

Benjamin Thompson

CS 312 Section 1

Lab 2

Code commented with complexity:

```
/**
 * Main function for solving the outer hull problem
 *
 */
public void Solve(List<System.Drawing.PointF> pointList)
{
    // Lambda function for sorting by X
    // https://msdn.microsoft.com/en-us/library/b0zbh7b6.aspx --  $O(n \log n)$ 
    // average,  $O(n^2)$  worst case
    pointList.Sort((a, b) => a.X.CompareTo(b.X));

    // Convex Hull -  $O()$ 
    pointList = ConvexHull(ref pointList, 0, pointList.Count-1);

    // Draw Convex Hull
    g.DrawPolygon(new Pen(Color.Red), pointList.ToArray());
}
```

```

/**
 * Convex Hull algorithm. Finds the smallest encompassing hull of a group of
vertices
 *
 * Complexity =>
 *               worst    O(n^2)
 *               avg      O(nlogn)
 *               best     O(nlogn)
 */
public List<PointF> ConvexHull(ref List<System.Drawing.PointF> pointList, int
left, int right)
{
    List<PointF> convex = new List<PointF>();
    List<PointF> convex1 = new List<PointF>();
    // Check if grouping is too big
    if ((right - left) > 2)
    {
        convex = ConvexHull(ref pointList, left, (right - left + 1) / 2 - 1 +
left);
        convex1 = ConvexHull(ref pointList, (right - left + 1) / 2 + left,
right);
        //return convex;
    }

    List<System.Drawing.PointF> points = new List<PointF>();
    for (int i = left; i < right + 1; i++) // O(n)
    {
        points.Add(pointList[i]);
    }

    if (convex.Count > 0)
    {
        convex = Merge(ref convex, ref convex1); // O(n)
    }
    else
    {
        convex = SortCC(points); // O(n^2)
    }
    return convex;
}

```

```

/**
 *
 * Complexity =>
 *          worst    O(n*m) // where m is the number of changes made to the
 *                          tangents (including both up and bottom)
 *          avg      O(n*m)
 *          best     O(n) // Assuming only one change
 */
public List<PointF> Merge (ref List<PointF> c_left, ref List<PointF> c_right)
{
    int right_pos = Rightmost(ref c_left);
    int top_right = 0, top_left = right_pos, bottom_right = 0, bottom_left =
right_pos;
    bool isChanged = true;

    while (isChanged) // Top Tangent // O(m)
    {
        isChanged = false;

        // Find Top Right
        int p = top_right;
        top_right = TopRight(c_left[top_left], ref c_right, top_right); // O(n)
        if (p != top_right)
        {
            isChanged = true;
        }

        // Find Top Left
        p = top_left;
        top_left = TopLeft(c_right[top_right], ref c_left, top_left); // O(n)
        if (p != top_left)
        {
            isChanged = true;
        }

        // Find Bottom Right
        p = bottom_right;
        bottom_right = BottomRight(c_left[bottom_left], ref c_right,
bottom_right); // O(n)
        if (p != bottom_right)
        {
            isChanged = true;
        }

        // Find Bottom Left
        p = bottom_left;
        bottom_left = BottomLeft(c_right[bottom_right], ref c_left, bottom_left);
// O(n)
        if (p != bottom_left)
        {
            isChanged = true;
        }
    }
}

```

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```

List<PointF> convex = new List<PointF>();

if (top_left == 0)
{
    convex.Add(c_left[0]); // O(1)
}
else
{
    convex = c_left.GetRange(0, top_left+1); // O(n)
}

List<PointF> temp_c;
if (bottom_right == 0)
{
    temp_c = c_right.GetRange(top_right, c_right.Count - top_right); //
O(n)
    temp_c.Add(c_right[0]); // O(1)
}
else
{
    temp_c = c_right.GetRange(top_right, bottom_right-top_right+1); //
O(n)
}
temp_c.ForEach(v => convex.Add(v)); // O(n)

if (bottom_left == c_left.Count-1)
{
    convex.Add(c_left[bottom_left]); // O(n)
}
else if (bottom_left != 0)
{
    if (c_left.Count - 1 == bottom_left)
    {
        convex.Add(c_left[bottom_left]); // O(1)
    }
    else
    {
        temp_c = c_left.GetRange(bottom_left, c_left.Count - bottom_left);
// O(n)
        temp_c.ForEach(v => convex.Add(v)); // O(n)
    }
}

return convex;
}

```

```

/**
 * Calculations for finding the top right point of the Tangent
 *
 * Complexity =>
 *          worst    O(n)
 *          avg      O(n)
 *          best     O(1)
 */
public int TopRight(PointF pt, ref List<PointF> convex, int top_right)
{
    double prev = CalcSlope(pt, convex[top_right]);
    for (int i = top_right+1; i < convex.Count; i++)
    {
        double curr = CalcSlope(pt, convex[i]);
        if (prev > curr)
        {
            return i - 1;
        }
        prev = curr;
    }
    return convex.Count - 1;
}

/**
 * Calculations for finding the top left point of the Tangent
 *
 * Complexity =>
 *          worst    O(n)
 *          avg      O(n)
 *          best     O(1)
 */
public int TopLeft(PointF pt, ref List<PointF> convex, int right)
{
    double prev = CalcSlope(pt, convex[right]);

    for (int i = right - 1; i > -1; i--)
    {
        double curr = CalcSlope(pt, convex[i]);
        if (prev < curr)
        {
            return i + 1;
        }
        prev = curr;
    }
    if (right == convex.Count-1)
        return 0;
    return 0;
}

```

```

/**
 * Calculations for finding the top bottom point of the Tangent
 *
 * Complexity =>
 *          worst    O(n)
 *          avg      O(n)
 *          best     O(1)
 */
public int BottomRight(PointF pt, ref List<PointF> convex, int pos)
{
    double prev = CalcSlope(pt, convex[pos]);
    if (pos == 0)
    {
        pos = convex.Count;
    }
    for (int i = pos - 1; i > -1; i--)
    {
        double curr = CalcSlope(pt, convex[i]);
        if (prev < curr)
        {
            if (i == convex.Count-1) { return 0; }
            return i + 1;
        }
        prev = curr;
    }
    if (pos != convex.Count - 1) {return 0;}
    return convex.Count - 1;
}

/**
 * Calculations for finding the bottom left point of the Tangent
 *
 * Complexity =>
 *          worst    O(n)
 *          avg      O(n)
 *          best     O(1)
 */
public int BottomLeft(PointF pt, ref List<PointF> convex, int pos)
{
    double prev = CalcSlope(pt, convex[pos]);
    if (pos == convex.Count-1)
    {
        pos = -1;
    }
    for (int i = pos+1; i < convex.Count; i++)
    {
        double curr = CalcSlope(pt, convex[i]);
        if (prev > curr)
        {
            if (i == 0) { return convex.Count - 1; }
            return i - 1;
        }
        prev = curr;
    }
    if (pos != convex.Count - 1) {return convex.Count - 1;}
    return 0;
}

```

```

/**
 * Finds the rightmost value of a clockwise sorted array of PointF's
 * Complexity =>
 *          worst    O(n)
 *          avg      O(n)
 *          best     O(n)
 */
public int Rightmost (ref List<PointF> convex)
{
    PointF prev = new PointF { X = 0, Y = 0 };
    int right = -1;
    for (int i = 0; i < convex.Count; i++) // O(n)
    {
        if (right == -1) { right = 0; }
        else
        {
            right++;
            if (prev.X > convex[i].X) { return i-1; }
            // The previous value is the rightmost
        }
        prev = convex[i];
    }
    return convex.Count-1; // For the case where there is only 1 edge
}

/**
 * Sorts a given list Counter Clockwise starting at the first point,
 * the furthest left in our case
 * Complexity =>
 *          worst    O(n^2)
 *          avg      O(nlogn)
 *          best     O(nlogn)
 */
private List<PointF> SortCC(List<PointF> pointList)
{
    List<PointF> sortedList = new List<PointF>();
    for (int i = 1; i < pointList.Count; i++) // O(n)
    {
        sortedList.Add(pointList[i]);
    }
    sortedList.Sort((a, b) => // O(nlogn) at best, O(n^2) at worst // generally a
quick sort
        ((b.Y - pointList[0].Y) / (b.X - pointList[0].X)).CompareTo(
            (a.Y - pointList[0].Y) / (a.X - pointList[0].X)
        ));
    sortedList.Insert(0, pointList[0]); // Insert the first point at the beginning
// O(n)
    return sortedList;
}

/** Used for calculating the slope of two vertices */
private Double CalcSlope(PointF p1, PointF p2)
{
    return (p2.Y - p1.Y) / (p2.X - p1.X);
}

```

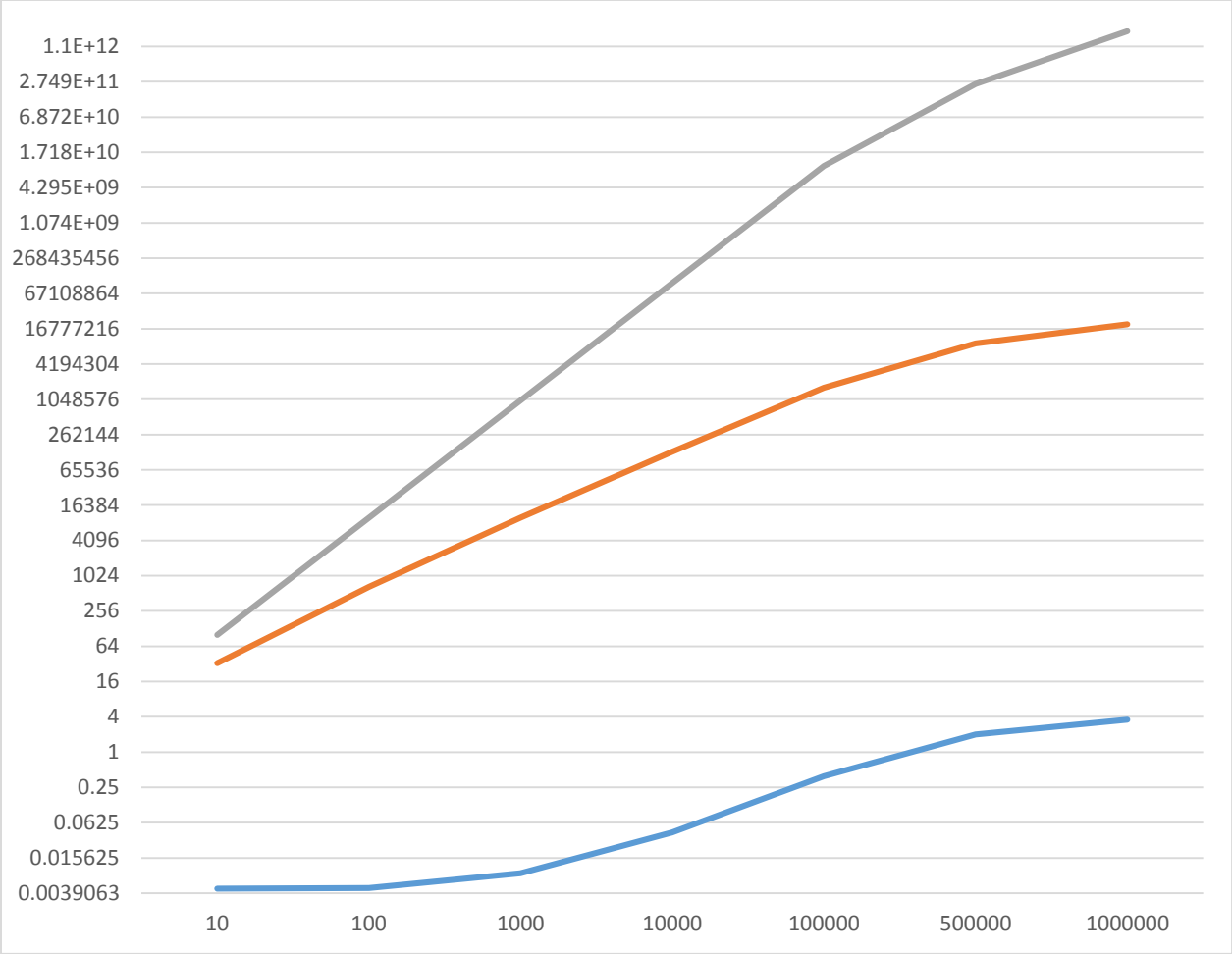


The convex ( $O(n^2)$  at worst) hull does seem to fit my data. The data (blue line graph) seems to fit a similar pattern to my thoughts. The graph I used is a logarithmic one, meaning that if the data set is logarithmic, then it will be linear on the plot. As you can see, my data starts out linear ( $O(\log n)$ ), the graph strays from it a bit, showing more of a  $O(n \log n)$  pattern (compare to the orange line graph). The orange line graph is an  $n \log n$  graphing of the same initial values, the gray is an  $O(n^2)$  graphing. Variations between the two graphs (blue and yellow) are slight.

The recurrence relation for this algorithm is  $T(n) = 2T(n/2) + O(n)$  for the average case, and  $T(n) = 2T(n/2) + O(n^2)$  for the worst case, which is extremely unlikely as the data shows. This solves to  $O(n \log n)$  and  $O(n^2)$ .

Data :

10	0.0047104
100	0.0047735
1000	0.0086178
10000	0.0426202
100000	0.3960408
500000	2.0077503
1000000	3.5943062



Screenshots

