## KVik metrics

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t := | \operatorname{Hot} \ \operatorname{Files} | / (| \operatorname{Hot} \ \operatorname{Files} | + | \operatorname{Data} \ \operatorname{File} |)
S_k - \operatorname{size} \ \operatorname{of} \ \operatorname{key} \ (\operatorname{in} \ \operatorname{bytes})
S_v - \operatorname{size} \ \operatorname{of} \ \operatorname{value} \ (\operatorname{in} \ \operatorname{bytes})
N - \operatorname{number} \ \operatorname{of} \ \operatorname{records} \ \operatorname{in} \ \operatorname{KVik}
M - \operatorname{number} \ \operatorname{of} \ \operatorname{keys} \ \operatorname{Log} \ \operatorname{Hash} \ \operatorname{Table} \ \operatorname{can} \ \operatorname{hold} \ (\operatorname{capacity})
f - \operatorname{size} \ \operatorname{of} \ \operatorname{offset} \ (\operatorname{in} \ \operatorname{bytes})
U_w - \operatorname{speed} \ \operatorname{of} \ \operatorname{write} \ \operatorname{requests}
E - \operatorname{number} \ \operatorname{of} \ \operatorname{write} \ \operatorname{cycles} \ (\operatorname{parameter} \ \operatorname{of} \ \operatorname{your} \ \operatorname{SSD})
F - \operatorname{size} \ \operatorname{of} \ \operatorname{SSD}
w := S_k + 2 \cdot f
c - \operatorname{Log} \ \operatorname{Hash} \ \operatorname{Table} \ \operatorname{fill} \ \operatorname{factor} \ (\operatorname{by} \ \operatorname{default}, \ \operatorname{it} \ \operatorname{is} \ \operatorname{assumed} \ \operatorname{to} \ \operatorname{be} \ 0.9)
s - \operatorname{Index} \ \operatorname{Hash} \ \operatorname{Table} \ \operatorname{fill} \ \operatorname{factor} \ (\operatorname{by} \ \operatorname{default}, \ \operatorname{it} \ \operatorname{is} \ \operatorname{assumed} \ \operatorname{to} \ \operatorname{be} \ 0.6)
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Assumption: key and value can be placed on one memory page

$$RAF = c \cdot M/N + (N - c \cdot M)/N \cdot (1 + 1) = 2 - c \cdot M/N$$

Assuming, requests are distributed uniformly and will hit Log Hash Table with probability  $c \cdot M/N$ .

$$\text{WAF} = 1 + 1/(c \cdot M) \cdot (2 \cdot c \cdot M) + 1/(N \cdot t) \cdot (2 \cdot N \cdot t) = 5$$

First term is write operation to Log File that happens every time. Second term stands for situations when Log Hash Table is filled and needs to be merged into Index Hash Table (in fact it should be better because before performing modification of Index Hash Table we presort all keys by hash function used in there, so ones that hit same bucket will be written using only one write operation). Since it happens once in  $c \cdot M$  operations we divide write operations it takes on it to take average. Each key will have two write operations (one for key written to Index Hash Table and one for value written to Hot Files). The third term is really similar to the second one but there we need one write operation to write value to Data File and one more write operation to update offset value in Index Hash Table (actually in case of update it can be done inplace so offset can remain the same but in general case it is so).

$$\begin{split} \text{SAF} &= (c \cdot M \cdot (S_k + S_v) + t \cdot N \cdot (S_v + S_k) + N \cdot S_v + 1.67 \cdot N \cdot w) / (N \cdot (S_v + S_k)) = \\ &= c \cdot M / N + t + 1 - S_k / (S_v + S_k) + 1.67 \cdot (S_k + 2f) / (S_v + S_k) = \\ &= t + 1 + c \cdot M / N + 0.67 \cdot S_k / (S_v + S_k) + 3.4 \cdot f / (S_v + S_k) \end{split}$$

Here everything is very simple: the first term in numerator is Log File, the second is Hot Files, the third is Data File, the fourth is Index Hash Table (fill factor was determined experimentally and  $0.6^{-1} = 1.(6)$ ).

$$\mathrm{MO} = \tfrac{5 \cdot N/8 + M \cdot (S_k + f)}{N} = 5/8 + M/N \cdot (S_k + f)$$

The first term in numerator is memory that we need to store Filter, the second term is Log Hash Table.