

Name: Michael Rogers

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CSCI 3104, Algorithms  
Problem Set 6b (40 points)

Profs. Hoenigman & Agrawal  
Fall 2019, CU-Boulder

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**Instructions for submitting your solution:**

- The solutions **should be typed** and we cannot accept hand-written solutions. Here's a short intro to Latex.
  - You should submit your work through **Gradescope** only.
  - If you don't have an account on it, sign up for one using your CU email. You should have gotten an email to sign up. If your name based CU email doesn't work, try the identikey@colorado.edu version.
  - Gradescope will only accept **.pdf** files (except for code files that should be submitted separately on Gradescope if a problem set has them) and **try to fit your work in the box provided**.
  - You cannot submit a pdf which has less pages than what we provided you as Gradescope won't allow it.
  - Verbal reasoning is typically insufficient for full credit. Instead, write a logical argument, in the style of a mathematical proof.
  - For every problem in this class, you must justify your answer: show how you arrived at it and why it is correct. If there are assumptions you need to make along the way, state those clearly.
  - You may work with other students. However, **all solutions must be written independently and in your own words**. Referencing solutions of any sort is strictly prohibited. You must explicitly cite any sources, as well as any collaborators.
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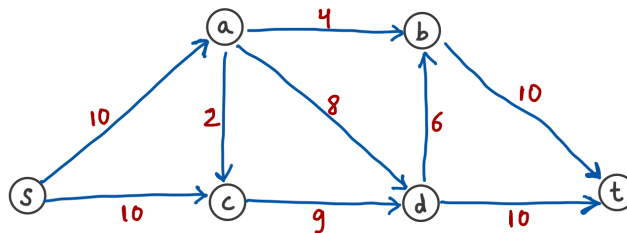
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1. (19 pts) Based on the following network and the given edge capacities answer the following.



- (a) (12 pts) Suppose we start the Ford-Fulkerson algorithm and **select the path  $s \rightarrow a \rightarrow c \rightarrow d \rightarrow t$  in the first iteration (Do not choose the first s-t path on your own)**. Complete all the iterations of Ford-Fulkerson to find the Max-Flow (including the first round that is incomplete). Clearly show each round with

- The path that you are selecting in that round.
- The bottleneck edge on this path.
- The additional flow that you push from the source by augmenting (pushing maximum allowed flow along) this selected augmenting path.
- The residual graph with the residual capacities (on both the forward and backward) edges.

Also, report the Max-Flow after the algorithm terminates.

*Solution.* Paths -

Path 1: s-a-c-d-t, Bottleneck: 2, Additional Flow: 2, Total Flow at iteration: 2

Path 2: s-c-d-t, Bottleneck: 7, Additional Flow: 7, Total Flow at iteration: 9

Path 3: s-c-a-b-t, Bottleneck: 2, Additional Flow: 2, Total Flow at iteration: 11

Path 4: s-a-d-b-t, Bottleneck: 6, Additional Flow: 6, Total Flow at iteration: 17

Path 5: s-a-b-t, Bottleneck: 2, Additional Flow: 2, Total Flow at iteration: 19

Algorithm Terminates. Max Flow = 19

Pictures Associated with each iteration listed below.

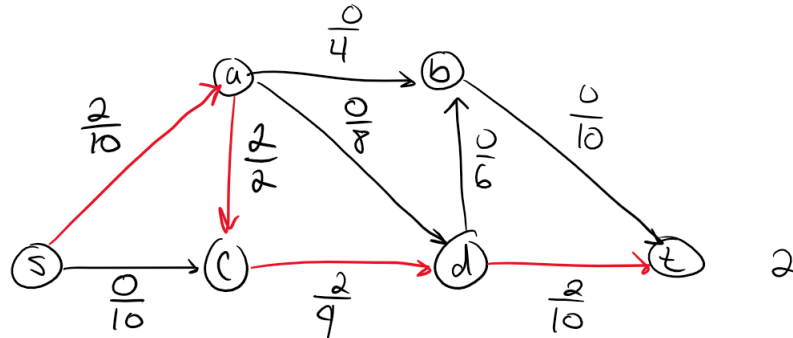
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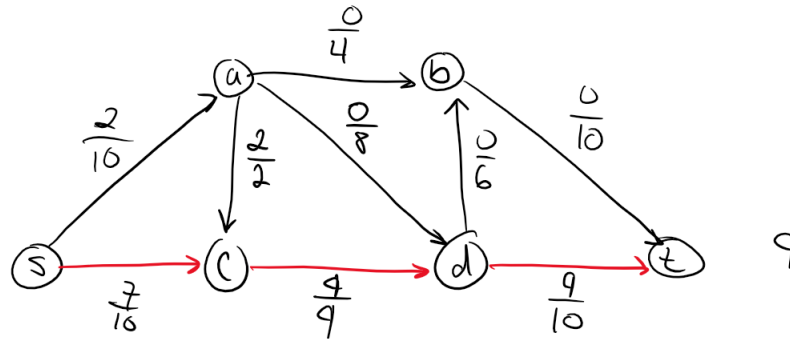
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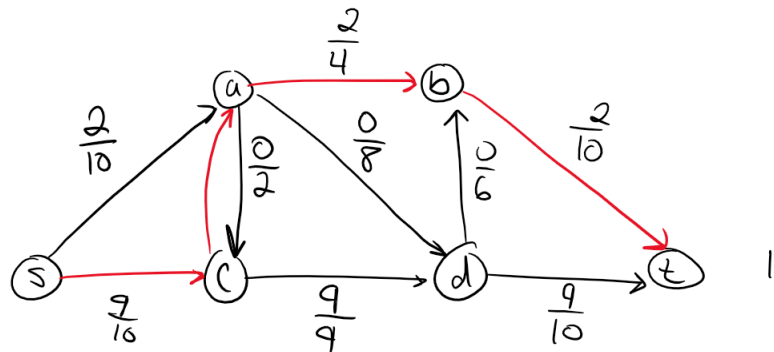
Path 1:



Path 2:



Path 3:



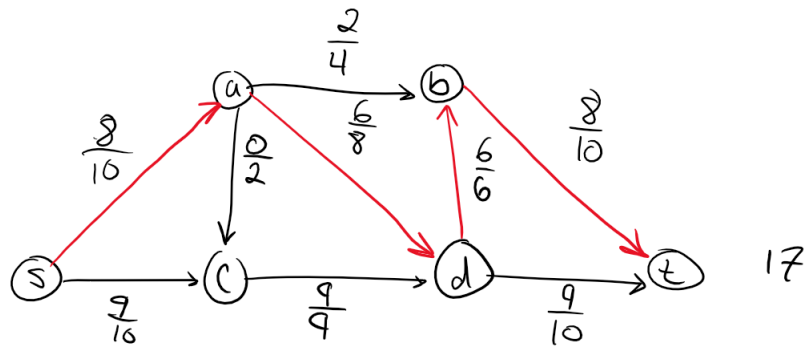
Path 4:

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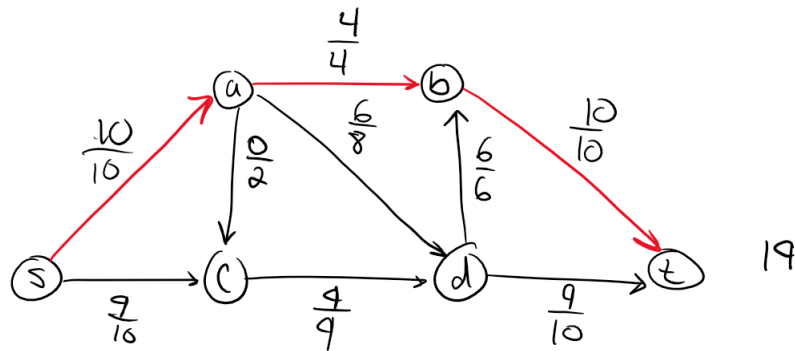
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Path 5:



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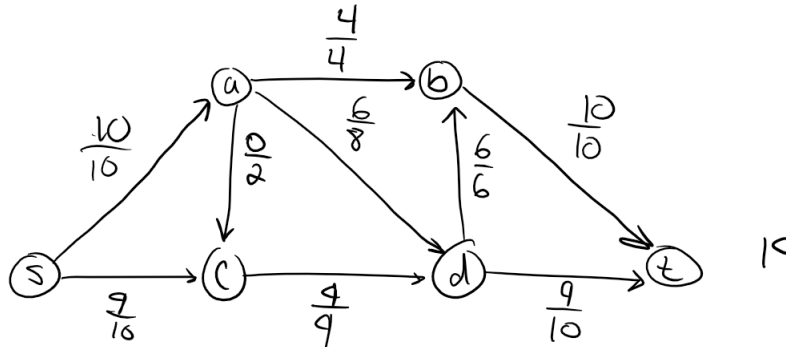
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- (b) (3 pts) Show the final flow  $f(e)$  for the edges of the original graph when the Ford-Fulkerson algorithm terminates.

*Solution.* Final Flow:



- (c) (4 pts) Find the minimum capacity cut with respect to the capacities on the original graph. Is this minimum capacity equal to the Max-Flow that you earlier identified? Justify your answer in a sentence. Also, report the crossing edges in this cut that are saturated (can't carry any more flow).

*Solution.* Min Cut:

The min cut would divide the graph into two disjoint sets:  $A = \{s\}$ ,  $B = \{a, b, c, d, t\}$ . This is the min cut because there is no other cut you can make in the graph, other than the same cut to separate  $t$ , that is less than 20. In this min cut, the edge  $s-a$  is saturated.

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2. (10 pts) Let  $(X, Y)$  be any s-t cut in the network  $G$  and  $a$  be any flow.

- (a) (5 pts) Prove that the value of the flow  $a$  equals the **net** flow that crosses the cut  $(X, Y)$ .

$$\text{i.e. } \text{value}(a) = \sum_{e \text{ out of } X} a(e) - \sum_{e \text{ in to } X} a(e)$$

You should use the flow conservation property to complete the proof.

**Hint :** Recollect the definition of a flow  $a$ .

$$\text{value}(a) = \sum_{e \text{ out of } s} a(e) - \sum_{e \text{ in to } s} a(e) \text{ where } s \text{ is the source.}$$

*Solution.* By the flow conservation property,

$$\sum_{e \text{ out of } V} f(e) = \sum_{e \text{ out of } V} f(e) \text{ , where } V \text{ is some vertex in } G$$

Since this property is an essential property of a valid flow network, it must be the case that  $\forall V \in (X, Y) \sum_{e \text{ out of } V} f(e)$  must be equal to  $\sum_{e \text{ into } V} f(e)$  on the other side of the cut. If this was not the case, it would not be a valid flow network.

- (b) (5 pts) Use the above proof (from part Q3a) to prove that the value of the flow  $a \leq \text{Capacity of the cut } (X, Y)$ .

*Solution.* Another fundamental property of a flow network is the fact that a flow cannot exceed the capacity of an edge. If the flow  $a$  exceeds that capacity of the cut, then it would be violating the rule that was just proved above. If

$$\sum_{e \text{ out of } V} f(e) > \sum_{\text{all edges in cut}} f(e)$$

then the the flow out of the cut must be equal to the flow into the cut by the flow conservation property. If this is the case, then it is violating an essential rule of flow networks. Either there would be vertices that are retaining flow, or the capacity would be broken, both of which constitute an invalid flow.

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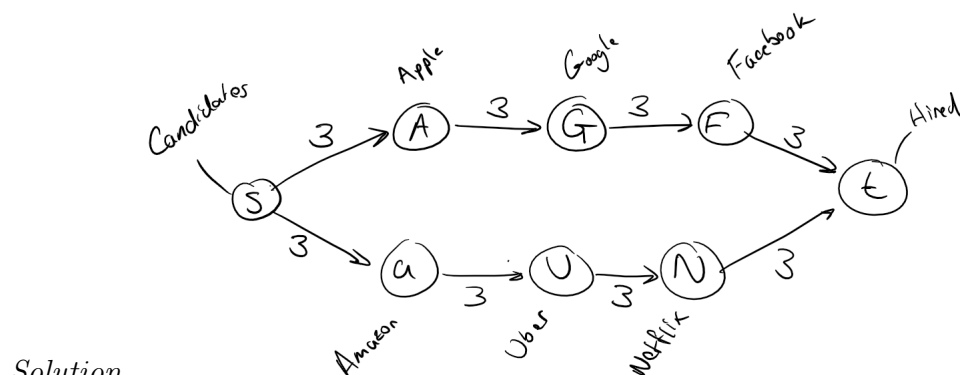
3. (11 pts) CU is organising a spot job fair where many companies participate and take tests to select students. After their final round of interviews, they all find a preference list of candidates that they would like to hire. All the companies just want to hire one student (because recession). All the companies sat together and they realised that if they extend offers to the same students, only one of them would get the student so they decide to run an algorithm to hire the maximum number of students they can together.

Example - Following is one such preference of each company after the final round of interviews. If they all give the offer to just their first preference only 3 students will get hired. But a better offer is Apple - Alice, Google - Dave, Facebook - Carol, Amazon - Eliza, Uber - Frank, Netflix - Bob and this gets 6 students hired.

Help them come up with an algorithm to find an offer set that gets the maximum students the job using Ford-Fulkerson.

Preference	Apple	Google	Facebook	Amazon	Uber	Netflix
1	Alice	Alice	Alice	Carol	Carol	Bob
2	Bob	Bob	Carol	Eliza	Eliza	Eliza
3	Dave	Dave	Frank		Frank	Frank

- (a) (6 pts) Draw a network  $G$  to represent this problem as a flow maximisation problem for the example given above. Clearly indicate the source, the edge directions, the sink and the capacities and label the vertices.



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- (b) (5 pts) Assume that you have access to Ford-Fulkerson sub-routine called **Ford-Fulkerson(G)** that takes a network and gives out max-flow in terms of  $f(e)$  for all the edges. How will you use this sub-routine to find the offer set that employs the maximum number of students. Clearly explain your solution.

*Solution.* I would use this sub-routine to get an idea of how many students can be pushed through a channel from source to sink. In other words, I would use this to find the bottlenecking value of the graph which will help with pushing flow through the graph. This would help the algorithm know how many companies already have suitable matches.