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MIT Digital Currency Initiative and the University of Brasilia presents

# Cryptocurrency Design and Engineering

Assignment 3: Background  
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# Required Materials

- Computer with Python3 installed
  - <https://www.python.org/downloads/>
- Python's ecdsa library
  - `python -m pip install ecdsa`
  - `pip install ecdsa`
- Instructions and Starter Code
  - <https://github.com/mit-dci-cde-2025/mit-dci-cde-2025-classroom-students-assignment-2-cde2025.2-digital-signatures-makerere>

# Goals

- Have a basic working understanding of ECDSA's usage in bitcoin
- A completed assignment 3
  - Create a signature
  - Sign a message
  - Forge 2 messages
  - Alter a valid signature

# Hashing

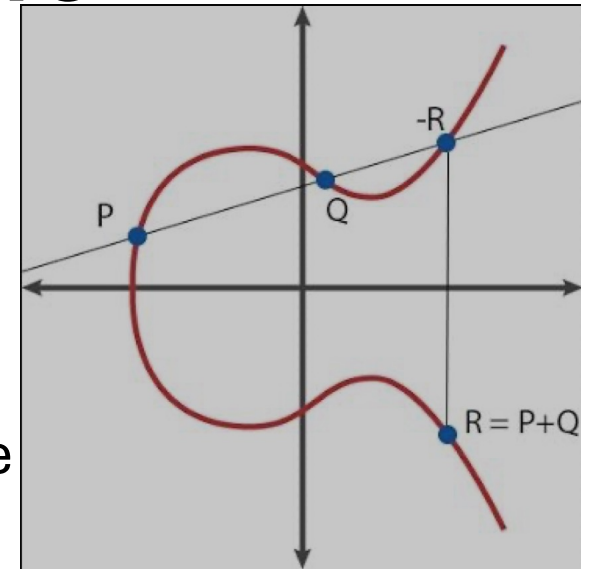
- Takes an arbitrary length input
- Outputs fixed length output
- The output appears random but is deterministic
- Allows for data integrity
- Allows for the commitment of information
  - It will rain tomorrow
    - aa6937cfa2c147206b40c2fe45e8182174bad1d705eae2d91efb8e773e22d36f
  - It will not rain tomorrow
    - 725763ae3279cc027377cf04cce9a5fa14f9850852be6435040b9c16083e7d13

# Digital Signatures

- Method by which a data can be verified
- Consists of two key components
  - Private Key
  - Public Key
- Used to verify information came from a specific source
- Relies on complex mathematics
  - RSA
  - ECDSA

# Elliptic Curve Digital Signature Algorithm

- Smaller keys
- Faster computation
- Patent free
- Capital letters will represent a Point on the curve
- Lower Case letters represent a scalar
- $P+Q+R = 0$
- $P+Q = -R$



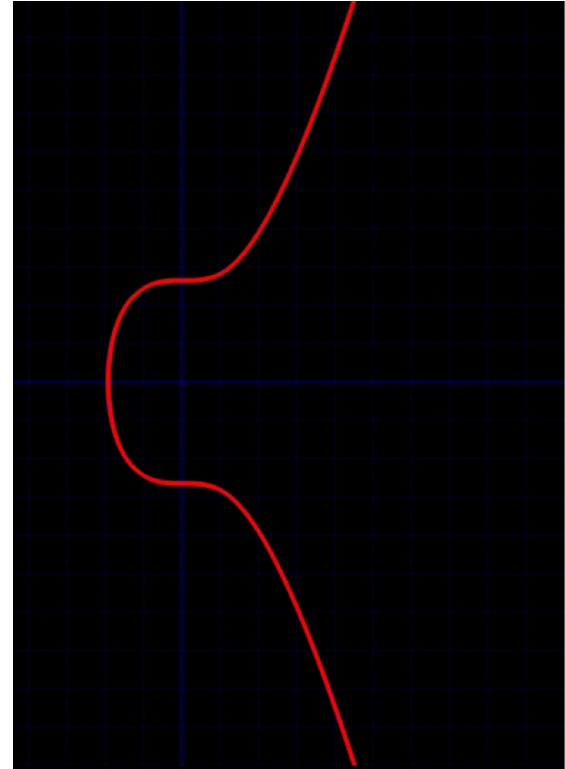
*Tadge Dryja. Cryptocurrency Engineering and Design.*  
2018. Massachusetts Institute of Technology: MIT OpenCourseWare,  
<https://ocw.mit.edu/>. License: Creative Commons BY-NC-SA.

# Elliptic Curve Operations

- $A+B$
- $A-B$
- $A*b$
- $A/b$

# Elliptic Curve Digital Signature Algorithm

- SECP256k1
- $y^2 = x^3 + 7$
- G and n are constants
- Private key is random integer a
- Public key is derived from the private key
  - $E = e * G$
- $R = k * G$
- $r = x \bmod n$
- $s = k^{(-1)} * (\text{hash}(m) + r * e) \bmod n$



Tadge Dryja. *Cryptocurrency Engineering and Design*. 2018. Massachusetts Institute of Technology: MIT OpenCourseWare, <https://ocw.mit.edu/>. License: Creative Commons BY-NC-SA.



# Assignment 2

- Directories
  - data
  - graders
  - implementation
  - Solutions
- <https://ecdsa.readthedocs.io/en/latest/modules.html>

# Exercise 1

- Create a Public/Private key pair
- Sign a message
  - `hashlib.sha256(b"").digest().hex()`
- Output your public key and signed message
  - All hex characters
  - Compressed format public key

# Exercise 2

- Given a known  $k$  forge a message
- $s = k^{-1} * (\text{hash}(m) + r * e) \bmod n$ 
  - $\rightarrow e = (s * k - h) * r_{\text{inv}} \% n$
- $r_{\text{inv}} = \text{pow}(r, -1, n)$

# Exercise 3

- Given  $k$  reuse forge a message
- Sony PlayStation 5 Hack
- $k = (\text{hash}(m1) - \text{hash}(m2)) / (s1 - s2) \bmod n$

# Exercise 4

- Signature malleability
- $(r, s) \equiv (r, n-s)$
- Vulnerability was patched

# Exercise 5

- Modify a block
- $(r, s) \equiv (r, n-s)$
- $r = \text{tx\_data}[47:79]$
- $s = \text{tx\_data}[81:113]$