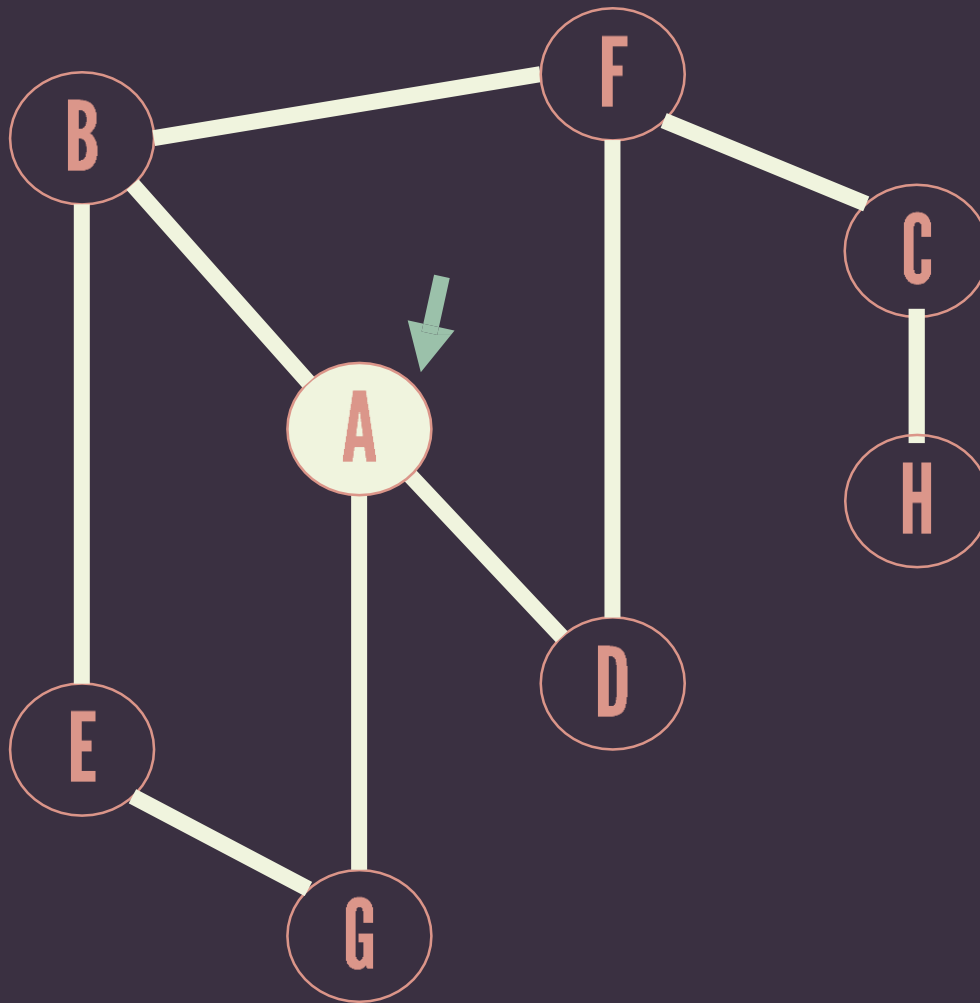
Starting at some vertex  $v$ , process  $v$  and recursively traverse all vertices adjacent to  $v$ .

```
void DFS( vertex v, graph G ){  
    print v;  
    visited[v] = TRUE;  
    for each w adjacent to v:  
        if( !visited[w] )  
            DFS(w);  
}
```

```
void DFS( vertex v, graph G ){  
    stack S;  
    push(v,S);  
    while stack S is not empty{  
        v = pop(S);  
        if ( !visited[v] ){ print v;  
            visited[v] = TRUE;  
            for each w adjacent to v:  
                if( !visited[w] )  
                    push(w, S);  
        }  
    }  
}
```

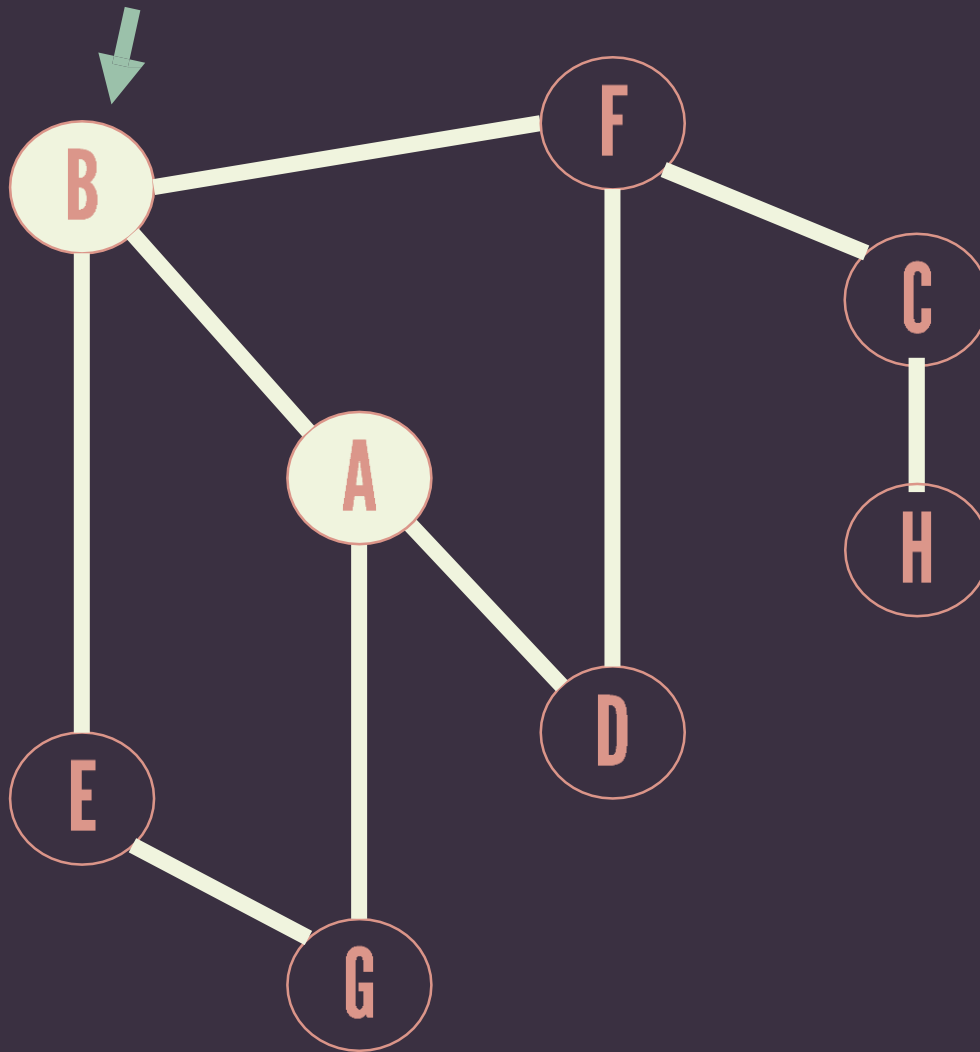



Result:



A

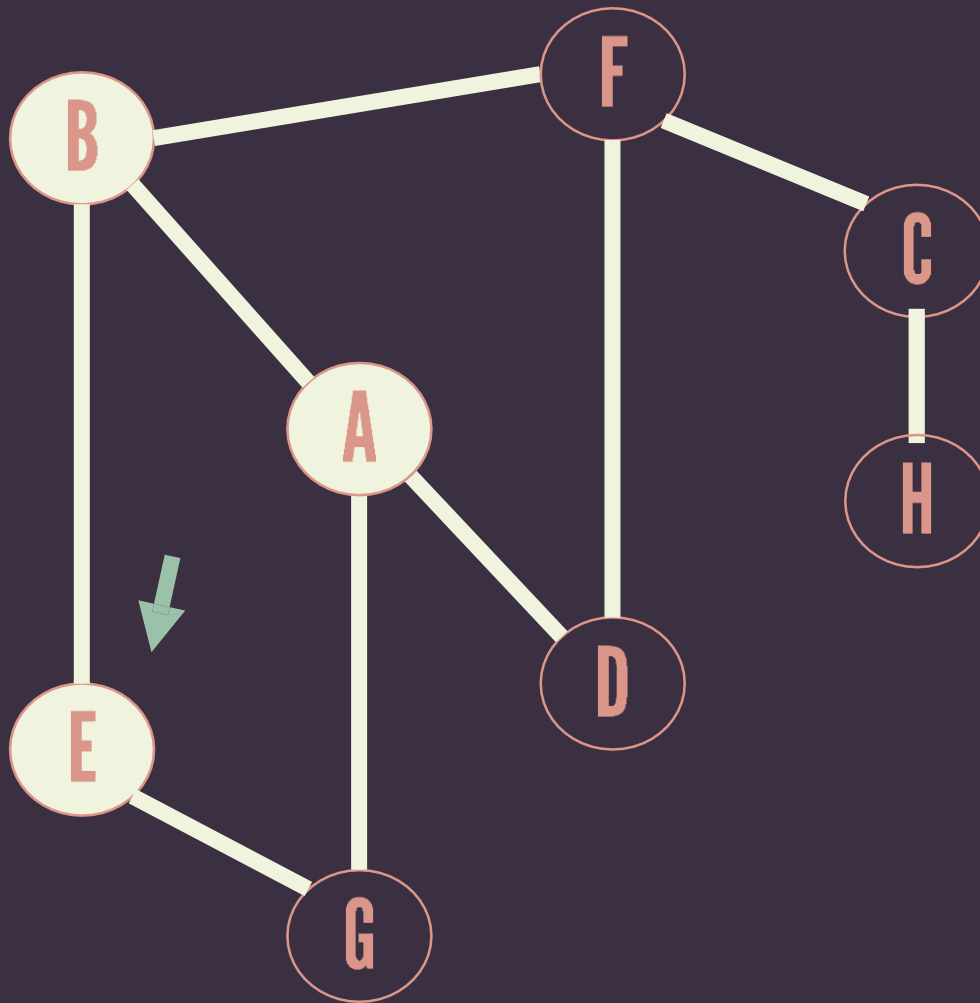
Result: A



B
A

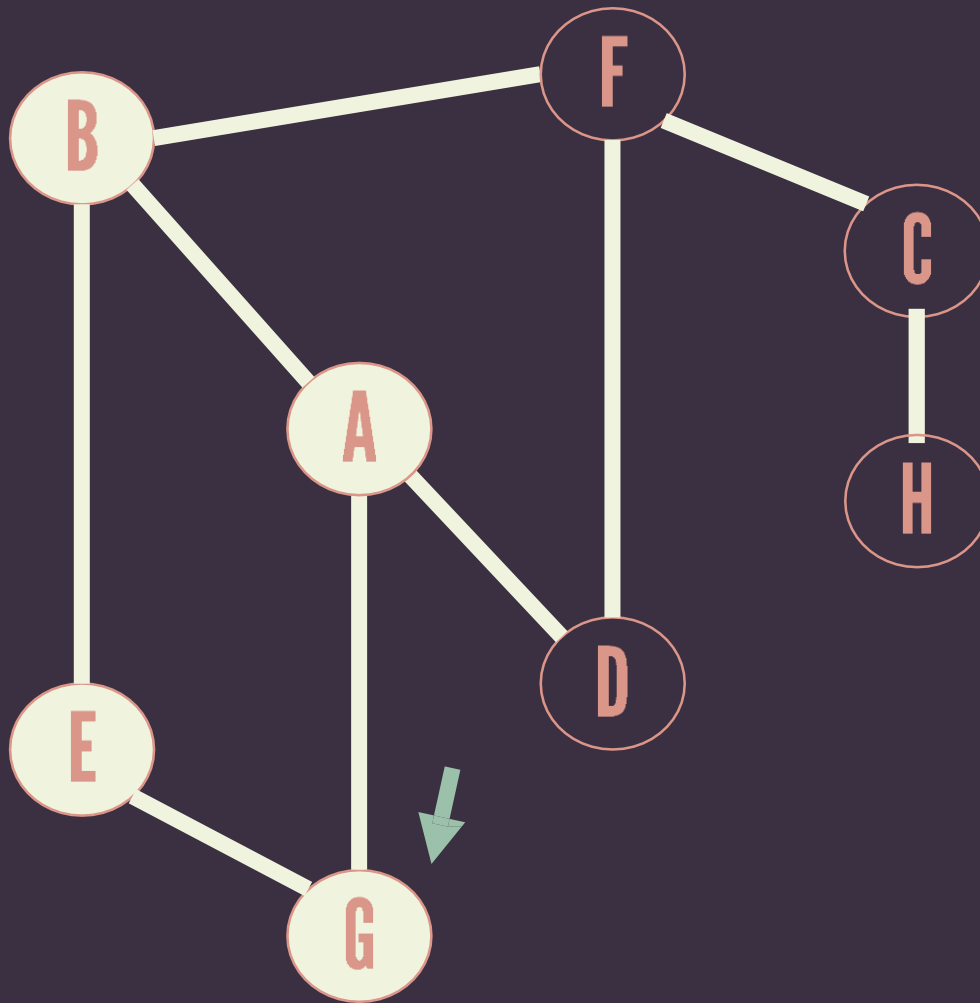
Result: AB





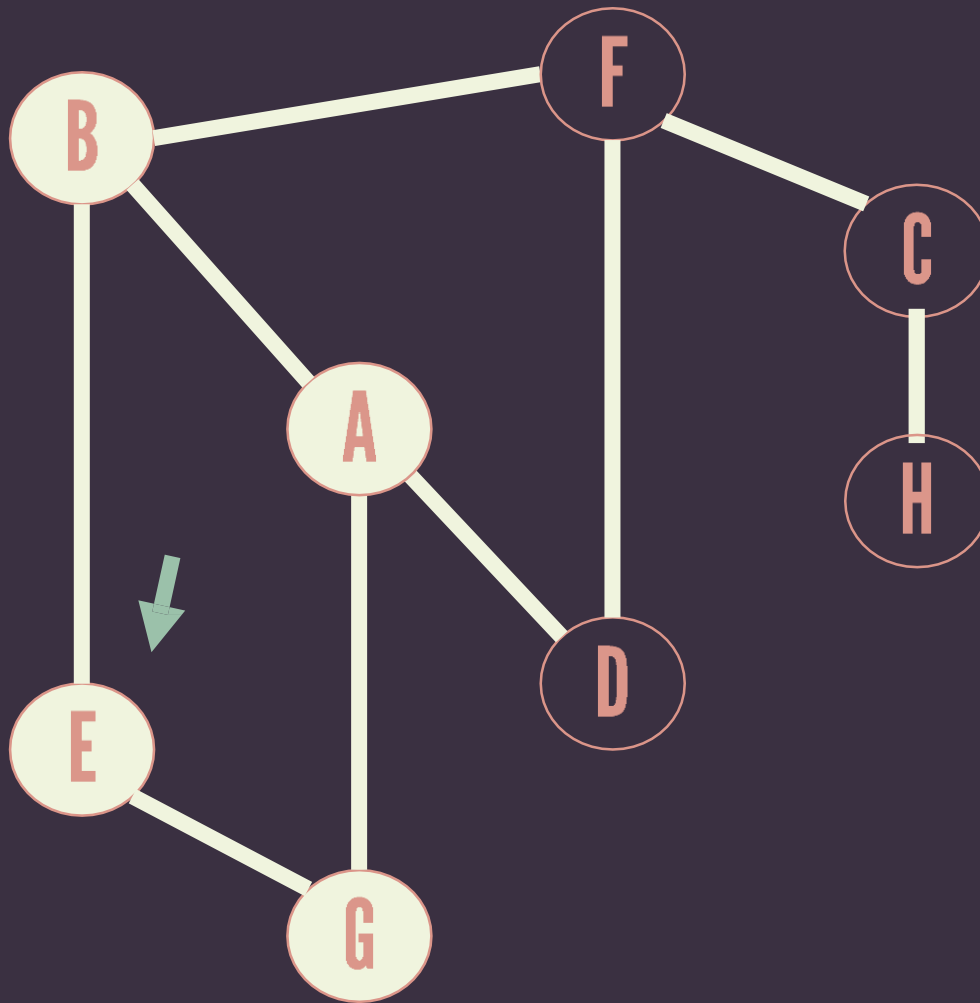
E
B
A

Result: A B E



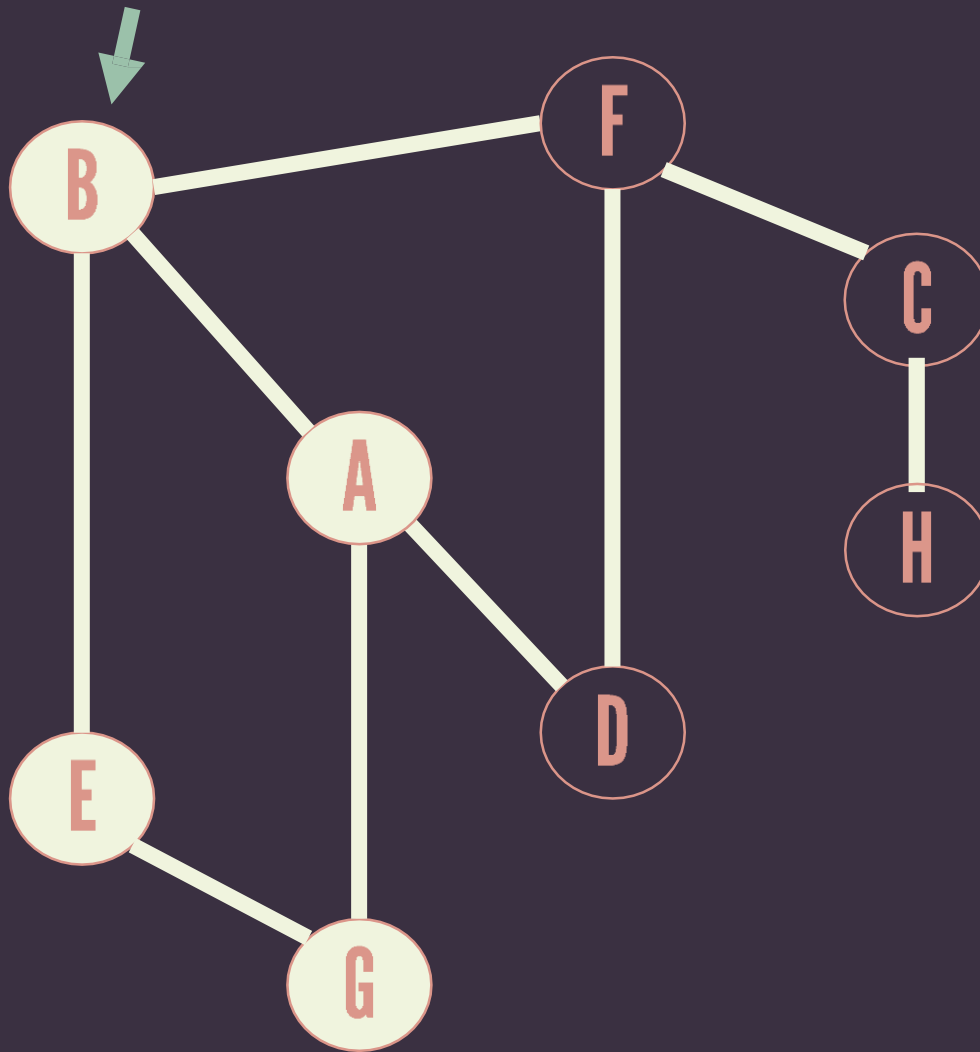
G
E
B
A

Result: A B E G



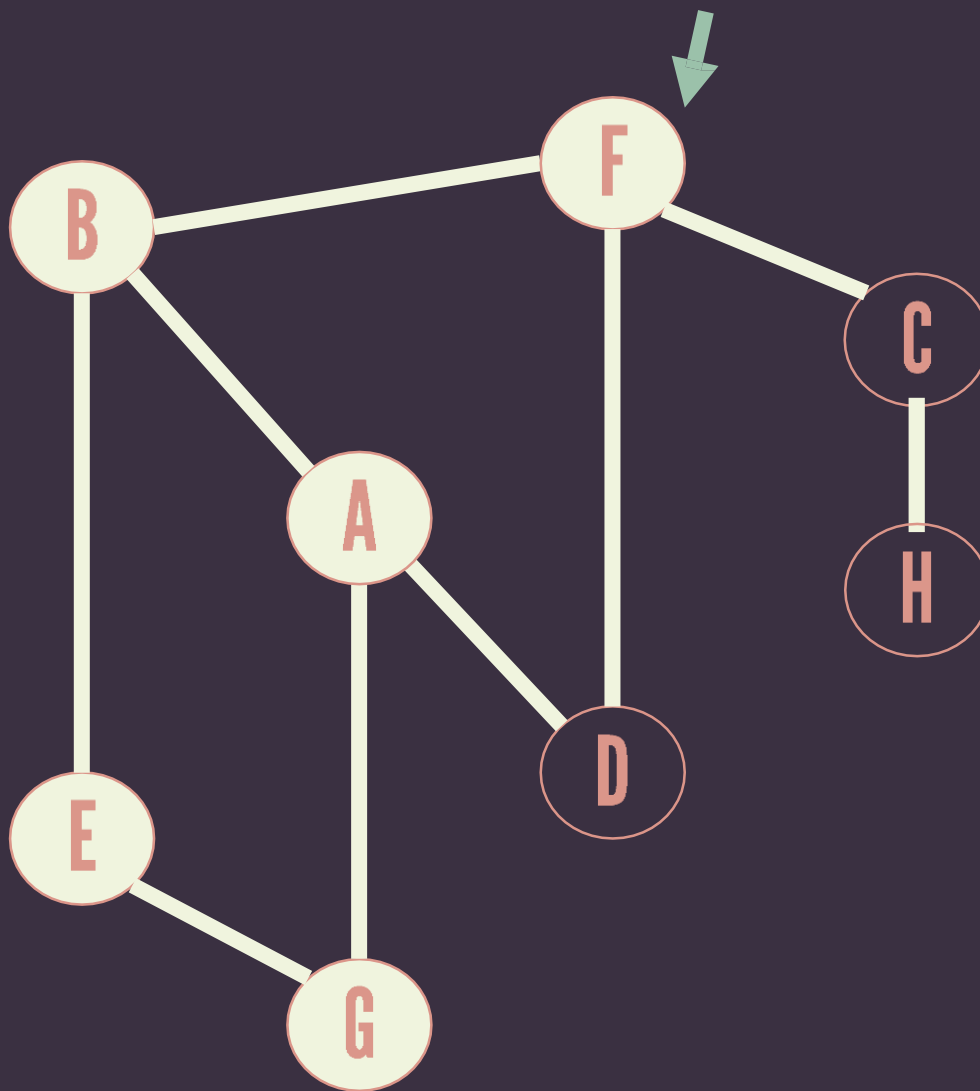
E
B
A

Result: A B E G



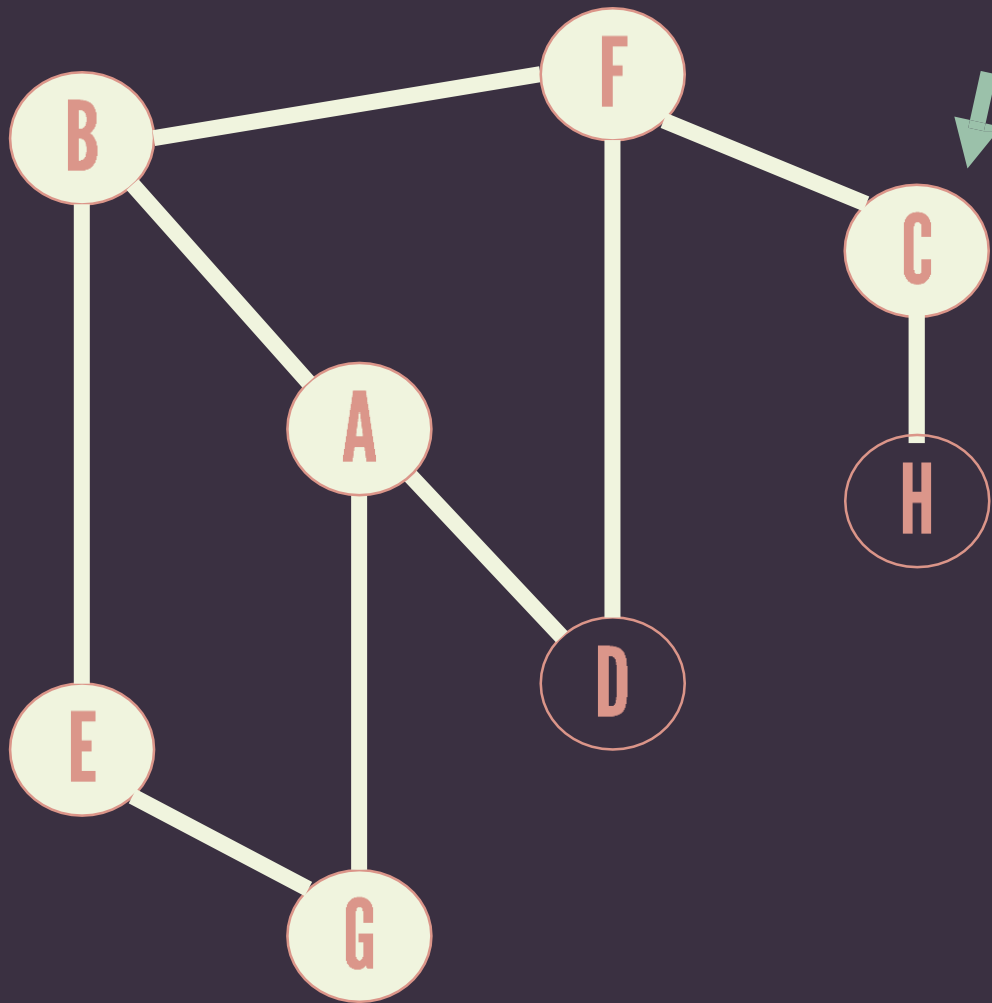
B
A

Result: A B E G



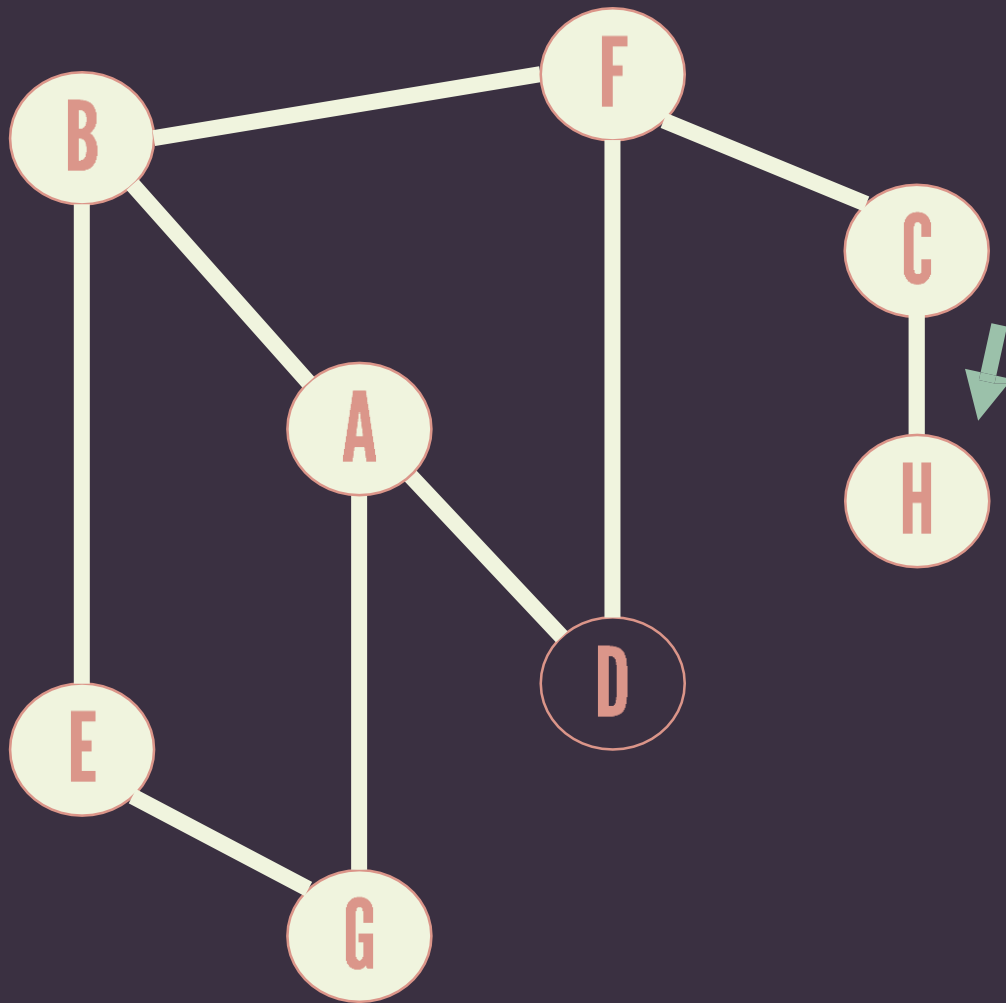
F
B
A

Result: A B E G F



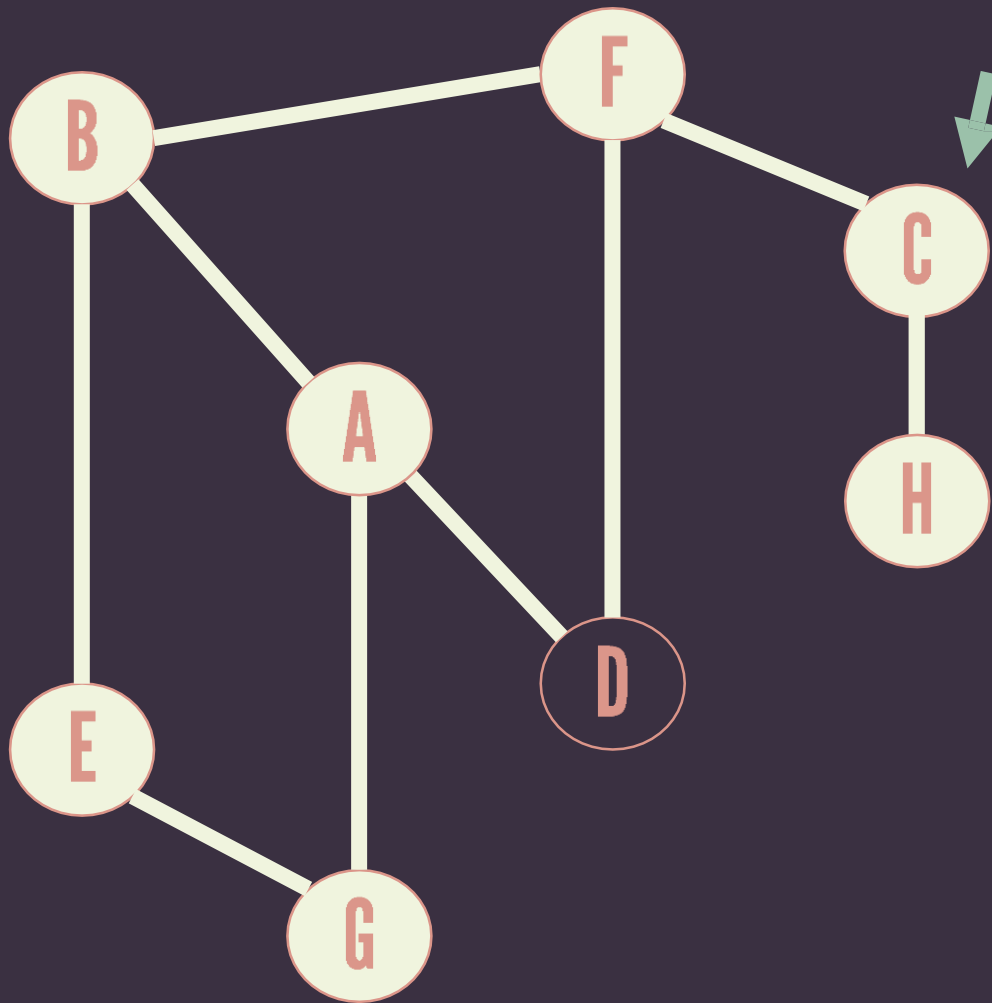
C
F
B
A

Result: A B E G F C



H
C
F
B
A

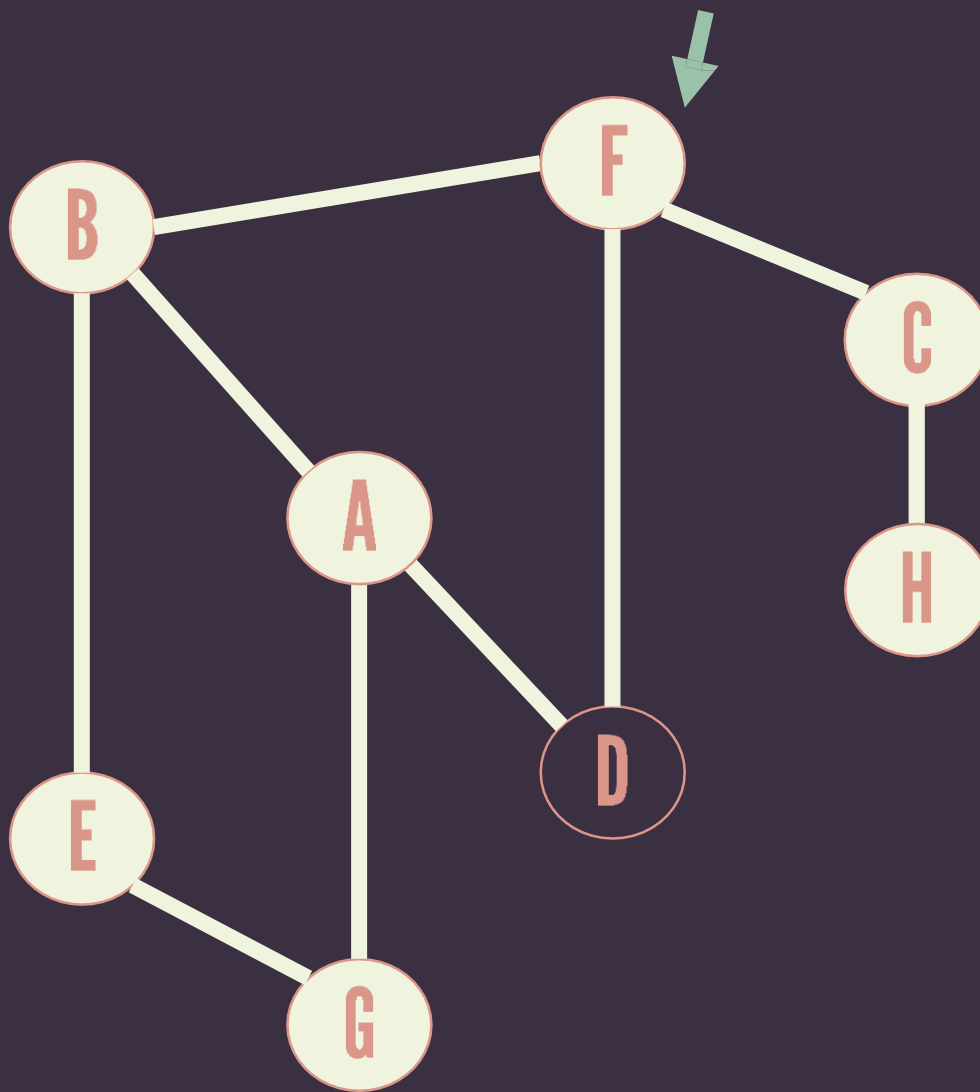
Result: A B E G F C H



C
F
B
A

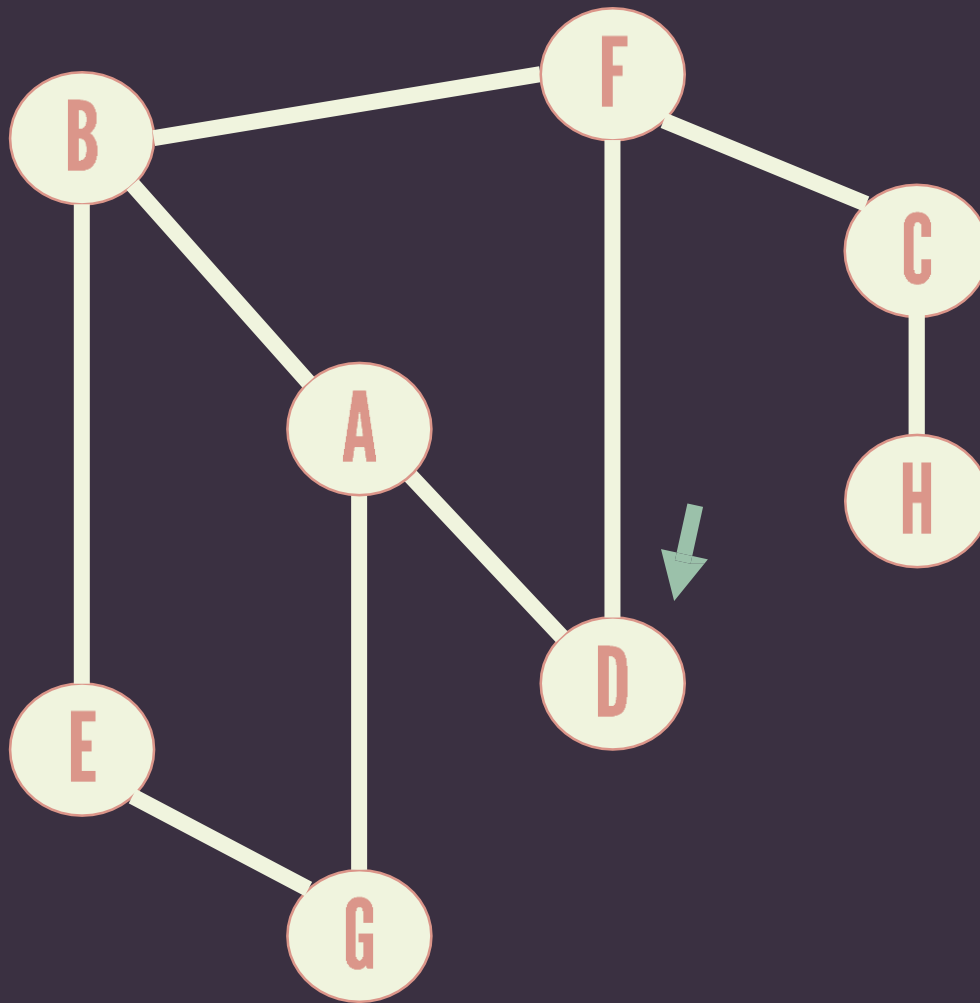
Result: A B E G F C H





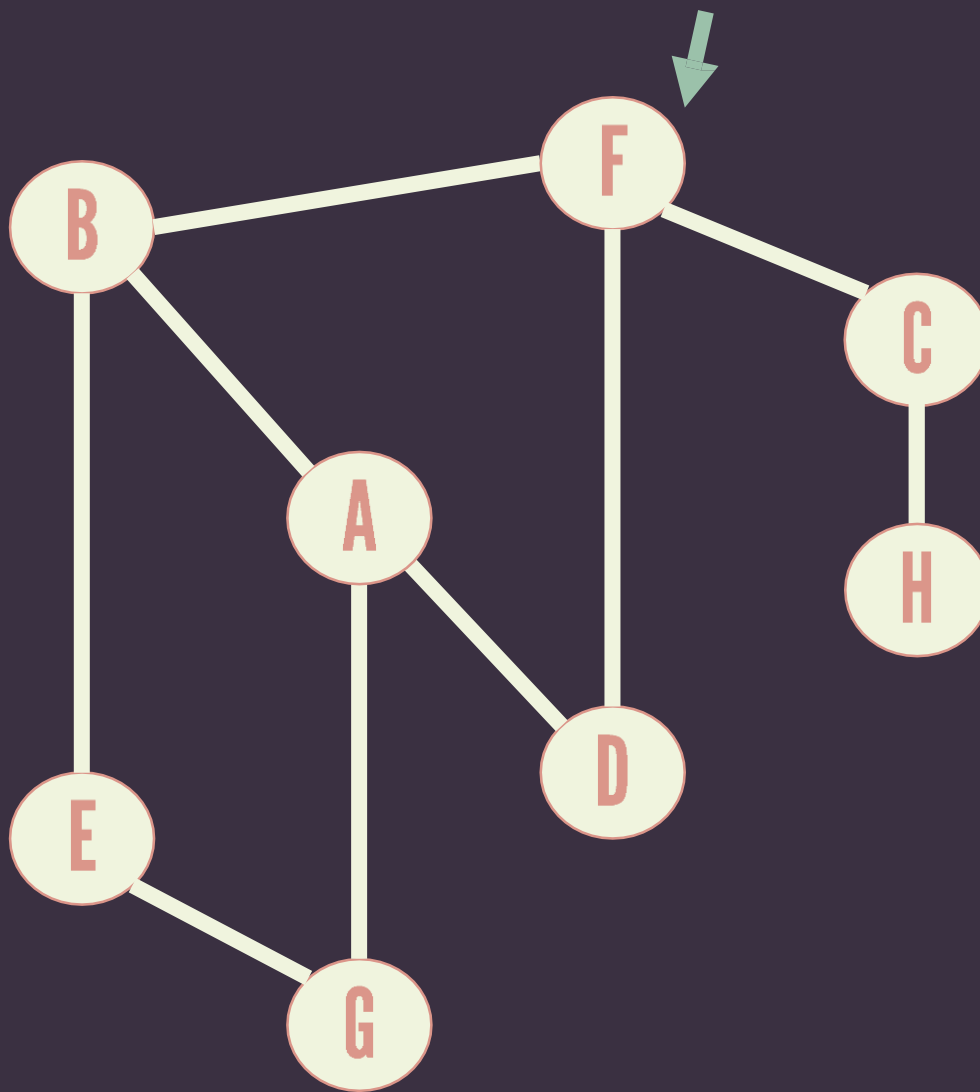
F
B
A

Result: A B E G F C H



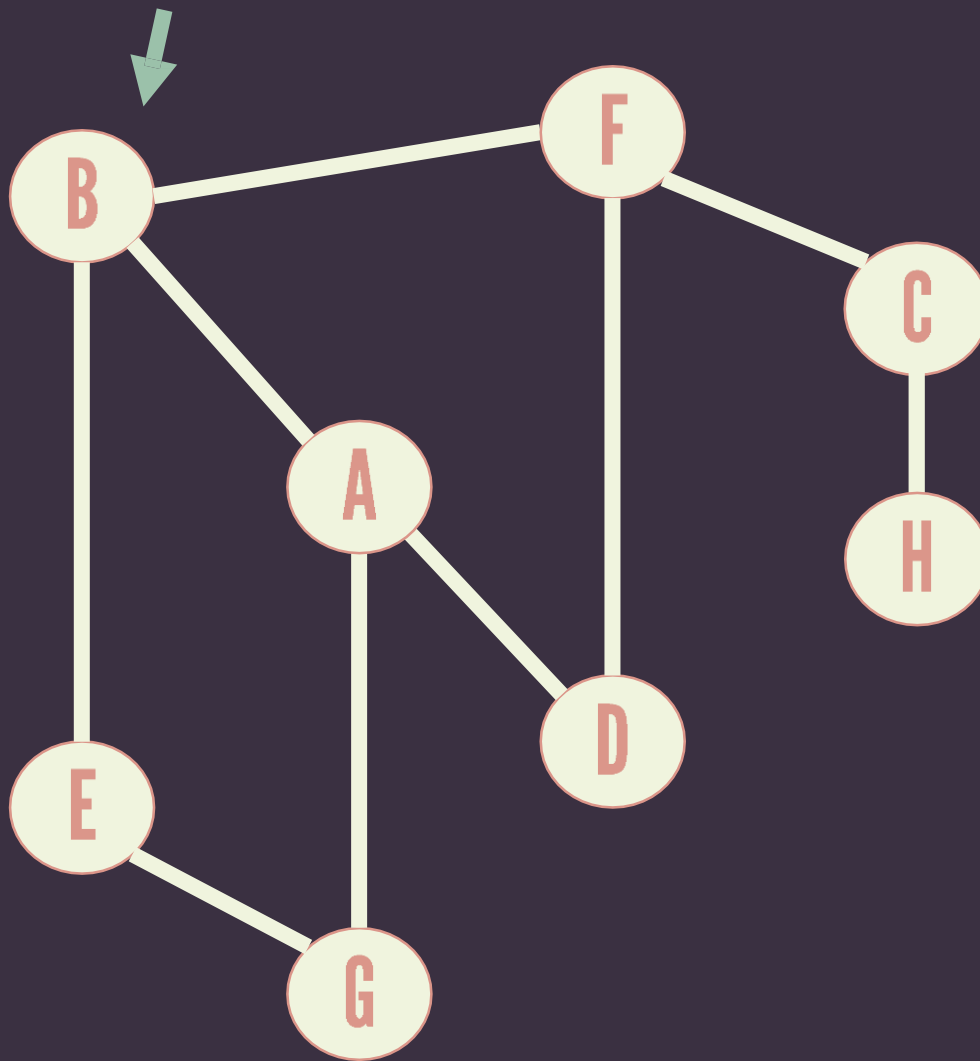
D
F
B
A

Result: A B E G F C H D



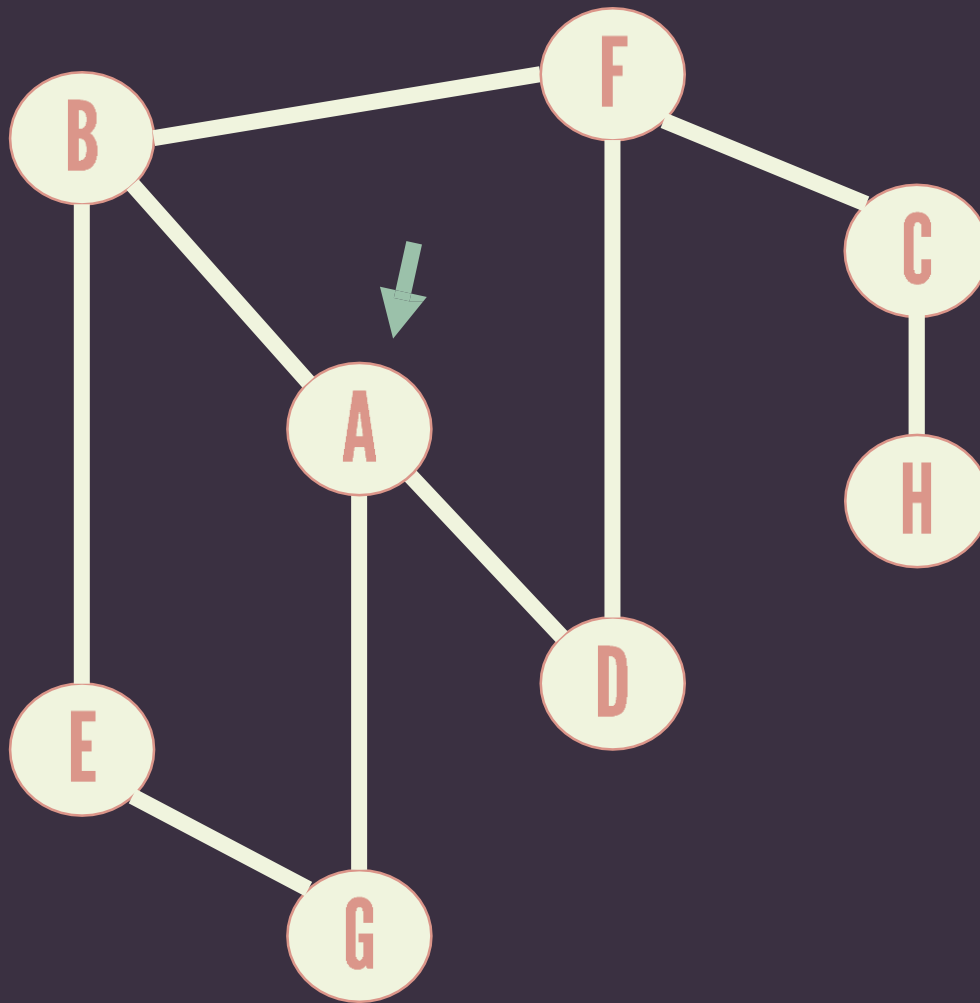
F
B
A

Result: A B E G F C H D



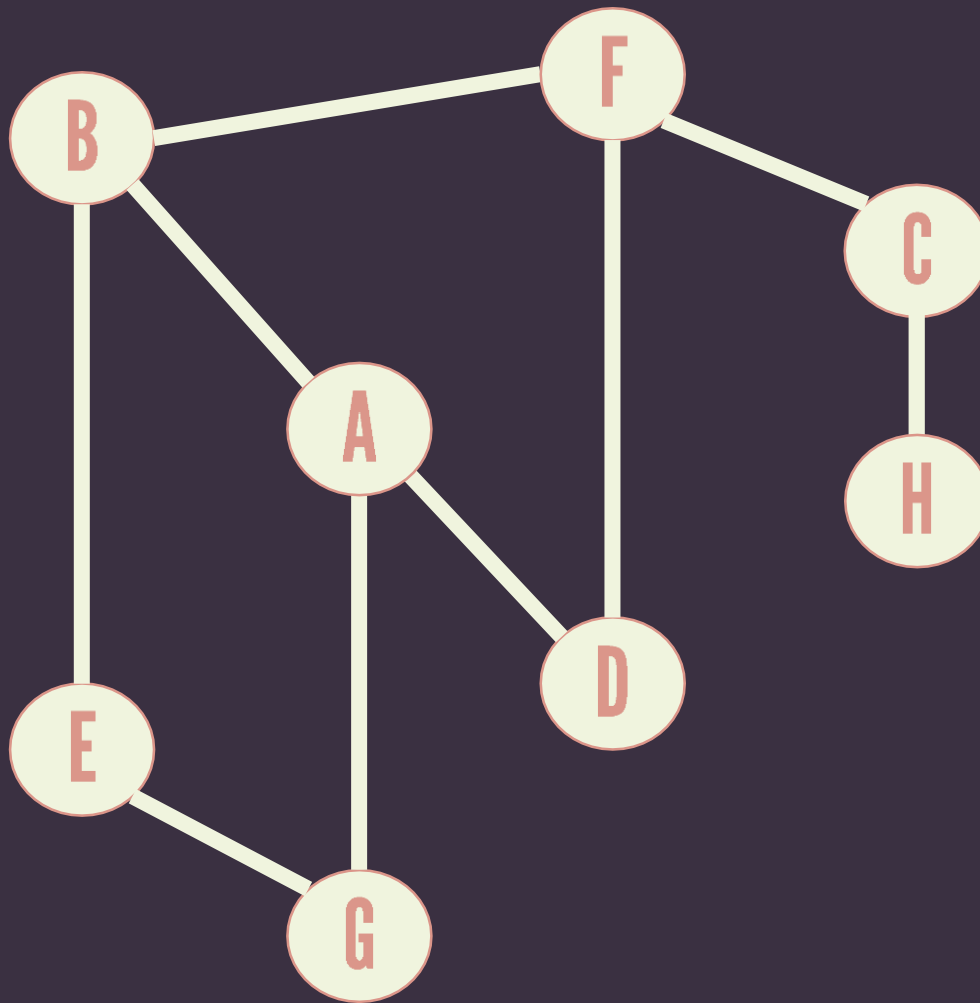
B
A

Result: A B E G F C H D



A

Result: A B E G F C H D




Result: A B E G F C H D

# DFS

## applications

Detecting cycle in a graph.

Topological sorting.

Path finding.

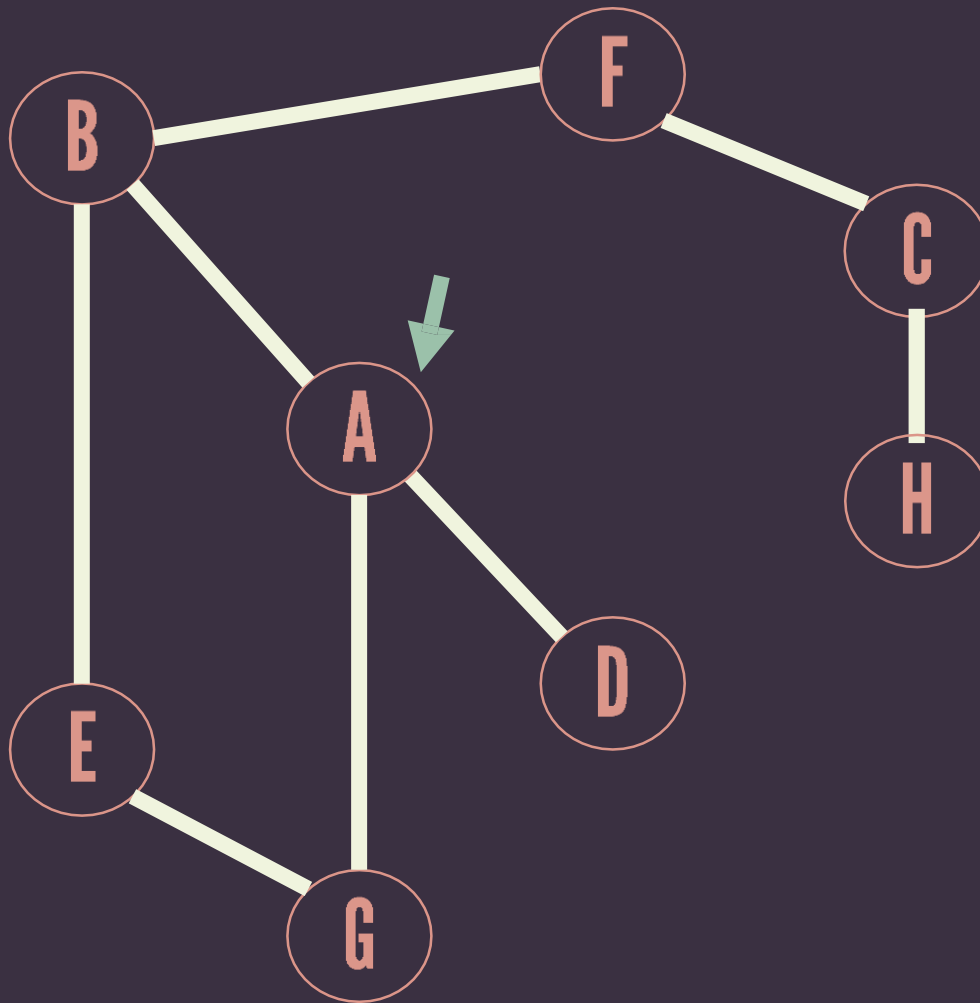
Spanning trees.

# DFS

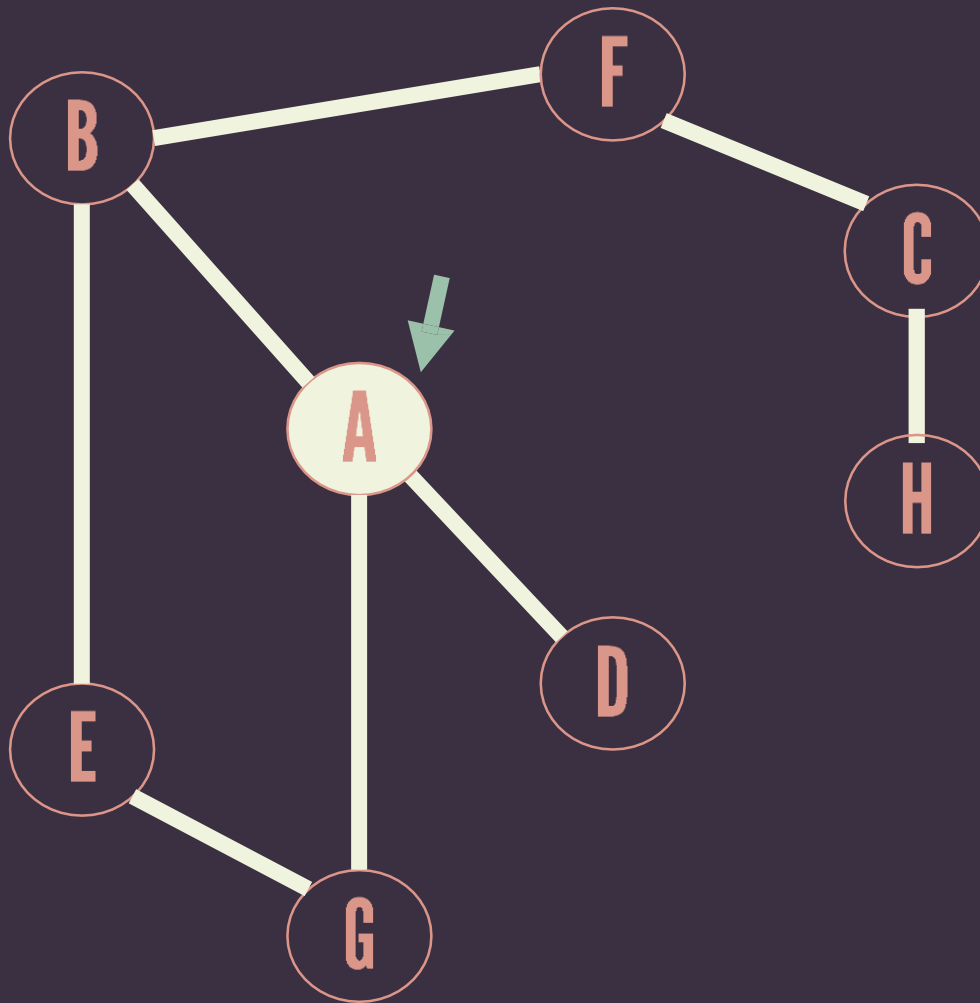
detecting  
cycle in a  
graph

A graph has a cycle if  
and only if we can find a  
back edge during DFS.



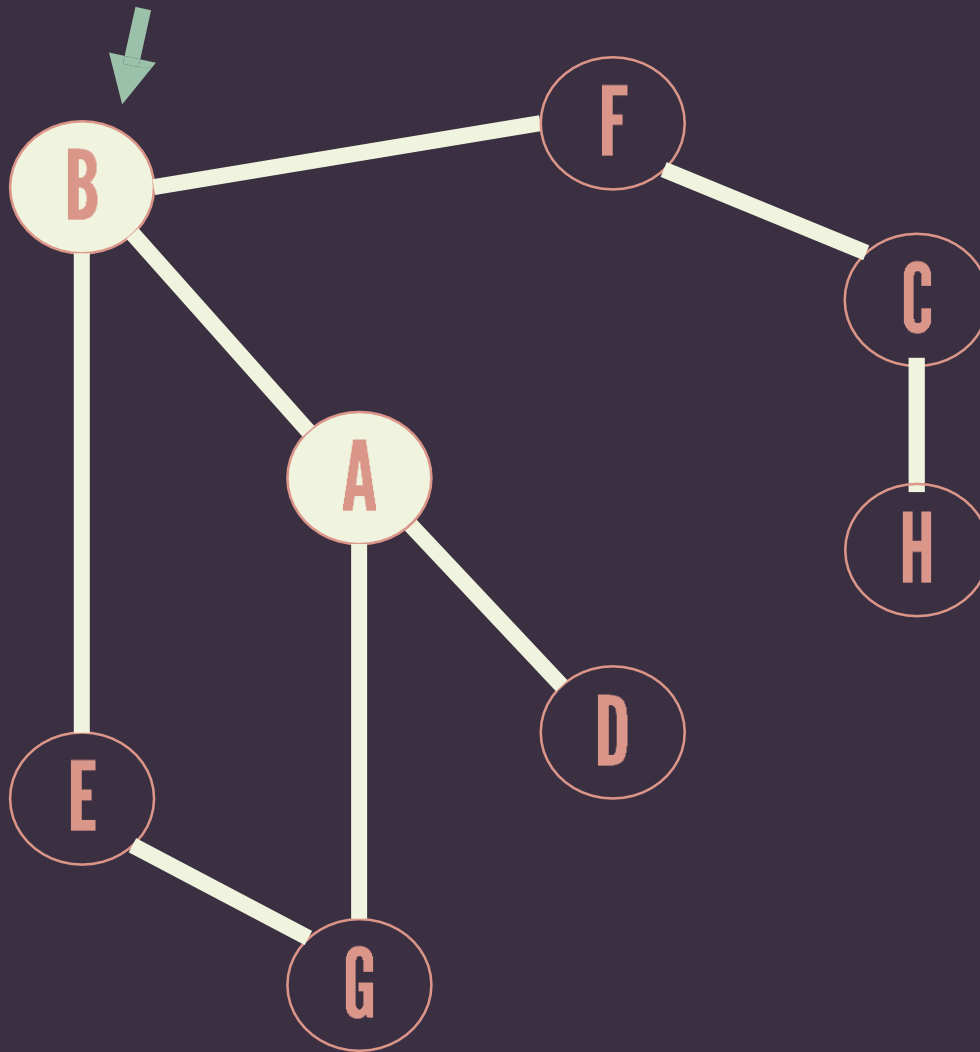



Result/Visited:



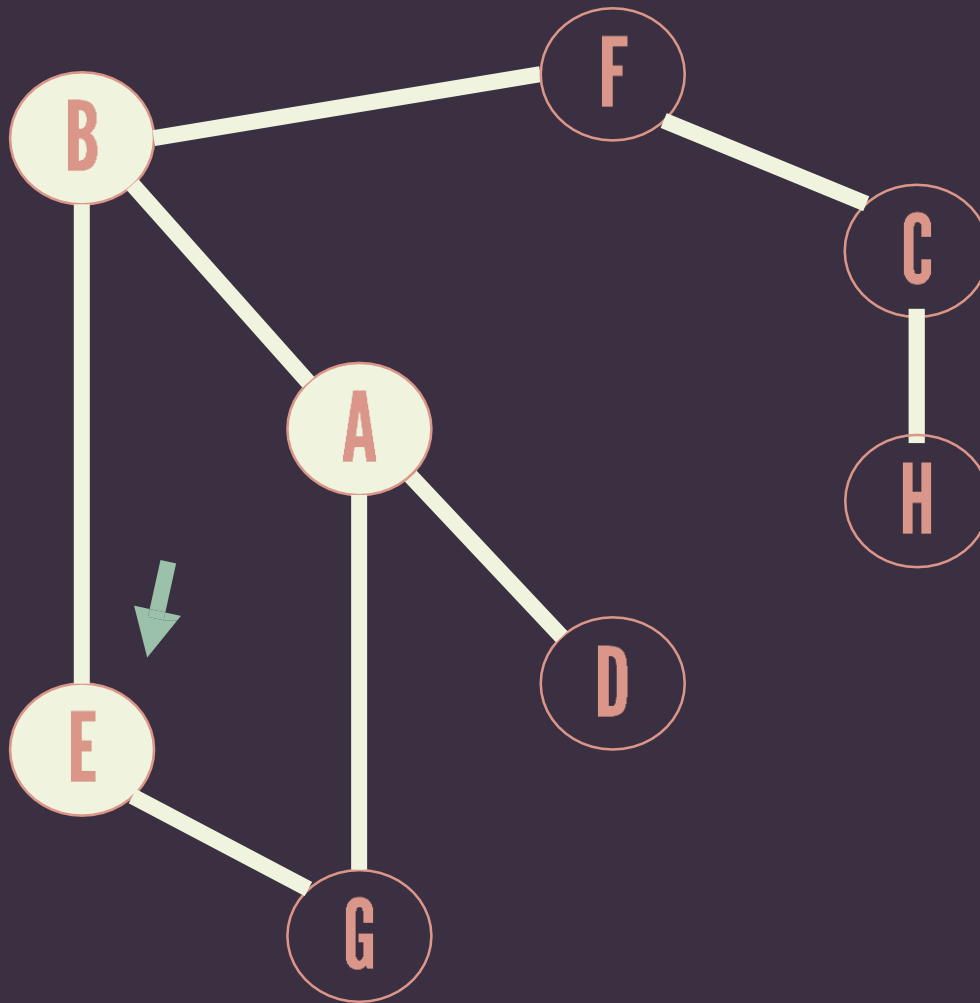
A

Result/Visited: A



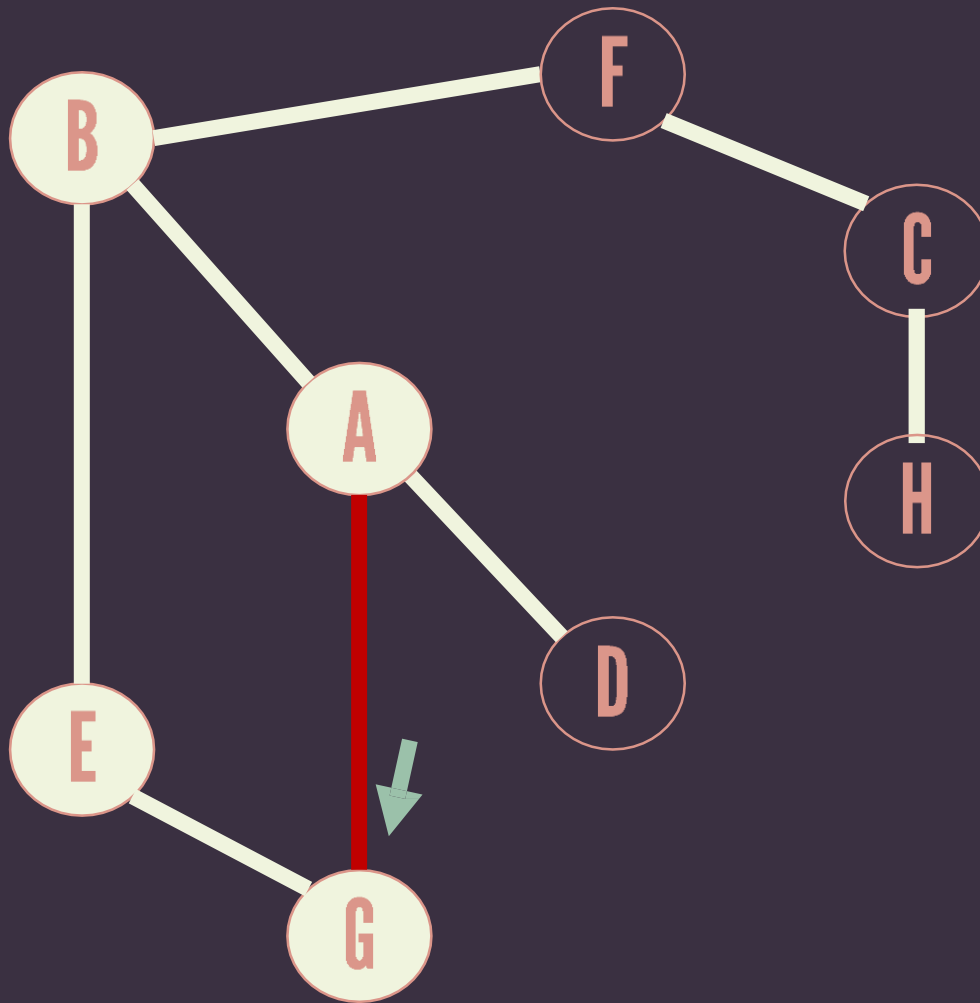
B
A

Result/Visited: AB



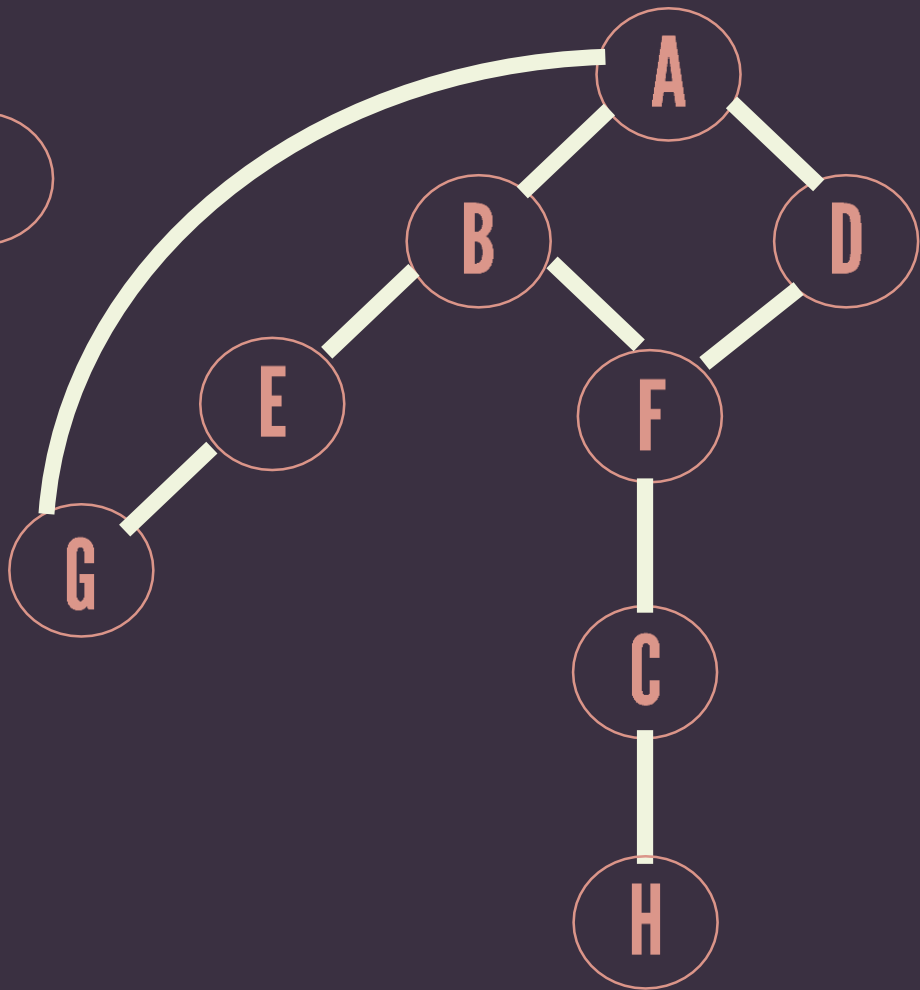
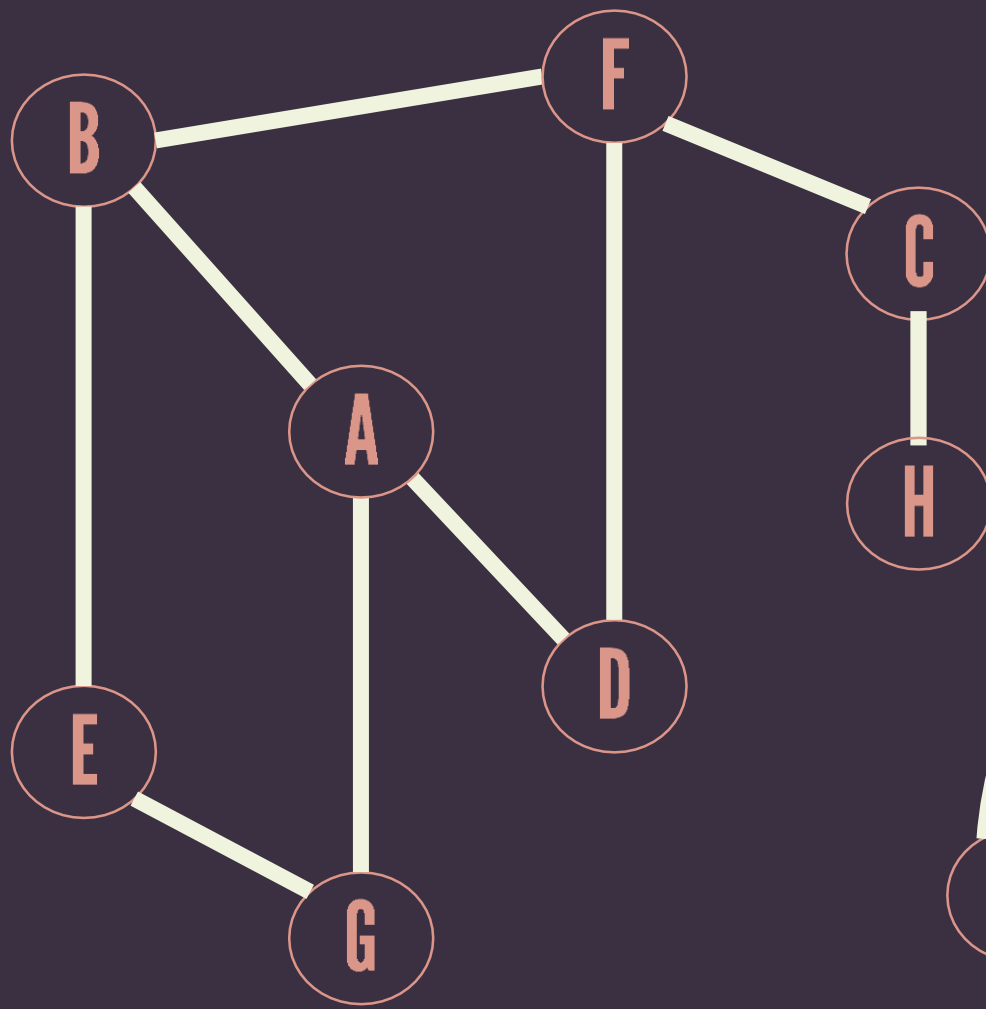
E
B
A

Result/Visited: A B E



G
E
B
A

Result/Visited: A B E G



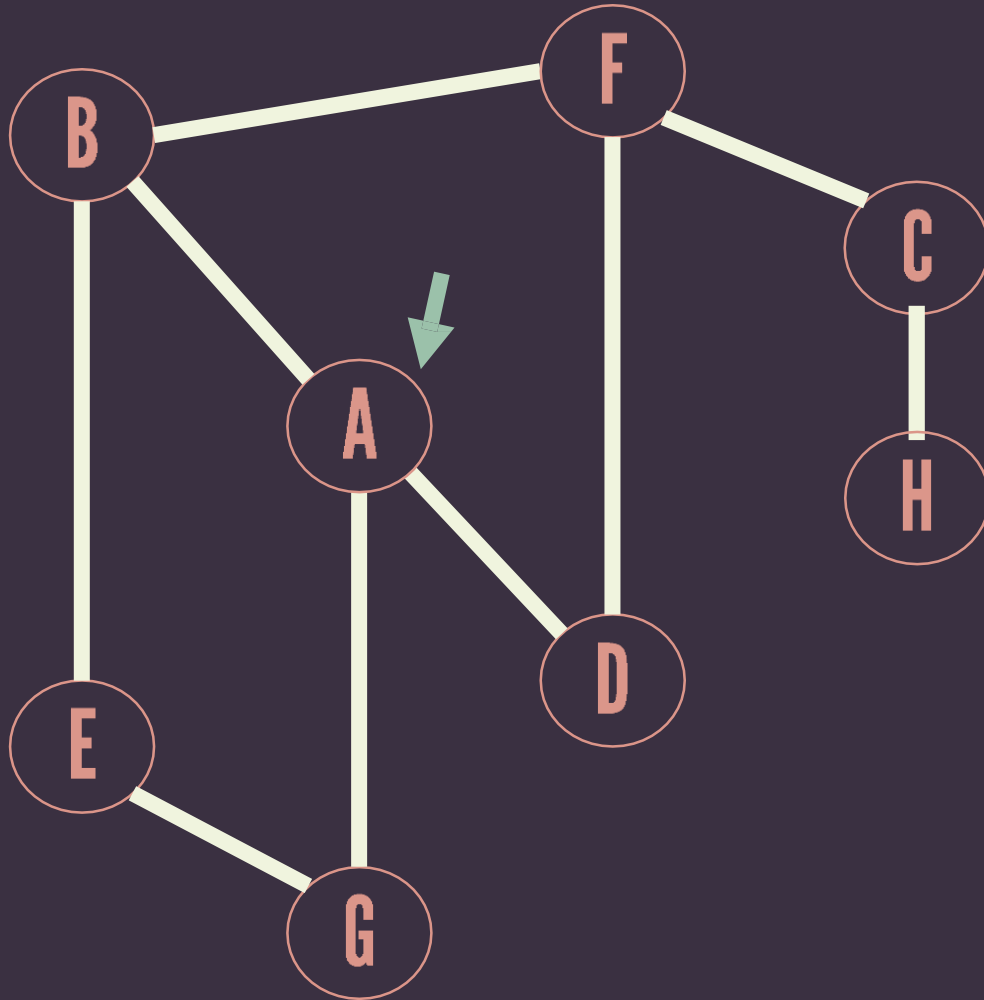
# BREADTH FIRST search

Generalization of  
level-order traversal  
of trees.

```
void BFS( vertex v, graph G ){
    queue Q;
    enqueue(v,Q);
    while queue Q is not empty{
        v = dequeue(Q);
        if ( !visited[v] ){
            print v;
            visited[v] = TRUE;
            for each w adjacent to v:
                if( !visited[w] )
                    enqueue(w, Q);
        }
    }
}
```

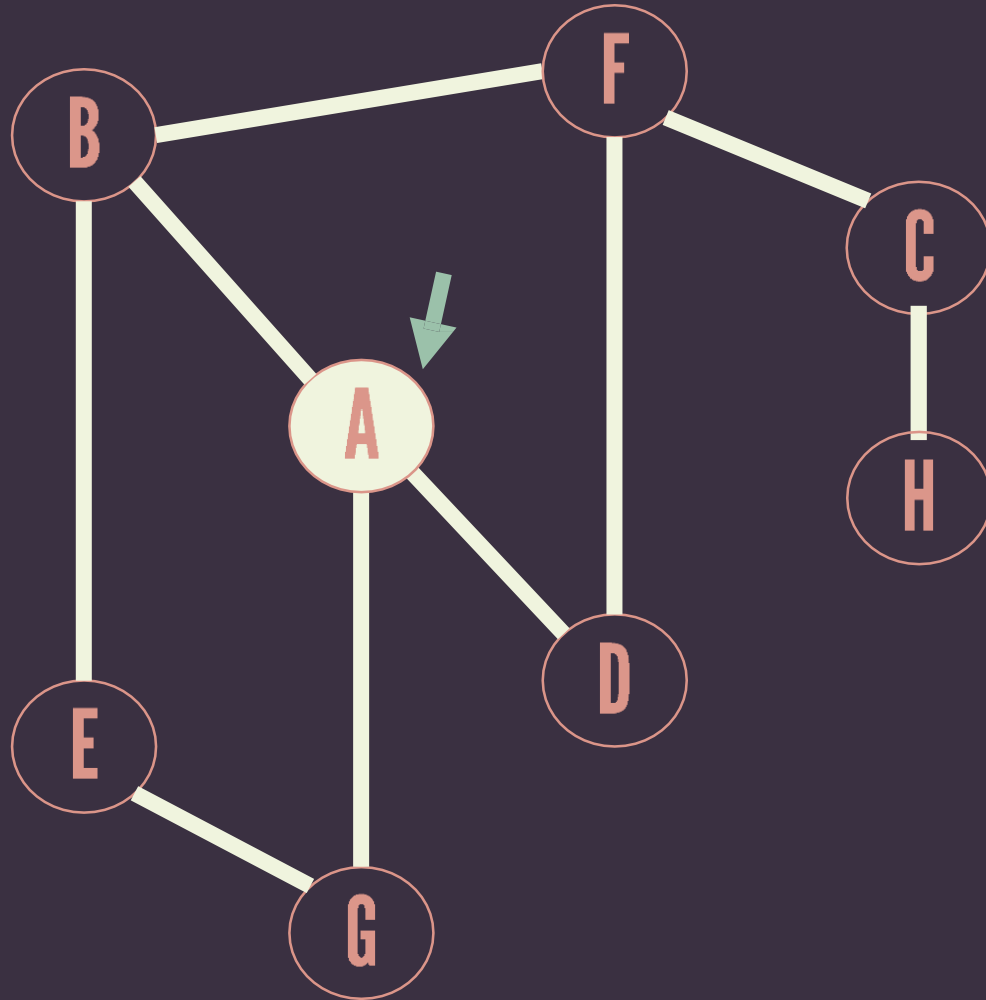


--	--	--	--	--	--	--



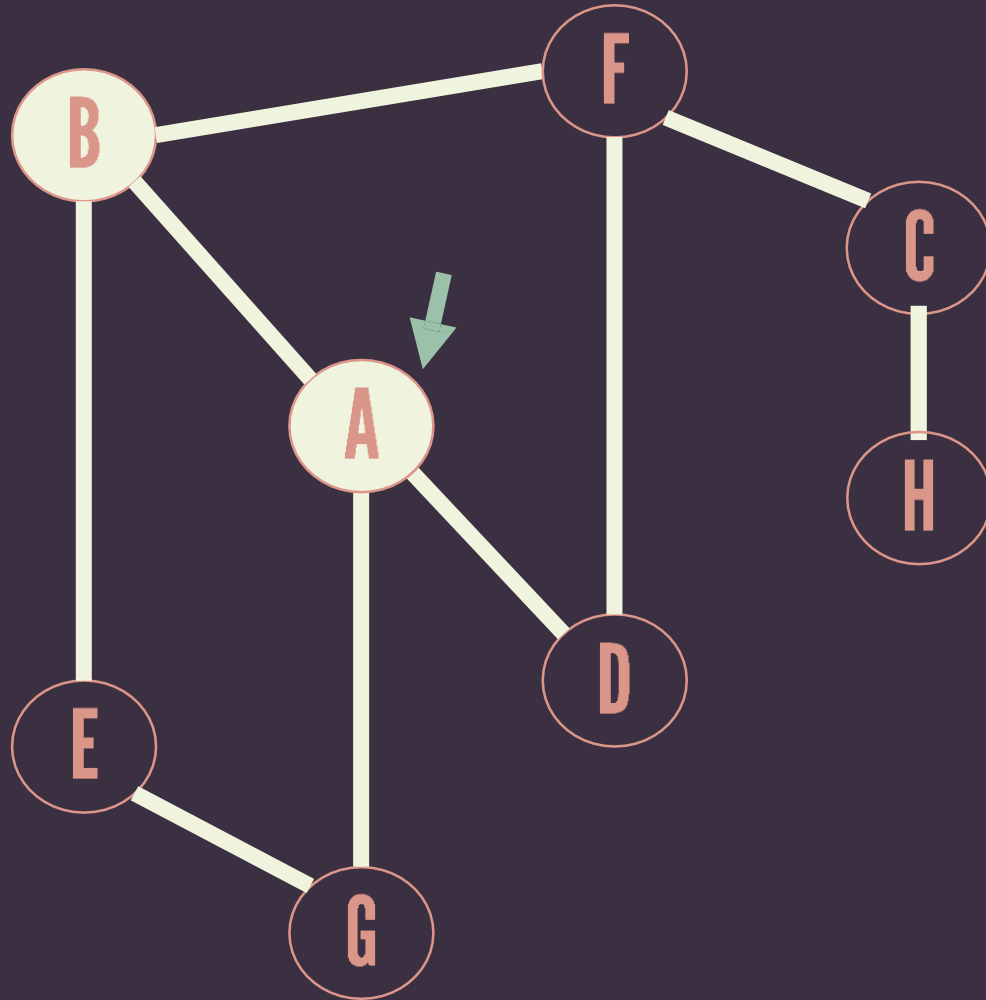
Result: A

A						
---	--	--	--	--	--	--



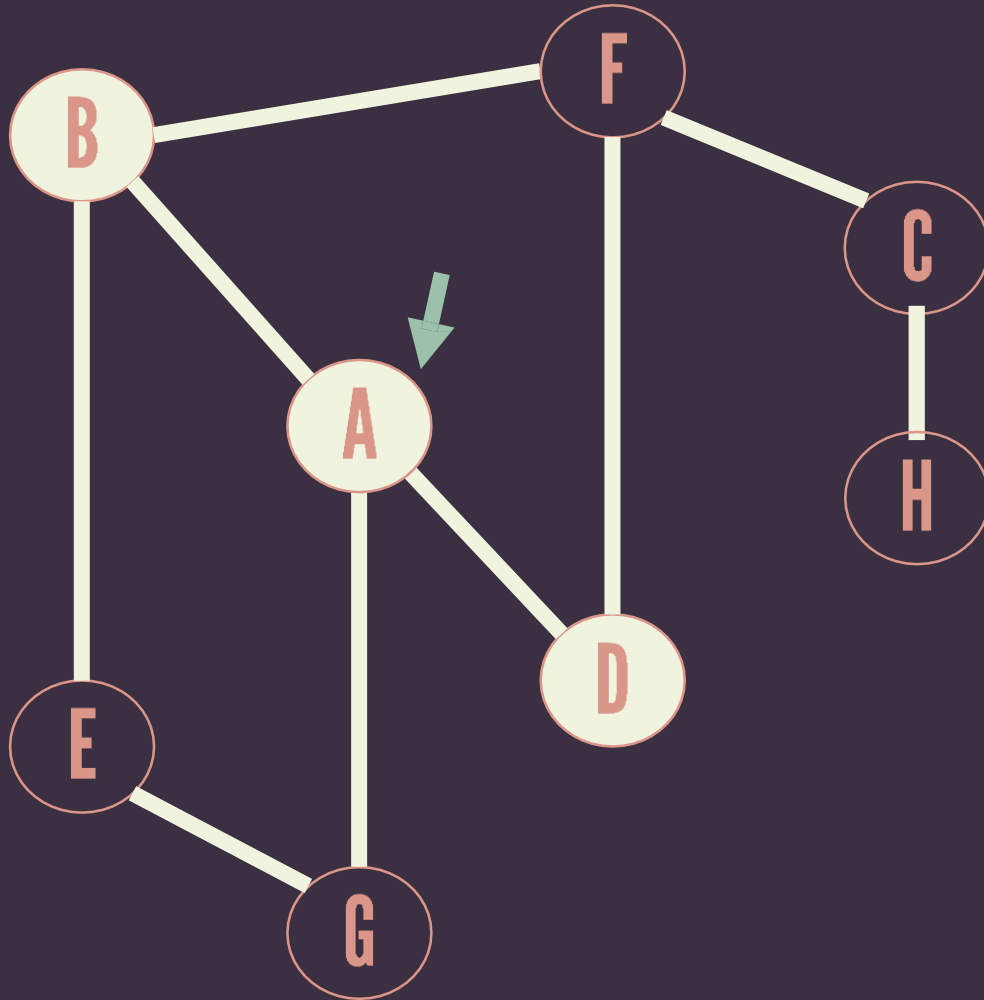
Result: A

A	B					
---	---	--	--	--	--	--



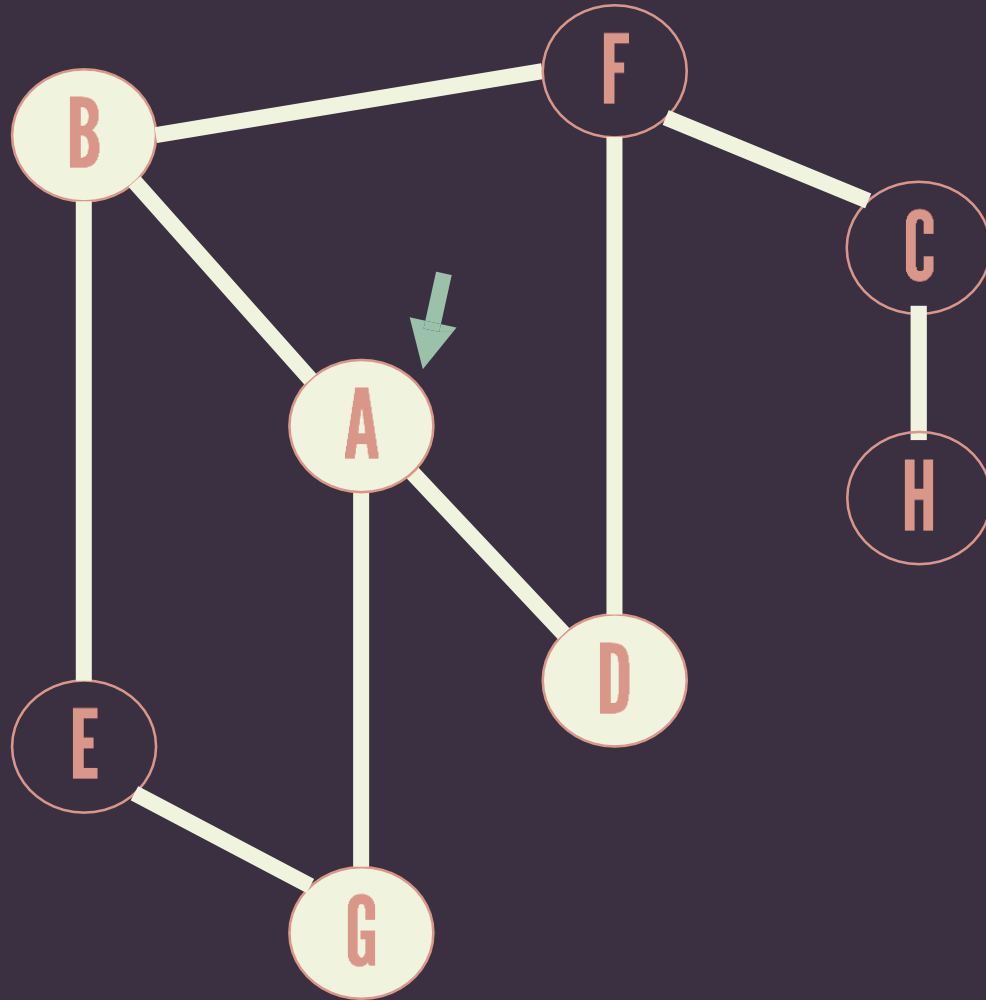
Result: A B

A	B	D				
---	---	---	--	--	--	--



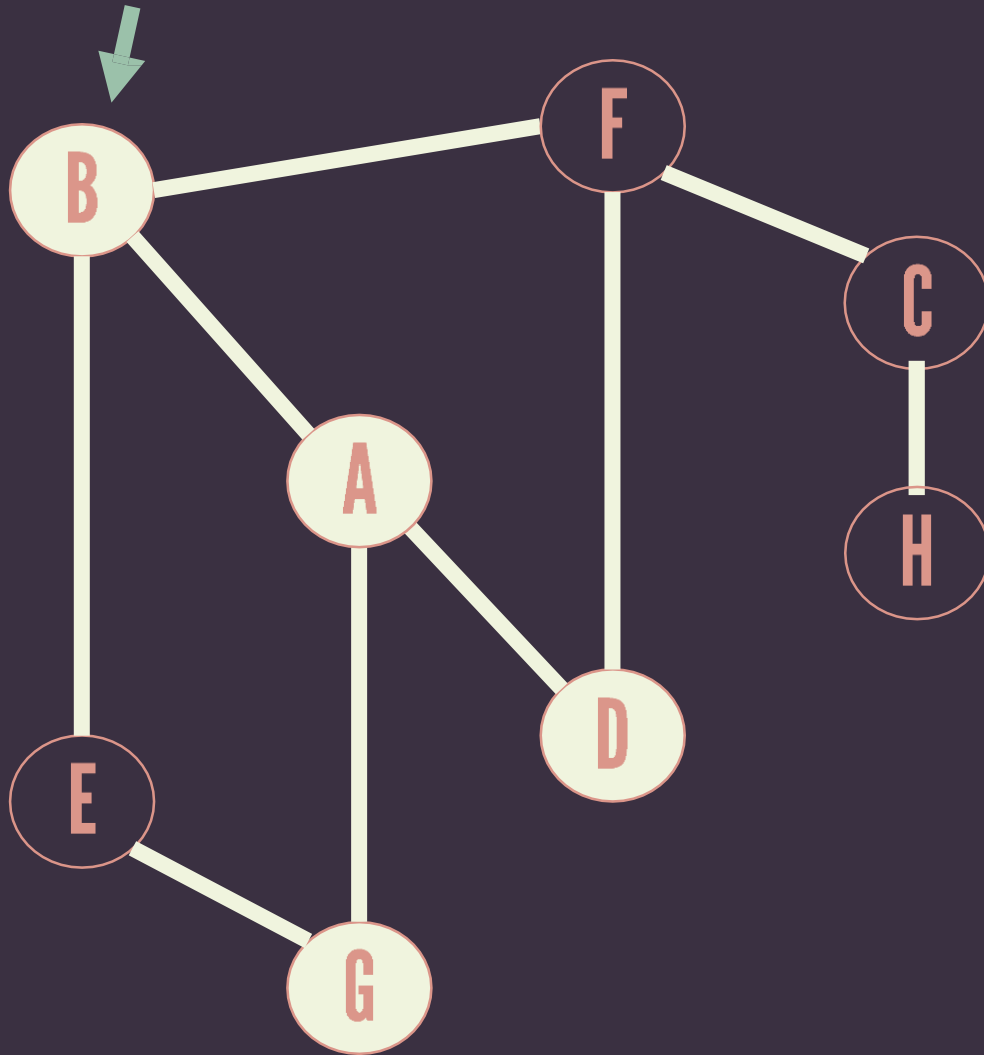
Result: A B D

A	B	D	G			
---	---	---	---	--	--	--



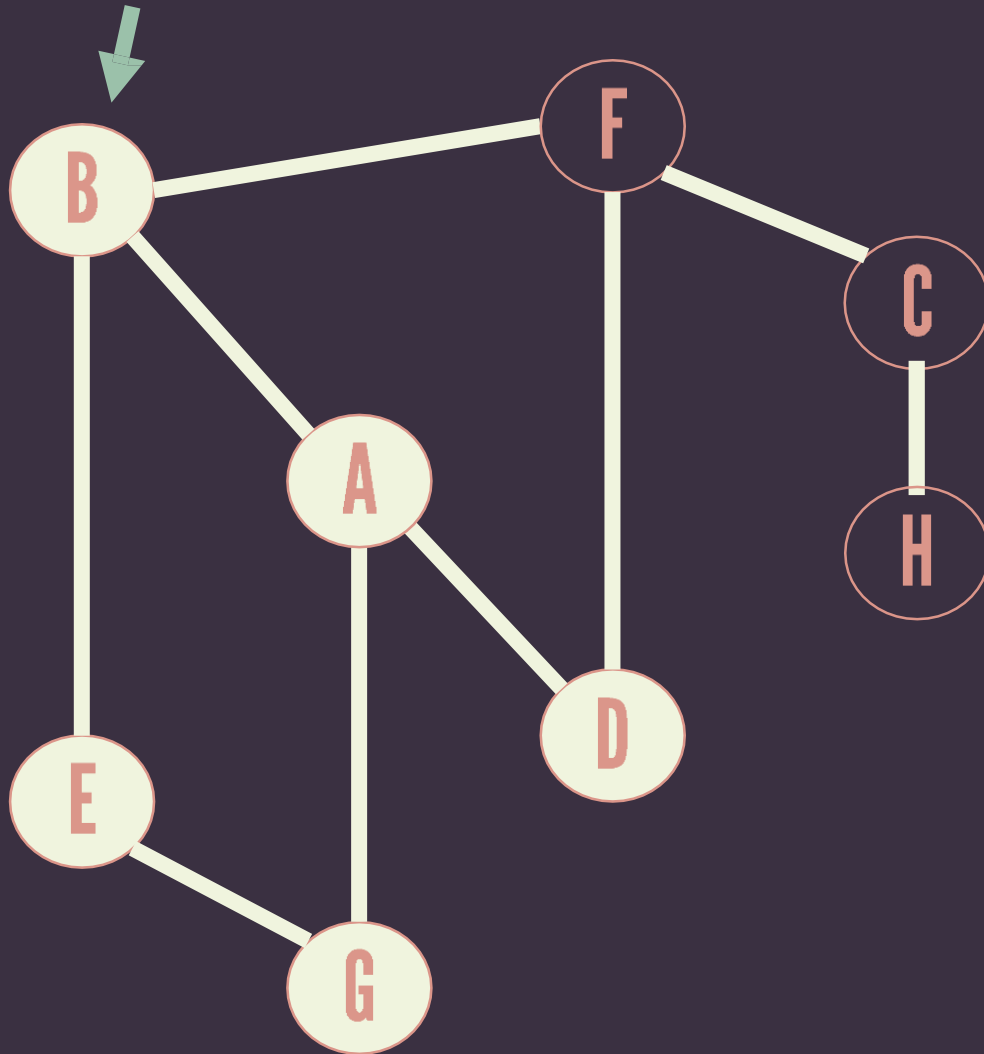
Result: A B D G

B	D	G				
---	---	---	--	--	--	--



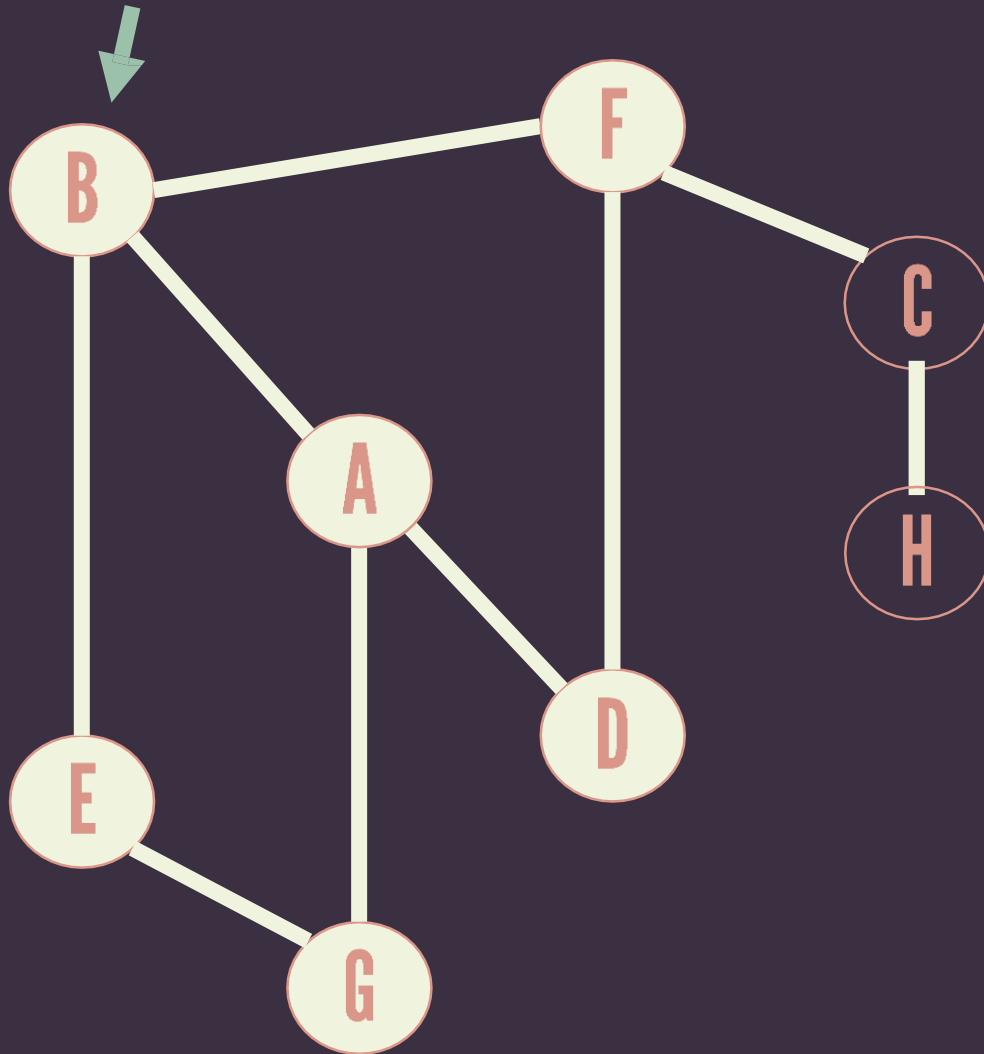
Result: A B D G

B	D	G	E			
---	---	---	---	--	--	--



Result: A B D G E

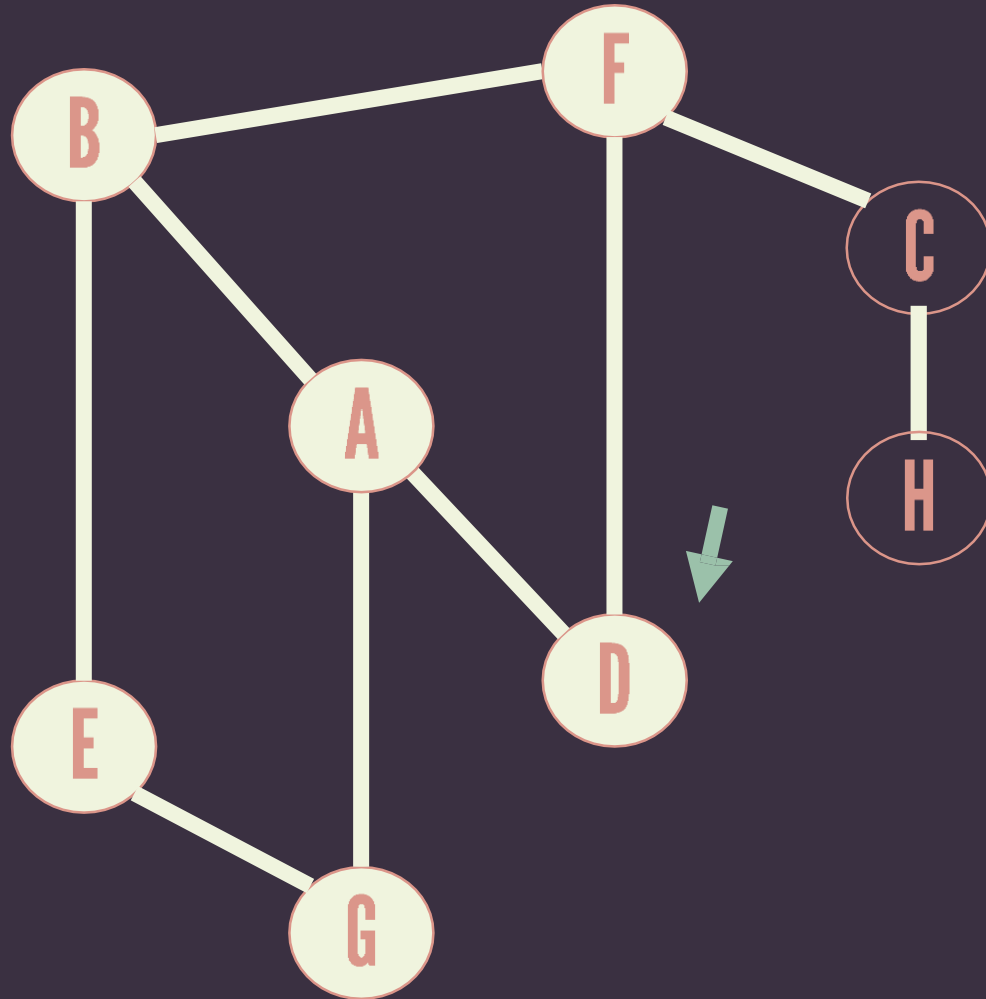
B	D	G	E	F		
---	---	---	---	---	--	--



Result: A B D G E F

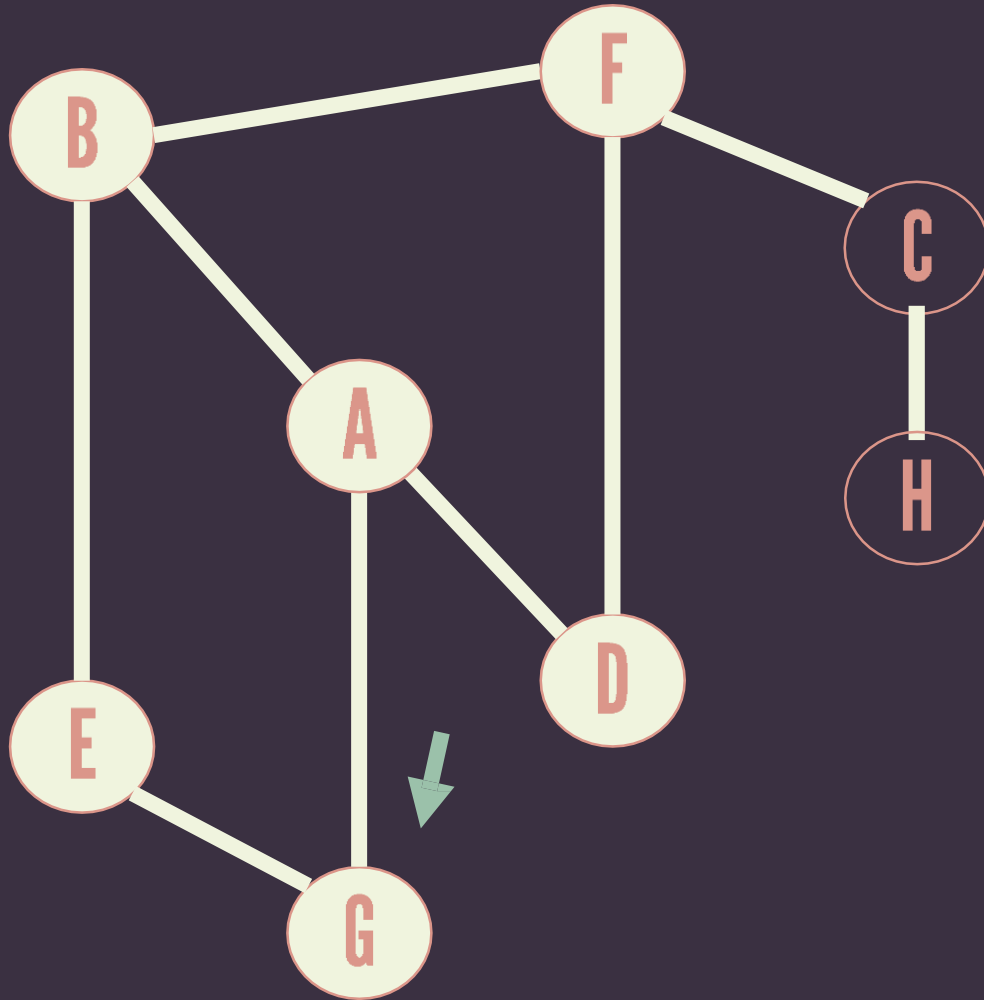


D	G	E	F			
---	---	---	---	--	--	--



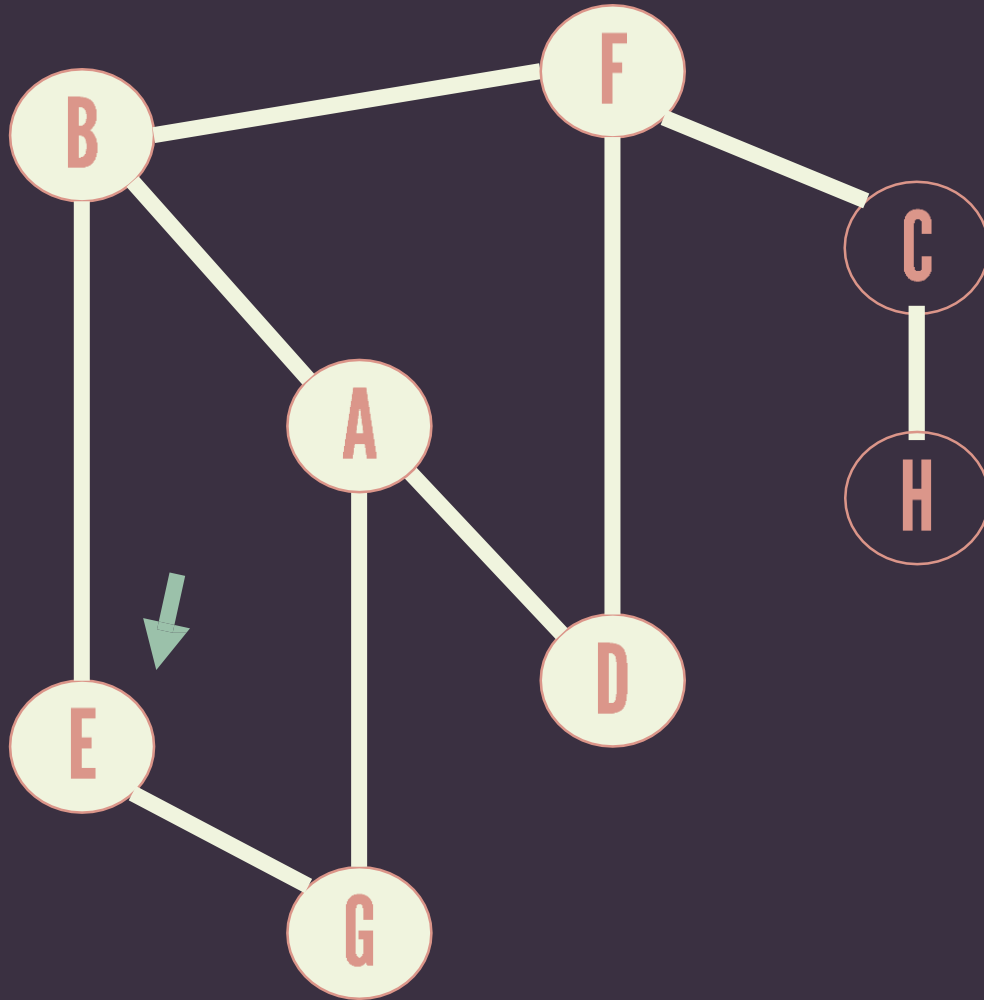
Result: A B D G E F

G	E	F				
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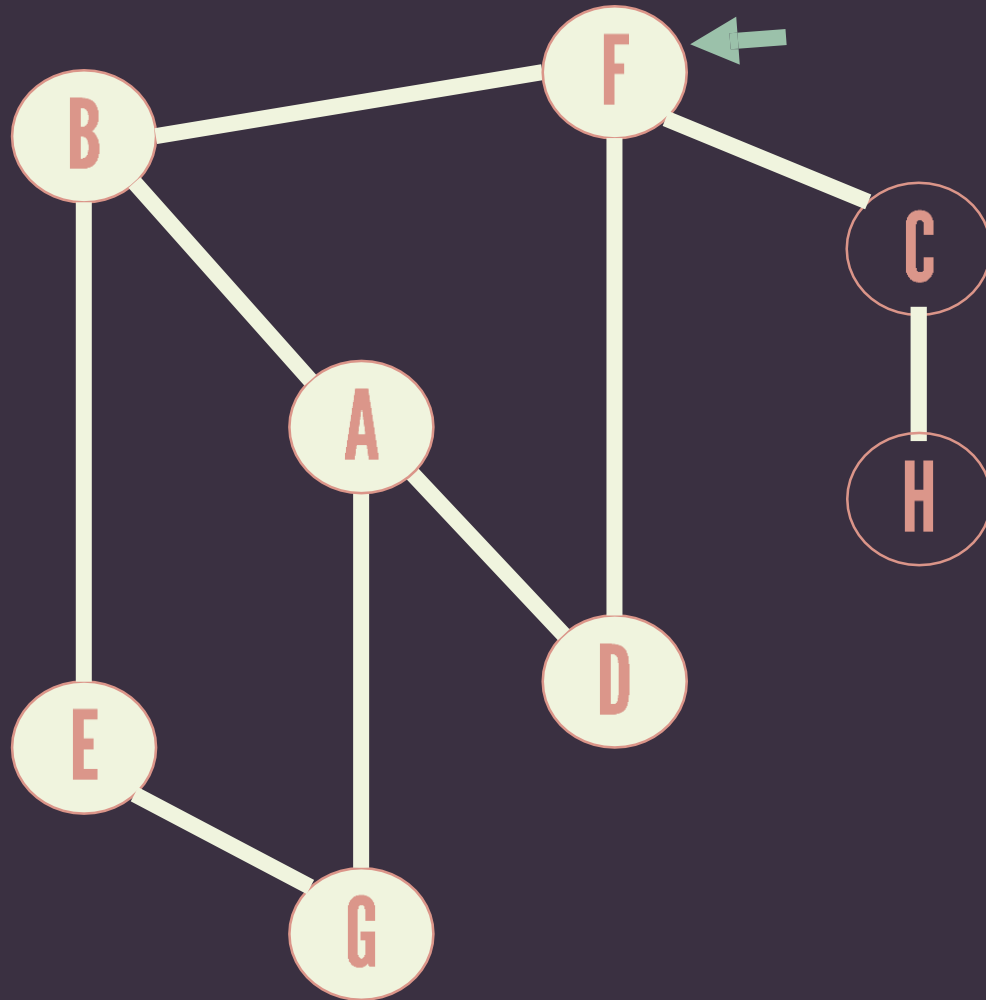
Result: A B D G E F

E	F					
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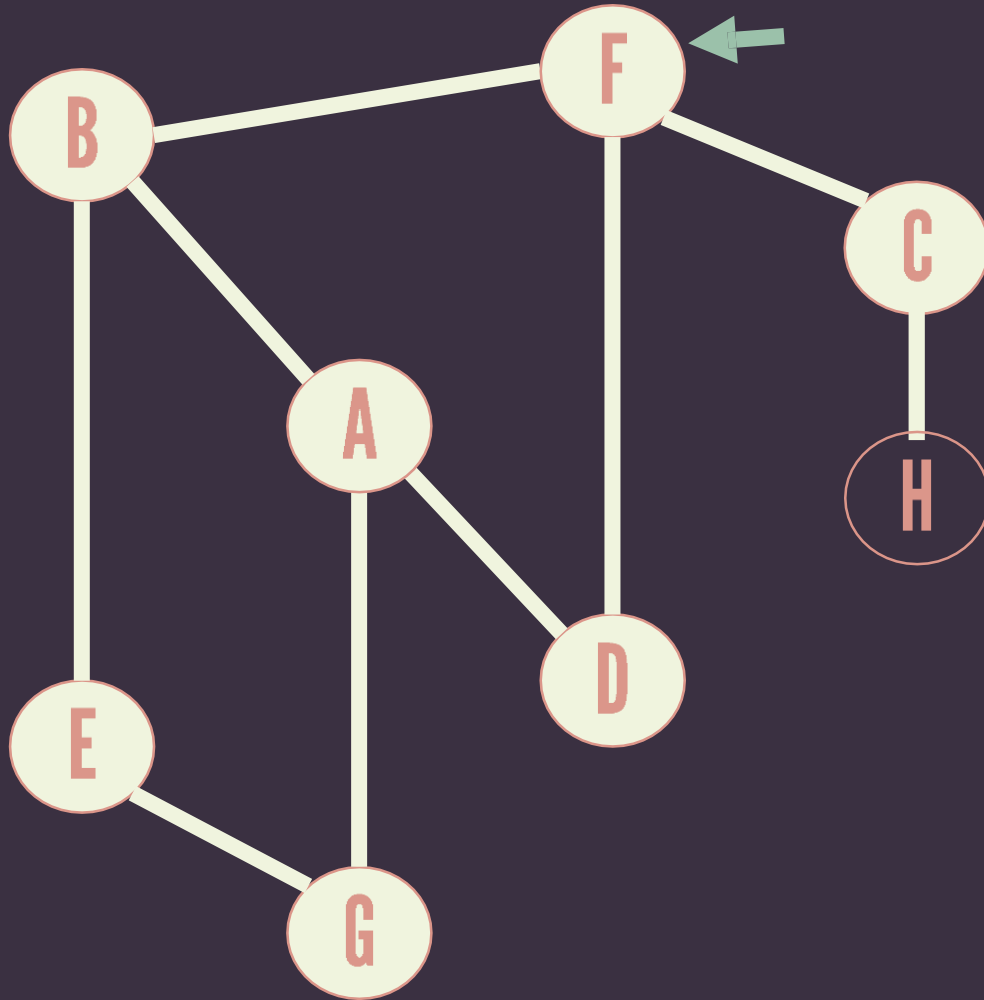
Result: A B D G E F

F						
---	--	--	--	--	--	--



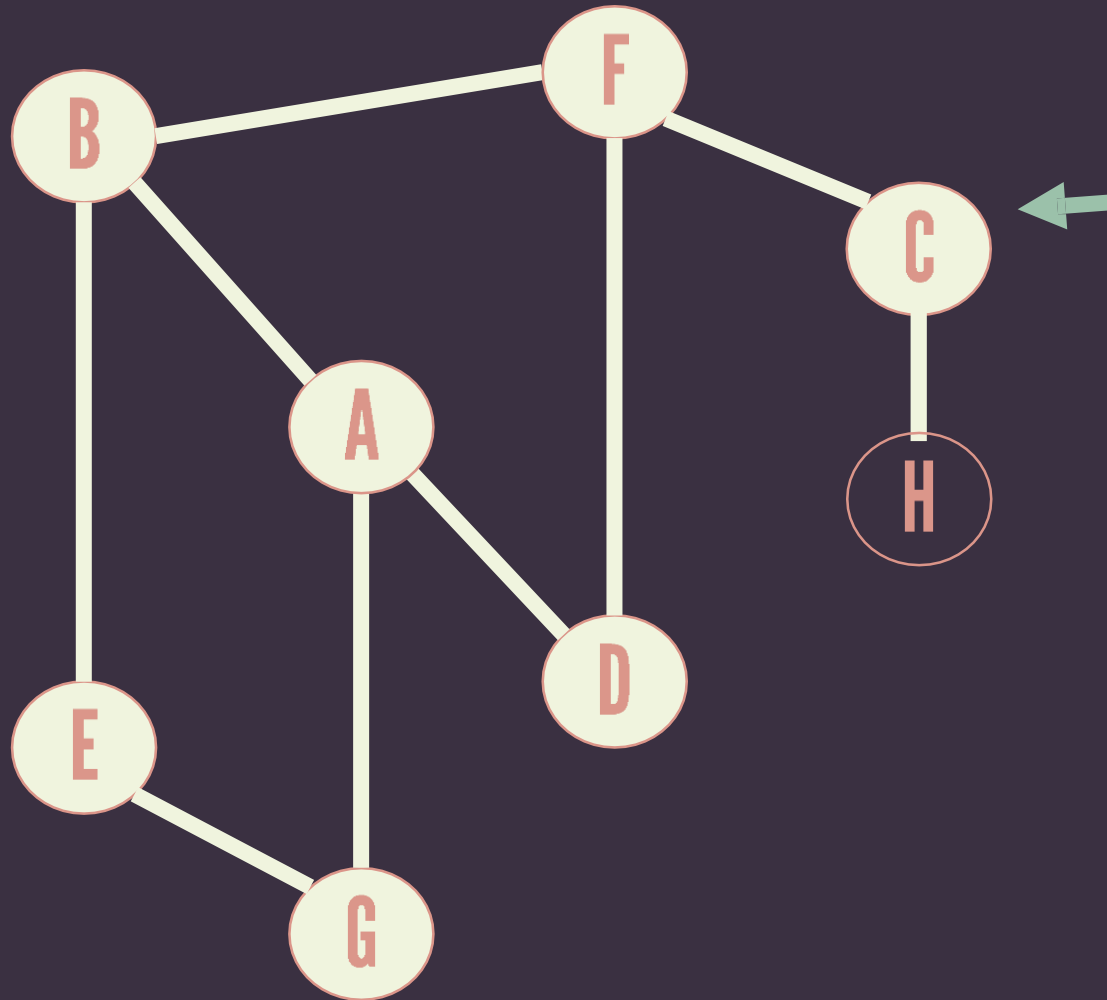
Result: A B D G E F

F	C					
---	---	--	--	--	--	--



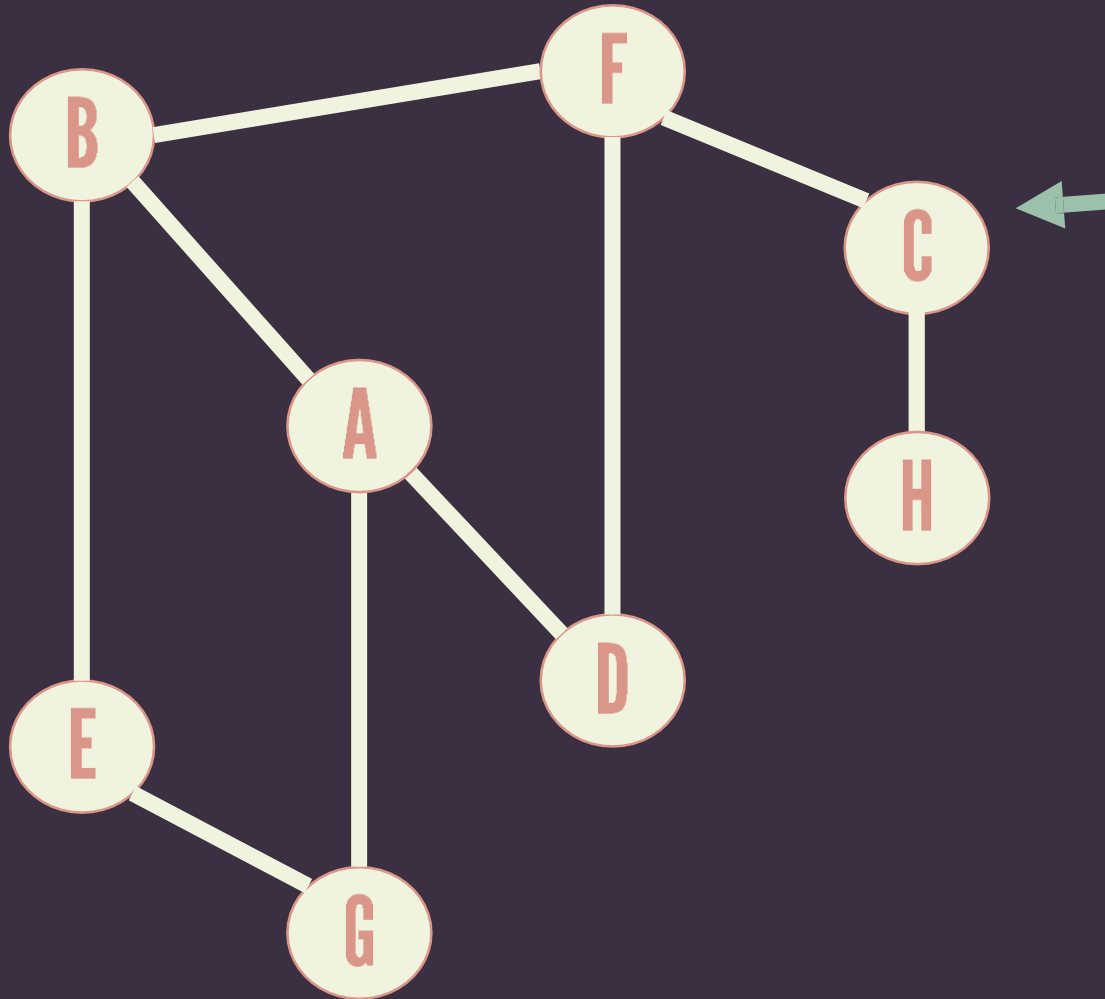
Result: A B D G E F C

C						
---	--	--	--	--	--	--



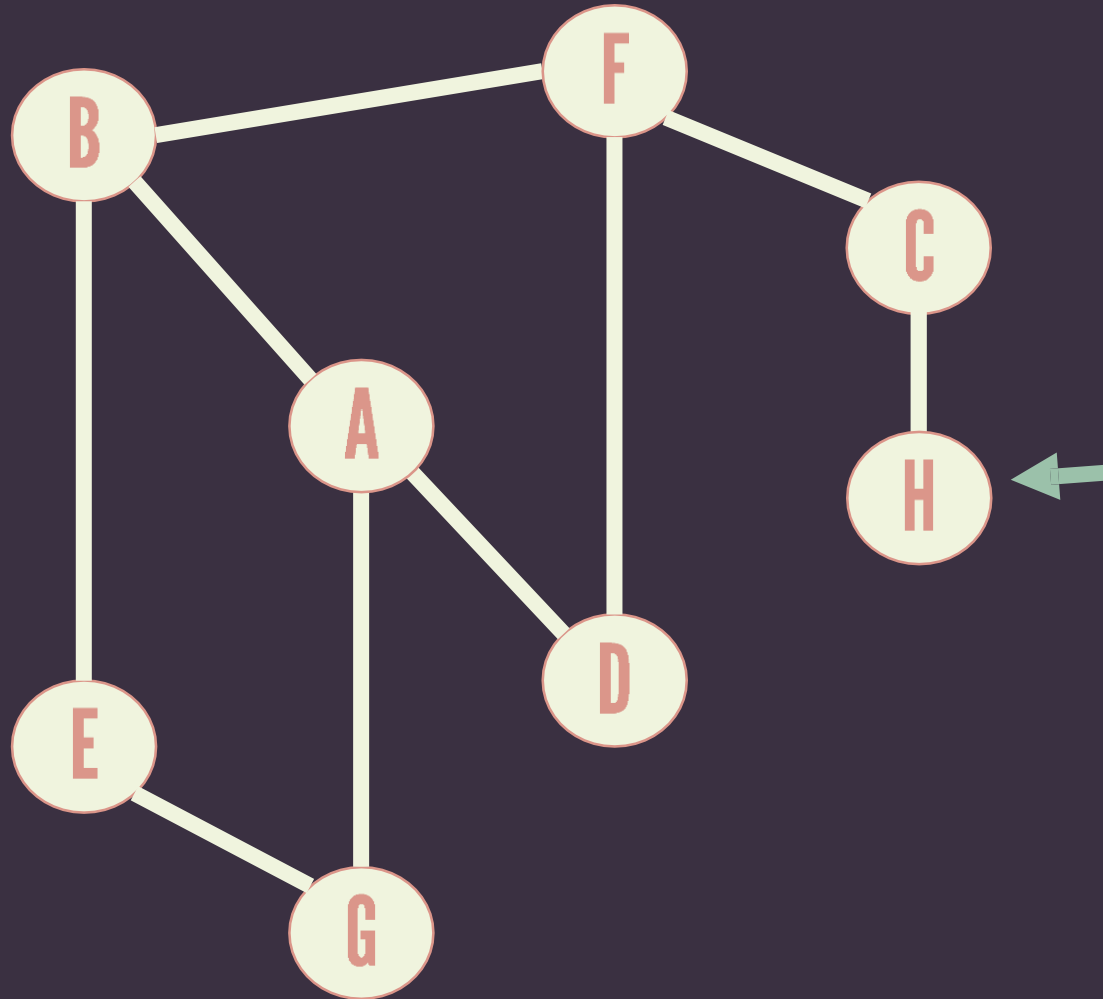
Result: A B D G E F C

C	H					
---	---	--	--	--	--	--



Result: A B D G E F C H

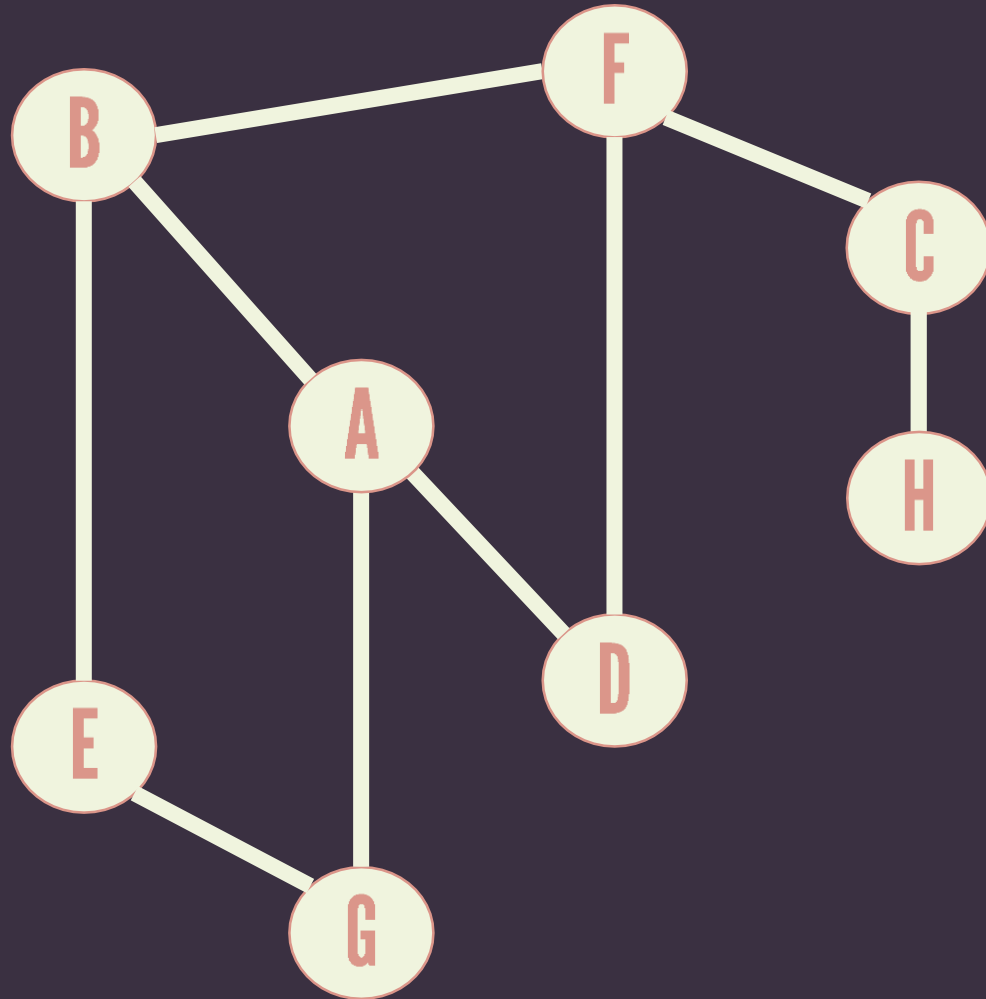
H						
---	--	--	--	--	--	--



Result: A B D G E F C H



--	--	--	--	--	--	--



Result: A B D G E F C H

# BFS

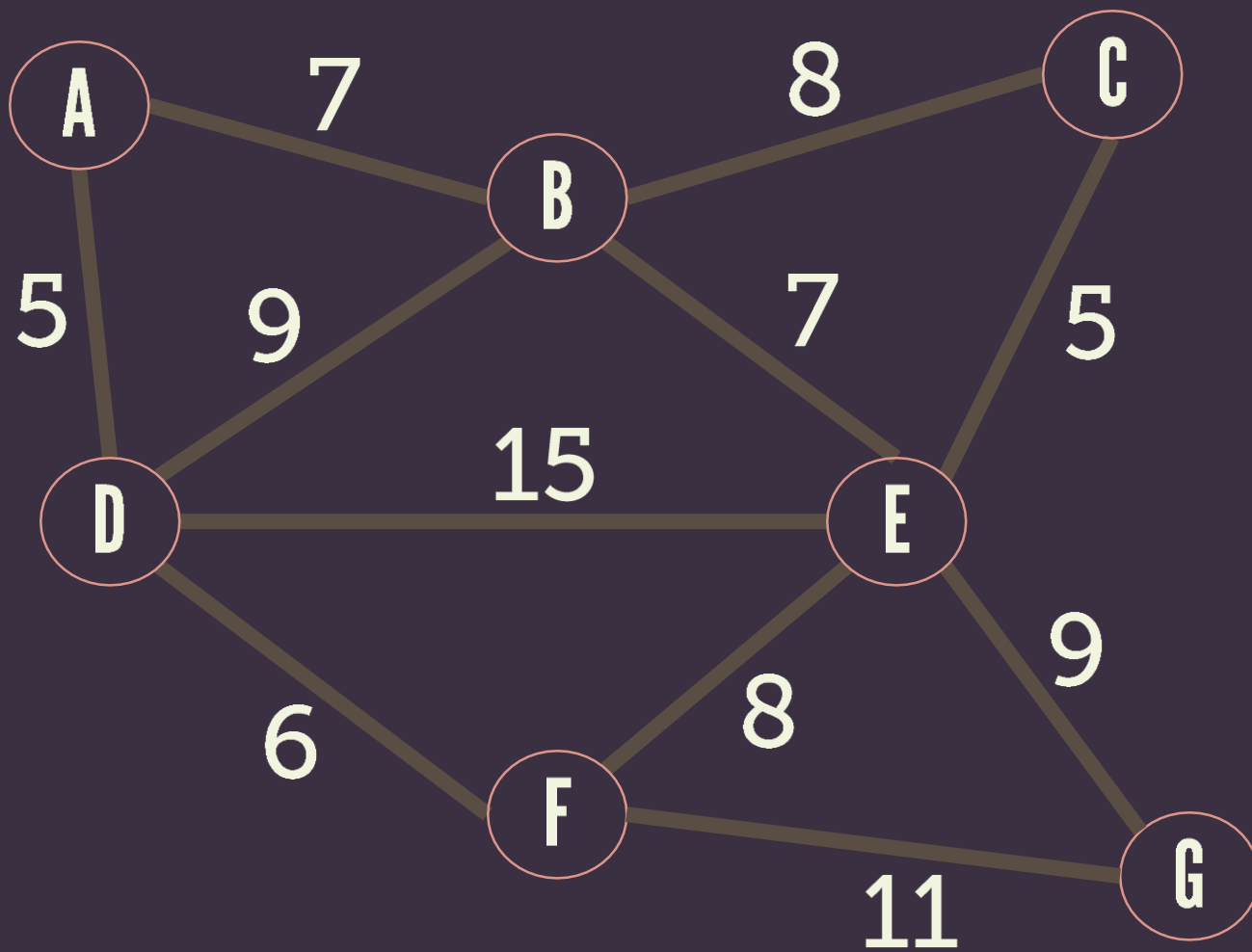
## applications

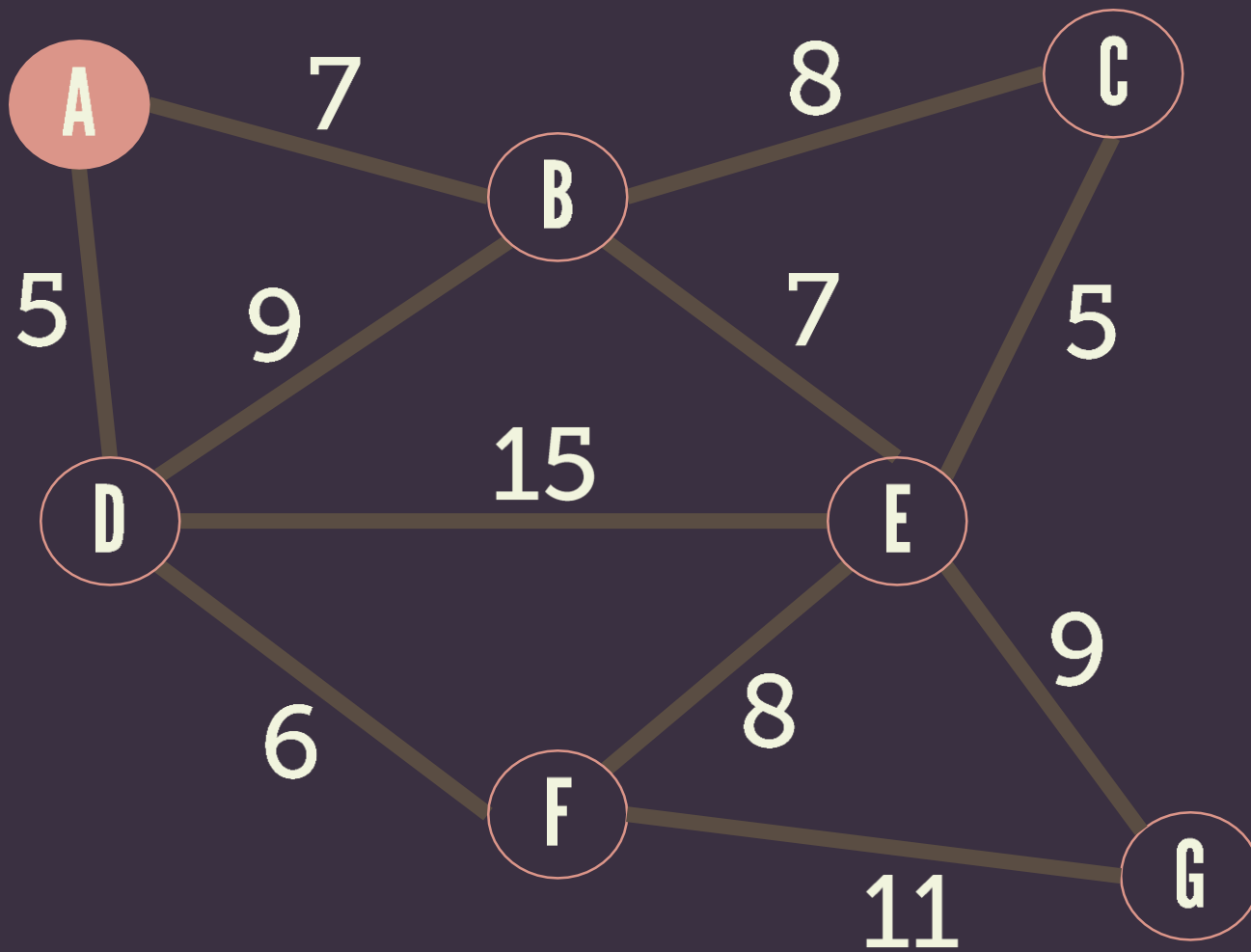
Finding nodes within one.  
connected components.  
Testing for bipartiteness.  
Finding shortest paths.

minimum  
**SPANNING TREES**

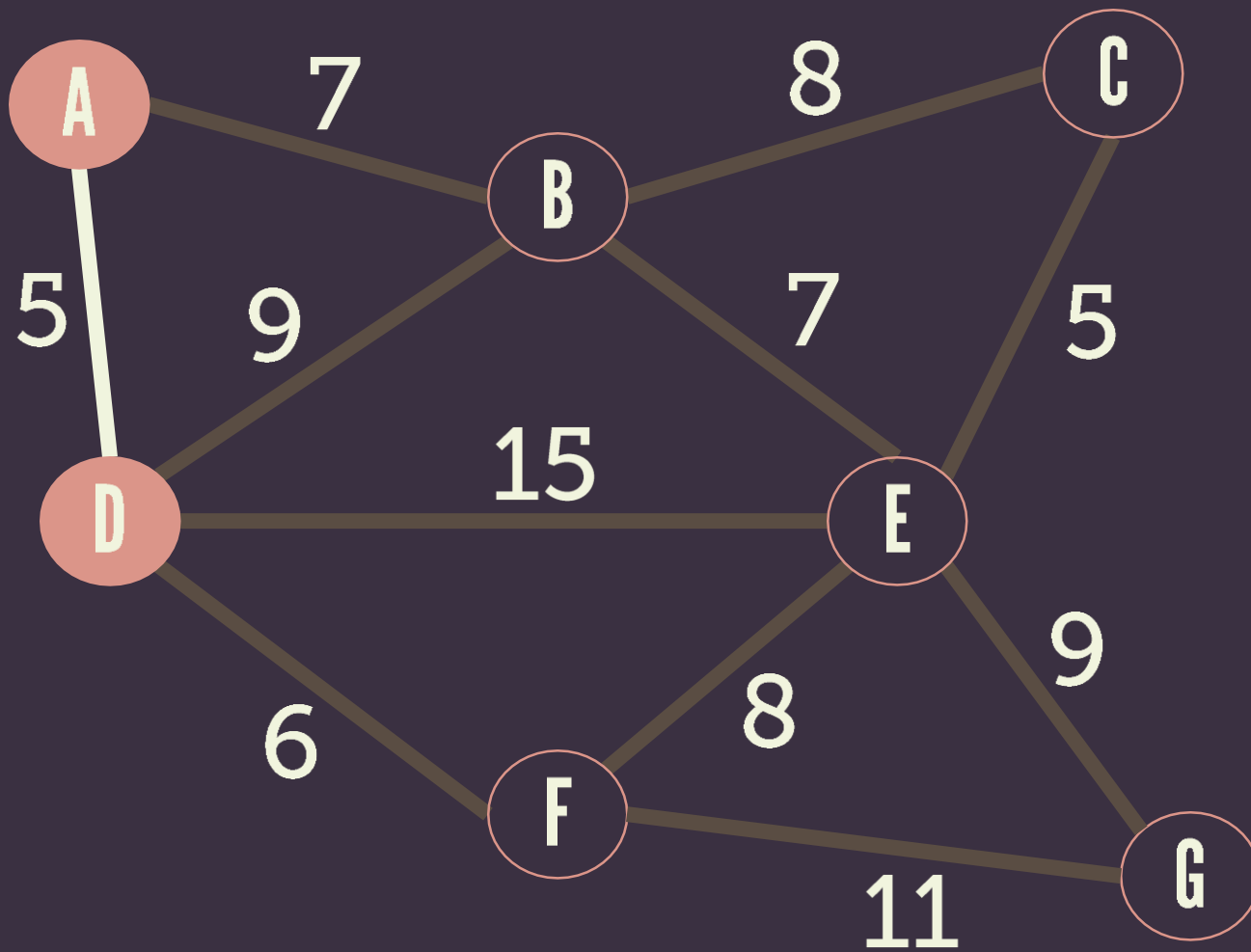
minimum  
**SPANNING**  
**TREE**

A tree formed from graph edges that connects all the vertices of the graph at lowest cost.

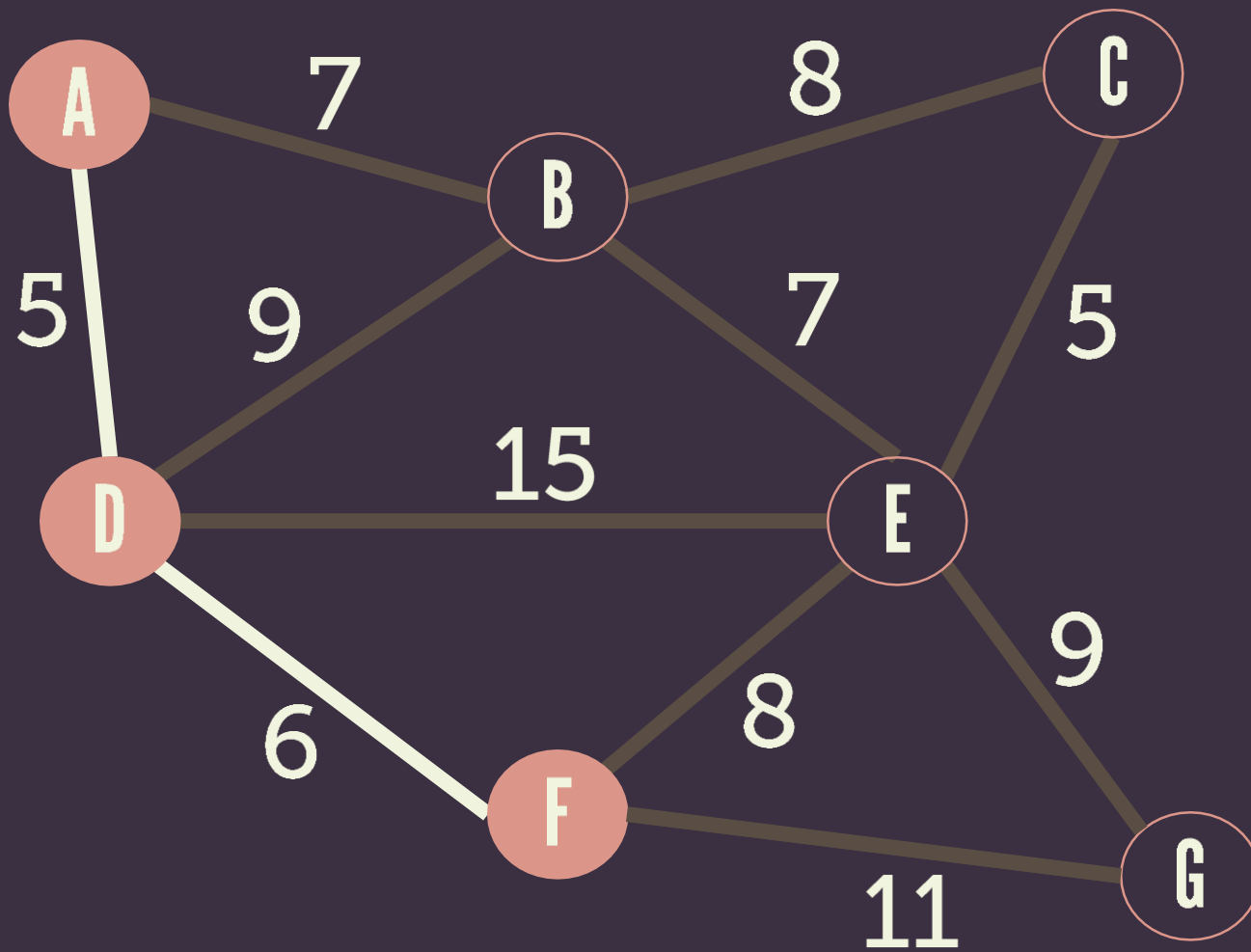




cost = ?

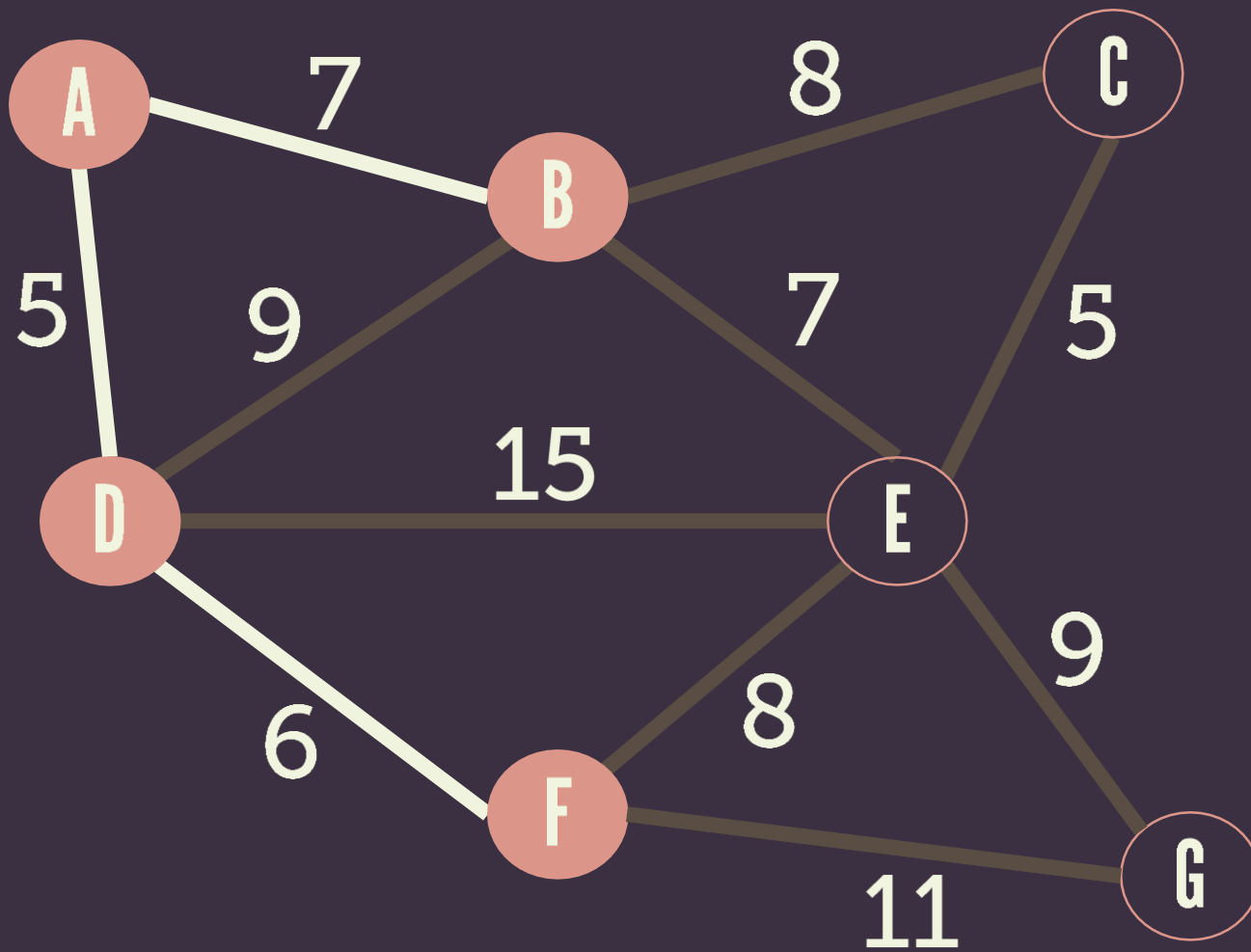


cost = 5

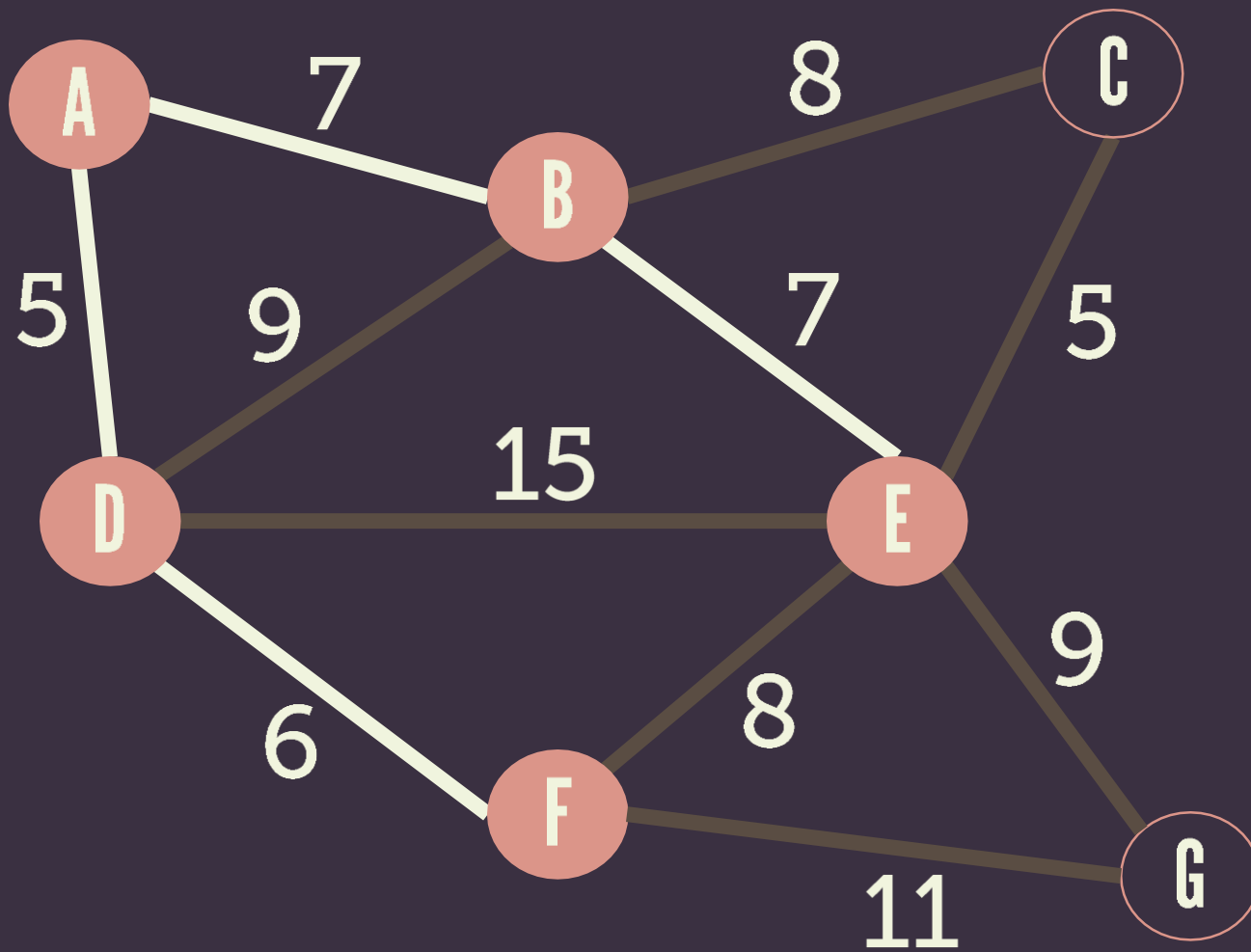


cost = 11

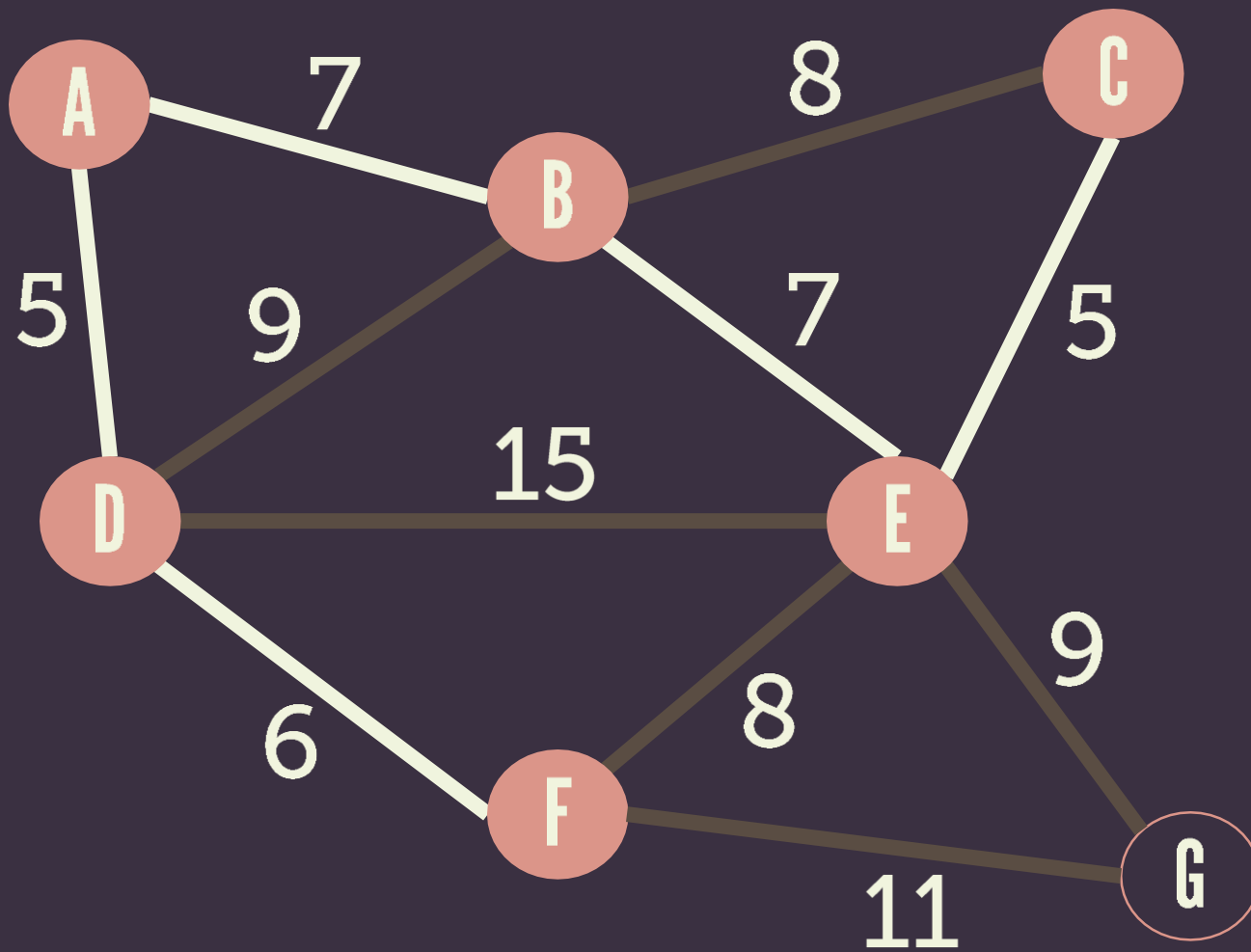




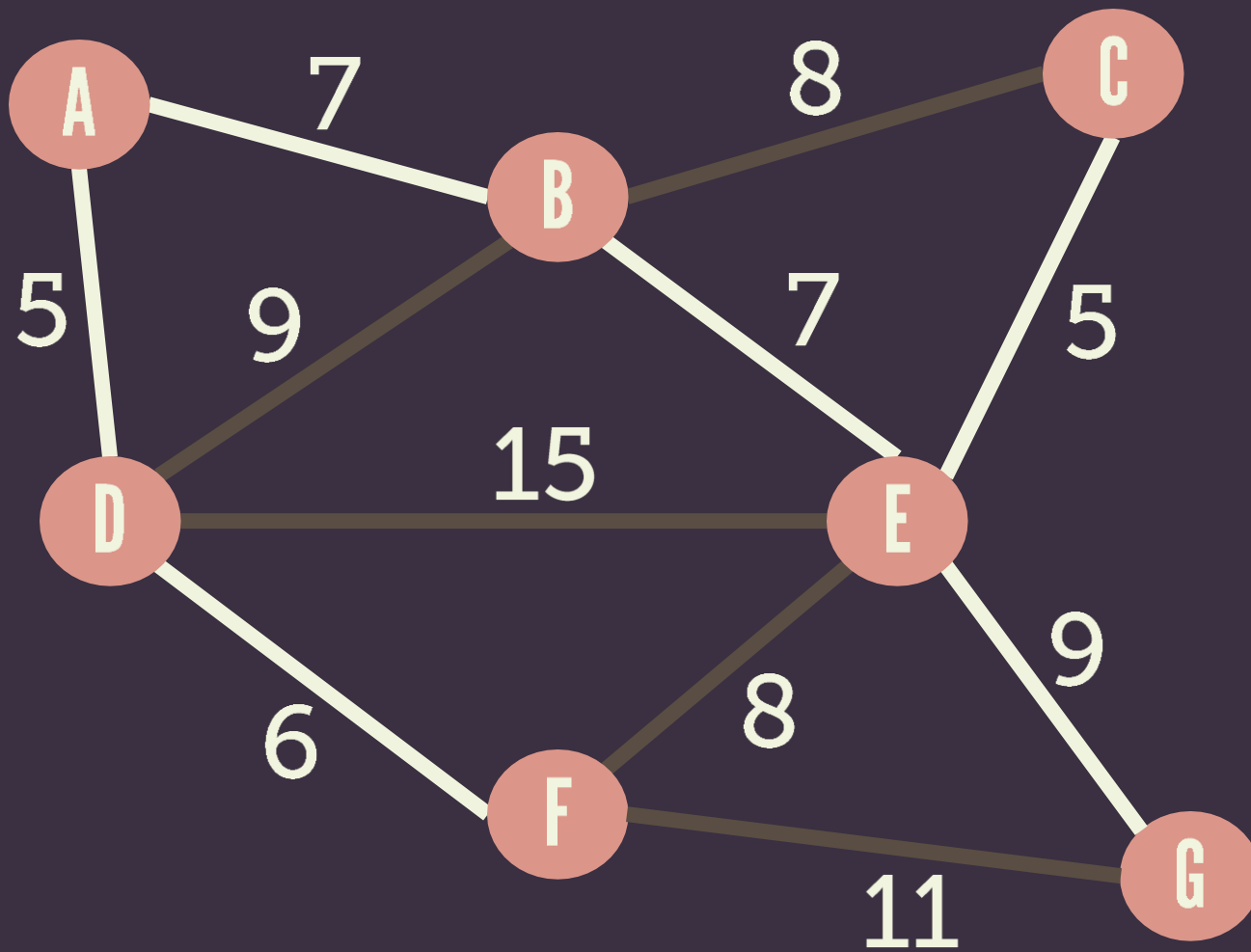
cost = 18



cost = 25



cost = 30



cost = 39