6.033 Spring 2017Lecture #6

- Monolithic kernels vs. Microkernels
- Virtual Machines

Enforcing Modularity via Virtualization

in order to enforce modularity + build an effective operating system

 programs shouldn't be able to refer to (and corrupt) each others' memory

virtual memory

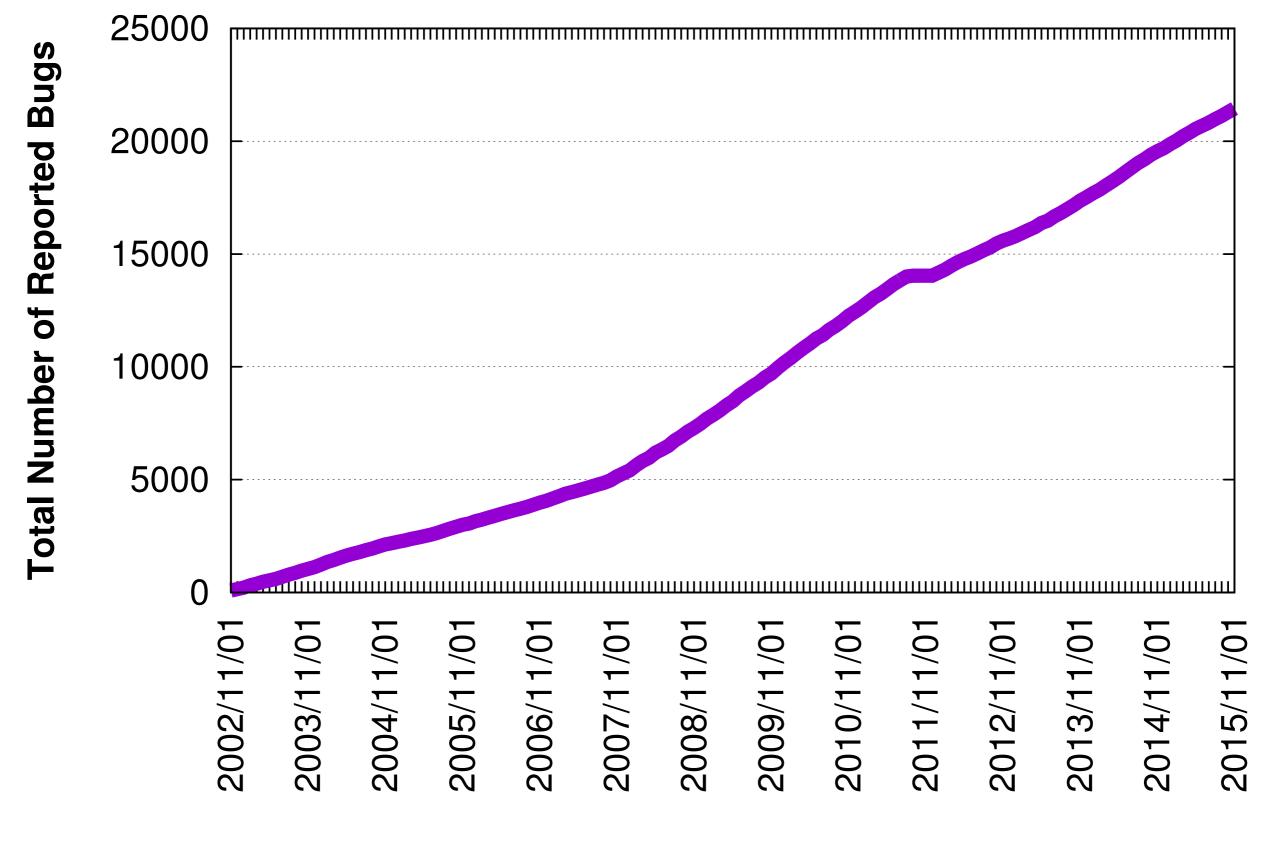
programs should be able to communicate

bounded buffers
(virtualize communication links)

3. programs should be able to **share a CPU** without one program halting the progress of the others

threads (virtualize processors)

today: can we rely on the kernel to work properly?



Date

source: bugzilla.kernel.org, count of all bugs currently marked NEW, ASSIGNED, REOPENED, RESOLVED, VERIFIED, or CLOSED, by creation date

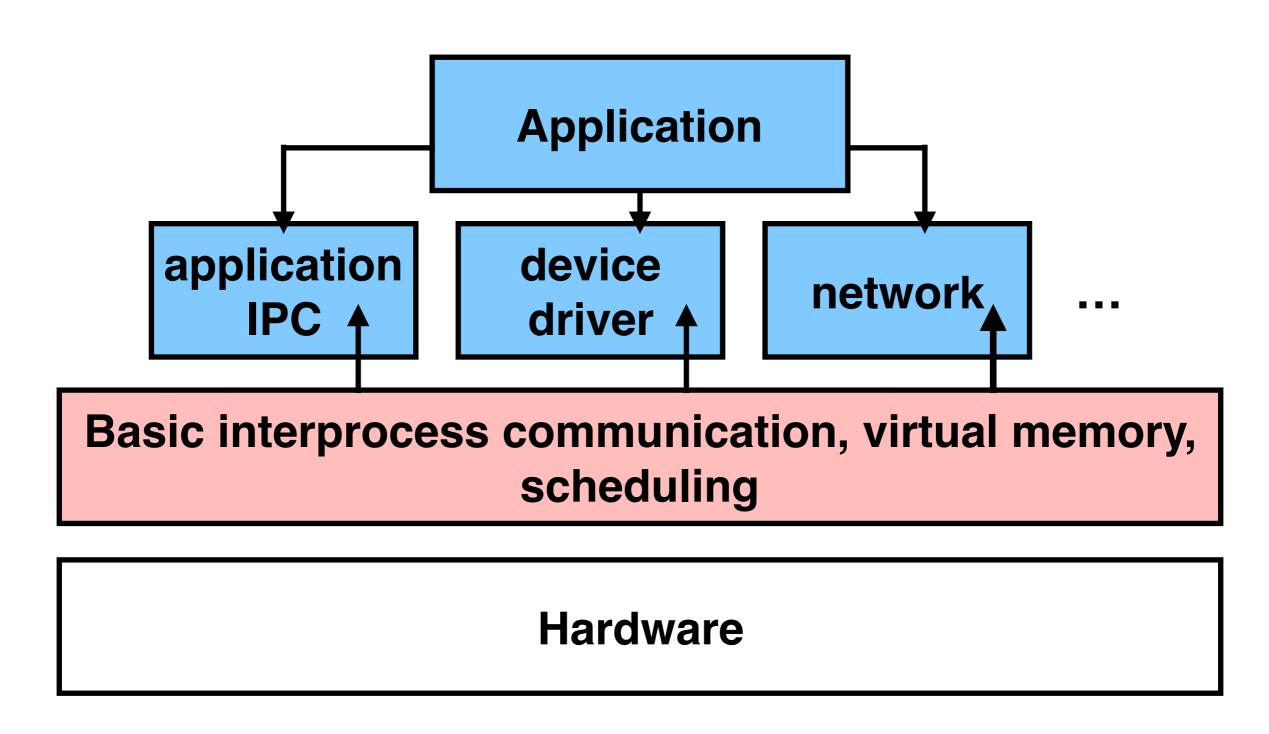
monolithic kernels: no enforced modularity within the kernel itself

Application

Basic interprocess communication, virtual memory, scheduling, file server, device drivers, network, ...

Hardware

microkernels: enforce modularity by putting subsystems in user programs



problem: how do we deal with bugs in the Linux kernel without redesigning Linux from scratch?

Virtual Machines

virtual machine running guest OS

virtual machine running guest OS

physical hardware

problem: if guest OSes run in kernel mode, we haven't solved our original program; if they run in user mode, they can't execute privileged instructions

Virtual Machines

VMM runs in kernel-mode on hardware

virtual machine running guest OS

virtual machine running guest OS

virtual machine monitor (VMM)

physical hardware

guest OS

guest OS

virtual hardware

U/K PTR page table

virtual hardware

U/K PTR page table

virtual machine monitor (VMM)

physical hardware

U/K, PTR, page table, ...

VMM's goal: virtualize hardware

guest OS

guest OS

guest virtual address

virtual hardware

```
U/K
PTR
page table
```

virtual hardware

```
U/K
PTR
page table
```

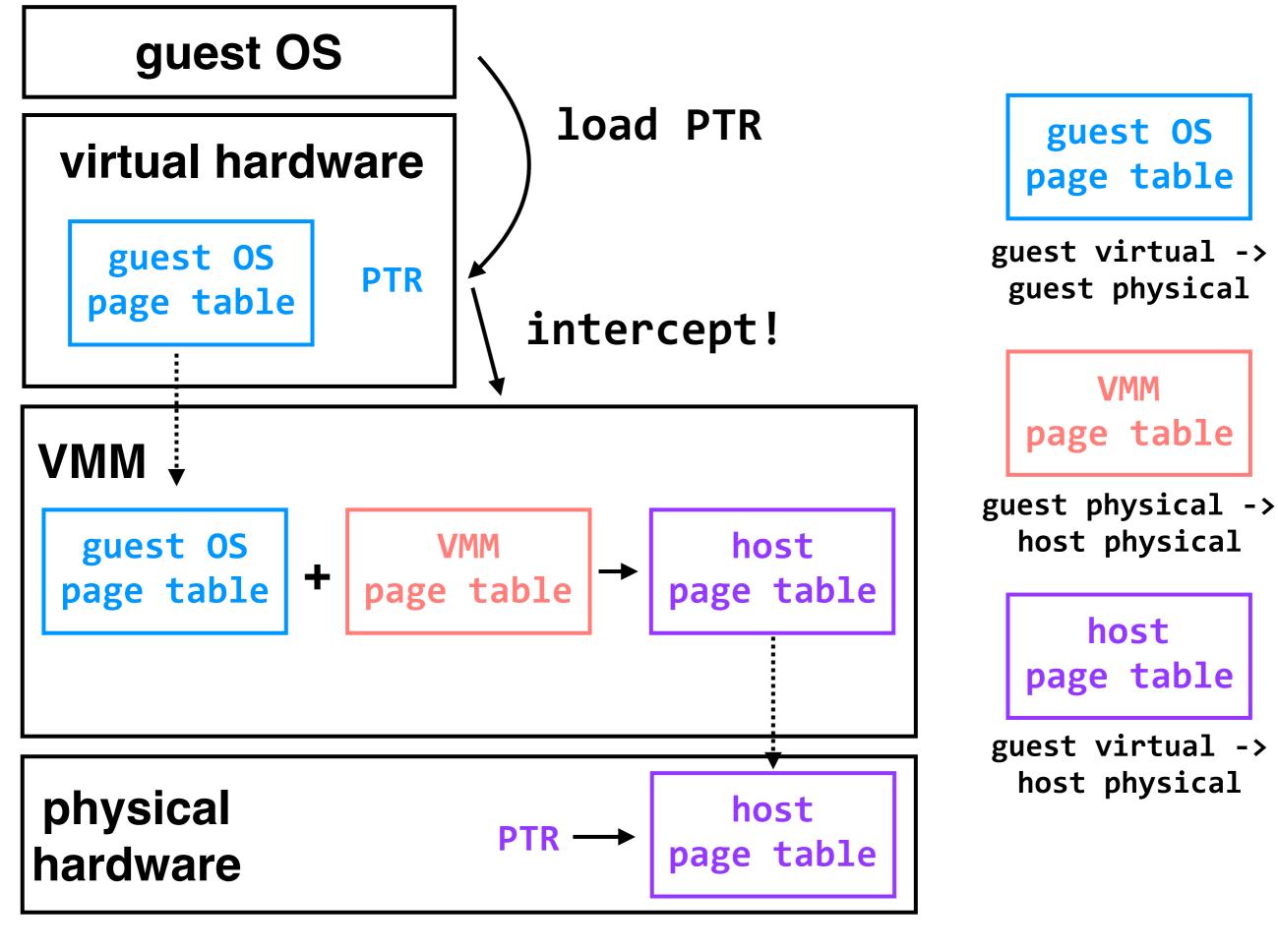
guest physical
 address

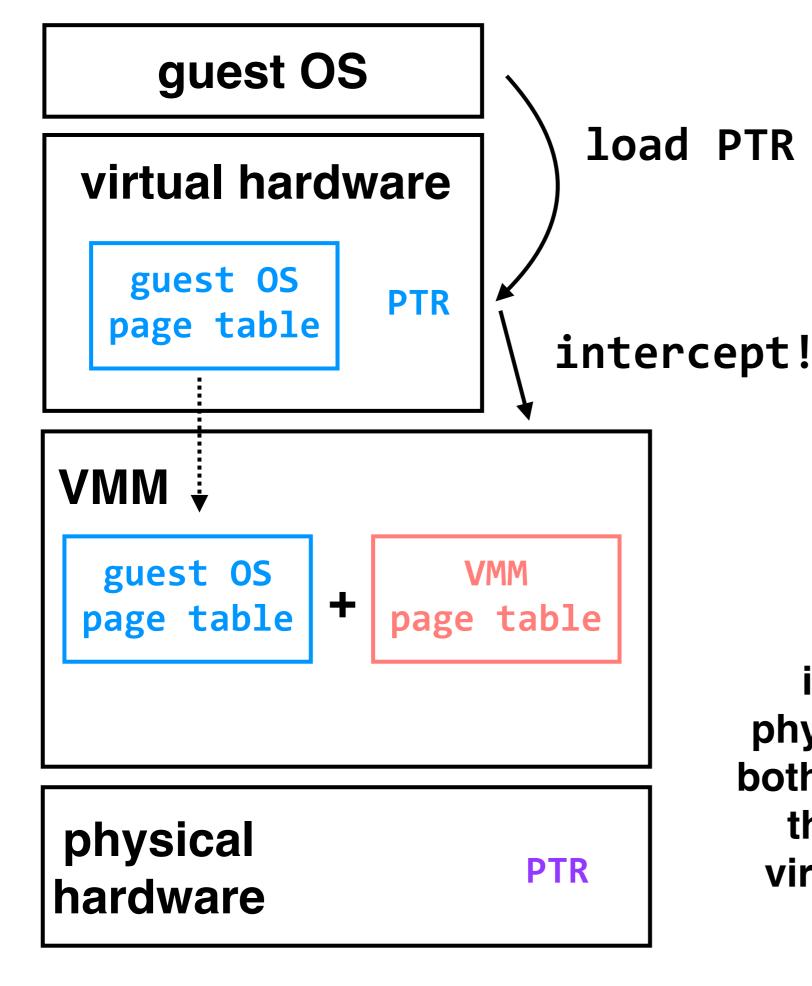
virtual machine monitor (VMM)

physical hardware

U/K, PTR, page table, ...

host physical address





guest OS
page table

guest virtual ->
 guest physical

VMM page table

guest physical ->
 host physical

in modern hardware, the physical hardware is aware of both page tables, and performs the translation from guest virtual to host physical itself

guest OS

guest OS

virtual hardware

```
U/K
PTR
page table
```

virtual hardware

```
U/K
PTR
page table
```

virtual machine monitor (VMM)

physical hardware

```
U/K, PTR, page table, ...
```

VMM's goal: virtualize hardware

Kernel Structure

Monolithic kernels provide no enforced modularity within the kernel. **Microkernels** do, but redesigning monolithic kernels as microkernels is challenging.

Virtual Machines

Virtual machines allow us to run multiple **isolated**OSes on a single physical machine, similar to how we used an OS to run multiple programs on a single CPU.
VMs must handle the challenges of virtualizing the hardware (examples: virtualizing memory, the U/K bit, and disk).