

Introduction to OpenMP

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Outline

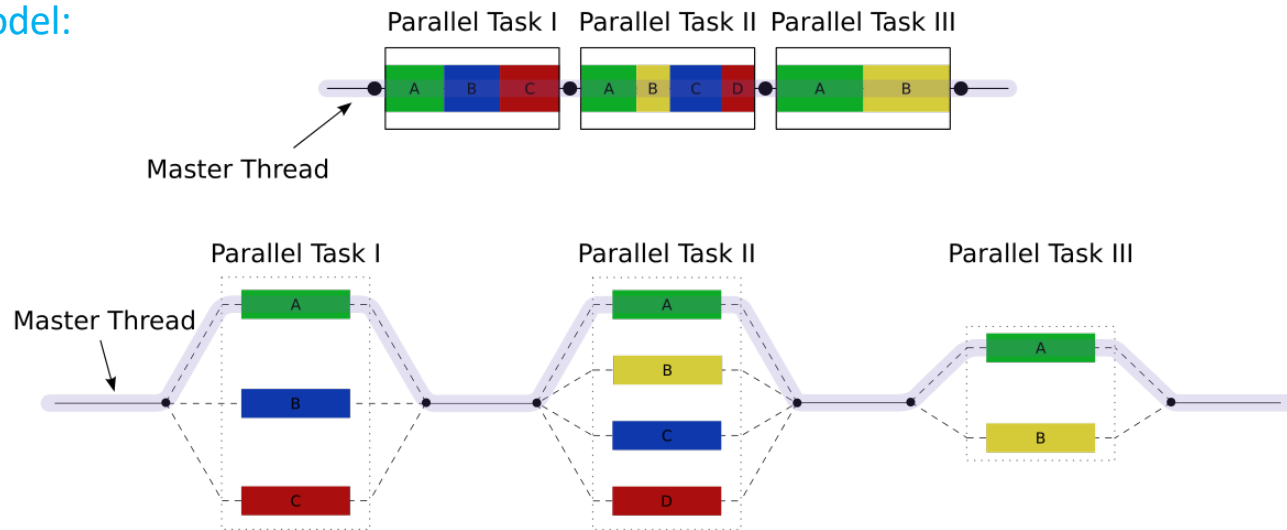
- A brief overview of OpenMP
- Constructs and clauses
- Parallel loops
- Data race
- OpenMP in Numpy

OpenMP

- ❑ OpenMP (Open Multi-Processing) is an API (application programming interface) that supports multi-platform **shared memory multiprocessing** programming.
- ❑ Supporting languages: **C, C++, and Fortran.**
- ❑ Consists of a set of **compiler directives, library routines, and environment variables** that influence run-time behavior.
- ❑ OpenMP supports accelerators.

Parallelism of OpenMP

- **Multithreading:** a master thread forks a specified number of slave threads and the system divides a task among them. The threads then run concurrently, with the runtime environment allocating threads to different processors (or cores).
- **Fork-join model:**



- Reference: <http://en.wikipedia.org/wiki/OpenMP>

The first OpenMP program: Hello world!

- Hello world in C

```
#include <omp.h>
int main() {
    int id;
    #pragma omp parallel private(id)
    {
        id = omp_get_thread_num();
        if (id%2==1)
            printf("Hello world from thread %d, I am odd\n", id);
        else
            printf("Hello world from thread %d, I am even\n", id);
    }
}
```

- Hello world in Fortran

```
program hello
  use omp_lib
  implicit none
  integer i
  !$omp parallel private(i)
  i = omp_get_thread_num()
  if (mod(i,2).eq.1) then
    print *, 'Hello from thread', i, ', I am odd!'
  else
    print *, 'Hello from thread', i, ', I am even!'
  endif
  !$omp end parallel
end program hello
```

OpenMP directive syntax

- In C/C++ programs

```
#pragma omp directive-name [clause[[,] clause]. . . ]
```

- In Fortran programs

```
!$omp directive-name [clause[[,] clause]. . . ]
```

- **Directive-name** is a specific keyword, for example *parallel*, that defines and controls the action(s) taken.
- **Clauses**, for example *private*, can be used to further specify the behavior.

Compile and run OpenMP programs

Compile C/Fortran codes

```
gcc -fopenmp name.c -o name  
gfortran -fopenmp name.f90 -o name
```

Run OpenMP programs

```
export OMP_NUM_THREADS=8    # set number of threads  
./name  
time ./name                 # run and measure the time.
```


OpenMP programming

- Parallel Construct
- Work-Sharing Constructs
- Clauses
- Data race: atomic, reduction

➤ **Construct** : An OpenMP executable directive and the associated statement, loop, or structured block, not including the code in any called routines.

Parallel construct

- Syntax in C/C++ programs

```
#pragma omp parallel [clause [, clause]. . . ]  
..... code block .....
```

- Syntax in Fortran programs

```
!$omp parallel [clause [, clause]. . . ]  
..... code block .....
```

```
!$omp end parallel
```

- Parallel construct is used to specify the computations that should be executed in parallel.
- A team of threads is created to execute the associated parallel region.
- The work of the region is replicated for every thread.
- At the end of a parallel region, there is an implied barrier that forces all threads to wait until the computation inside the region has been completed.

Work-sharing constructs

Functionality	Syntax in C/C++	Syntax in Fortran
Distribute iterations	#pragma omp for	!\$omp do
Distribute independent works	#pragma omp sections	!\$omp sections
Use only one thread	#pragma omp single	!\$omp single
Parallelize array syntax	N/A	!\$omp workshare

- Many applications can be parallelized by using just a parallel region and one or more of work-sharing constructs, possibly with clauses.

- The parallel and work-sharing (except single) constructs can be combined.
- Following is the syntax for combined parallel and work-sharing constructs,

Combine parallel construct with ...	Syntax in C/C++	Syntax in Fortran
Loop construct	#pragma omp parallel for	!\$omp parallel do
Sections construct	#pragma omp parallel sections	!\$omp parallel sections
Workshare construct	N/A	!\$omp parallel workshare

Loop construct

- The loop construct causes the iterations of the loop immediately following it to be executed in parallel.

- Syntax in C/C++ programs

```
#pragma omp for [clause[[,] clause]. . . ]  
..... for loop .....
```

- Syntax in Fortran programs

```
!$omp do [clause[[,] clause]. . . ]  
..... do loop .....
```

```
[!$omp end do]
```

- The terminating **!\$omp end do** directive in Fortran is optional but recommended.

- Distribute iterations in a parallel region

```
#pragma omp parallel for shared(n,a) private(i)
  for (i=0; i<n; i++)
    a[i] = i + n;
```

- **shared clause:** All threads can read from and write to a shared variable.
- **private clause:** Each thread has a local copy of a private variable.
- The maximum iteration number n is shared, while the iteration number i is private.
- Each thread executes **a subset** of the total iteration space $i = 0, \dots, n - 1$
- The mapping between iterations and threads can be controlled by the schedule clause.

- Two work-sharing loops in one parallel region

```
#pragma omp parallel shared(n,a,b) private(i)
{
    #pragma omp for
    for (i=0; i<n; i++) a[i] = i+1;
    // there is an implied barrier

    #pragma omp for
    for (i=0; i<n; i++) b[i] = 2 * a[i];
} /*-- End of parallel region --*/
```

- The distribution of iterations to threads could be different for the two loops.
- The **implied barrier** at the end of the first loop ensures that all the values of `a[i]` are updated before they are used in the second loop.

Exercise 1

- SAXPY in OpenMP:

The SAXPY program is to add a scalar multiple of a real vector to another real vector:

$$s = a * x + y.$$

1. Provided a serial SAXPY code, parallelize it using OpenMP directives.
2. Compare the performance between serial and OpenMP codes.

Sections construct

- Syntax in C/C++ programs

```
#pragma omp sections [clause[[,] clause]. . . ]  
{  
  [#pragma omp section ]  
  ..... code block 1 .....  
  [#pragma omp section  
  ..... code block 2 ..... ]  
  . . .  
}
```

- Syntax in Fortran programs

```
!$omp sections [clause[[,] clause]. . . ]  
[!$omp section ]  
..... code block 1 .....  
[!$omp section  
..... code block 2 ..... ]  
.  
.  
.  
!$omp end sections
```

- The work in each section must be **independent**.
- Each section is distributed to one thread.

- Example of parallel sections

```
#pragma omp parallel sections
{
    #pragma omp section
    funcA();
    #pragma omp section
    funcB();
} /*-- End of parallel region --*/
```

- The most common use of the sections construct is probably to execute function or subroutine calls in parallel.
- There is a **load-balancing problem**, if the amount of work in different sections are not equal.

Single construct

- Syntax in C/C++ programs

```
#pragma omp single [clause[,] clause]. .  
.  
..... code block .....
```

- Syntax in Fortran programs

```
!$omp single [clause[,] clause]. . . ]  
..... code block .....
```

- The code block following the single construct is executed by one thread only.
- The executing thread could be any thread (not necessary the master one).
- The other threads wait at a barrier until the executing thread has completed.

- An example of the single construct

```
#pragma omp parallel shared(a,b) private(i)
{
    #pragma omp single
    {
        a = 10;
    }
    /* A barrier is automatically inserted here */
    #pragma omp for
    for (i=0; i<n; i++)
        b[i] = a;
} /*-- End of parallel region --*/
```

- Only one thread initializes the shared variable *a*.
- If the single construct is omitted here, multiple threads could assign the value to *a* at the same time, potentially resulting in a memory problem.
- The **implicit barrier** at the end of the single construct ensures that the correct value is assigned to the variable *a* before it is used by all threads.

Data race

- A data race condition arises **when multithreads read or write the same shared data simultaneously**.
- **Example:** two threads each increases the value of a shared integer variable by one.

Correct sequence

Thread 1	Thread 2		value
			0
read value		←	0
Increase value			0
write back		→	1
	read value	←	1
	increase value		1
	write back	→	2

Incorrect sequence

Thread 1	Thread 2		value
			0
read value		←	0
	read value	←	0
increase value			0
	increase value		0
write back		→	1
	write back	→	1

- Example of data race: sums up elements of a vector

Different threads read and write the shared data *sum* simultaneously.

A data race condition arises!

The final result of *sum* could be incorrect!

```
sum = 0;
#pragma omp parallel for shared(sum,a,n) private(i)
for (i=0; i<n; i++)
{
    sum = sum + a[i];
} /*-- End of parallel for --*/
printf("Value of sum after parallel region: %f\n",sum);
```

Atomic construct

- The atomic construct allows multiple threads to safely update a shared variable.
- The memory update (such as write) in the next instruction will be performed atomically. It does not make the entire statement atomic. **Only the memory update is atomic.**
- It is applied **only to the (single) assignment statement** that immediately follows it.

C/C++ programs

- Syntax

```
#pragma omp atomic  
..... a single statement .....
```

- Supported operators

```
+, *, -, /, &, ^, |, <<, >>.
```

Fortran programs

```
!$omp atomic  
..... a single statement .....
```

```
!$omp end atomic
```

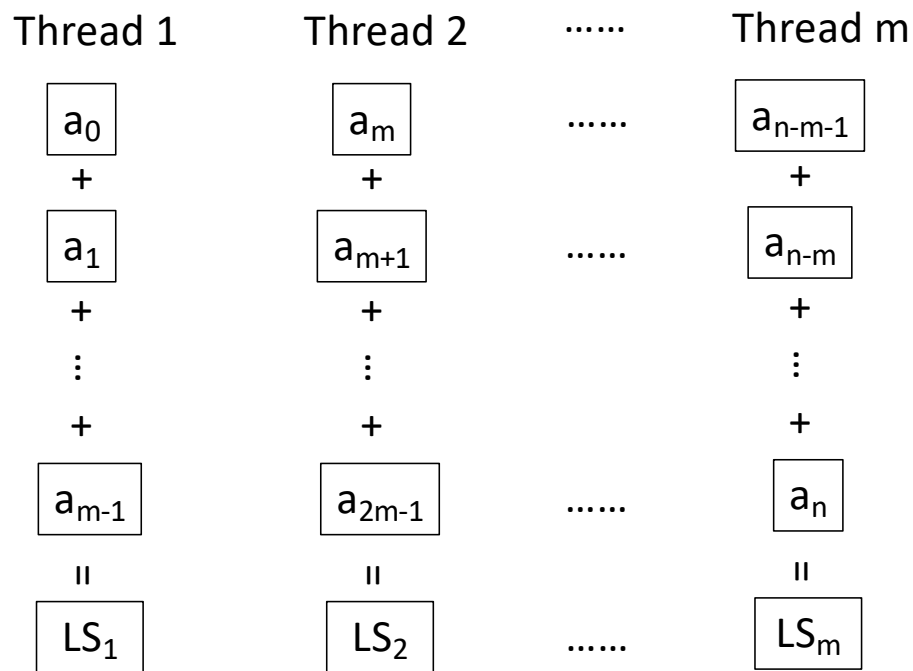
```
+, *, -, /, .AND., .OR., .EQV., .NEQV. .
```

- The first try to solve the data-race problem: use *atomic* (correct but slow)
- The atomic construct avoids the data racing condition. Therefore, the result is correct.
- But all elements are added *sequentially*, and there is *performance penalty for using atomic*, because the system coordinates all threads.
- This code is *even slower than a serial code*!

```
sum = 0;
#pragma omp parallel for shared(n,a,sum) private(i) // Optimization: use reduction instead
of atomic
for (i=0; i<n; i++)
{
    #pragma omp atomic
    sum += a[i];
} /*-- End of parallel for --*/
printf("Value of sum after parallel region: %d\n",sum);
```


- A partially parallel scheme to avoid data race

Step 1: Calculate local sums in parallel



m: number of
threads

n: array length

LS: local sum

Step 2: Update total sum sequentially

Thread 1	Thread 2	Thread m
Read initial S			
$S = S + LS_1$			
Write S			
	Read S		
	$S = S + LS_2$		
	Write S		
		
			Read S
			$S = S + LS_m$
			Write S

m: number of threads

LS: local sum

S: total sum

- The second try to solve the data-race problem: use *atomic* (correct and fast)
- Each thread adds up its local sum.
- The *atomic* is only applied for adding up local sums to obtain the total sum.

```
sum = 0;
#pragma omp parallel shared(n,a,sum) private(sumLocal)
{
    sumLocal = 0;
    #pragma omp for
    for (i=0; i<n; i++) sumLocal += a[i];
    #pragma omp atomic
    sum += sumLocal;
} /*-- End of parallel region --*/
printf("Value of sum after parallel region: %d\n",sum);
```

Reduction clause

- The third try to solve the data-race problem: use *reduction* (correct, fast and simple)

```
#pragma omp parallel for default(none) shared(n,a) private(i) reduction(+:sum)
for (i=0; i<n; i++)
    sum += a[i];
/*-- End of parallel reduction --*/
```

- The reduction variable is protected to *avoid data race*.
- The *partially parallel scheme* is applied behind the scene.
- The reduction variable is shared by default and it is unnecessary to specify it explicitly.

- Operators and statements supported by the reduction clause

	C/C++	Fortran
Typical statements	$x = x \text{ op } \text{expr}$ $x \text{ binop } = \text{expr}$ $x = \text{expr} \text{ op } x$ (except for subtraction) $x++$ $++x$ $x--$ $--x$	$x = x \text{ op } \text{expr}$ $x = \text{expr} \text{ op } x$ (except for subtraction) $x = \text{intrinsic } (x, \text{expr_list})$ $x = \text{intrinsic } (\text{expr_list}, x)$
<i>op</i> could be	+, *, -, &, ^, , &&, or	+, *, -, .and., .or., .eqv., or .neqv.
<i>binop</i> could be	+, *, -, &, ^, or	N/A
<i>Intrinsic</i> function could be	N/A	max, min, iand, ior, ieor

Exercise 2

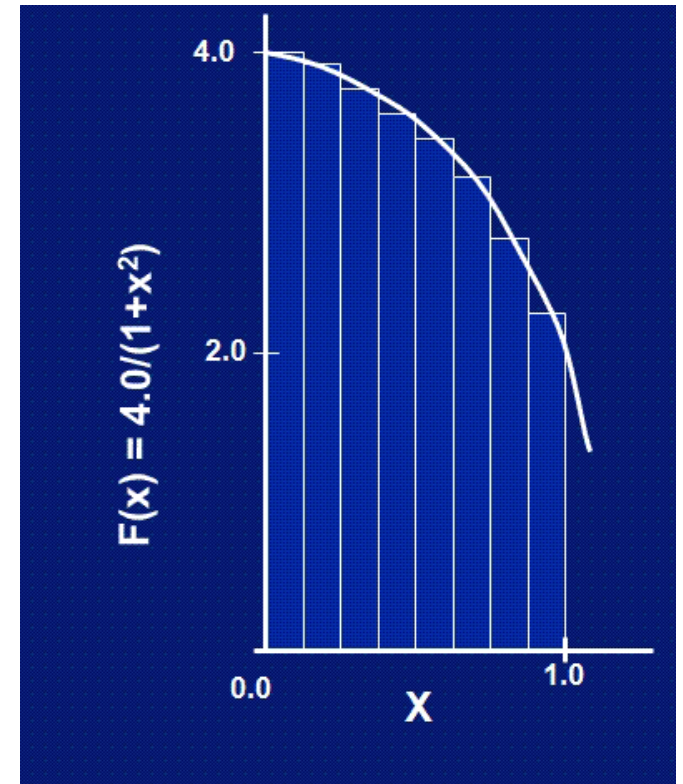
- Compute the value of pi:

$$\int_0^1 \frac{4.0}{(1+x^2)} dx = \pi$$

The integration can be numerically approximated as the sum of a number of rectangles.

$$\sum_{i=0}^N F(x_i) \Delta x \approx \pi$$

1. Provided the serial code for computing the value of pi, parallelize it using OpenMP directives.
2. Compare the performance between serial and OpenMP codes.



Using OpenMP in Numpy

❑ OpenMP-enabled Libraries

- **OpenBLAS**: BLAS = Basic Linear Algebra Subprograms
- **MKL**: Linear algebra libraries optimized for intel CPUs.

❑ Numpy

- Install Numpy with precompiled OpenBLAS or MKL
- Numpy calls OpenMP-enabled routines under the hood.
- Run a python program with multiple threads

```
pip install numpy
```

```
import numpy as np
n = 10000
a = np.random.rand(n,n)
b = np.random.rand(n,n)
c = np.matmul(a, b)
```

```
export OMP_NUM_THREADS=8
python name.py
```

Appendix A: OpenMP built-in functions

- Enable the usage of OpenMP functions:

C/C++ program: include `omp.h` .

Fortran program: include `omp_lib.h` or use `omp_lib` module.

- List of OpenMP functions:

`omp_set_num_threads(integer)` : set the number of threads

`omp_get_num_threads()`: returns the number of threads

`omp_get_thread_num()`: returns the number of the calling thread.

`omp_set_dynamic(integer|logical)`: dynamically adjust the number of threads

`omp_get_num_procs()`: returns the total number of available processors when it is called.

`omp_in_parallel()`: returns true if it is called within an active parallel region. False otherwise.

Appendix B: OpenMP runtime variables

OMP_NUM_THREADS : the number of threads (**=integer**)

OMP_SCHEDULE : the schedule type (**=kind,chunk** . Kind could be static, dynamic or guided)

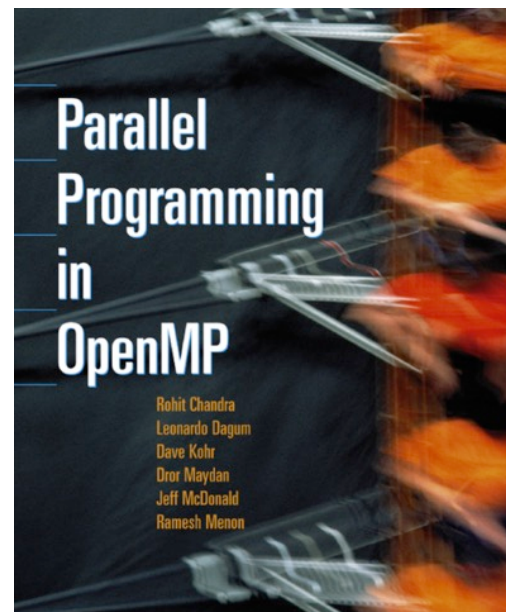
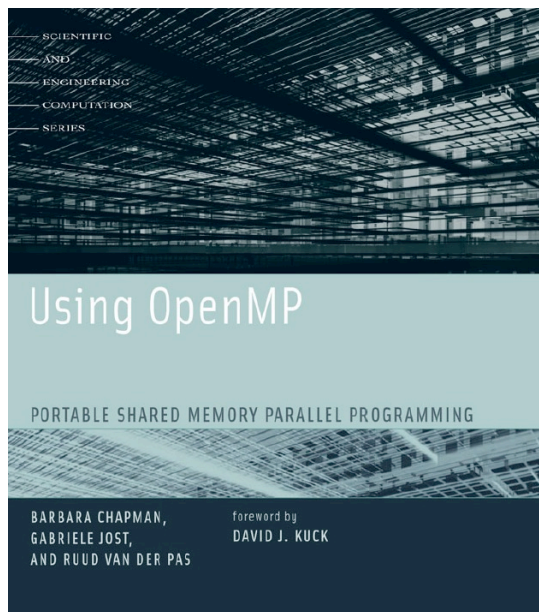
OMP_DYNAMIC : dynamically adjust the number of threads (**=true | =false**).

KMP_AFFINITY : only for intel compiler, to bind OpenMP threads to physical processing units.
(**=compact | =scatter | =balanced**).

Example usage: **export KMP_AFFINITY= compact,granularity=fine,verbose** .

Further information

References



OpenMP official website: <http://openmp.org/wp/>