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CSCI 3104, Algorithms
Problem Set 1a (10 points)

Profs. Hoenigman & Agrawal
Fall 2019, CU-Boulder

Advice 1: For every problem in this class, you must justify your answer: show how you arrived at it and why it is correct. If there are assumptions you need to make along the way, state those clearly.

Advice 2: Verbal reasoning is typically insufficient for full credit. Instead, write a logical argument, in the style of a mathematical proof.

Instructions for submitting your solution:

- The solutions **should be typed** and we cannot accept hand-written solutions. Here's a short intro to Latex.
 - You should submit your work through **Gradescope** only.
 - If you don't have an account on it, sign up for one using your CU email. You should have gotten an email to sign up. If your name based CU email doesn't work, try the identikey@colorado.edu version.
 - Gradescope will only accept **.pdf** files (except for code files that should be submitted separately on Gradescope if a problem set has them) and **try to fit your work in the box provided**.
 - You cannot submit a pdf which has less pages than what we provided you as Gradescope won't allow it.
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1. (3 pts) *What are the three components of a loop invariant proof? Write a one-sentence description for each one.*

Initialization - Loop invariant true prior to first iteration.

Maintenance - Loop invariant is true immediately before or after the loop body executes.

Termination - Loop invariant is true after termination of the loop.

2. (6 pts total) *Identify the loop invariant in the following algorithms.*

(a) FindMaxElement(A) : //suppose array A is not empty
 ret = A[0]
 for i = 1 to length(A)-1 {
 if A[i] > ret{
 ret = A[i]
 }
 }
 return ret

ret stores the largest A[i] at the start of the ith iteration.

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(b) FindElement(A, n) : //suppose no duplicates in array A and array A is not empty
 ret = -1 //index -1 implies the element haven't been found yet
 for i = 0 to length(A)-1 {
 if A[i] == n{
 ret = i
 }
 }
 return ret

ret stores either -1 or i at the start of the i^{th} iteration. Where i is the index of the found number.

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(c) SumArray(A) : //suppose array A is not empty
    sum = 0
    for i = 0 to length(A)-1 {
        sum += A[i]
    }
    return sum
```

sum is storing $A[0]+A[1]+A[2]+ \dots + A[i]$ at the start of the i^{th} iteration.

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3. (1 pt) If r is a real number not equal to 1, then for every $n \geq 0$,

$$\sum_{i=0}^n r^i = \frac{(1 - r^{n+1})}{(1 - r)}.$$

Provide the first two steps of a proof by induction i.e. base case and the inductive hypothesis. You will be asked to complete this proof later in **PS1b**.

Base Case - $i = 0$: $\sum_{i=0}^0 r^0 = 1 = \frac{(1-r^{0+1})}{1-r} = 1$ and thus proves our base case.

Inductive Hypothesis - Assume that for some integer k , $\sum_{i=0}^k r^i = \frac{(1-r^{k+1})}{(1-r)}$