

1 2.1  $N, \sqrt{N}, N^{1.5}, N^2, N \log N, N \log \log N, N \log^2 N, N \log(N^2),$   
 $\frac{2}{N}, 2^N, 2^{\frac{N}{2}}, 37, N^2 \log N, N^3$

$\frac{2}{N}, 37, \sqrt{N}, N, N \log \log N, N \log N = N \log(N^2), N \log^2 N,$   
 $N^{1.5}, N^2, N^2 \log N, N^3, 2^{\frac{N}{2}} = 2^N$

2 2.6 day 1 = 2, day 2 = 4, day 3 = 16, day 4 = 256, day 5 = 65536  
 a 1 2 4 16 256 65536

start at day  $N=1$  with a fine of \$2,  
 then the fine on day  $N = 4^{2^{(N-2)}}$

b  $D = 4^{2^{(N-2)}}$   $2^N = 4^{\frac{1}{2}N}$   
 $D = 4^{\frac{2^N}{2}} \Rightarrow D = 4^{2^{\frac{N}{2}}}$   
 $\log_4 D = \log_4 2^N = \log_4 4^{\frac{N}{2}}$   
 $\log_4 D = \frac{N}{2}$

$$N = 2 \log_4 D$$

$$N = \log D$$

3 2.7a 1. sum = 0;  
 for (i = 0; i < n; i++) = n  $\rightarrow O(n)$   
 sum++;

2. sum = 0;  
 for (i = 0; i < n; i++) = n  
 for (j = 0; j < n; j++) = n  $\rightarrow O(n^2)$   
 sum++;

3. sum = 0;  
 for (i = 0; i < n; i++) = n  $\rightarrow O(n^3)$   
 for (j = 0; j < n \* n; j++) =  $n^2$   
 sum++;

4. sum = 0;  
 for (i = 0; i < n; i++) = n  $\rightarrow O(n^2)$   
 for (j = 0; j < i; j++) 0, 1, 2, ..., n-1  $\rightarrow$   
 sum++;  $\frac{(n-1)n}{2} \rightarrow n^2$



5.  $sum = 0;$   
 for  $i = 0; i < n; i++ = n \rightarrow O(n^5)$   
 for  $j = 0; j < i; j++$   
 for  $k = 0; k < j; k++ \rightarrow \frac{(n^2-1)(n^2)}{2} = n^4$   
 $i=0 \quad i=1 \quad i=2 \quad i=3 \quad \dots \quad i=n-1$   
 $0 \quad 0 \quad 0+1+2+3=6 \quad 0+1+2+3+4+\dots+8=36 \quad 0 \text{ to } n^2-1$

6  $sum = 0;$   
 for  $i = 1; i < n; i++ = n$   
 for  $j = 1; j < i; j++ = n^2 \rightarrow O(n^5)$   
 if  $(j \% i == 0)$   
 for  $k = 0; k < j; k++ = n^2$  assume this case always runs for 0  
 $sum++;$

4 2.10 a add 2 N digit numbers  
 $sum = 0, carry = 0$   
 for  $(i = n-1 \text{ to } i = 0) = n \rightarrow O(N)$   
 $sum = (n1i + n2i + carry) \% 10 \times 10^{\text{power}}$   
 if  $(n1i + n2i + carry > 9)$   
 $carry = 1$

5 2.11 0.5 ms for  $N=100$ , time for  $N=500$   
 a  $0.5(5) = 2.5 \text{ ms}$  for linear  
 b  $0.5(5) \left( \frac{\log 500}{\log 100} \right) = 3.37 \text{ ms}$  for  $O(N \log N)$   
 c  $0.5(5^2) = 12.5 \text{ ms}$  for quadratic  
 d  $0.5(5^3) = 62.5 \text{ ms}$  for cubic

6 2.15  $A_i = i$  for an array  $A_1 < A_2 < A_3 < \dots < A_N$

1	2	5	7	8	$\frac{4+0}{2} = 2$	$A_i > i$
0	1	2	3	4		$A_i \leq i$
1	2	3	4	5		

0 1 2 3 4 5  
 1 2 3 4 5

```
boolean hasMatchingIndexEl(int[] arr) {
    int beg = 1; int end = arr.length;
    int mid;
    while (beg <= end) {
        mid = beg + (end - beg) / 2;
```



```

    if (arr[mid-1] > mid)
        end = mid-1;
    else if (arr[mid-1] < mid)
        beg = mid+1;
    else if (arr[mid-1] == mid)
        return true;
}

```

```

return false;
}

```

→ each time, splits the portion of the array to consider in half (like binary search)

$O(\log N)$

3.1 void printLots(L, P) {

~~int lIdx = 0; int lIter = L.iterator();~~

~~int pIdx = 0; int pIter = P.iterator();~~

int lIdx = 0;

Iterator<T> lIter = L.iterator();

Iterator<Integer> pIter = P.iterator();

int toPrint;

while (pIter.hasNext()) {

toPrint = pIter.next();

while (lIter.hasNext()) {

if (lIdx == toPrint) {

System.out.println(lIter.next());

lIdx++;

break;

}

lIdx++;

}

}

if (N == L.size())  $O(N)$