



Modular Algorithms in Quadratic Algebraic Number Fields

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Computer Algebra









Theorem (Chinese Remainder Theorem (CRT))

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Essentially...there is a u satisfying u \equiv u_1 \pmod{m_1}, \vdots u \equiv u_n \pmod{m_n}, when m_i and m_i are coprime.
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Proof sketch.

$$\phi: \mathbb{Z}_{m_1\cdots m_n} \to \mathbb{Z}_{m_1} \times \cdots \times \mathbb{Z}_{m_n}$$

$$u \mapsto (u \mod m_1, \ldots, u \mod m_n)$$

is an isomorphism.





Find *u* with

$$u \equiv 49 \pmod{32771},$$

 $u \equiv -21 \pmod{65537},$
 $u \equiv -30 \pmod{131101}.$

$$u = v_0 + v_1(m_1) + v_2(m_1m_2) + \dots + v_{n-1}(\prod_{i=1}^{n-1} m_i)$$

$$-104905043721354 = 49 - 28(32771) - 48845(32771 \cdot 65537)$$







Modular Algorithms





$$22x + 44y + 74z = 1,$$

$$15x + 14y - 10z = -2,$$

$$-25x - 28y + 20z = 34.$$





Fraction-free Gaussian elimination:

$$1257315840x = 7543895040,$$

$$-57150720y = 314328960,$$

$$162360z = 243540.$$

$$\left\{x = 6, \quad y = -\frac{11}{2}, \quad z = \frac{3}{2}\right\}$$



Solve
$$A\vec{x} = \vec{b}$$
 [TI61],[CL77].

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1: p_1, \ldots, p_k \leftarrow \text{distinct primes}

2: for i = 1, \ldots, n do

3: d_i \leftarrow \det A \pmod{p_i}

4: \vec{x_i} \leftarrow \text{solve } A\vec{x} = \vec{b} \pmod{p_i}

5: D \leftarrow \text{CRT on } \{d_1 \pmod{p_1}, \ldots, d_n \pmod{p_n}\}

6: \vec{y} \leftarrow \text{pointwise CRT on } \{d_1 \cdot \vec{x_1} \pmod{p_1}, \ldots, d_n \cdot \vec{x_n} \pmod{p_n}\}

7: return \frac{1}{D} \cdot \vec{y}
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Modular linear solver:

$$\begin{bmatrix} x \\ y \\ z \end{bmatrix} = \frac{1}{-7380} \begin{bmatrix} -44280 \\ 40590 \\ -11070 \end{bmatrix} = \begin{bmatrix} 6 \\ -11/2 \\ 3/2 \end{bmatrix}$$











$$\mathbb{Q}(\sqrt{M})$$



?









Quadratic Integers







Let $M \in \mathbb{Z}$ be squarefree. The quadratic integers of $\mathbb{Q}(\sqrt{M})$ are

$$R = \begin{cases} \mathbb{Z}[\sqrt{M}] & M \equiv 2, 3 \pmod{4} \\ \mathbb{Z}[\frac{1+\sqrt{M}}{2}] & M \equiv 1 \pmod{4} \end{cases}$$

[Mar77]. Write $R = \mathbb{Z}[\gamma]$.

- ▶ What are the primes?
- ▶ Are any primes $p \in \mathbb{Z}$ also prime in R (inert)?









Chebotarëv Density Theorem implies infinitely many inert primes.

Algorithm: ([IR13] Proposition 5.2.2)

▶ Input: Set of odd primes *P*.

Output: An odd inert prime not contained in P.







Proposition

Let $p \in \mathbb{Z}$ be an inert prime in $R = \mathbb{Z}[\gamma]$. Then

$$R/pR \simeq egin{cases} \mathbb{Z}_p[x]/\langle x^2-M
angle & M \equiv 2,3 \pmod{4} \ \mathbb{Z}_p[x]/\langle x^2-x-rac{M-1}{4}
angle & M \equiv 1 \pmod{4} \end{cases}$$

Moreover, the isomorphism is given by

$$\phi([a+b\gamma])=[a+bx].$$







Modular Algorithms in $\mathbb{Q}(\sqrt{M})$



- ▶ Use Garner's integer CRT algorithm in $\mathbb{Z}[\gamma]$ (for inert primes).
- ▶ Use the same modular linear solver!
- ▶ No unique factorisation in *R*.
- ► No GCD.
- ▶ Rational reconstruction $\mathbb{Z}_k[\gamma] \leadsto \mathbb{Q}[\gamma] = \mathbb{Q}(\sqrt{M})$.



Solve Ax = b, where:

$$A = \begin{pmatrix} 81 - 19\sqrt{122} & 78 - 89\sqrt{122} & -81 - 80\sqrt{122} \\ 22 - 53\sqrt{122} & -8 + 66\sqrt{122} & -43 - 19\sqrt{122} \\ 50 - 30\sqrt{122} & -90 + 154\sqrt{122} & -2 - 124\sqrt{122} \end{pmatrix}, \quad b = \begin{pmatrix} 26851 - 2700\sqrt{122} \\ -41098 + 883\sqrt{122} \\ -67029 + 1076\sqrt{122}. \end{pmatrix}$$

$$x = \begin{pmatrix} \frac{5\sqrt{122}}{2} \\ \frac{33}{2} - 3\sqrt{122} \\ 15 \end{pmatrix}$$

- ▶ Fraction-free Gaussian elimination: 105-bit integers
- ▶ Modular solver: 57-bit integers (seven 16-bit primes)











- [CL77] S Cabay and TPL Lam. Congruence techniques for the exact solution of integer systems of linear equations. *ACM Transactions on Mathematical Software (TOMS)*, 3(4):386–397, 1977.
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- [Mar77] D.A. Marcus. *Number Fields*. Springer-Verlag, 1977.
 - [TI61] H Takahasi and Y Ishibashi. A new method for 'exact computation' by a digital computer. In *Inf. Processing in Japan*, pages 28–42, 1961.



