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CPTS 260 Homework 3
WSU
Part 1: Intro To Assembly
1.
   f = $s0
   g = $s1
   h = s2
   i = $s3
   i = $s4
   Base address for arrays A = \$s6, and B = \$s7
   t0 = f * 4;
   t0 = A[f];
   t1 = g * 4;
   t1 = B[g];
   f = A[f];
   t2 = A[f + 1];
   t0 = A[f + 1];
   t0 = A[f + 1] + A[f];
   B[g] = A[f+1] + A[f];
2.
A.
   $s0 = 0x80000000 in Hex
   $s0 = 2147483648 in Dec
   s1 = 0xD0000000 in Hex
   $s1 = 3489660928 \text{ in Dec}
   2147483648 + 3489660928 = 5637144576
    1000 0000 0000 0000 0000 0000 0000 0000
    1101 0000 0000 0000 0000 0000 0000 0000
  1 0101 0000 0000 0000 0000 0000 0000 = 5637144576
```

There is overflow because the carry in and carry out bits are not equal.

```
В.
   $s0 = 0x80000000 in Hex
   $s0 = 2147483648 in Dec
   $s1 = 0xD0000000 in Hex
   $s1 = 3489660928 in Dec
   2147483648 - 3489660928 = -1342177280
   1000 0000 0000 0000 0000 0000 0000 0000
   0011 0000 0000 0000 0000 0000 0000 0000
                                  2's compliment
   There is no overflow because there were no extra bits that went into the carry in or carry out.
C.
   First MiPs line
   $s0 = 0x800000000 in Hex
   $s0 = 2147483648 in Dec
   $s1 = 0xD0000000 in Hex
   $s1 = 3489660928 in Dec
   2147483648 + 3489660928 = 5637144576
   1000 0000 0000 0000 0000 0000 0000 0000
   1101 0000 0000 0000 0000 0000 0000 0000
  1 0101 0000 0000 0000 0000 0000 0000 =
                                 5637144576
   Second MiPs line
   $s0 = 0x80000000 in Hex
   $s0 = 2147483648 in Dec
   t0 = 0x150000000 in Hex
   t0 = 5637144576 in Dec
   2147483648 + 5637144576 = 7784628224
   0000 1000 0000 0000 0000 0000 0000 0000 0000
   0001 0101 0000 0000 0000 0000 0000 0000 0000
   0001 1101 0000 0000 0000 0000 0000 0000 = 7784628224
```

There is no overflow because there were no extra bits that went into the carry in or carry out.

```
3.
```

## Part 2: Intro to MIPS

## 1. A.

ор	rs	rt	rd	shamt	funct
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

op is the ADDi instruction rs is 9 which is \$t1 rt is 9 which is \$t1

The rest of the 16 bits once translated to dec is 1 which is used as the constant for the addi command.

The instruction is ADDi \$t1, \$t1, 1

NAME	NUMBER		
\$zero	0		
\$at	1		
\$v0-\$v1	2-3		
\$a0-\$a3	4-7		
\$t0-\$t7	8-15		
\$s0-\$s7	16-23		
\$t8-\$t9	24-25		
\$k0-\$k1	26-27		
\$gp	28		
\$sp	29		
\$fp	30		
\$ra	31		

B.

```
sw $t5, 32($t2)
```

sw according to the reference sheet has a hex of 2b. \$t5 = 21Dec 32 = \$t2 = 10Dec

- The offset is 6 bits so (31 through 26) is 101011 for the sw instruction
- The following next five bits is \$t2 or 10 or 01010 in binary
- The following next five bits is \$t5 or 21 or 01101 in binary
- The following 16 bits are used for the offset and it's 32 so it's 0000 0000 0010 0000 in binary

The whole string is 101011 01010 01101 0000 0000 0010 0000 When translated to hexadecimal it's AD5D0020

C.

ор	rs	rt	rd	shamt	funct
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

$$op = 0$$
,  $rs = 3$ ,  $rt = 2$ ,  $rd = 13$ ,  $Shamt = 0$ ,  $funct = 36$ 

op is 6 bits so it's 000000 rs is 5 bits and 3 so it's 00011 rt is 5 bits and 2 so it's 00010 Shamt is 5 bits and 0 so it's 00000 funky is 6 bits and 36 so it's 100100

all together it's: 000000 00011 00010 01101 00000 100100 once translated to the MIPS command it's AND \$t4, \$v1, \$v0

2.

A.

I think the I-format would be the best because there are arrays in the command and this instruction format has space set to keep the address of the array available.

В.

loop: slt \$t1, \$t2, 1 beg #t1, \$zero, Done addi \$t2, \$t2, -1 j loop:

3.

x = \$s1 y = \$s2 A = \$s3 add \$t0, \$s3, \$t0 lw \$t0, 5(\$t0)