

## IV. Constructing the Real Numbers

UBC M320 Lecture Notes by Philip D. Loewen

11 Oct 2023

Introduce the following notation:

- $CS(\mathbb{Q})$ : the set of all Cauchy sequences with rational elements.  
 $x, y, z$ : typical symbols for elements of  $CS(\mathbb{Q})$ . Thus, e.g.,  $x = (x_1, x_2, \dots)$ .  
 $R[x]$ : the subset of  $CS(\mathbb{Q})$  associated with a given  $x \in CS(\mathbb{Q})$  as follows:

$$R[x] = \left\{ x' \in CS(\mathbb{Q}) : \lim_{n \rightarrow \infty} |x'_n - x_n| = 0 \right\}.$$

- $\Phi$ : the function that takes each rational number  $q$  into the subset of  $CS(\mathbb{Q})$  containing the corresponding constant sequence, i.e.,

$$\Phi(q) = R[(q, q, \dots)] \quad \forall q \in \mathbb{Q}.$$

Let  $\mathcal{R} = \{R[x] : x \in CS(\mathbb{Q})\}$ . This will be our model for  $\mathbb{R}$ . Note that each set  $R[x]$  contains *infinitely many Cauchy sequences*.

### A. Equality

For  $x, x' \in CS(\mathbb{Q})$ , define “ $x \sim x'$ ” to mean  $(x'_n - x_n) \rightarrow 0$ , i.e.,

$$\forall \varepsilon > 0 \ (\varepsilon \in \mathbb{Q}), \ \exists N \in \mathbb{N} : \forall n \geq N, \ |x'_n - x_n| < \varepsilon.$$

This is an equivalence relation ( $x \sim x$ ;  $x \sim y$  iff  $y \sim x$ ;  $x \sim y$  and  $y \sim z$  imply  $x \sim z$ ). For each  $x \in CS(\mathbb{Q})$ ,  $R[x]$  is the equivalence class. This means that  $R[x] = \{x' \in CS(\mathbb{Q}) : x' \sim x\}$ , and it implies (see HW05) that, for any  $x, y \in CS(\mathbb{Q})$ ,

$$\begin{aligned} x &\in R[x], \quad \text{and} \\ R[x] \cap R[y] \neq \emptyset &\iff R[x] = R[y]. \end{aligned}$$

### B. Addition

Extend “+” to  $\mathcal{R}$  by first extending it to  $CS(\mathbb{Q})$ . Let’s say that for  $x = (x_n)$  and  $y = (y_n)$  in  $CS(\mathbb{Q})$ ,

$$x + y = (x_1 + y_1, x_2 + y_2, \dots).$$

The result belongs to  $CS(\mathbb{Q})$  [HW04 Q6(a)].

**Proposition.** *The following binary operation on  $\mathcal{R}$  is well-defined:*

$$R[x] + R[y] = R[x + y], \quad x, y \in CS(\mathbb{Q}).$$

*That is, whenever  $R[x'] = R[x]$  and  $R[y'] = R[y]$ , one has  $R[x' + y'] = R[x + y]$ .*

*Proof.* For  $x, x', y, y'$  as in the setup, we have  $x' \in R[x]$  and  $y' \in R[y]$ , so

$$x'_n - x_n \rightarrow 0, \quad y'_n - y_n \rightarrow 0.$$

The sum rule for limits implies that

$$(x'_n + y'_n) - (x_n + y_n) \rightarrow 0.$$

This shows that  $(x' + y') \in R[x + y]$ . But of course  $(x' + y') \in R[x' + y']$  too, so  $R[x + y] \cap R[x' + y'] \neq \emptyset$ . Therefore  $R[x + y] = R[x' + y']$ .  
 ////

12 Oct 2023

**Proposition.**  $(\mathcal{R}, +)$  is an Abelian group, i.e., for ever  $x, y, z \in CS(\mathbb{Q})$ ,

- (A1)  $R[x] + R[y]$  is a well-defined element of  $\mathcal{R}$ .
- (A2)  $R[x] + R[y] = R[y] + R[x]$ .
- (A3)  $(R[x] + R[y]) + R[z] = R[x] + (R[y] + R[z])$ .
- (A4) The element  $\Phi(0) = R[(0, 0, \dots)]$  obeys  $\Phi(0) + R[x] = R[x]$ .
- (A5)  $\mathcal{R}$  contains a set, denoted “ $-R[x]$ ”, such that

$$R[x] + (-R[x]) = \Phi(0).$$

*Proof.* (A1) This takes some thinking. We did it in the previous proposition.

- (A2) For the underlying Cauchy sequences  $x, y \in CS(\mathbb{Q})$ ,  $x + y = y + x$  follows from the properties of addition for rational numbers. (Apply these component-by-component.)
- (A3) For the underlying Cauchy sequences  $x, y, z \in CS(\mathbb{Q})$ ,  $(x + y) + z = (x + y) + z$  follows from the properties of addition for rational numbers. (Apply these component-by-component.)
- (A4) For the underlying Cauchy sequences  $x \in CS(\mathbb{Q})$  and  $(0, 0, 0, \dots)$ ,  $x + (0, 0, \dots) = x$  follows from the properties of addition for rational numbers. (Apply these component-by-component.)
- (A5) Given any  $x \in CS(\mathbb{Q})$ , invent the related sequence “ $-x$ ” by defining

$$-x = (-x_1, -x_2, -x_3, \dots).$$

Clearly  $-x \in CS(\mathbb{Q})$ , and  $x + (-x) = (0, 0, \dots)$  in  $CS(\mathbb{Q})$ . Therefore  $R[x] + R[-x] = R[(0, 0, \dots)] = \Phi(0)$ . So the set  $R[-x]$  (which clearly lies in  $\mathcal{R}$ ) can be used to define  $-R[x]$ .  
 ////

## C. Multiplication

Extend “ $\cdot$ ” to  $\mathcal{R}$  by first extending it to  $CS(\mathbb{Q})$ . Let’s say that if  $x = (x_n)$  and  $y = (y_n)$  then

$$x \cdot y = (x_1 y_1, x_2 y_2, \dots).$$

The result belongs to  $CS(\mathbb{Q})$  [HW04 Q6(b)].

**Proposition.** *The following binary operation on  $\mathcal{R}$  is well-defined:*

$$R[x] \cdot R[y] = R[x \cdot y], \quad x, y \in CS(\mathbb{Q}).$$

*That is, whenever  $R[x'] = R[x]$  and  $R[y'] = R[y]$ , one has  $R[x' \cdot y'] = R[x \cdot y]$ .*

*Proof.* For  $x, x', y, y'$  as in the setup, we have  $x' \in R[x]$  and  $y' \in R[y]$ , so

$$x'_n - x_n \rightarrow 0, \quad y'_n - y_n \rightarrow 0.$$

Let's rearrange

$$x'_n y'_n - x_n y_n = [(x'_n - x_n) + x_n] y'_n - x_n y_n = [x'_n - x_n] y'_n + x_n [y'_n - y_n].$$

Every Cauchy sequence is bounded, so there are numbers  $M_{y'}$  and  $M_x$  such that

$$|x'_n y'_n - x_n y_n| \leq M_{y'} |x'_n - x_n| M_x |y'_n - y_n|, \quad \forall n \in \mathbb{N}.$$

The right side converges to 0, so the left side does too. This shows that  $(x' \cdot y') \in R[x \cdot y]$ . Of course  $(x' \cdot y') \in R[(x' \cdot y')]$  also, so  $R[x \cdot y] \cap R[x' \cdot y'] \neq \emptyset$ , and this gives  $R[x \cdot y] = R[x' \cdot y']$ . ////

**Proposition.**  *$(\mathcal{R} \setminus \{\Phi(0)\}, \cdot)$  is an Abelian group, i.e., for any  $x, y, z \in CS(\mathbb{Q})$  such that  $\Phi(0) \notin \{R[x], R[y], R[z]\}$ , ...*

(M1)  $R[x] \cdot R[y]$  is a well-defined element of  $\mathcal{R}$ .

(M2)  $R[x] \cdot R[y] = R[y] \cdot R[x]$ .

(M3)  $(R[x] \cdot R[y]) \cdot R[z] = R[x] \cdot (R[y] \cdot R[z])$ .

(M4) The element  $\Phi(1) = R[(1, 1, \dots)]$  obeys  $\Phi(1) \neq \Phi(0)$  and  $\Phi(1) \cdot R[x] = R[x]$ .

(M5)  $\mathcal{R}$  contains a set, denoted “ $1/R[x]$ ”, such that

$$R[x] \cdot (1/R[x]) = \Phi(1).$$

*Proof.* (M1) This takes some thinking. We did it in the previous proposition.

(M2) For the underlying Cauchy sequences  $x, y \in CS(\mathbb{Q})$ ,  $x \cdot y = y \cdot x$  follows from the properties of multiplication for rational numbers. (Apply these component-by-component.)

(M3) For the representative Cauchy sequences  $x, y, z \in CS(\mathbb{Q})$ ,  $(x \cdot y) \cdot z = (x \cdot y) \cdot z$  follows from the properties of multiplication for rational numbers. (Apply these component-by-component.)

(M4) For the representative Cauchy sequences  $x \in CS(\mathbb{Q})$  and  $(1, 1, 1, \dots)$ , the identity  $x \cdot (1, 1, \dots) = x$  follows from the properties of multiplication for rational numbers. (Apply these component-by-component.)

(M5) This takes some thinking. It's HW05 Question 6. ////

Often we simplify notation by assuming the “ $\cdot$ ” between elements of  $\mathcal{R}$  written side-by-side.

## D. Distribution

**Proposition.** For any  $x, y, z \in CS(\mathbb{Q})$ , we have

$$R[x] \cdot (R[y] + R[z]) = (R[x] \cdot R[y]) + (R[x] \cdot R[z]).$$

*Proof.* Carefully apply the definitions for  $\cdot$  and  $+$  to condense both sides:

$$\begin{aligned} R[x] \cdot (R[y] + R[z]) &= R[x] \cdot R[y + z] = R[x \cdot (y + z)], \\ (R[x] \cdot R[y]) + (R[x] \cdot R[z]) &= R[x \cdot y] + R[x \cdot z] = R[x \cdot y + x \cdot z]. \end{aligned}$$

So a Cauchy sequence belonging to the set on the left has terms  $x_n(y_n + z_n)$ . These terms, taken one by one, equal  $(x_n y_n + x_n z_n)$ , so this same Cauchy sequence belongs to  $R[x \cdot y + x \cdot z]$ . Therefore the sets on the left and right in the desired identity have an element (which is a Cauchy sequence) in common. Therefore the sets must be identical. (HW05 Q3) ////

## E. Order

**Definition.** For  $x, y$  in  $\mathcal{R}$ , say  $R[x] < R[y]$  if and only if

$$\exists r > 0 \ (r \in \mathbb{Q}), \ \exists N \in \mathbb{N} : \forall n > N, \ x_n + r < y_n.$$

**Proposition.** The relation “ $<$ ” is well-defined on  $\mathcal{R}$ . Also, ...

(i) For every  $x, y \in CS(\mathbb{Q})$ , exactly one of the following holds:

$$R[x] < R[y], \quad R[x] = R[y], \quad R[y] < R[x].$$

(ii) Whenever  $x, y, z \in CS(\mathbb{Q})$ , one has

$$(R[x] < R[y] \text{ and } R[y] < R[z]) \implies R[x] < R[z].$$

*Proof.* HW05, Problems 4–5. ////

13 Oct 2023

**Proposition OC (Order Components).** Let  $a, b \in CS(\mathbb{Q})$ .

- (a) If  $R[a] > R[b]$ , then there exists  $N \in \mathbb{N}$  such that  $R[a] > \Phi(b_N)$ .
- (b) If  $R[a] \leq \Phi(b_k)$  for all  $k \in \mathbb{N}$ , then  $R[a] \leq R[b]$ .

*Proof.* It suffices to prove (a), since (b) is the contrapositive.

If  $R[a] > R[b]$ , then there exist  $N_0 \in \mathbb{N}$  and  $r > 0$  in  $\mathbb{Q}$  such that

$$a_n > b_n + r \quad \forall n \geq N_0.$$

Also, the sequences  $a$  and  $b$  are Cauchy, so there exist  $N_a, N_b \in \mathbb{N}$  such that

$$\begin{aligned} |a_m - a_n| &< \frac{r}{3} & \forall n \geq N_a, \\ |b_m - b_n| &< \frac{r}{3} & \forall n \geq N_b. \end{aligned}$$

Let  $N = \max \{N_0, N_a, N_b\}$ : for every  $m \geq N$ ,

$$b_N = b_m + (b_N - b_m) < b_m + \frac{r}{3} < (a_m - r) + \frac{r}{3} = a_m - \frac{2r}{3}.$$

This shows  $\Phi(b_N) < R[a]$ , as required. ////

**Pause for Reflection.** Note that  $\Phi: \mathbb{Q} \rightarrow \mathcal{R}$  provides a working copy of  $\mathbb{Q}$  in the set of equivalence classes of rational-valued Cauchy sequences. For any rational numbers  $p, q$ , the corresponding constant sequences satisfy (among other things)

$$\begin{aligned} \Phi(p + q) &= \Phi(p) + \Phi(q), & \Phi(-q) &= -\Phi(q), \\ \Phi(pq) &= \Phi(p)\Phi(q), & \Phi(1/p) &= 1/\Phi(p) \text{ [for } p \neq 0], \\ p < q &\iff \Phi(p) < \Phi(q). \end{aligned}$$

The promise and potential of this whole project is that the set  $\mathcal{R}$  contains an equivalence class for *every* Cauchy sequence with rational entries, not only the constant sequences. This is a definite step forward.

## F. Completeness

As suggested above, the embedding  $\Phi: \mathbb{Q} \rightarrow \mathcal{R}$  falls far short of being surjective. Many of the equivalence classes in  $\mathcal{R}$  contain no constant sequence at all. Indeed, there are enough additional equivalence classes in  $\mathcal{R} \setminus \Phi(\mathbb{Q})$  to make  $\mathcal{R}$  order-complete.

**Notation Shift.** Let's start using single Greek letters for the elements of  $\mathcal{R}$ , with typical correspondences like  $\alpha = R[a]$ ,  $\beta = R[b]$ , ..., for  $a, b, \dots \in CS(\mathbb{Q})$ .

**Theorem.** *Let  $A \neq \emptyset$  be a subset of  $\mathcal{R}$  with an upper bound, i.e., suppose there exists  $\mu \in \mathcal{R}$  such that*

$$\forall \alpha \in A, \quad \alpha \leq \mu.$$

*Then there exists an element  $\beta \in \mathcal{R}$  such that*

- (i)  $\forall \alpha \in A, \alpha \leq \beta$  (i.e.,  $\beta$  is an upper bound for  $A$ ), and
- (ii) for each  $\gamma \in \mathcal{R}$  with  $\gamma < \beta$ , there exists  $\alpha \in A$  such that  $\alpha > \gamma$  (i.e.,  $\beta$  is the least upper bound for  $A$ ).

*Proof.* For each  $n$ , consider the rational number

$$b_n = \min(S_n), \quad \text{where} \quad S_n = \left\{ \frac{k}{2^n} : k \in \mathbb{Z}, \alpha \leq \Phi\left(\frac{k}{2^n}\right) \forall \alpha \in A \right\}.$$

To see that the minimum on the right makes sense, start with the assumption that  $\mu \in \mathcal{R}$  is an upper bound for the set  $A$ . According to HW05 Q5(b), there must be some rational number  $r$  such that  $\mu < \Phi(r)$ , and this implies that the set of upper bounds for  $A$  in  $\mathcal{R}$  includes  $\Phi(s)$  for every rational  $s \geq r$ . So indeed the set  $S_n$  is not empty. On the other hand, for any  $\alpha \in A$  (note  $A \neq \emptyset$ ), HW05 Q5(b) provides a rational number  $q$  such that  $\Phi(q) < \alpha$ . This  $q$  provides a lower bound for  $S_n$ . As a nonempty discrete set with a lower bound,  $S_n$  must have a smallest element.

For each  $n$ , the definition of  $b_n$  implies two things:

$$\forall \alpha \in A, \alpha \leq \Phi(b_n) \quad \text{and} \quad \exists \alpha_n \in A : \alpha_n > \Phi\left(b_n - \frac{1}{2^n}\right). \quad (*)$$

Further, since  $S_n \subseteq S_{n+1}$  for all  $n$ , we have

$$b_n \geq b_{n+1} \geq b_n - \frac{1}{2^{n+1}}, \quad \text{so} \quad 0 \leq b_n - b_{n+1} \leq \frac{1}{2^{n+1}}.$$

Iterating this leads to an inequality valid for all  $p, n \in \mathbb{N}$ :

$$\begin{aligned} 0 \leq b_n - b_{n+p} &= (b_n - b_{n+1}) + (b_{n+1} - b_{n+2}) + \dots + (b_{n+p-1} - b_{n+p}) \\ &\leq \frac{1}{2^{n+1}} + \frac{1}{2^{n+2}} + \dots + \frac{1}{2^{n+p}} < \frac{1}{2^n}. \end{aligned}$$

It follows that the sequence  $b = (b_1, b_2, b_3, \dots)$  is Cauchy. (Given  $\varepsilon > 0$ , choose  $N \in \mathbb{N}$  so that  $2^{-N} < \varepsilon$ . Then any  $m \geq n \geq N$  make  $|b_m - b_n| = b_n - b_m < 2^{-n} < \varepsilon$ .)

Let  $\beta = R[b]$ . Recall from (\*) that every  $\alpha \in A$  obeys  $\alpha \leq \Phi(b_n)$  for every  $n \in \mathbb{N}$ . According to Proposition OC above, it follows that  $\alpha \leq R[b] = \beta$ . This shows (i).

To confirm (ii), fix any  $\gamma = R[c]$  such that  $\gamma < \beta$ . There must be some rational  $r > 0$  and index  $N \in \mathbb{N}$  such that  $c_n + r < b_n$  for all  $n > N$ . By increasing  $N$  if necessary, we may assume also that  $1/2^N < r/2$ . Recalling that the sequence  $(b_n)$  is nonincreasing, we have

$$c_n + \frac{r}{2} = c_n + r - \frac{r}{2} < b_n - \frac{r}{2} < b_n - \frac{1}{2^N} \quad \forall n > N.$$

This element-by-element inequality gives  $R[c] < \Phi\left(b_N - \frac{1}{2^N}\right)$ . The second assertion in (\*) above (use  $n = N$  and transitivity of  $>$ ) now establishes (ii). /////

Our construction is finished. The set  $\mathcal{R}$  with the operations of  $+$  and  $\cdot$  and  $<$  is an ordered field, in the sense defined by Rudin in Definitions 1.5–1.8 and 1.12. Further, it is order-complete, in the sense of Rudin's Definition 1.10. This proves Rudin's Theorem 1.19.

## G. Consequences

**Theorem (Archimedes).** *The set  $\mathbb{N}$  has no upper bound in  $\mathbb{R}$ . In particular, for every  $x \in \mathbb{R}$  there exists  $n \in \mathbb{N}$  satisfying  $n > x$ .*

*Proof.* Suppose, seeking a contradiction, that  $\mathbb{N}$  has an upper bound. Let  $\alpha \stackrel{\text{def}}{=} \sup(\mathbb{N})$ . Then  $\alpha - \frac{1}{2}$  is not an upper bound for  $\mathbb{N}$ , so there exists  $n \in \mathbb{N}$  such that  $n > \alpha - \frac{1}{2}$ . But then  $n + 1$  is an integer satisfying  $n + 1 > \alpha + \frac{1}{2} > \alpha$ , and this contradicts the defining property of  $\alpha$ .  
 ////

*Remark.* The Archimedean property of  $\mathbb{Q}$  was never in doubt. But not every ordered field containing a work-alike copy of  $\mathbb{Q}$  inherits the property. See Rechnitzer's M319 notes for a reasonably simple example in which the elements are rational functions with rational coefficients. (<https://arechnitzer.gitlab.io/m319text/>)