

ECE 471 Lab 1

Secret-Key Encryption Lab

Mitchell Dzurick

2/3/2020

Github with all documentation - <https://www.github.com/mitchdz/ECE471>

Contents

1	Overview	2
2	Lab Tasks	3
2.1	Task 1: Frequency c Against Monoalphabetic Substitution Cipher	3
2.1.1	Task 1: solution	3
2.2	Task 2: Encryption using Different Ciphers and Modes	21
2.2.1	Task2: solution	21
2.3	Task 3: Encryption Mode – ECB vs. CBC	26
2.3.1	Task 3: Solution	27
2.4	Task 4: Padding	31
2.4.1	Task 4: Solution	32
2.5	Task 5: Error Propagation – Corrupted Cipher Text	32
2.5.1	Task 5: Solution	32
2.6	Task 6: Initial Vector (IV)	32
2.6.1	Task 6: Solution	32

Secret Key Encryption Lab

Copyright © 2018 Wenliang Du, Syracuse University. The development of this document was partially funded by the National Science Foundation under Award No. 1303306 and 1718086. This work is licensed under a Creative Commons Attribution-NonCommercial-ShareAlike 4.0 International License. A human-readable summary of (and not a substitute for) the license is the following: You are free to copy and redistribute the material in any medium or format. You must give appropriate credit. If you remix, transform, or build upon the material, you must distribute your contributions under the same license as the original. You may not use the material for commercial purposes.

1 Overview

The learning objective of this lab is for students to get familiar with the concepts in the secret-key encryption. After finishing the lab, students should be able to gain a first-hand experience on encryption algorithms, encryption modes, paddings, and initial vector (IV). Moreover, students will be able to use tools and write programs to encrypt/decrypt messages. This lab covers the following topics:

- Secret-key encryption
- Substitution cipher and frequency analysis
- Encryption modes and paddings
- Programming using the crypto library

Lab Environment. This lab has been tested on our pre-built Ubuntu 12.04 VM and Ubuntu 16.04 VM, both of which can be downloaded from the SEED website.

2 Lab Tasks

2.1 Task 1: Frequency c Against Monoalphabetic Substitution Cipher

It is well-known that monoalphabetic substitution cipher (also known as monoalphabetic cipher) is not secure, because it can be subjected to frequency analysis. In this lab, you are given a cipher-text that is encrypted using a monoalphabetic cipher; namely, each letter in the original text is replaced by another letter, where the replacement does not vary (i.e., a letter is always replaced by the same letter during the encryption). Your job is to find out the original text using frequency analysis. It is known that the original text is an English article.

2.1.1 Task 1: solution

Shown below is the linux command

```
$ ls -l
```

in the lab1/task1 directory of my github repository. Inside that repository, the following command is used to copy the contents of the ciphertext into my X clipboard:

```
$ xclip -sel clip -i ciphertext.txt
```

```
mitch@light: ~/git/ECE471/lab1/task1 (master) $ ls -l
total 16
-rw-rw-r-- 1 mitch mitch 4759 Jan 29 18:44 ciphertext.txt
-rw-rw-r-- 1 mitch mitch 4759 Jan 29 18:49 plaintext.txt
mitch@light: ~/git/ECE471/lab1/task1 (master) $ xclip -sel clip -i ciphertext.txt
```

Figure 1: listing of lab1/task1 directory

The contents of ciphertext.txt is copied below:

```
ytn xqavhq yzhu xu qzupvd ltmat qnncq vgxzy hmrtv ybynh ytmq ixur qyhurn
vlvhpq yhme ytn gvrrnh bnniq imsn v uxuvrnuvhmvu yxx
```

```
ytn vlvhpq hvan lvq gxxsnupnp gd ytn pncmqn xb tvhfnd lnmuqynmu vy myq xzyqny
vup ytn veevhnu ymceixqmxu xb tmq bmic axcevud vy ytn nup vup my lvq qtvenp gd
ytn ncncnruan xb cnyxx ymcnq ze givasrxlu eximymaq vhcaupd vaymfmqc vup
v uvymxuvi axufnhqvymxu vq ghmnb vup cvp vq v bnfnh phnvc vgxzy ltnytnh ytnhn
xzrty yx gn v ehnqmpnuy lmubhnd ytn qnvqxu pmpuy ozqy qnnc nkyhv ixur my lvq
nkyhv ixur gnavzqn ytn xqavhq lnhn cxfnp yx ytn bmhqv lnnsnup mu cvhat yx
vfxmp axubimaymur lmyt ytn aixqmur anhncxud xb ytn lmuynh xidcemaq ytvusq
ednxuratvur
```

```
xun gmr jznqymxu qzhhxzupmur ytmq dnvhq vavpncl vlvhpq mq txl xh mb ytn
```

anhncxud lmii vpphnqq cnyxx nqenamviid vbynh ytn rxipnu rixgnq ltmat gnavcn
v ozgmivuy axcmurxzy evhyd bxh ymcnq ze ytn cxfncnuy qenvhtnvpnp gd
exlnhbzi txiidlxxp lxcnu ltx tniemp hvmqn cmiimxuq xb pxivhq yx bmrty qnkzvi
tvhvqqcnuy vhxzup ytn axzuyhd

qmruvimur ytnmh qzeexhy rxipnu rixgnq vyynupnnq qlvytvp ytn cqnifnq mu givas
qexhypn iveni emuq vup qxzupnp xbb vgxzy qnkqy exlnh mcgvivuanq bhxc ytn hnp
avheny vup ytn qyvrn xu ytn vmh n lvq aviinp xzy vgxzy evd munjzmyd vbynh
myq bxhcnh vuatxh avyy qvpinh jzmy xuan qtn invhunp ytv ytn lvq cvsmur bvh
inqq ytvu v cvin axtxqy vup pzhmur ytn anhncxud uvyyvimm exhycvu yxxs v gizuy
vup qvymqbdmur pmr vy ytn viicvin hxqynh xb uxcmuvynp pmhnayxhq txl axzip
ytvy gn yxeenp

vq my yzhuq xzy vy invqy mu ynhcq xb ytn xqavhq my ehxgvgid lxuy gn

lxcnu mufxifnp mu ymcnq ze qvmp ytv yiytxzrt ytn rixgnq qmrumbmnp ytn
mumymvymfnq ivzuat ytnd unfnh muynupnp my yx gn ozqy vu vlvhq qnvqxu
avcevmru xh xun ytv gnavcn vqqxamvynp xuid lmyt hnpavheny vaymxuq muqynvp
v qexsnqlxvcu qvmp ytn rhxze mq lxhsmur gntmup aixqnp pxxhq vup tvq qmuau
vcvqqnp cmiimxu bxh myq inrvi pnbnuqn bzup ltmat vbynh ytn rixgnq lvq
bixxpn np lmyt ytxzqvupq xb pxuvymxuq xb xh inqq bhxc enxein mu qxcn
axzuyhmnp

ux avii yx lnvh givas rxluq lnuy xzy mu vpfvuan xb ytn xqavhq ytxzrt ytn
cxfncnuy lmii vicxqy anhyvmuid gn hnbnhnuanp gnbxhn vup pzhmur ytn anhncxud
nqenamviid qmuau fxavi cnyxx qzeexhynhq imsn vqtind ozpp ivzhv pnhu vup
umaxin smpcvu vhn qatnpzinp ehnqnuynhq

vuxytnh bnvyzhn xb ytmq qnvqxu ux xun hnviid suxlp ltx mq rxmlur yx lmu gnqy
emayznhn vhrzvgid ytmq tveenuq v ixy xb ytn ymcn muvhrzvgid ytn uvmigmynh
uvhhvymfn xuid qnhfnq ytn vlvhq tden cvatmun gzy xbynu ytn enxein bxhnavqymur
ytn hvan qxaviinp xqavhxixrmqyq avu cvsn xuid npzavynp rznqgnq

ytn lvd ytn vavpncl yvgzivynq ytn gmr lmuunh pnxnqy tnie mu nfnhd xytnh
avynrxhd ytn uxcmunn lmyt ytn cxqy fxynq lmuq gzy mu ytn gnqy emayznhn
avynrxhd fxynhq vhn vqsnp yx imqy ytnmh yxe cxfmnq mu ehnbnhnuymvi xhpnh mb v
cxfmn rnyq cxhn ytvu enhanuy xb ytn bmhqyeivan fxynq my lmuq ltnu ux
cxfmn cvu vrnq ytv ytn xun lmyt ytn bnlqy bmhqyeivan fxynq mq nimcmuvynp vup
myq fxynq vhn hnpmqyhmgzynp yx ytn cxfmnq ytv rvhunhnp ytn nimcmuvynp gviixyq
qanaxupeivan fxynq vup ytmq axuymuznq zuymi v lmuunh ncncrnq

my mq vii ynhmgid axubzqmur gzy veevhnu yid ytn axuqnuqzq bvfxhmy ytn
vtnvp mu ytn nup ytmq cnvuq ytv ytn nupxbqnvqxu vlvhq atvyy ytn
mufxifnp yxhzhnp qenazivymxu vgxzy ltmat bmic lxzip cxqy imsnid gn fxynh

qanaxup xh ytmhp bvxhmyv vup ytnu njzviid yxhyzhnp axuaizqmxuq vgxzy ltmat
bmic cmrty ehnfvmi

mu my lvq v yxqqze gnylnnu gxdtxxp vup ytn nfnuyzvi lmuunh gmhpcvu
mu lmyt ixyq xb nkenhyq gnyymur xu ytn hnfuvuy xh ytn gmr qtxhy ytn
ehmwn lnuy yx qexyimrty ivqy dnvh unvhid vii ytn bxhnavqynhq pnaivhnp iv
iv ivup ytn ehnqzceymfn lmuunh vup bxh ylx vup v tvib cmuzynq ytnd lnhn
axhhnay gnbxhn vu nufnixen quvbz lvq hnfnvinp vup ytn hmrtbybzi lmuunh
cxxuimrty lvq ahxlunp

ytmq dnvh vlvhpq lvyatnhq vhn zunjzviid pmfmnpn gnylnnu ythnn gmiigxvhpq
xzyqmpn nggmur cmqqxzhm ytn bvxhmyv vup ytn qtven xb lvynh ltmat mq
ytn gvrrnhq ehnpmaymxu lmyt v bnl bxhnavqymur v tvmi cvhd lmu bxh rny xzy

gzy vii xb ytxqn bmicq tvfn tmqyxhnavi xqavhxymur evyynhuq vrvmuqy ytnc ytn
qtven xb lvynh tvq uxcmuvymxuq cxhn ytvu vud xytnh bmic vup lvq viqx
ucvnp ytn dnvhq gnqy gd ytn ehxpzanhq vup pmhnayxhq rzmidpq dny my lvq uxy
uxcmuvynp bxh v qahnnu vayxhq rzmidp vlvhp bxh gnqy nuqncgin vup ux bmic tvq
lxu gnqy emayzhn lmytxzy ehnfmxzqid ivupmur vy invqy ytn vayxhq uxcmuvymxu
qmuu ghvfntnvhy mu ytmq dnvh ytn gnqy nuqncgin qvr nupnp ze rxmur yx
ythnn gmiigxvhpq ltmat mq qmrumbmavuy gnatzqn vayxhq cvsn ze ytn vavpncdq
ivhrnqy ghvuat ytv ymicl min pmfmqmfn viqx lxu ytn gnqy phvcv rxipnu rixgn
vup ytn gvbyv gzy myq bmiccvsnh cvhymu capxuvrt lvq uxy uxcmuvynp bxh gnqy
pmhnayxh vup vevhy bhxc vhrx cxfmnq ytv ivup gnqy emayzhn lmytxzy viqx
nvhumur gnqy pmhnayxh uxcmuvymxuq vhn bnl vup bv gnylnnu

The related plaintext is shown below using the final (substitution) key ‘abcdefghijklmnoprstu-
vwxyz’ to ‘cfmvpbrlqxwiejdsgkhnazotu’. The process to derive this key is outlined below
with accompanying pictures.

```
mitch@light: ~/git/ECE471/lab1/task1 (master) $ tr 'abcdefghijklmnopqrstuvwxyz' 'cfmvpbrlqxwiejdsgkhnazotu' < ciphertext.txt > plaintext.txt
```

Figure 2: Command to convert ciphertext to plaintext

Figure 2 shows the linux command used to convert the ciphertext that was shown into the
accompanying plaintext.

the oscars turn on sunday which seems about right after this long strange
awards trip the bagger feels like a nonagenarian too

the awards race was bookended by the demise of harvey weinstein at its outset
and the apparent implosion of his film company at the end and it was shaped by
the emergence of metoo times up blackgown politics armcandy activism and
a national conversation as brief and mad as a fever dream about whether there
ought to be a president winfrey the season didnt just seem extra long it was

extra long because the oscars were moved to the first weekend in march to avoid conflicting with the closing ceremony of the winter olympics thanks pyeongchang

one big question surrounding this years academy awards is how or if the ceremony will address metoo especially after the golden globes which became a jubilant comingout party for times up the movement spearheaded by powerful hollywood women who helped raise millions of dollars to fight sexual harassment around the country

signaling their support golden globes attendees swathed themselves in black sported lapel pins and sounded off about sexist power imbalances from the red carpet and the stage on the air e was called out about pay inequity after its former anchor catt sadler quit once she learned that she was making far less than a male cohost and during the ceremony natalie portman took a blunt and satisfying dig at the allmale roster of nominated directors how could that be topped

as it turns out at least in terms of the oscars it probably wont be

women involved in times up said that although the globes signified the initiatives launch they never intended it to be just an awards season campaign or one that became associated only with redcarpet actions instead a spokeswoman said the group is working behind closed doors and has since amassed million for its legal defense fund which after the globes was flooded with thousands of donations of or less from people in some countries

no call to wear black gowns went out in advance of the oscars though the movement will almost certainly be referenced before and during the ceremony especially since vocal metoo supporters like ashley judd laura dern and nicole kidman are scheduled presenters

another feature of this season no one really knows who is going to win best picture arguably this happens a lot of the time inarguably the nailbiter narrative only serves the awards hype machine but often the people forecasting the race socalled oscarologists can make only educated guesses

the way the academy tabulates the big winner doesnt help in every other category the nominee with the most votes wins but in the best picture category voters are asked to list their top movies in preferential order if a movie gets more than percent of the firstplace votes it wins when no movie manages that the one with the fewest firstplace votes is eliminated and its votes are redistributed to the movies that garnered the eliminated ballots

secondplace votes and this continues until a winner emerges

it is all terribly confusing but apparently the consensus favorite comes out ahead in the end this means that endofseason awards chatter invariably involves tortured speculation about which film would most likely be voters second or third favorite and then equally tortured conclusions about which film might prevail

in it was a tossup between boyhood and the eventual winner birdman in with lots of experts betting on the revenant or the big short the prize went to spotlight last year nearly all the forecasters declared la la land the presumptive winner and for two and a half minutes they were correct before an envelope snafu was revealed and the rightful winner moonlight was crowned

this year awards watchers are unequally divided between three billboards outside ebbing missouri the favorite and the shape of water which is the baggers prediction with a few forecasting a hail mary win for get out

but all of those films have historical oscarvoting patterns against them the shape of water has nominations more than any other film and was also named the years best by the producers and directors guilds yet it was not nominated for a screen actors guild award for best ensemble and no film has won best picture without previously landing at least the actors nomination since braveheart in this year the best ensemble sag ended up going to three billboards which is significant because actors make up the academys largest branch that film while divisive also won the best drama golden globe and the bafta but its filmmaker martin mcdonagh was not nominated for best director and apart from argo movies that land best picture without also earning best director nominations are few and far between

After the ciphertext is copied, a program is utilized that is supplied by cryptoclub.org. The domain is shown below.

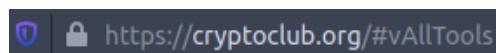


Figure 3: cryptoclub website

Below is the page you should see upon opening the URL in Figure 3



Figure 4: cryptoclub main page

The contents of the ciphertext are copied into the program as shown in Figure 5. On this page, the ciphertext shows the frequency of letters and relates the frequency of letters in the ciphertext to the frequency of letters in the English alphabet.



Figure 5: Crack Substitution Tool main page

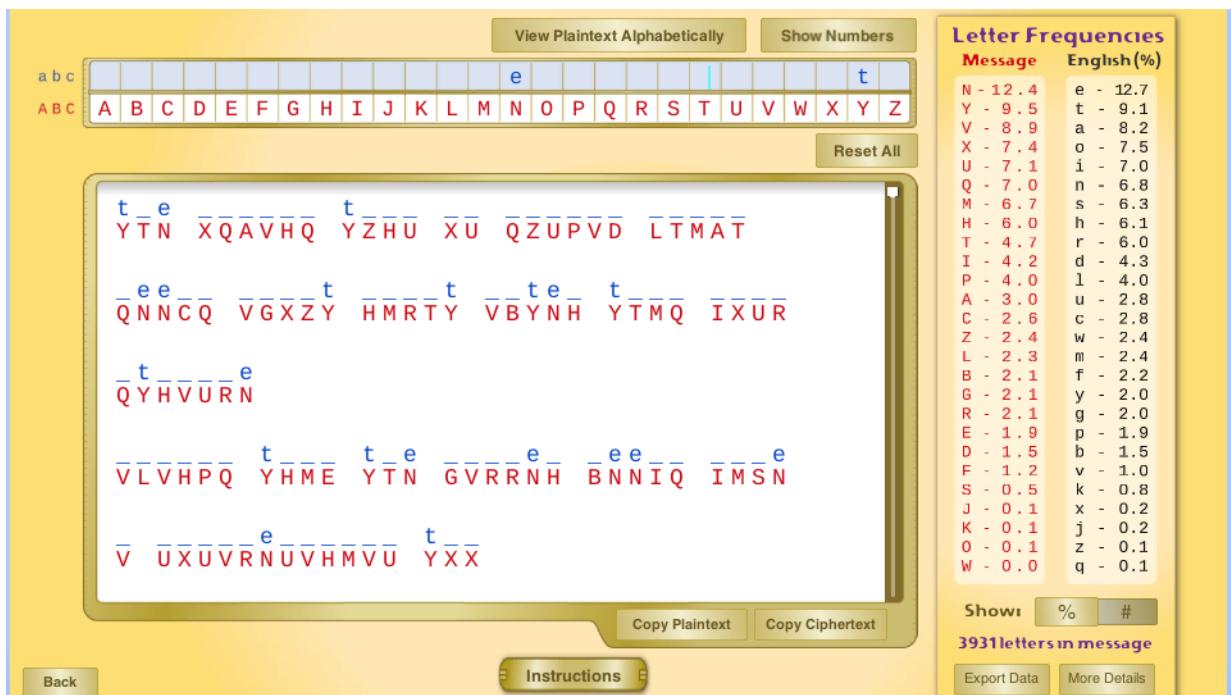


Figure 6: Mapping N to e and Y to t

In Figure 6 N is mapped to ‘e’ and Y is mapped to ‘t’ due to frequency analysis (You can see that N matches e and Y matches t in frequency on the right). It is very easy then to notice that T maps to ‘h’ shown in Figure 7 to create the Trigram ‘the’.

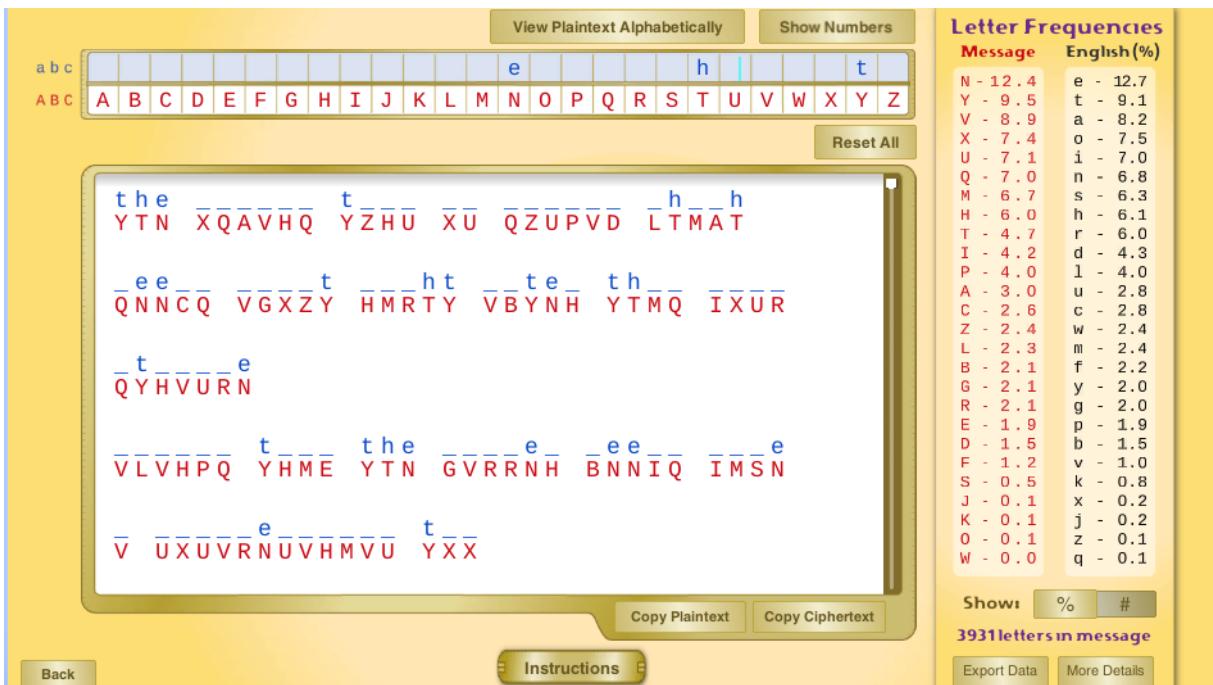


Figure 7: Mapping T to h

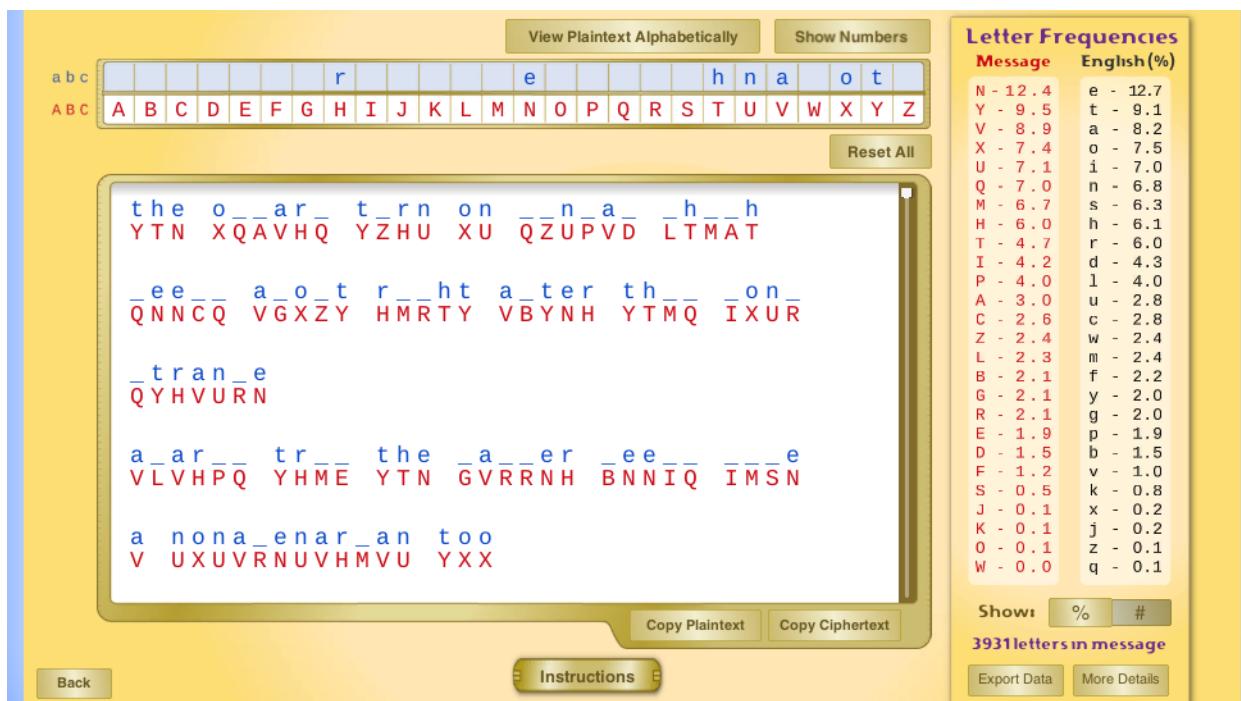


Figure 8: Mapping U to n, V to a, X to o, and H to r

Figure 8 shows how the ciphertext letters 'UVXH' are mapped to 'naor'. This mapping was done by looking at the ciphertext frequency and mapping the frequency to the closest related

frequency in the english alphabet. Words do not particularly make sense yet, so it is just assumed for now that this is the correct mapping.

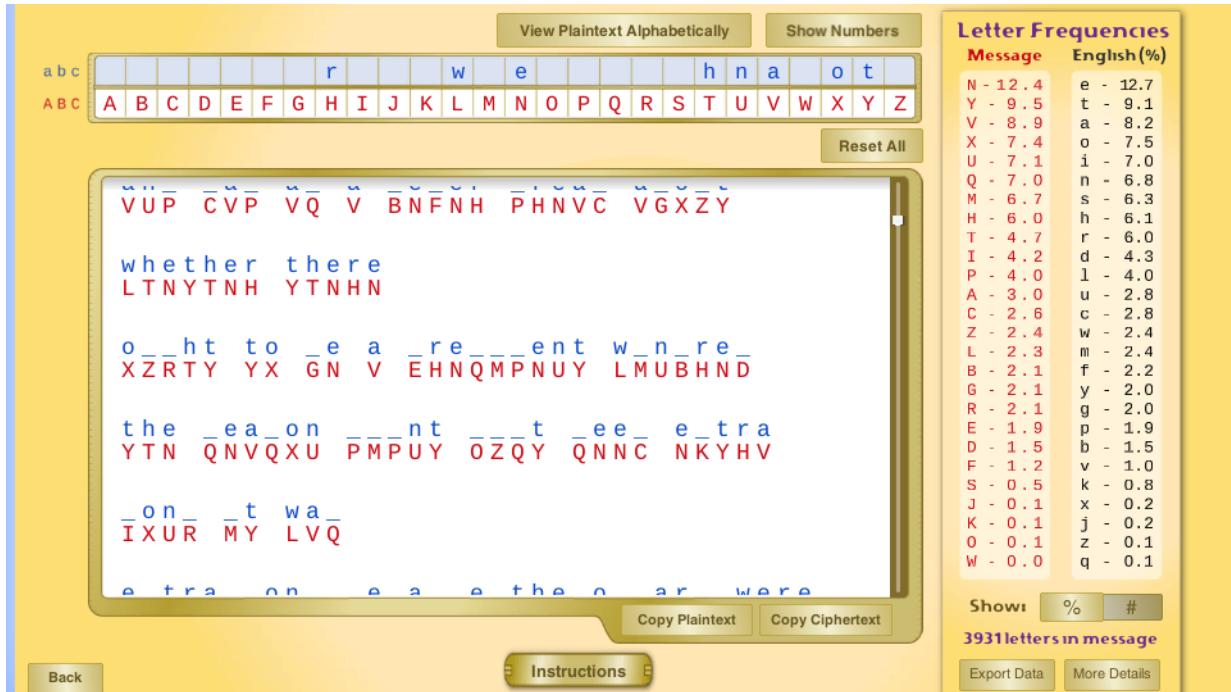


Figure 9: mapping L to w

In Figure 9 the string ‘hether’ was shown, which clearly meant to spell ‘whether’. Therefore, L maps to w.

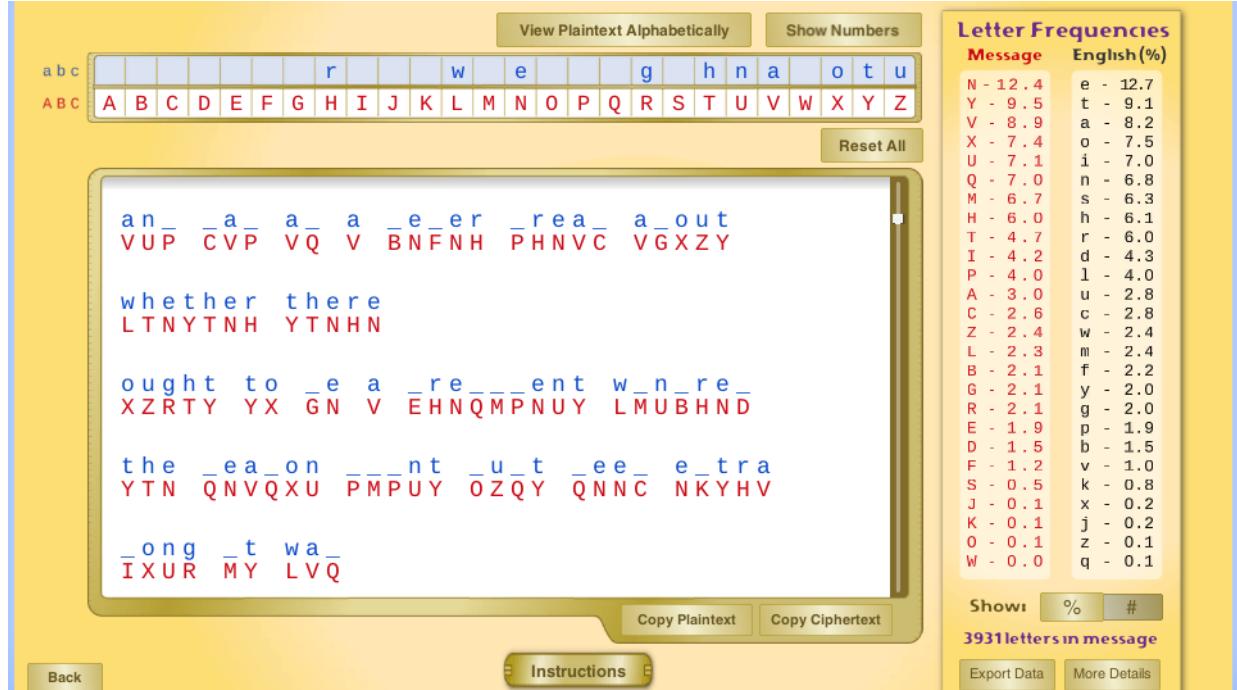


Figure 10: Mapping R to g

In Figure 10 the string 'ou ht' is shown, which is meant to say 'ought'. Therefore, R maps to g.

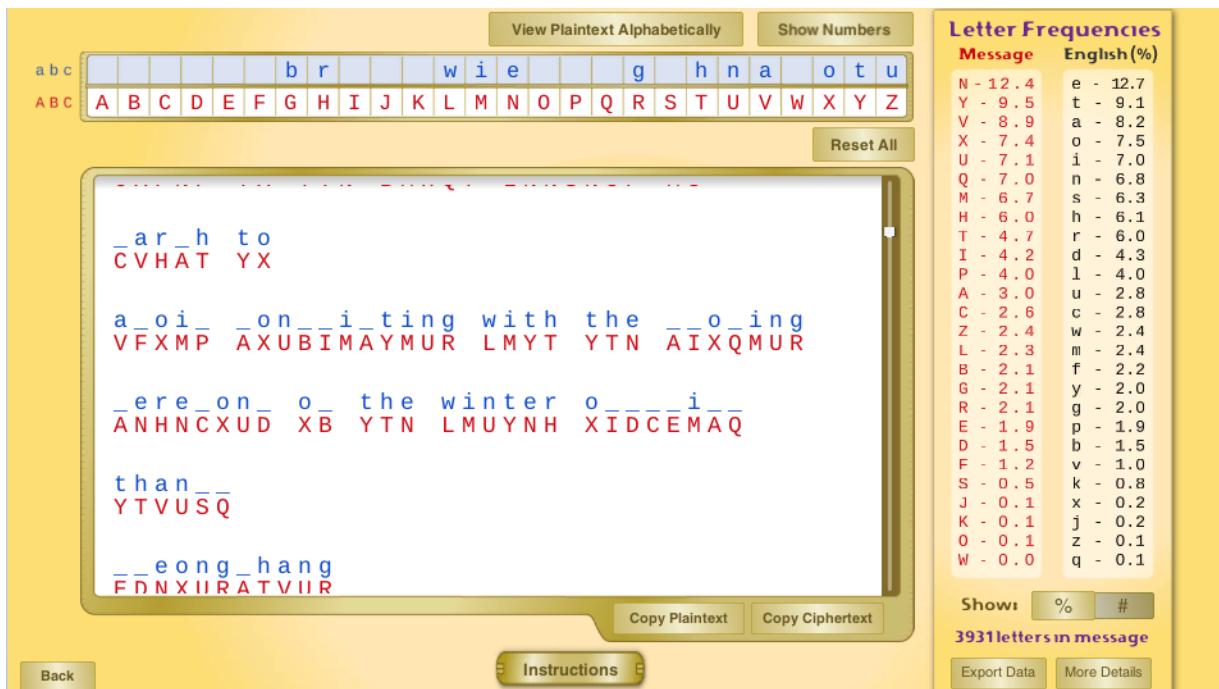


Figure 11: Mapping G to b and M to i

In Figure 11 The string 'w-nter' was shown, which should mean 'winter'. Thus M maps to i. The mapping for G to b is not shown in this picture, but was done in a very similar fashion with another word.

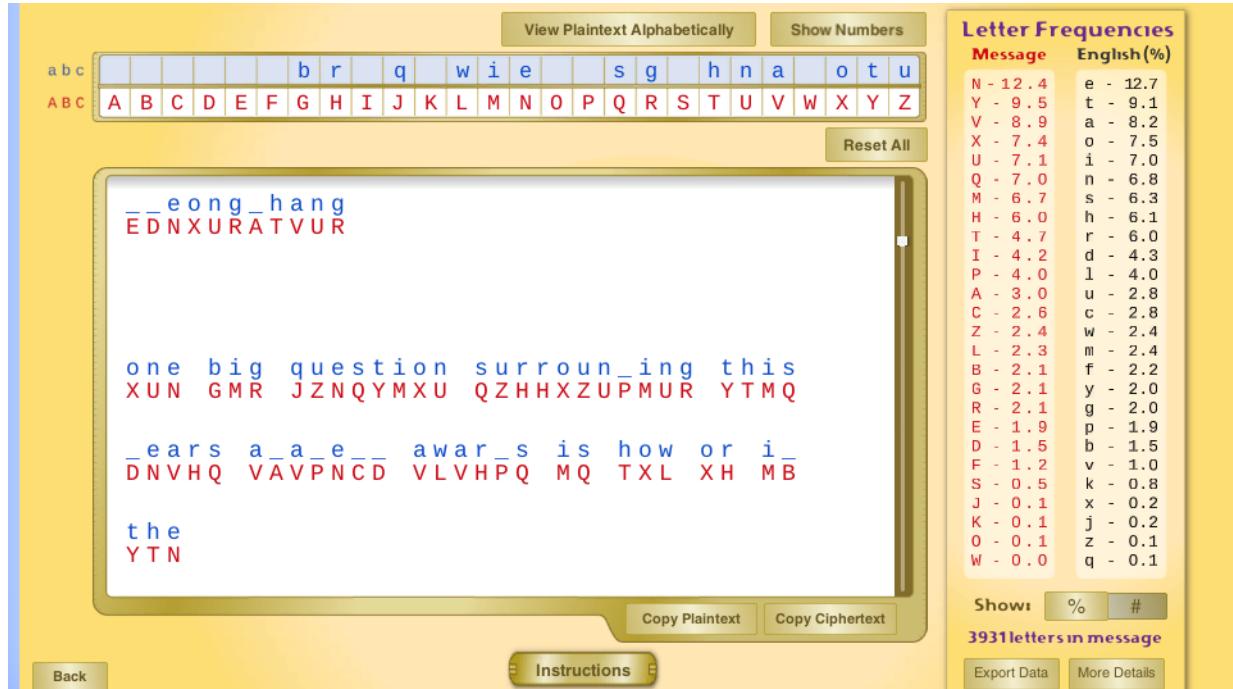


Figure 12: Mapping J to q and Q to s

In Figure 12 The string '-ue-tion' was present, which obviously meant 'question', therefore J maps to q Q maps to s.

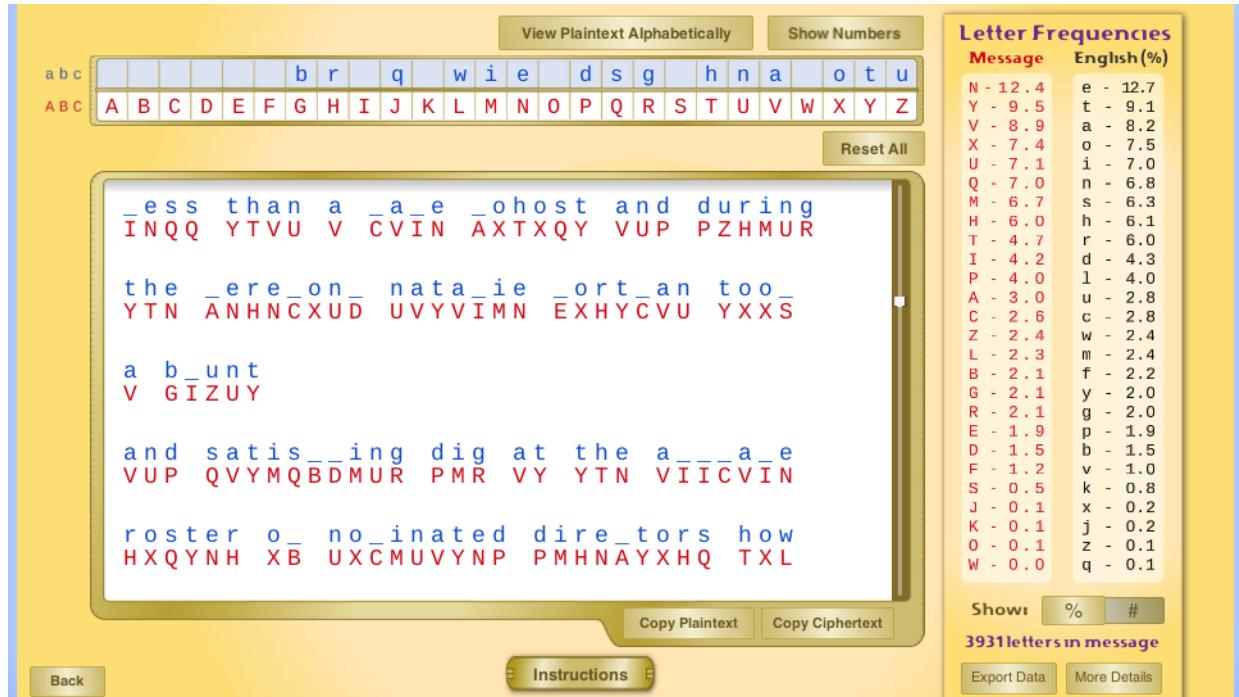


Figure 13: Mapping P to d

In Figure 13 The string 'an-' and '-ig' and '-uring' was present. Therefore, it is clear that P maps to d.

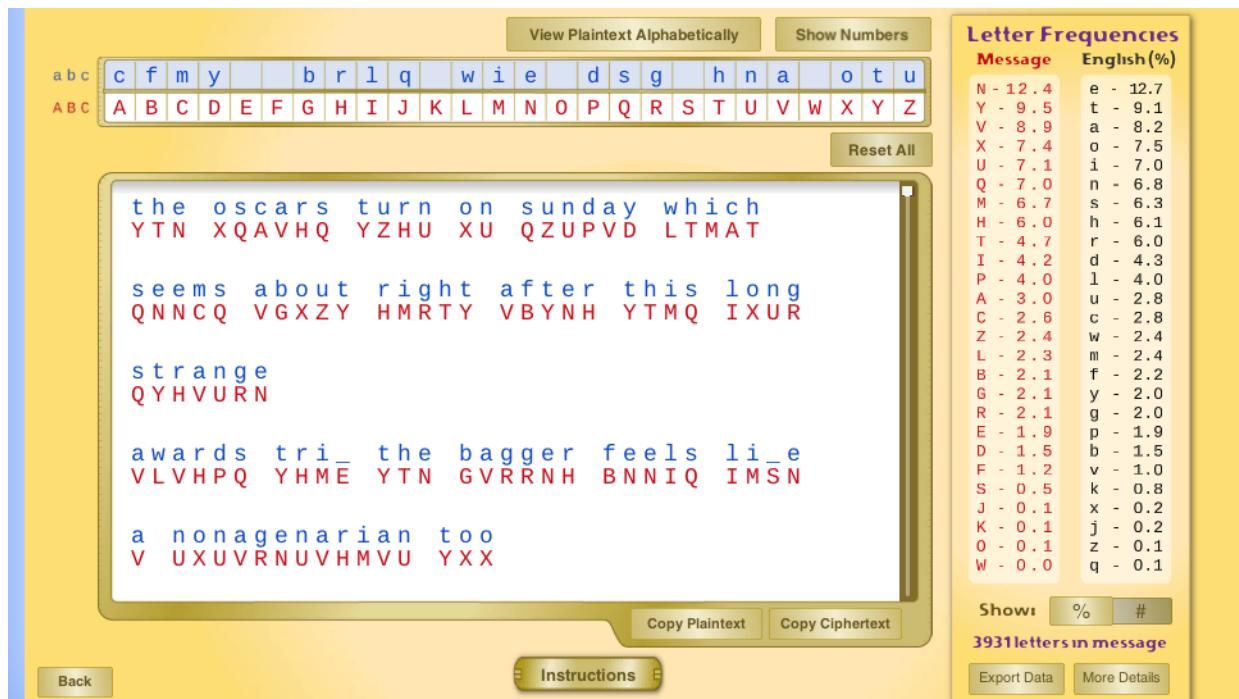


Figure 14: Mapping A to c, B to f, C to m, D to y, and I to l

In Figure 14 there are a lot of substitutions made. In particular, substitution A to c, B to f, C to m, D to y, and I to l. 'sunda-' was present which lead D to map to y, 'fee-s' which mapped I to l, 'os-ars' which mapped A to c, and the mapping from C to m is not present in Figure 14, but was done in a much similar fashion.

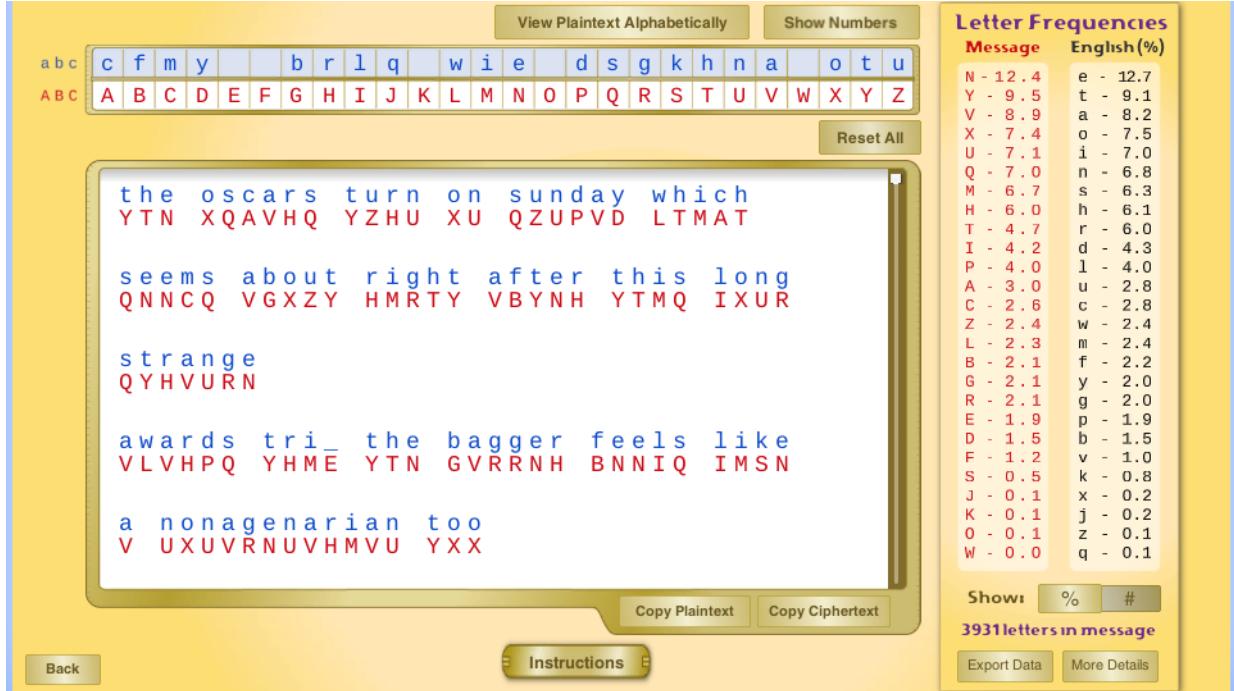


Figure 15: Mapping S to k

In figure 15 the mapping from S to k was done. This is due to the work 'li-e' being present, which is *likely* to be 'like'. Therefore, S maps to K. (get the pun)

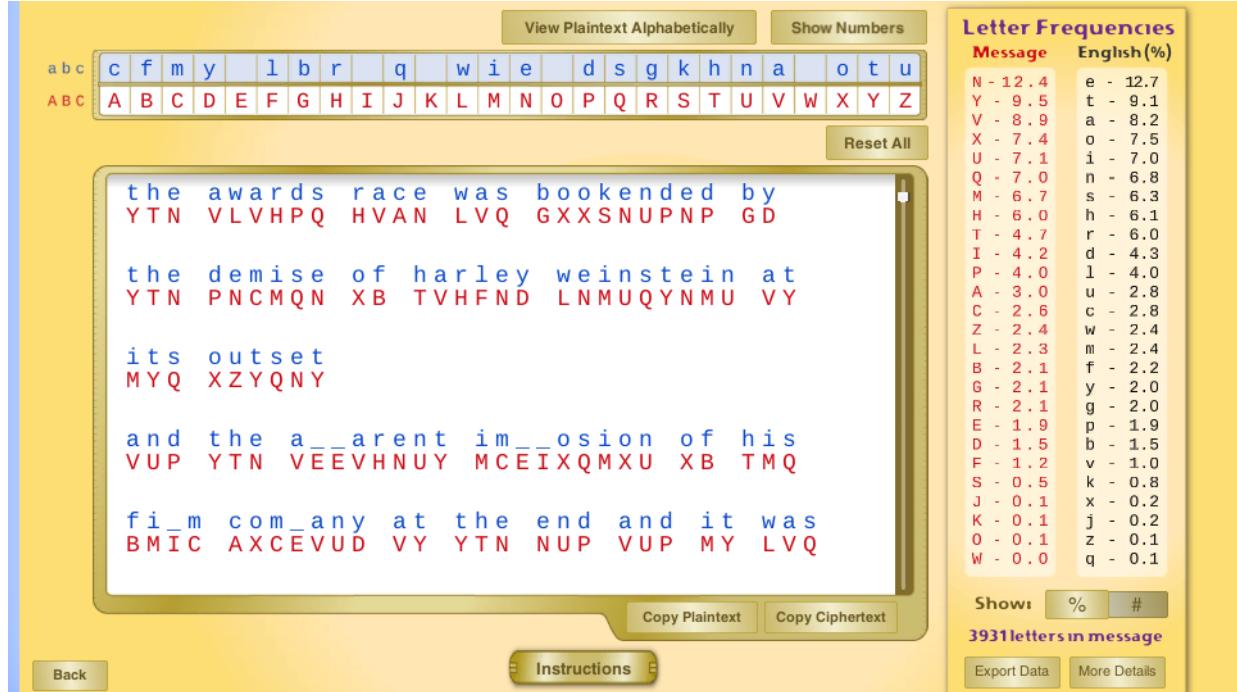


Figure 16: mapping F to l, removing I to l

In Figure 16 F is now mapped to l because the assumption was that the string 'harley' needed to be made. This removed the mapping from I to l.

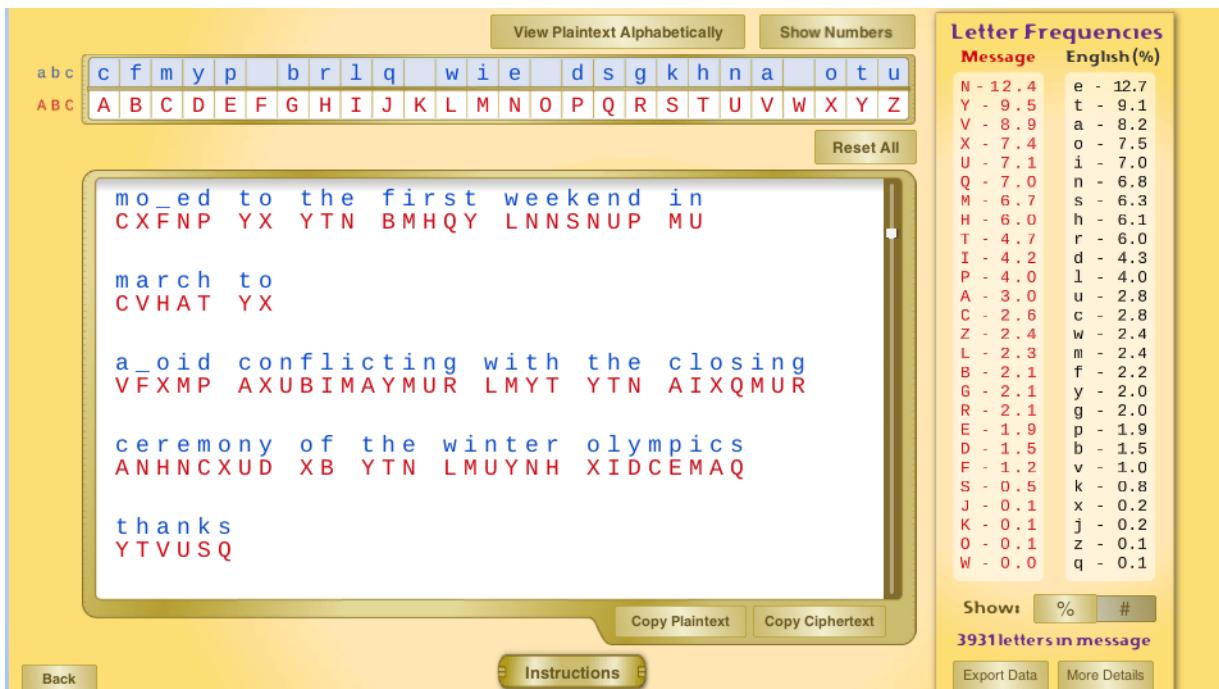


Figure 17: Mapping E to p, and re-mapping I to l

In Figure 17 E is mapped to p because 'olymp-ics' was present, which definitely means olympics. It was also noted that I should definitely be L because the string 'with the c-osing ceremony' was present.

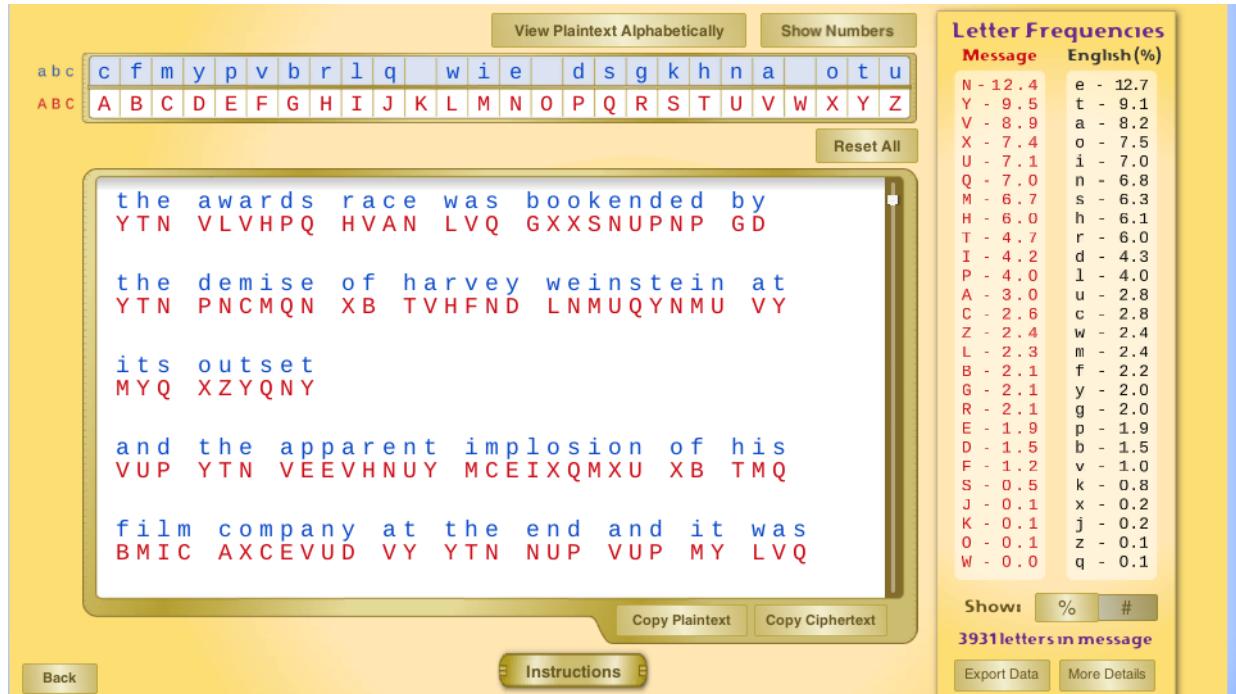


Figure 18: Mapping F to v

In Figure 18 it is clear that F should be mapped to v now, and that 'har-ey' should be 'harvey' to complete the name Harvey Weinstein.

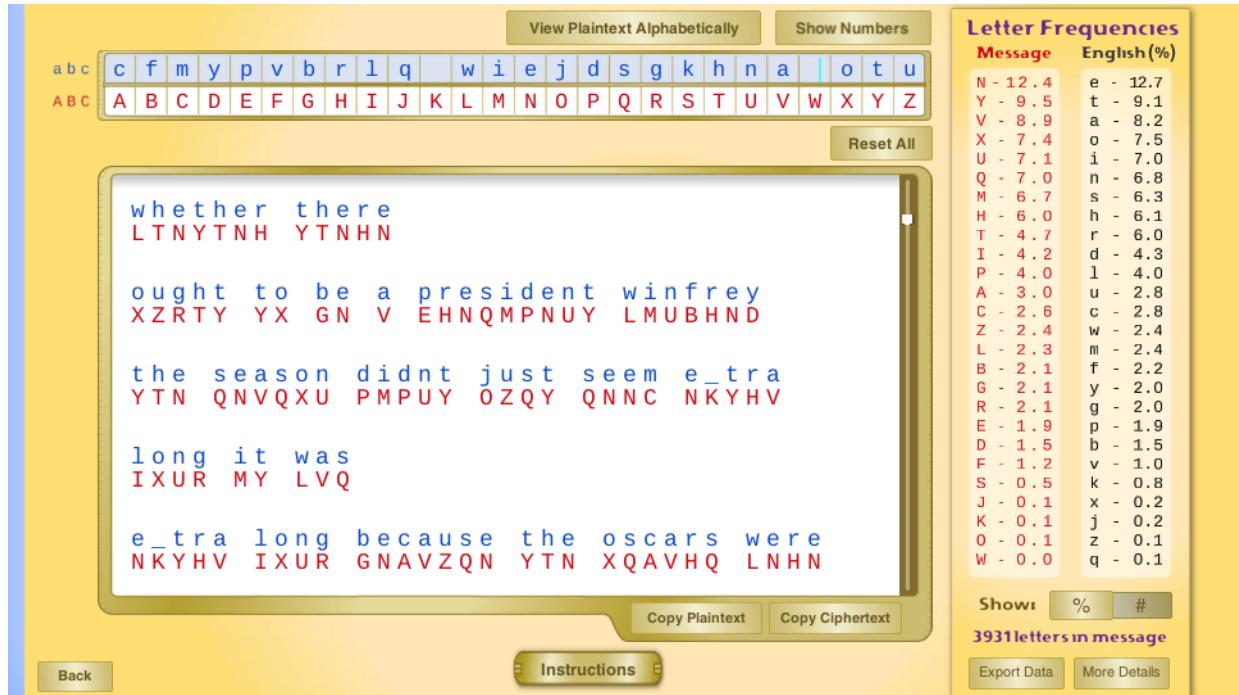


Figure 19: Mapping O to j

In Figure 19 it is clear that O should be mapped to j because the phrase 'the season didnt -ust seem' which obviously meant to say just.

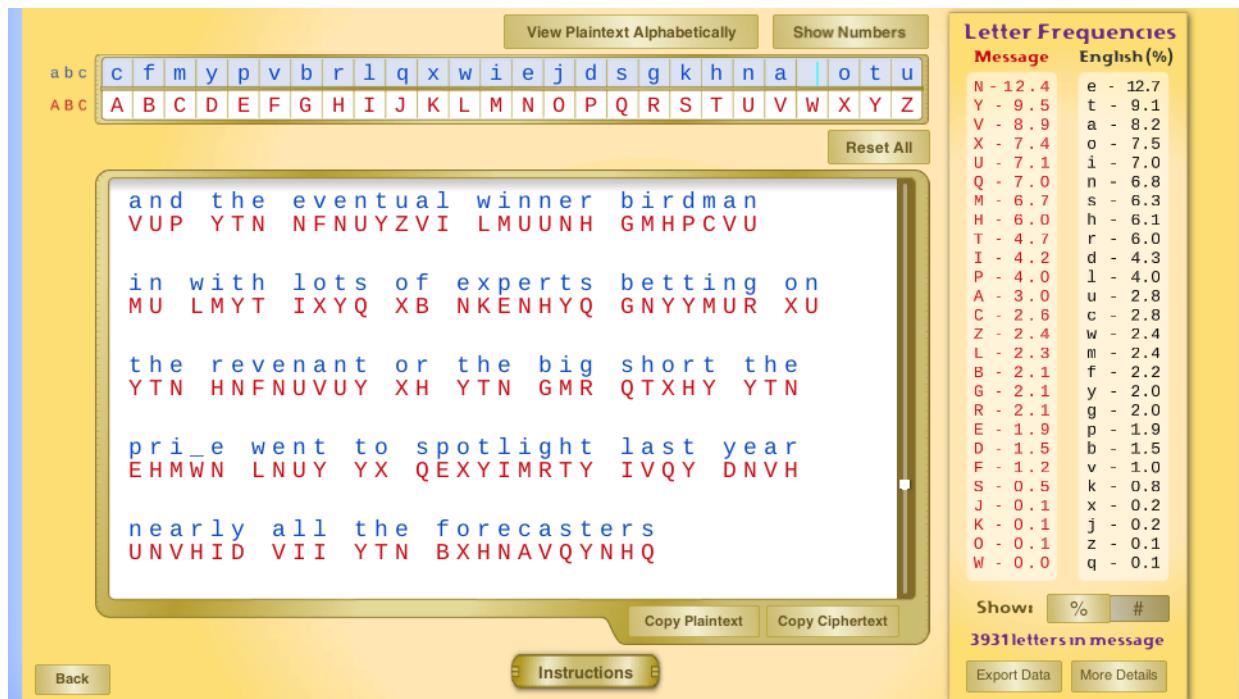


Figure 20: Mapping K to x

In Figure 20 it is clear that K should map x because the string 'e-perts' is present which should map to 'experts'.

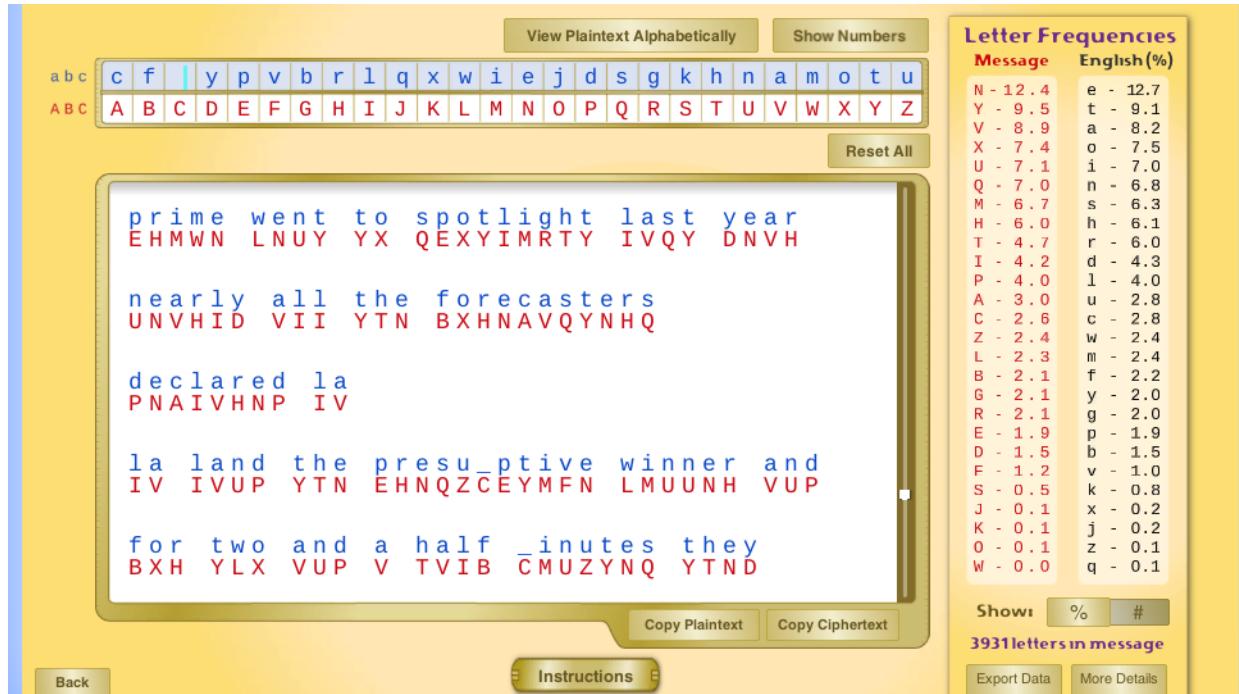


Figure 21: Mapping W to m, removing C to m

In Figure 21 it is clear that W should map to 'm' because the string 'pri-e' is present which is referring to 'prime'

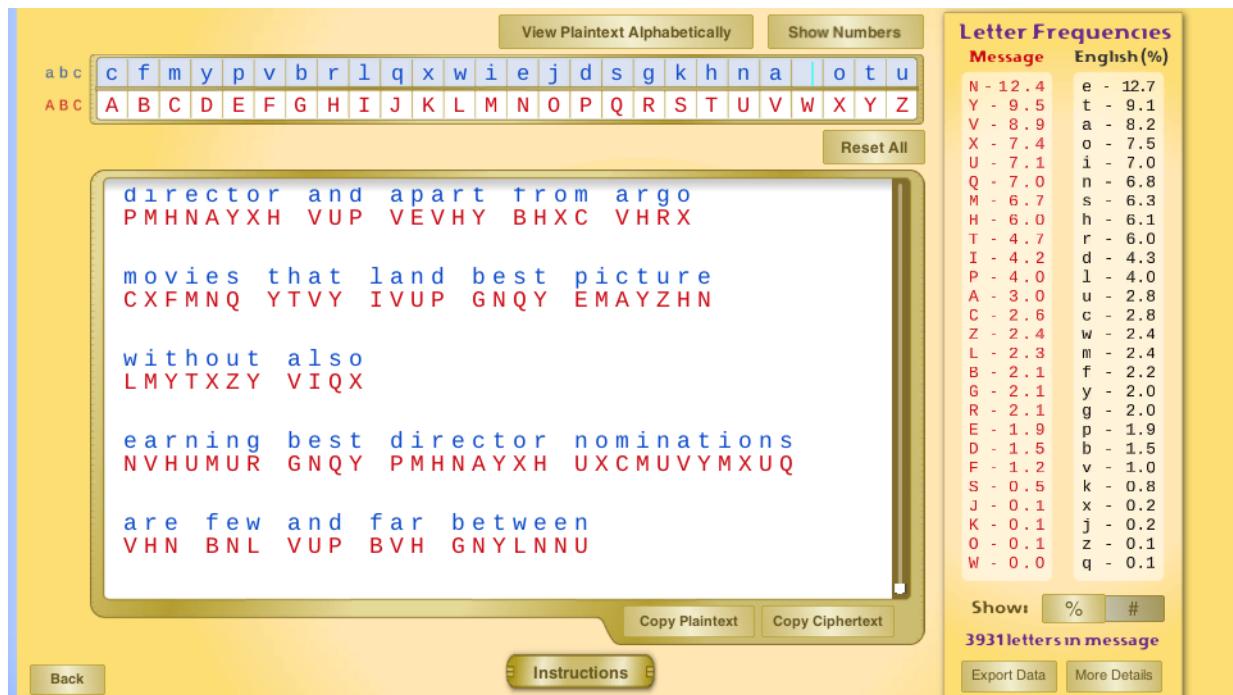


Figure 22: mapping C to m and removing W to m

In Figure 22 It is now really clear that mapping W to m was a mistake because the string 'fro- argo -ovies that' which maps to 'from argo movies that'.

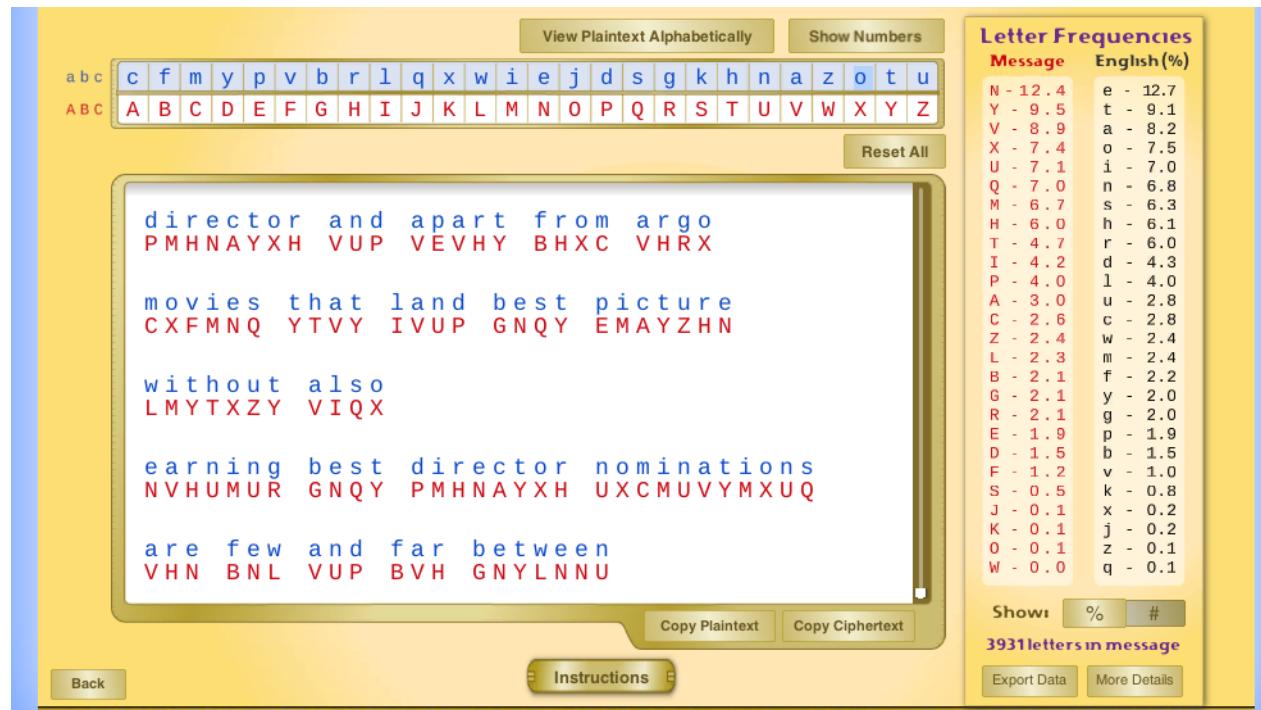


Figure 23: Mapping W to z

In Figure 23 W maps to z due to the process of elimination, because W does not show up, and neither does Z.

Finally, the key ‘abcdefghijklmnopqrstuvwxyz’ to ‘cfmvpbrlqxwiejdsgkhnazotu’ is obtained, where the first string ‘abcdefghijklmnopqrstuvwxyz’ is the ciphertext letters.

2.2 Task 2: Encryption using Different Ciphers and Modes

In this task, we will play with various encryption algorithms and modes. You can use the following openssl enc command to encrypt/decrypt a file. To see the manuals, you can type man openssl and man enc.

```
openssl enc -ciphertext -e -in plain.txt -out cipher.bin \
-K 00112233445566778889aabbcdddeeff \
-iv 0102030405060708
```

Please replace the ciphertext with a specific cipher type, such as -aes-128-cbc, -bf-cbc, -aes-128-cfb, etc. In this task, you should try at least 3 different ciphers. You can find the meaning of the command-line options and all the supported cipher types by typing ”man enc”. We include some common options for the openssl enc command in the following:

-in <file>	input file
-out <file>	output file
-e	encrypt
-d	decrypt
-K/-iv	key/iv in hex is the next argument
-[pP]	print the iv/key (then exit if -P)

2.2.1 Task2: solution

In order to complete this task, we must utilize the command man as shown in Figure 24 to find what all our possible encryption schemes are.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ man enc
```

Figure 24: using man on enc

base64	Base 64
bf-cbc	Blowfish in CBC mode
bf	Alias for bf-cbc
blowfish	Alias for bf-cbc
bf-cfb	Blowfish in CFB mode
bf-ecb	Blowfish in ECB mode
bf-ofb	Blowfish in OFB mode
cast-cbc	CAST in CBC mode
cast	Alias for cast-cbc
cast5-cbc	CAST5 in CBC mode
cast5-cfb	CAST5 in CFB mode
cast5-ecb	CAST5 in ECB mode
cast5-ofb	CAST5 in OFB mode
chacha20	ChaCha20 algorithm
des-cbc	DES in CBC mode
des	Alias for des-cbc
des-cfb	DES in CFB mode
des-ofb	DES in OFB mode
des-ecb	DES in ECB mode
des-edede-cbc	Two key triple DES EDE in CBC mode
des-edede	Two key triple DES EDE in ECB mode
des-edede-cfb	Two key triple DES EDE in CFB mode
des-edede-ofb	Two key triple DES EDE in OFB mode
des-edede3-cbc	Three key triple DES EDE in CBC mode
des-edede3	Three key triple DES EDE in ECB mode
des3	Alias for des-edede3-cbc
des-edede3-cfb	Three key triple DES EDE CFB mode
des-edede3-ofb	Three key triple DES EDE in OFB mode
desx	DESX algorithm.
gost89	GOST 28147-89 in CFB mode (provided by ccgost engine)
gost89-cnt	GOST 28147-89 in CNT mode (provided by ccgost engine)
idea-cbc	IDEA algorithm in CBC mode
idea	same as idea-cbc
idea-cfb	IDEA in CFB mode
idea-ecb	IDEA in ECB mode
idea-ofb	IDEA in OFB mode

Figure 25: encryption schemes from enc using man part 1

rc2-cbc	128 bit RC2 in CBC mode
rc2	Alias for rc2-cbc
rc2-cfb	128 bit RC2 in CFB mode
rc2-ecb	128 bit RC2 in ECB mode
rc2-ofb	128 bit RC2 in OFB mode
rc2-64-cbc	64 bit RC2 in CBC mode
rc2-40-cbc	40 bit RC2 in CBC mode
rc4	128 bit RC4
rc4-64	64 bit RC4
rc4-40	40 bit RC4
rc5-cbc	RC5 cipher in CBC mode
rc5	Alias for rc5-cbc
rc5-cfb	RC5 cipher in CFB mode
rc5-ecb	RC5 cipher in ECB mode
rc5-ofb	RC5 cipher in OFB mode
seed-cbc	SEED cipher in CBC mode
seed	Alias for seed-cbc
seed-cfb	SEED cipher in CFB mode
seed-ecb	SEED cipher in ECB mode
seed-ofb	SEED cipher in OFB mode
sm4-cbc	SM4 cipher in CBC mode
sm4	Alias for sm4-cbc
sm4-cfb	SM4 cipher in CFB mode
sm4-ctr	SM4 cipher in CTR mode
sm4-ecb	SM4 cipher in ECB mode
sm4-ofb	SM4 cipher in OFB mode
aes-[128 192 256]-cbc	128/192/256 bit AES in CBC mode
aes[128 192 256]	Alias for aes-[128 192 256]-cbc
aes-[128 192 256]-cfb	128/192/256 bit AES in 128 bit CFB mode
aes-[128 192 256]-cfb1	128/192/256 bit AES in 1 bit CFB mode
aes-[128 192 256]-cfb8	128/192/256 bit AES in 8 bit CFB mode
aes-[128 192 256]-ctr	128/192/256 bit AES in CTR mode
aes-[128 192 256]-ecb	128/192/256 bit AES in ECB mode
aes-[128 192 256]-ofb	128/192/256 bit AES in OFB mode

Figure 26: encryption schemes from enc using man part 2

```
aria-[128|192|256]-cbc 128/192/256 bit ARIA in CBC mode
aria[128|192|256]       Alias for aria-[128|192|256]-cbc
aria-[128|192|256]-cfb 128/192/256 bit ARIA in 128 bit CFB mode
aria-[128|192|256]-cfb1 128/192/256 bit ARIA in 1 bit CFB mode
aria-[128|192|256]-cfb8 128/192/256 bit ARIA in 8 bit CFB mode
aria-[128|192|256]-ctr 128/192/256 bit ARIA in CTR mode
aria-[128|192|256]-ecb 128/192/256 bit ARIA in ECB mode
aria-[128|192|256]-ofb 128/192/256 bit ARIA in OFB mode

camellia-[128|192|256]-cbc 128/192/256 bit Camellia in CBC mode
camellia[128|192|256]       Alias for camellia-[128|192|256]-cbc
camellia-[128|192|256]-cfb 128/192/256 bit Camellia in 128 bit CFB mode
camellia-[128|192|256]-cfb1 128/192/256 bit Camellia in 1 bit CFB mode
camellia-[128|192|256]-cfb8 128/192/256 bit Camellia in 8 bit CFB mode
camellia-[128|192|256]-ctr 128/192/256 bit Camellia in CTR mode
camellia-[128|192|256]-ecb 128/192/256 bit Camellia in ECB mode
camellia-[128|192|256]-ofb 128/192/256 bit Camellia in OFB mode
```

Figure 27: encryption schemes from enc using man part 3

Figure 25, Figure 26 and Figure 27 show the results of the command shown in Figure 24 which is an impressive amount of encryption schemes that could be used, most of which have 3 levels of bits that can be used to effectively make the algorithm more resistant to certain attacks such as brute force attacks.

This task requires us to try at least 3 different ciphers on the ciphertext. Just for fun, the following ciphers are used:

1. bf-cbc
 2. aria-192-ecb
 3. camellia-192-ofb

First off is blowfish-cbc! There is no reason this was picked, it just sounds fun. blowfish is a symmetric-key block cipher which was designed in 1993. Blowfish was created as an alternative to the, at the time, aging DES standard. Blowfish became popular because most other ciphers were proprietary and required money and licensing to use. Blowfish was, and will be, public domain!

Figure 28: using bf-cbc to encrypt

Figure 28 shows the results of encrypting a simple plaintext file using bf-cbc (blowfish-cbc), and the resulting cipher.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat cipher.bin && echo ""
F♦5o(v2♦jbj♦W♦
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ openssl enc -bf-cbc -d -in cipher.bin -out out.txt -K 00112233445566778899aabbccddeeff -iv 0102030405060708
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat out.txt
hello, I am plaintext!
```

Figure 29: using bf-cbc to decrypt

Figure 29 shows the process of decrypting blowfish-cbc using openssl. It is extremely similar to encrypting, the only difference is swapping the in/out file, and adding -d instead of -e. Just for fun, the time of encrypting and decrypting the program was done using the linux time command, which is shown below.

```
time openssl enc -bf-cbc -d -in cipher.bin -out out.txt \
-K 00112233445566778889aabbccddeeff \
-iv 0102030405060708
```

Encryption time: 0.004s

Decryption Time: 0.004s

In the end, blowfish-cbc is a very fun algorithm with a fun background. It's a shame it's not seen as much.

Now on to aria-192-ecb!

The process for aria-192-ecb is *extremely* similar to blowfish-cbc as shown in Figure 28. Therefore, for this section no pictures will be shown, but rather the commands will be simply printed as text.

```
$ openssl enc -aria-192-ecb -d -in plain.txt -out aria-192-ecb_cipher.bin \
-K 00112233445566778889aabbcdddeeff \
-iv 0102030405060708
warning: iv not use by this cipher
hex string is too short, padding with zero bytes to length
```

Here, openssl outputs some interesting comments on the program, stating no IV is used, and the hex string is too short. Let's try that again

```
$ openssl enc -aria-192-ecb -d -in plain.txt -out aria-192-ecb_cipher.bin \
      -K 00112233445566778889aabbcdddeefffffffefffffefff
$ cat plain.txt
hello, I am plaintext!
$ cat aria-192-ecb_cipher.bin
6&eI-R^:&
```

There we go, that is much more like it!

And of course, the process to decrypt is very similar to the encryption process.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat aria-192-ecb_cipher.bin && echo ""  
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ time openssl enc -aria-192-ecb -d -in aria-192-ecb_cipher.bin -out out.txt -K 00112233445566778889aabccddeeffffffeffffff  
  
real      0m0.006s  
user      0m0.005s  
sys       0m0.001s  
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat out.txt  
hello, I am plaintext!
```

Figure 30: using aria-192-ecb to decrypt

Figure 30 shows the process of decrypting with aria-192-ecb. time was utilized again in order to determine how long aria-192-ecb takes.

Encryption Time: 0.010s

Decryption Time: 0.006s

Now for the final algorithm to check out, camellia-192-ofb!

```
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ openssl enc -camellia-192-ofb -e -in plain.txt -out camellia-256-ofb_cipher.bin -K 0011223344556677889aabbcdddeeffffffefffffff -iv 01020304050607081234132123432145
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat plain.txt
hello, I am plaintext!
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat camellia-256-ofb_cipher.bin && echo ""
98M99959999
D-9999
```

Figure 31: using camellia-256-ofb to encrypt

Figure 31 shows the process of encrypting with camellia-192-ofb.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ openssl enc -camellia-192-ofb -d -in cam  
ellia-256-ofb_cipher.bin -out out.txt -K 00112233445566778899aabbcdddeeffffffeffffffeff -iv 0  
1020304050607081234132123432145  
mitch@light: ~/git/ECE471/lab1/task2-6/t2 (master) $ cat out.txt  
hello. I am plaintext!
```

Figure 32: using camellia-256-ofb to decrypt

Figure 32 shows the process of decrypting with camellia-192-ofb. And of course, ztime was used for both. The results are shown below.

Encrypting: 0.003s

Decrypting: 0.008s

2.3 Task 3: Encryption Mode – ECB vs. CBC

The file *pic_original.bmp* can be downloaded from my github github.com/mitchdz/ECE471, and it contains a simple picture. We would like to encrypt this picture, so people without

the encryption keys cannot know what is in the picture. Please encrypt the file using the ECB (Electronic Code Book) and CBC (Cipher Block Chaining) modes, and then do the following:

1. Let us treat the encrypted picture as a picture, and use a picture viewing software to display it. However, For the .bmp file, the first 54 bytes contain the header information about the picture, we have to set it correctly, so the encrypted file can be treated as a legitimate .bmp file. We will replace the header of the encrypted picture with that of the original picture. We can use the bless hex editor tool (already installed on our VM) to directly modify binary files. We can also use the following commands to get the header from p1.bmp, the data from p2.bmp (from offset 55 to the end of the file), and then combine the header and data together into a new file.

```
$ head -c 54 p1.bmp > header  
$ tail -c +55 p2.bmp > body  
$ cat header body > new.bmp
```

2. Display the encrypted picture using a picture viewing program (we have installed an image viewer program called eog on our VM). Can you derive any useful information about the original picture from the encrypted picture? Please explain your observations.

Select a picture of your choice, repeat the experiment above, and report your observations.

2.3.1 Task 3: Solution



Figure 33: Original image that is being encrypted

Figure 33 shows the original file. The file is a red oval in the center, with a teal square on the bottom right that has a blue border.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ openssl enc -aes-128-cbc -e -in pic_original.bmp -out cbc_cipher.bmp -K 00112233445566778899aabccddeeff -iv 01020304050607081234123412341212  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ head -c 54 pic_original.bmp > header  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ tail -c +55 cbc_cipher.bmp > body  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ cat header body > cbc_cipher.bmp
```

Figure 34: using cbc to encrypt image

Figure 34 shows the commands used in order to encrypt the image using cbc.



Figure 35: results of cbc on image

Figure 35 shows the results of encrypting the image using cbc. There is not much data that is able to be extracted from this image.

Similar to above, ecb (electronic codebook is now used)

```
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ openssl enc -aes-128-ecb -e -in pic_original.bmp -out ecb_cipher.bmp -K 00112233445566778889aabccddeff -iv 01020304050607081234123412341212
warning: iv not use by this cipher
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ head -c 54 pic_original.bmp > header
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ tail -c +55 ecb_cipher.bmp > body
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ cat header body > ecb_cipher.bmp
```

Figure 36: encrypting image using ECB

Figure 36 shows the command used to encrypt the image using ECB. Note the comment saying that the iv is not used by this cipher. Unlike CBC, ECB does not need an initial value.

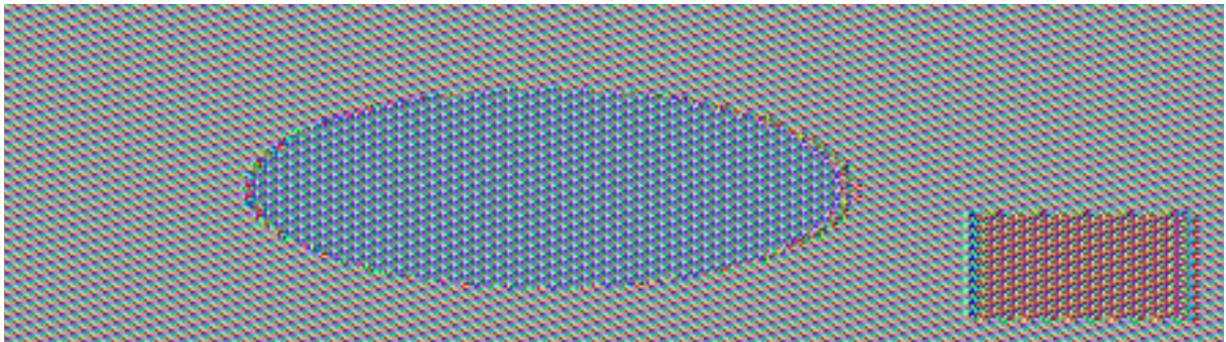


Figure 37: results of ECB on image

Figure 37 shows the resulting ciphertext after encrypting the image with ECB. It is very easy to identify what the original image is. Although some details such as the blue border

are gone, and the colors of the oval and the square have changed. If this was a document of a person, the person would most likely be easily identifiable.

Now it is time to try this experiment on my own image!



Figure 38: Personal image to be encrypted

Figure 38 shows the personal image used to do this experiment on. This image is the background that I use for my laptop.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ openssl enc -aes-128-cbc -e -in t3_personal_image.bmp -out personal_cbc_cipher.bmp -K 00112233445566778889aabbccddeeff  
iv undefined  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ head -c 54 t3_personal_image.bmp > header  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ tail -c +55 personal_cbc_cipher.bmp > body  
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ cat header body > t3_personal_cbc.bmp
```

Figure 39: Encrypting personal image using CBC

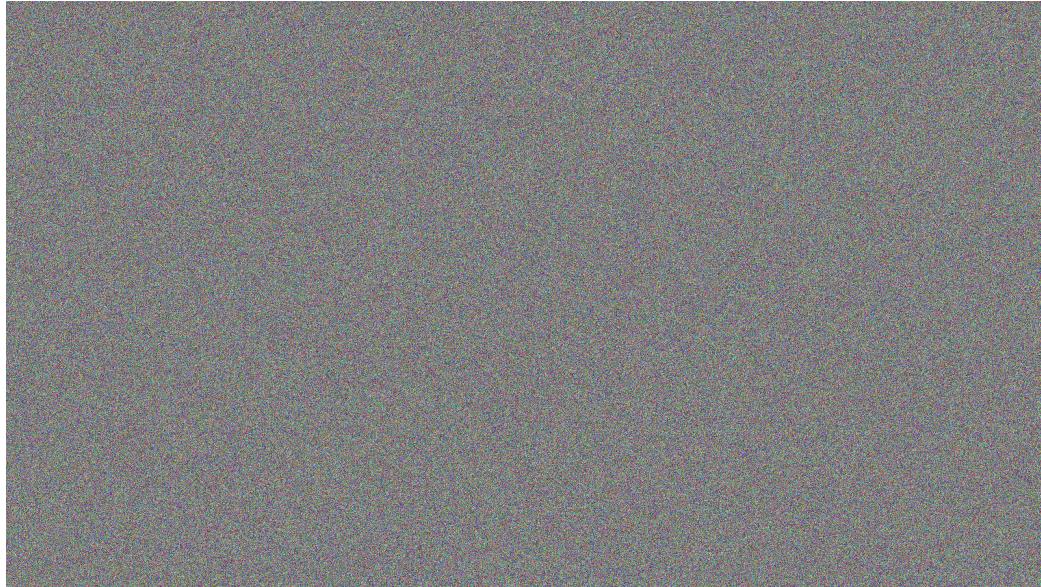


Figure 40: results of encrypting Figure 38 with CBC

Figure 40 shows the results of encrypting Figure 38 with CBC. The result is, as expected, seemingly random. There is nothing to be able to determine about the picture from viewing this image.

```
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ openssl enc -aes-128-ecb -e -in t3_personal_image.bmp -out personal_ecb_cipher.bmp -K 00112233445566778889aabccddeff
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ head -c 54 t3_personal_image.bmp > header
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ tail -c +55 personal_ecb_cipher.bmp > body
mitch@light: ~/git/ECE471/lab1/task2-6/t3 (master) $ cat header body > new.bmp
```

Figure 41: Encrypting personal image using ECB

Figure 41 shows the process of encrypting the personal image using ECB. This process is very similar to Figure 36 and Figure 34.

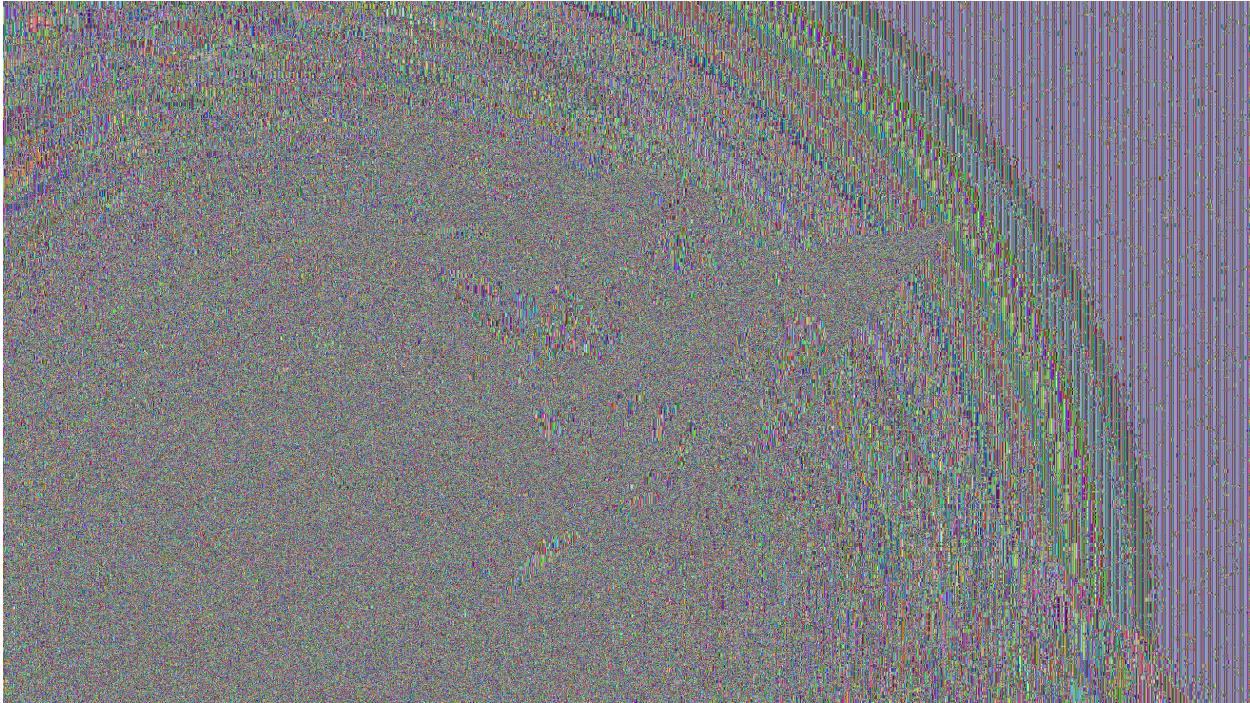


Figure 42: Result of encrypting with ECB

Figure 42 shows the results of encrypting Figure 38 with ECB, which honestly is not too easy to identify what the image should be. This leads me to believe that ECB only reveals the contents via inspection if there is a large enough contrast between objects and colors.

2.4 Task 4: Padding

For block ciphers, when the size of a plaintext is not a multiple of the block size, padding may be required. All the block ciphers normally use PKCS5 padding, which is known as standard block padding. We will conduct the following experiments to understand how this type of padding works:

1. Use ECB, CBC, CFB, and OFB modes to encrypt a file (you can pick any cipher). Please report which modes have paddings and which ones do not. For those that do not need paddings, please explain why.
2. Let us create three files, which contain 5 bytes, 10 bytes, and 16 bytes, respectively. We can use the following "echo -n" command to create such files. The following example creates a file f1.txt with length 5 (without the -n option, the length will be 6, because a newline character will be added by echo):

```
$ echo -n "12345" > f1.txt
```

We then use "openssl enc -aes-128-cbc -e" to encrypt these three files using 128-bit AES with CBC mode. Please describe the size of the encrypted files.

We would like to see what is added to the padding during the encryption. To achieve this goal, we will decrypt these files using "openssl enc -aes-128-cbc -d". Unfortunately, decryption by default will automatically remove the padding, making it impossible for us to see the padding. However, the command does have an option called "-nopad", which disables the padding, i.e., during the decryption, the command will not remove the padded data. Therefore, by looking at the decrypted data, we can see what data are used in the padding. Please use this technique to figure out what paddings are added to the three files.

It should be noted that padding data may not be printable, so you need to use a hex tool to display the content. The following example shows how to display a file in the hex format:

```
$ hexdump -C p1.txt
00000000 31 32 33 34 35 36 37 38 39 49 4a 4b 4c 0a  |123456789IJKL.|
```

```
$ xxd p1.txt
00000000: 3132 3334 3536 3738 3949 4a4b 4c0a          123456789IJKL.
```

2.4.1 Task 4: Solution

2.5 Task 5: Error Propagation – Corrupted Cipher Text

2.5.1 Task 5: Solution

2.6 Task 6: Initial Vector (IV)

2.6.1 Task 6: Solution