a.)

Say P(T) is the probability of hitting the target, and P(W) is the probability of it being windy. We are given that P(T|W) = 0.4 and  $P(T|W^C) = 0.7$ , as well as that P(W) = 0.3. We note that conditional probability theory states that in general,  $P(A|B) = \frac{P(A \cap B)}{P(B)}$ .

(i) on a given shot there is a gust of wind and she hits her target;  $P(T \cap W)$ 

$$P(T|W) = \frac{P(T \cap W)}{P(W)}$$

$$P(T \cap W) = P(T|W) \cdot P(W)$$

$$P(T \cap W) = (0.4)(0.3)$$

$$\boxed{P(T \cap W) = 0.12}$$

(ii) she hits the target with her first shot; P(T)

$$P(T) = P(T \cap W) + P(T \cap W^{C})$$

$$P(T) = P(T|W) \cdot P(W) + P(T|W^{C}) \cdot P(W^{C}), \quad P(W^{C}) = 1 - P(W) = 0.7$$

$$P(T) = (0.4)(0.3) + (0.7)(0.7)$$

$$\boxed{P(T) = 0.61}$$

(iii) she hits the target exactly once in two shots;  $P(T) \cdot P(T^C) + P(T^C) \cdot P(T)$ 

$$P(T) \cdot P(T^C) + P(T^C) \cdot P(T) = 2 \cdot P(T) \cdot P(T^C), \quad P(T^C) = 1 - P(T) = 0.39$$

$$P(T) \cdot P(T^C) + P(T^C) \cdot P(T) = 2(0.61)(0.39)$$

$$\boxed{P(T) \cdot P(T^C) + P(T^C) \cdot P(T) = 0.4758}$$

(iv) there was no gust of wind on an occasion when she missed;  $\frac{P(T^C \cap W^C)}{P(T^C)}$ 

$$\frac{P(T^C \cap W^C)}{P(T^C)} = \frac{P(T^C | W^C)P(W^C)}{P(T^C)}, \quad P(T^C | W^C) = 1 - P(T | W^C) = 0.3$$

$$\frac{P(T^C \cap W^C)}{P(T^C)} = \frac{(0.3)(0.7)}{(0.39)}$$

$$\frac{P(T^C \cap W^C)}{P(T^C)} = 0.538$$

**b.**)

We are given

$$P(A|B,C) = P(A|B \cap C) > P(A|B)$$

and from the properties of conditional probability we find

$$P(A|B \cap C) = \frac{P(A \cap B \cap C)}{P(B \cap C)} > \frac{P(A \cap B)}{P(B)}.$$
 (1)

We are looking to show

$$P(A|B, C^C) = P(A|B \cap C^C) < P(A|B).$$

or equivalently

$$P(A|B \cap C^C) = \frac{P(A \cap B \cap C^C)}{P(B \cap C^C)} < \frac{P(A \cap B)}{P(B)}.$$
 (2)

Returning to equation 1, we find through algebraic manipulation that

$$P(A \cap B \cap C) > \frac{P(A \cap B)P(B \cap C)}{P(B)}$$

$$-P(A \cap B \cap C) < -\frac{P(A \cap B)P(B \cap C)}{P(B)}$$

$$P(A \cap B) - P(A \cap B \cap C) < P(A \cap B) - \frac{P(A \cap B)P(B \cap C)}{P(B)}$$

$$P(A \cap B) - P(A \cap B \cap C) < \frac{P(A \cap B)}{P(B)} (P(B) - P(B \cap C))$$

$$\frac{P(A \cap B) - P(A \cap B \cap C)}{P(B) - P(B \cap C)} < \frac{P(A \cap B)}{P(B)}.$$
(3)

Since in general,

$$P(X\cap Y) + P(X\cap Y^C) = P(X),$$
 (also expressed as 
$$P(X\cap Y^C) = P(X) - P(X\cap Y))$$

we can express equation 3 as

$$\frac{P(A\cap B\cap C^C)}{P(B\cap C^C)}<\frac{P(A\cap B)}{P(B)},$$

equivalent to equation 2.

a.)

Assume a positive semidefinite matrix  $A \in \mathbb{R}^{n \times n}$  so that  $x^T A x \ge 0$ . By definition, we say  $A \succeq 0$  (i).

(i)⇒(ii):

Since, by definition,  $A \succeq 0 \Rightarrow x^T A x \geq 0$ , we could let x = By where  $B \in \mathbb{R}^{n \times n}$  is invertible. Then

$$x^T A x = (By)^T A (By) \ge 0$$

$$y^T B^T A B y \ge 0$$

This is of the form defining a semidefinite matrix, implying that  $B^TAB \succeq 0$ .

 $(ii) \Rightarrow (i)$ :

We assume now that  $B^TAB \succeq 0$  and  $B \in \mathbb{R}^{n \times n}$  is invertible. By the definition of semidefinite matrices,

$$x^T B^T A B x \ge 0.$$

Utilizing the properties of transpose matrices, we find

$$(Bx)^T A(Bx) \ge 0.$$

Let Bx = y, then this becomes

$$y^T A y \ge 0$$

and equivalent statement to  $A \succeq 0$ .

(i)⇒(iii):

The eigenvalues of A are defined as  $\lambda$  when  $Ax = \lambda x$ . Using this equality in the definition of semidefinite matrix  $A (A \succeq 0)$ , we find

$$x^T A x = x^T \lambda x > 0.$$

As a constant, we can reexpress this as

$$\lambda x^T x > 0$$

$$\lambda |x|^2 \ge 0$$

We note  $|x|^2 = \sum_i x_i^2$  and  $x_i^2 \ge 0$  .:  $|x|^2 \ge 0$ . With  $|x|^2 \ge 0$ ,  $\lambda \ge 0$  in order for  $\lambda |x|^2 \ge 0$  to be true.

 $(iii) \Rightarrow (iv)$ :

Since matrix A is symmetric (given), we can apply the spectral theorem for symmetric matrices (as stated by Mark Gockenbach on his website. This theorem states that for symmetric  $A \in \mathbb{R}^{n \times n}$  there exists a diagonal matrix  $D \in \mathbb{R}^{n \times n}$  and an orthogonal matrix  $P \in \mathbb{R}^{n \times n}$  such that  $A = PDP^T$ . Furthermore, the diagonal entries of D are the eigenvalues of A.

Any nonnegative diagonal matrix  $E \in \mathbb{R}^{n \times n}$  can be equivalently represented as  $F^2$  where  $F \in \mathbb{R}^{n \times n}$  is another diagonal matrix with elements  $F_i = \sqrt{E_i}$ . We showed in (iii) that all eigenvalues of A are nonnegative. Therefore, if all entries of D are the eigenvalues of A, then D is a nonnegative diagonal matrix and  $D = F^2$ . Since all diagonal matrices are symmetrical,  $D = F^2 = FF^T$ .

Together with the spectral theorem, we find

$$A = PDP^T = PFF^TP^T$$

$$A = PF(PF)^T$$

If we let  $U = PF (U \in \mathbb{R}^{n \times n})$ , then

$$A = UU^T$$
.

#### (iv)⇒(i):

Assume that  $\exists$  matrix  $U \in \mathbb{R}^{n \times n}$  such that  $A = UU^T$ . We want to show that  $A = UU^T \succeq 0$ .

$$A = UU^T$$

$$x^T A x = x^T U U^T x$$

$$x^T A x = (U^T x)^T U^T x$$

We can say  $U^T x = v$ , where  $v \in \mathbb{R}^n$ , so

$$x^T A x = v^T v$$

$$x^T A x = |v|^2$$

 $|v|^2 = \sum_i v_i^2$  and for all real numbers,  $v_i^2 \ge 0$   $\therefore$   $|v|^2 \ge 0$ .

$$x^T A x > 0$$

$$A \succ 0$$
.

### **b.**)

 $A \in \mathbb{R}^{n \times n}$  is a positive definite matrix, such that  $A \succ 0$ .

(i)

We are given that every  $\lambda > 0$  and want to prove that  $A + \lambda I > 0$ . By definition, this is true iff

$$x^T(A + \lambda I)x > 0$$

$$x^T A x + x^T \lambda I x > 0$$

By definition, A is positive definite, so  $x^TAx > 0$ . Furthermore,  $x^T\lambda Ix = x^T\lambda x = \lambda |x|^2$ .  $|x|^2 > 0$  and it is given that  $\lambda > 0$ , so the product  $\lambda |x|^2 > 0$ . Additionally, if  $\lambda |x|^2 > 0$  and  $x^TAx > 0$ , then their sum must also be greater than 0.

$$x^T A x + x^T \lambda I x > 0$$
 :  $A + \lambda I > 0$ .

(ii)

We are attempting to prove  $\exists \gamma > 0$  such that  $A - \gamma I \succ 0$ .

$$A - \gamma I \succ 0 \implies x^T (A - \gamma I)x > 0.$$

$$x^T A x - x^T \gamma I x > 0$$

Like in 2.b.iii, we note that the eigenvalues of A are given by  $\lambda$ , when  $Ax = \lambda x$ .

$$x^T \lambda x - x^T \gamma Ix > 0$$

$$\lambda x^T x - \gamma x^T x > 0$$

$$(\lambda - \gamma)x^Tx > 0$$

$$(\lambda - \gamma) \cdot |x|^2 > 0$$

Since  $|x|^2 > 0$ , dividing it out yields

$$\lambda - \gamma > 0$$
.

Now, using the definition of positive definite matrices,

$$x^T A x = x^T \lambda x > 0$$

$$\lambda x^T x > 0$$

$$\lambda |x|^2 > 0.$$

Since  $|x|^2 > 0$ ,  $\lambda > 0$  for the product  $\lambda |x|^2 > 0$ .

Now, since we have shown  $\lambda > 0$ , then when  $\lambda - \gamma > 0$ ,  $\{\gamma \in \mathbb{R} \mid 0 < \gamma < \lambda\}$ .

(iii)

First, we note that the procedure for calculating  $a = x^T A x$  is  $a = \sum_{j=1}^n \sum_{k=1}^n x_j x_k A_{jk}$ .

By definition,  $A \succ 0 \implies x^T A x > 0 \ \forall \ x \in \mathbb{R}^n - \{0\}$ . Since all  $x \in \mathbb{R}^n - \{0\}$  must satisfy this equation, showing that the diagonal entries of A must be greater than zero for any set of x vectors is sufficient to prove this case.

With this in mind, consider the basis of unit vectors  $e_1, \dots e_n$ , where  $e_1^T = (1 \ 0 \ \dots \ 0), \ e_2^T = (0 \ 1 \ \dots \ 0), \ e_n^T = (0 \ 0 \ \dots \ 1),$  etc.

Letting each of these unit vectors  $e_i = x$ ,

$$x^{T}Ax > 0$$

$$e_{i}^{T}Ae_{i} > 0$$

$$\sum_{j=1}^{n} \sum_{k=1}^{n} x_{j}x_{k}A_{jk} > 0$$

Since in  $e_i$ ,  $x_i = 1$  if i = k, otherwise  $x_i = 0$ , we can rewrite this summation as

$$x_i^2 A_{ii} > 0$$

$$x_i = 1 : x_i^2 = 1$$
, so

$$A_{ii} > 0$$
.

(iv)

By definition of a positive definite matrix, A > 0, we have  $x^T A x > 0$ . As noted in the proof of (iii), the procedure for calculating  $a = x^T A x$  is  $a = \sum_{i=1}^{n} \sum_{j=1}^{n} x_i x_j A_{ij}$ .

Again, since  $A \succ 0 \implies x^T A x > 0 \ \forall \ x \in \mathbb{R}^n - \{0\}$ , it is sufficient to show that  $\sum_{i=1}^n \sum_{j=1}^n A_{ij} > 0$  for any vector.

Consider the vector  $v^T = (1 \ 1 \ ... \ 1)$ , a vector in  $\mathbb{R}^n$  consisting of all ones.

$$v^T A v > 0$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} x_{i} x_{j} A_{ij} > 0, \text{ where } x_{i} = 1, x_{j} = 1 \ \forall i, j.$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} A_{ij} > 0. \quad \blacksquare$$

a.)

Let  $x, a \in \mathbb{R}^n$ . We can express the gradient of a function f(x) as the vector  $\nabla_x f(x)$  where the  $i^{\text{th}}$  element is given by  $\frac{df}{dx_i}$ . This gives

$$\nabla_x(a^T x) = \left(\frac{d}{dx_1}(a^T x) \quad \frac{d}{dx_2}(a^T x) \quad \dots \quad \frac{d}{dx_n}(a^T x)\right)^T$$

Noting that  $a^T x = \sum_{j=1}^n a_j x_j$ ,

$$\nabla_x(a^T x) = \left(\frac{d}{dx_1} \sum_{i=1}^n a_i x_i \quad \frac{d}{dx_2} \sum_{i=1}^n a_i x_i \quad \dots \quad \frac{d}{dx_n} \sum_{i=1}^n a_i x_i\right)^T$$

. For any summation,

$$\frac{d}{dx_i} \sum_{j=1}^{n} C_j x_j = \sum_{j=1}^{n} \frac{d}{dx_i} C_j x_j = \sum_{j=1}^{n} (C_j \delta_{ij}) = C_i$$

where  $C_j$  is a constant and  $\delta_{ij}$  is the Kronecker delta. Then,

$$\nabla_x(a^T x) = \begin{pmatrix} a_1 & a_2 & \dots & a_n \end{pmatrix}^T$$
$$\nabla_x(a^T x) = a$$

*b.*)

Let  $A \in \mathbb{R}^{n \times n}$ ,  $x \in \mathbb{R}^n$ . We note that  $x^T A x = \sum_{i=1}^n \sum_{j=1}^n x_i x_j A_{ij}$ . Then,

$$\nabla_x(x^T A x) = \left(\frac{d}{dx_1}(x^T A x) \quad \frac{d}{dx_2}(x^T A x) \quad \dots \quad \frac{d}{dx_n}(x^T A x)\right)^T$$

The  $i^{\text{th}}$  row of  $x^T A x$  is given by

$$\frac{d}{dx_i}(x^T A x) = \frac{d}{dx_i} \sum_{j=1}^n \sum_{k=1}^n A_{jk} x_j x_k$$

For any summation,

$$\frac{d}{dx_i} \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_j x_k = \sum_{j=1}^n \sum_{k=1}^n \frac{d}{dx_i} C_{jk} x_j x_k$$

$$\frac{d}{dx_i} \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_j x_k = \sum_{j=1}^n \sum_{k=1}^n C_{jk} (x_j \delta_{ij} + x_k \delta_{ik})$$

$$\frac{d}{dx_i} \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_j x_k = \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_j \delta_{ij} + \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_k \delta_{ik}$$

$$\frac{d}{dx_i} \sum_{j=1}^n \sum_{k=1}^n C_{jk} x_j x_k = \sum_{k=1}^n C_{ik} x_i + \sum_{j=1}^n C_{ji} x_i$$

Furthermore, for any matrix  $C \in \mathbb{R}^{n \times n}$ .

$$Cx = \left(\sum_{l=1}^{n} C_{l1}x_{1} \sum_{l=1}^{n} C_{l2}x_{2} \dots \sum_{l=1}^{n} C_{ln}x_{n}\right)^{T}$$
$$(Cx)_{i} = \sum_{l=1}^{n} C_{l1}x_{1}$$

Then,

$$\frac{d}{dx_i} \sum_{j=1}^{n} \sum_{k=1}^{n} C_{jk} x_j x_k = (Cx)_i + (C^T x)_i$$

We can use this conclusion in our above expression for  $\nabla_x(x^TAx)$ 

$$\nabla_{x}(x^{T}Ax) = \begin{bmatrix} (Ax)_{1} + (A^{T}x)_{1} & (Ax)_{2} + (A^{T}x)_{2} & \dots & (Ax)_{n} + (A^{T}x)_{n} \end{bmatrix}^{T}$$

$$\nabla_{x}(x^{T}Ax) = (Ax + A^{T}x)$$

$$\nabla_{x}(x^{T}Ax) = (A + A^{T})x$$

In the case that A is symmetric,  $A = A^T$  so

$$\nabla_x(x^T A x) = (A + A^T)x = (A + A)x$$
$$\nabla_x(x^T A x) = 2Ax$$

c.)

Let  $A, X \in \mathbb{R}^{n \times n}$ . Explicitly,  $(A^T X)_{ij} = \sum_{k=1}^n A_{ki} X_{kj}$ . The diagonals of  $(A^T X)$  are

$$(A^T X)_{ii} = \sum_{k=1}^n A_{ki} X_{ki}.$$

As the sum of the diagonals,

$$\operatorname{tr}(A^{T}X) = \sum_{i=1}^{n} (A^{T}X)_{ii}$$
$$\operatorname{tr}(A^{T}X) = \sum_{i=1}^{n} \sum_{i=1}^{n} A_{ki}X_{ki}$$

We can further show that

$$\frac{d}{dX_{lm}}\operatorname{tr}\left(A^{T}X\right) = \frac{d}{dX_{lm}}\sum_{i=1}^{n}\sum_{k=1}^{n}A_{ki}X_{ki}$$

$$\frac{d}{dX_{lm}}\operatorname{tr}\left(A^{T}X\right) = \sum_{i=1}^{n}\sum_{k=1}^{n}\frac{d}{dX_{lm}}A_{ki}X_{ki}$$

$$\frac{d}{dX_{lm}}\operatorname{tr}\left(A^{T}X\right) = \sum_{i=1}^{n}\sum_{k=1}^{n}A_{ki}\delta_{(ki)(lm)}$$

$$\frac{d}{dX_{lm}}\operatorname{tr}\left(A^{T}X\right) = A_{lm}$$

Then, since

$$\nabla_{X}(\operatorname{tr}\left(A^{T}X\right)) = \begin{pmatrix} \frac{d}{dX_{11}}(\operatorname{tr}\left(A^{T}X\right)) & \frac{d}{dX_{12}}(\operatorname{tr}\left(A^{T}X\right)) & \dots & \frac{d}{dX_{1n}}(\operatorname{tr}\left(A^{T}X\right)) \\ \frac{d}{dX_{21}}(\operatorname{tr}\left(A^{T}X\right)) & \frac{d}{dX_{22}}(\operatorname{tr}\left(A^{T}X\right)) & \dots & \frac{d}{dX_{2n}}(\operatorname{tr}\left(A^{T}X\right)) \\ \vdots & \vdots & \ddots & \vdots \\ \frac{d}{dX_{n1}}(\operatorname{tr}\left(A^{T}X\right)) & \frac{d}{dX_{12}}(\operatorname{tr}\left(A^{T}X\right)) & \dots & \frac{d}{dX_{nn}}(\operatorname{tr}\left(A^{T}X\right)) \end{pmatrix}$$

we can can say

$$\nabla_X(\operatorname{tr}(A^T X)) = \begin{pmatrix} A_{11} & A_{12} & \dots & A_{1n} \\ A_{21} & A_{22} & \dots & A_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ A_{n1} & A_{n2} & \dots & A_{nn} \end{pmatrix}$$

$$\nabla_X(\operatorname{tr}(A^T X)) = A$$

#### d.

To be a norm, the distance metric  $\delta(x,y) = f(x-y)$  must satisfy the triangle inequality,  $\delta(x,z) \leq \delta(x,y) + \delta(y,z)$ . In this case, it follows that

$$f(x-z) \le f(x-y) + f(y-z). \tag{4}$$

Here, x + y = z. For vectors  $x \in \mathbb{R}^2$ , we can test if  $f(x) = (\sqrt{|x_1|} + \sqrt{|x_2|})^2$  is a norm using equation 4.

$$(\sqrt{|x_1 - z_1|} + \sqrt{|x_2 - z_2|})^2 \le (\sqrt{|x_1 - y_1|} + \sqrt{|x_2 - y_2|})^2 + (\sqrt{|y_1 - z_1|} + \sqrt{|y_2 - z_2|})^2$$

$$|x_1-z_1|+|x_2-z_2|+2\sqrt{|x_1-z_1|\cdot|x_2-z_2|}\leq |x_1-y_1|+|x_2-y_2|+2\sqrt{|x_1-y_1|\cdot|x_2-y_2|}+|y_1-z_1|+|y_2-z_2|+2\sqrt{|y_1-z_1|\cdot|y_2-z_2|}+|y_1-z_1|+|y_2-z_2|+2\sqrt{|y_1-z_1|\cdot|y_2-z_2|}+|y_1-z_1|+|y_2-z_2|+2\sqrt{|y_1-z_1|\cdot|y_2-z_2|}+|y_1-z_1|+|y_2-z_2|+2\sqrt{|y_1-z_1|\cdot|y_2-z_2|}+|y_1-z_1|+|y_1-z_1|+|y_2-z_2|+2\sqrt{|y_1-z_1|\cdot|y_2-z_2|}+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-z_1|+|y_1-$$

Since x + y = z,  $x_i + y_i = z_i$ . From this,

$$x_i - z_i = -y_i$$
 and  $y_i - z_i = -x_i$ 

$$|x_i - z_i| = |y_i|$$
 and  $|y_i - z_i| = |x_i|$ 

Then, we have

$$|y_1| + |y_2| + 2\sqrt{|y_1| \cdot |y_2|} \le |x_1 - y_1| + |x_2 - y_2| + 2\sqrt{|x_1 - y_1| \cdot |x_2 - y_2|} + |x_1| + |x_2| + 2\sqrt{|x_1| \cdot |x_2|}$$

$$|y_1| - |x_1| + |y_2| - |x_2| + 2(\sqrt{|y_1| \cdot |y_2|} - \sqrt{|x_1| \cdot |x_2|}) \le |x_1 - y_1| + |x_2 - y_2| + 2\sqrt{|x_1 - y_1| \cdot |x_2 - y_2|}$$

Since it is always true that for any  $a, b \in \mathbb{R}$ ,  $|a-b| = |b-a| \ge |a| - |b|$ , we know it is true that both  $|x_1 - y_1| \ge |y_1| - |x_1|$  and  $|x_2 - y_2| \ge |y_2| - |x_2|$ . This leaves it only necessary to further prove

$$\begin{split} 2(\sqrt{|y_1|\cdot|y_2|} - \sqrt{|x_1|\cdot|x_2|}) &\leq 2\sqrt{|x_1 - y_1|\cdot|x_2 - y_2|} \\ \sqrt{|y_1 y_2|} - \sqrt{|x_1 x_2|} &\leq \sqrt{|x_1 - y_1|\cdot|x_2 - y_2|} \\ \sqrt{|y_1 y_2|} - \sqrt{|x_1 x_2|} &\leq \sqrt{|x_1 x_2 + y_1 y_2 - x_1 y_2 - x_2 y_1|} \\ \sqrt{|y_1 y_2|} &\leq \sqrt{|x_1 x_2 + y_1 y_2 - x_1 y_2 - x_2 y_1|} + \sqrt{|x_1 x_2|} \end{split}$$

If we wish to provide a counter example, it must satisfy the condition that

$$\sqrt{|y_1y_2|} > \sqrt{|x_1x_2 + y_1y_2 - x_1y_2 - x_2y_1|} + \sqrt{|x_1x_2|}$$

To satisfy this condition, we could try x, y such that  $|y_1y_2|$  is large,  $|x_1x_2|$  is small, and  $(-x_1y_2 - x_2y_1)$  is large, but smaller than  $y_1y_2$ . Specifically, we try

$$x = \begin{pmatrix} 1 \\ 64 \end{pmatrix} \qquad \qquad y = \begin{pmatrix} 10 \\ 1000 \end{pmatrix}$$

Then,

$$\sqrt{|y_1y_2|} > \sqrt{|x_1x_2 + y_1y_2 - x_1y_2 - x_2y_1|} + \sqrt{|x_1x_2|}$$

$$\sqrt{10^4} > \sqrt{|64 + 10^4 - 1000 - 640|} + \sqrt{|64|}$$

$$10^2 > \sqrt{10064 - 1640} + 8$$

$$92 > 91.782...$$

Since this satisfies this condition, we can try it in our original equation,

$$(\sqrt{|x_1 - z_1|} + \sqrt{|x_2 - z_2|})^2 \le (\sqrt{|x_1 - y_1|} + \sqrt{|x_2 - y_2|})^2 + (\sqrt{|y_1 - z_1|} + \sqrt{|y_2 - z_2|})^2$$

Now 
$$z = \begin{pmatrix} 11\\1064 \end{pmatrix}$$
 and so

$$(\sqrt{|1-11|} + \sqrt{|64-1064|})^2 \le (\sqrt{|1-10|} + \sqrt{|64-1000|})^2 + (\sqrt{|10-11|} + \sqrt{|1000-1064|})^2 + (\sqrt{10} + \sqrt{1000})^2 \le (\sqrt{9} + \sqrt{936})^2 + (\sqrt{1} + \sqrt{64})^2$$

But, instead we find 1210 ≰ 1209.565, and so we have found a counterexample. The given function is **not** a norm.

e.)

Let  $x \in \mathbb{R}^n$ . We know that  $||x||_{\infty} = \max_i |x_i|$ , and  $||x||_2 = \sum_i^n |x_i|^2$ .

### Minimum $||x||_2$ :

The minimum (nontrivial) value of  $||x||_2$  would be given for any vector with only one single nonzero element (say this element is the  $j^{\text{th}}$  element). This can be shown as:

$$\sqrt{x_1^2} \le \sqrt{x_1^2 + \dots + x_n^2}.$$

In this case,

$$||x||_2 = \sqrt{\sum_{i=1}^{n} |x_i|^2} = \sqrt{\sum_{i=1}^{n} |x_i \delta_{ij}|^2}$$
  
 $||x||_2 = \sqrt{|x_j|^2} = |x_j|$ 

With all other elements of x being zero,  $|x_i| = \max_i |x_i|$ .

#### Maximum $||x||_2$

Similarly, the maximum value of  $||x||_2$  would be given when every  $|x_i| = \max_i |x_i| = |x_j|$ . This can be shown as:

$$||x||_2 = \sqrt{\sum_{i=1}^{n} |x_i|^2}$$

$$||x||_2=\sqrt{n|x_j|^2}=\sqrt{n}|x_j|$$

With all elements of x being nonzero,  $|x_j| = \max_i |x_i|$ .

In both cases, we have defined  $|x_i| = \max_i |x_i|$ . Using proper substitutions, we find that

$$\begin{split} |x_j| & \leq \sqrt{x_1^2 + \ldots + x_n^2} \leq \sqrt{n} |x_j| \\ \max_i |x_i| & \leq \sqrt{x_1^2 + \ldots + x_n^2} \leq \sqrt{n} \max_i |x_i| \\ ||x||_{\infty} & \leq ||x||_2 \leq \sqrt{n} ||x||_{\infty} \quad \blacksquare \end{split}$$

*f*.)

Let  $x \in \mathbb{R}^n$ . We know that  $||x||_1 = \sum_i^n |x_i|$ , and  $||x||_2 = \sum_i^n |x_i|^2$ .

## Minimum $||x||_1$ :

The minimum (nontrivial) value of  $||x||_1$  would be given for any vector with only one single nonzero element (say this element is the  $j^{\text{th}}$  element). This can be shown as:

$$|x_1| \le |x_1| + \dots + |x_n|.$$

In this case,

$$||x||_1 = \sum_{i=1}^{n} |x_i \delta_{ij}| = |x_j|$$

$$||x||_2 = \sqrt{\sum_{i=1}^{n} |x_i \delta_{ij}|^2} = |x_j|$$

## Maximum $||x||_1$

Similarly, the maximum value of  $||x||_1$  would be given when every  $|x_i| = \max_i |x_i| = |x_j|$ . This can be shown as:

$$||x||_1 = \sum_{i=1}^{n} |x_i|$$

$$||x||_1 = n|x_j|.$$

From the previous problem, we also showed that in this case

$$||x||_2 = \sqrt{n}|x_j|$$

$$||x||_1 = \sqrt{n}||x||_2$$

Combining these minimum an maximum values:

$$||x||_2 \le ||x||_1 \le \sqrt{n}||x||_2$$

Let  $A \in \mathbb{R}^{n \times n}$  be a symmetric matrix with  $A \succeq 0$ .

#### a.)

As discussed in problem 2.a, the spectral theorem for symmetric matrices states that for symmetric  $A \in \mathbb{R}^{n \times n}$  there exists a diagonal matrix  $D \in \mathbb{R}^{n \times n}$  and an orthogonal matrix  $P \in \mathbb{R}^{n \times n}$  such that  $A = PDP^T$ . Furthermore, the diagonal entries of D are the eigenvalues of A.

We can use this to show

$$\lambda_{\max}(A) = \max_{\|x\|_2 = 1} x^T A x.$$

or equivalently (using the spectral theorem)

$$\lambda_{\max}(A) = \max_{||x||_2 = 1} x^T P D P^T x.$$

Since  $x \in \mathbb{R}^n$  and  $P \in \mathbb{R}^{n \times n}$ ,

$$P^T x = y$$
, and  $(x^T P)^T = y^T$ 

where  $y \in \mathbb{R}^n$ . We then find

$$\lambda_{\max}(A) = \max_{||x||_2 = 1} y^T D y.$$

We note that since P is orthogonal, the columns of P are orthonormal, and it follows that

$$||(P^T x)||_2 = 1$$
 and  $||(x^T P)||_2 = 1$   
 $||y||_2 = 1$  and  $||y^T||_2 = 1$   
 $\lambda_{\max}(A) = \max_{||y||_2 = 1} y^T Dy$ .

Now,

$$\max_{||y||_2=1} y^T D y = \max_{||y||_2=1} \sum_{i=1}^n y_i^2 D_{ii}$$

Since  $\sum_i y_i^2 = 1$  (it is constant), we will maximize  $\max_{||y||_2=1} \sum_i^n y_i^2 D_{ii}$  by choosing the configuration of y that favors  $\max_i D_{ii}$ . If we let  $\max_i D_{ii} = D_{jj}$ , then this would be the vector where  $y_{i=j} = 1$  and  $y_{i\neq j} = 0$ . Since  $y_j = 1$ ,

$$\max_{||y||_2=1} y^T D y = \max_{||y||_2=1} \sum_{i}^{n} D_{ii} \delta_{ij}$$

$$\max_{\|y\|_2=1} y^T D y = \max_{\|y\|_2=1} D_{ii}$$

Then.

$$\lambda_{\max}(A) = \max_{||y||_2 = 1} D_{ii}.$$

Returning to the assertion of the spectral theorem that the elements of D are the eigenvalues of A, this statement is true.  $\blacksquare$ 

#### **b.**)

Using the procedure established in part (a) we can repeat to show

$$\lambda_{\min}(A) = \min_{\|x\|_2 = 1} x^T A x.$$

or equivalently (using the spectral theorem)

$$\lambda_{\min}(A) = \min_{||x||_2 = 1} x^T P D P^T x.$$

Since  $x \in \mathbb{R}^n$  and  $P \in \mathbb{R}^{n \times n}$ ,

$$P^T x = y$$
, and  $(x^T P)^T = y^T$ 

where  $y \in \mathbb{R}^n$ . We then find

$$\lambda_{\min}(A) = \min_{\|x\|_2 = 1} y^T D y.$$

We note that since P is orthogonal, the columns of P are orthonormal, and it follows that

$$||(P^T x)||_2 = 1$$
 and  $||(x^T P)||_2 = 1$   
 $||y||_2 = 1$  and  $||y^T||_2 = 1$   
 $\lambda_{\min}(A) = \min_{\|y\|_2 = 1} y^T Dy.$ 

Now,

$$\min_{||y||_2=1} y^T D y = \min_{||y||_2=1} \sum_{i=1}^n y_i^2 D_{ii}$$

Since  $\sum_i y_i^2 = 1$  (it is constant), we will minimize  $\min_{||y||_2=1} \sum_i^n y_i^2 D_{ii}$  by choosing the configuration of y that favors  $\min_i D_{ii}$ . If we let  $\min_i D_{ii} = D_{jj}$ , then this would be the vector where  $y_{i=j} = 1$  and  $y_{i\neq j} = 0$ . Since  $y_j = 1$ ,

$$\min_{||y||_2=1} y^T D y = \min_{||y||_2=1} \sum_{i}^n D_{ii} \delta_{ij}$$

$$\min_{||y||_2=1} y^T D y = \min_{||y||_2=1} D_{ii}$$

Then,

$$\lambda_{\min}(A) = \min_{||y||_2 = 1} D_{ii}.$$

Returning to the assertion of the spectral theorem that the elements of D are the eigenvalues of A, this statement is true.

#### c.)

The conditions which must be satisfied for a minimization (maximization) problem to be satisfied are:

- (1) the objective function to be a convex (concave) function
- (2) the feasible region to be a convex set

By definition, a set  $\mathcal{C} \subseteq \mathbb{R}^n$  is convex iff  $\forall x, y \in \mathcal{C}, \ \forall t \in [0, 1], \ tx + (1 - t)y \in \mathcal{C}$ .

Say we choose 2 vectors  $x = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$  and  $y = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ . Furthermore, choose t = 0.75. Note that  $||x||_2 = 1$  and  $||y||_2 = 1$ , where  $||\cdot||_2 = 1$  is the condition defining our set.

$$0.75 \begin{pmatrix} 1 \\ 0 \end{pmatrix} + (1 - 0.75) \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0.75 \\ 0.25 \end{pmatrix} \notin \mathcal{C}.$$

The set is not convex, and so neither program is convex.

#### d.

Again using the spectral theorem for symmetric matrices, we know  $A = PDP^T$  where A is symmetric, P is orthogonal, and D is diagonal and its diagonals are the eigenvalues of A. Then, we can see

$$AA = (PDP^T)(PDP^T)$$
 
$$A^2 = PD\mathbb{1}DP^T, \text{ since } P^T = P^{-1} \text{ for orthogonal matrices}$$
 
$$A^2 = PD^2P^T$$

And so, since  $A^2$  is also symmetric,  $D^2$  has diagonals  $\lambda^2$  which are the eigenvalues of  $A^2$ .

Using this fact, we know from (a) and (b) that

$$\lambda_{\max}(A^2) = \max(D_{ii}^2) \text{ and } \lambda_{\min}(A^2) = \min(D_{ii}^2)$$
$$\lambda_{\max}(A^2) = (\max D_{ii})^2 \text{ and } \lambda_{\min}(A^2) = (\min D_{ii})^2$$
$$\lambda_{\max}(A^2) = \lambda_{\max}(A)^2 \text{ and } \lambda_{\min}(A^2) = \lambda_{\min}(A)^2$$

e.)

Noting that  $||Ax||_2 = \sqrt{(Ax)^T(Ax)}$ , we find

$$||Ax||_2 = \sqrt{x^T A^T A x}$$

Since A is symmetric,  $AA^T = A^2$ , and

$$||Ax||_2 = \sqrt{x^T A^2 x}$$

It is by definition that

$$\min_{||x||_2=1} ||Ax||_2 \le ||Ax||_2 \le \max_{||x||_2=1} ||Ax||_2.$$
(5)

From this, we can deduce

$$\begin{split} \min_{||x||_2=1} ||Ax||_2 &= \min_{||x||_2=1} \sqrt{x^T A^2 x} \text{ and } \max_{||x||_2=1} ||Ax||_2 = \max_{||x||_2=1} \sqrt{x^T A^2 x} \\ \min_{||x||_2=1} ||Ax||_2 &= \sqrt{\min_{||x||_2=1} x^T A^2 x} \text{ and } \max_{||x||_2=1} ||Ax||_2 = \sqrt{\max_{||x||_2=1} x^T A^2 x}. \end{split}$$

By substituting our answer from parts (a) and (b) we have

$$\min_{||x||_2=1} ||Ax||_2 = \sqrt{\lambda_{\min}(A^2)} \text{ and } \max_{||x||_2=1} ||Ax||_2 = \sqrt{\lambda_{\max}(A^2)}$$

and by then substituting our answer from part (d) we have

$$\min_{||x||_2=1} ||Ax||_2 = \sqrt{\lambda_{\min}(A)^2} \text{ and } \max_{||x||_2=1} ||Ax||_2 = \sqrt{\lambda_{\max}(A)^2}$$

$$\min_{||x||_2=1} ||Ax||_2 = \lambda_{\min}(A) \text{ and } \max_{||x||_2=1} ||Ax||_2 = \lambda_{\max}(A).$$

When consiered with equation 5, we have

$$\lambda_{\min}(A) \le ||Ax||_2 \le \lambda_{\max}(A)$$

*f.*)

As in part (e) we know  $||Ax||_2 = \sqrt{x^T A x}$ ,  $\forall x \in \mathbb{R}^n$ . The result found in (d) is a special case for unit vectors, where

$$\lambda_{\min}(A) \le ||Ax||_2 \le \lambda_{\max}(A).$$

Instead, in the general case, we can force any vector x to be a unit vector by dividing it by its magnitude:  $\frac{x}{||x||_2}$ . Then for non-unit vectors, we have

$$||Ax||_2 = \sqrt{y^T Ay ||x||_2^2} = ||x||_2 \sqrt{y^T Ay}$$

where  $y = \frac{x}{||x||_2}$ . Using this in conjunction with the part (d), we find that in general,  $\forall x \in \mathbb{R}^n$ ,

$$\lambda_{\min}(A)||x||_2 \le ||Ax||_2 \le \lambda_{\max}(A)||x||_2.$$

a.)

First order optimality conditions require  $\nabla_x f(x) = 0$  (where f(x) is the objective function). When computing this gradient, the  $i^{\text{th}}$  element is given by  $\frac{df}{dx_i}$ .

The objective function in this case is

$$\frac{1}{2}x^T A x - b^T x,$$

so

$$\nabla_x \left( \frac{1}{2} x^T A x - b^T x \right) = 0.$$

From problem 3.a we have  $\nabla_x(a^Tx) = a$ , and from problem 3.b we have  $\nabla_x(x^TAx) = 2Ax$ . We then find for our current objective function:

$$\nabla_x \left( \frac{1}{2} x^T A x - b^T x \right) = \frac{1}{2} (2Ax^*) - b = 0$$

$$Ax^* - b = 0$$

$$Ax^* = b$$

$$x^* = A^{-1}b$$

*b.*)

The update rule for gradient descent is given by

$$w \leftarrow w - \epsilon \nabla_w R(w)$$

where R(w) is the risk function and  $\epsilon$  is the step size.

For this problem, the update rule becomes

$$x \leftarrow x - (1)\nabla_x \left(\frac{1}{2}x^T A x - b^T x\right)$$

where  $R(w) = \frac{1}{2}x^T A x - b^T x$  and  $\epsilon = 1$ . Then,

$$x^{(k)} = x^{(k-1)} - (Ax^{(k-1)} - b).$$

$$x^{(k)} = x^{(k-1)} - Ax^{(k-1)} + b.$$

c.)

$$x^{(k)} - x^* = (x^{(k-1)} - Ax^{(k-1)} + b) - x^*$$

$$x^{(k)} - x^* = x^{(k-1)} - Ax^{(k-1)} + Ax^* - x^*$$

$$x^{(k)} - x^* = x^{(k-1)} - Ax^{(k-1)} - x^* + Ax^*$$

$$x^{(k)} - x^* = (I - A)(x^{(k-1)} - x^*)$$

$$x^{(k)} - x^* = (I - A)(x^{(k-1)} - x^*)$$

#### d.

We know from (c) that  $x^{(k)} - x^* = (I - A)(x^{(k-1)} - x^*)$ . It follows that  $||x^{(k)} - x^*||_2 = ||(I - A)(x^{(k-1)} - x^*)||_2$ . If we let  $y = x^{(k-1)} - x^*$  then this equation becomes

$$||x^{(k)} - x^*||_2 = ||(I - A)y||_2$$

If we can show  $(I - A) \succeq 0$  then we can use the inequality found in problem 4.f. For  $(I - A) \succeq 0$ ,  $x^T(I - A)x \ge 0$ . By the definition of an eigenvalue, we know that

$$Ax = \lambda x$$
$$Ax - \lambda x = 0$$
$$(A - \lambda I)x = 0$$
$$(\lambda I - A)x = 0$$

We are told that all eigenvalues of A are on the interval (0,1), so we know that  $(\lambda I)_{ii} < I_{ii} \quad \forall i \in 1,...,n$ . Then  $(I-A) > (\lambda I - A)$  and if  $x^T(\lambda I - A)x = 0$  then  $x^T(I-A)x > 0$  (assuming  $x \neq 0$ ). It becomes evident that indeed  $(I-A) \succeq 0$ .

Now that we have shown that (I - A) is semi-positive definite, we can use the result of problem 4.f to show

$$||(I-A)y||_2 \le \lambda_{\max}(I-A)||y||_2.$$

Since we know that  $(I-A)x = \lambda x$  where  $\lambda$  represents the eigenvalues of (I-A), we can show

$$(I - A)x = \lambda x$$
$$Ax - x = -\lambda x$$
$$Ax = (1 - \lambda)x$$

(i.e.  $(1 - \lambda)$  are the eigenvalues of A). Since we are told  $0 \le \lambda_{\min}(A)$  and  $\lambda_{\max}(A) \le 1$  we can deduce that  $0 \le \lambda_{\min}(I - A)$  and  $\lambda_{\max}(I - A)$  as well. Let  $\rho$  represent the maximum eigenvalue (still,  $0 < \rho < 1$ ).

Now,

$$||(I - A)y||_2 \le \rho ||y||_2.$$

Substituting  $||(I - A)y||_2 = ||x^{(k)} - x^*||_2$  and  $y = x^{(k-1)} - x^*$ ,

$$||x^{(k)} - x^*||_2 \le \rho ||x^{(k-1)} - x^*||_2|.$$

#### e.)

Assuming the "worst" case scenario, where

$$||x^{(k')} - x^*||_2 = \rho ||x^{(k'-1)} - x^*||_2,$$

we find that

$$||x^{(1)} - x^*||_2 = \rho ||x^{(0)} - x^*||_2$$
$$||x^{(2)} - x^*||_2 = \rho ||x^{(1)} - x^*||_2 = \rho^2 ||x^{(0)} - x^*||_2$$

$$||x^{(k)} - x^*||_2 = \rho ||x^{(k-1)} - x^*||_2 = \rho^k ||x^{(0)} - x^*||_2$$

We want our solution  $||x^{(k')} - x^*||_2 \le \epsilon$ , so

$$||x^{(n)} - x^*||_2 = \rho^k ||x^{(0)} - x^*||_2 \le \epsilon$$

Solving for k-iterations,

$$\rho^k \leq \frac{\epsilon}{||x^{(0)} - x^*||_2}$$

$$k\log\rho \le \frac{\epsilon}{||x^{(0)} - x^*||_2}.$$

Since  $\log \rho < 0$ , convergence to tolerance  $\epsilon$  will occur for

$$k \ge \frac{\epsilon}{\log \rho ||x^{(0)} - x^*||_2}$$

Since we are dealing with the worst case scenario, we can be sure that our algorithm will converge in this many iterations.

# *f*.)

The iteration of gradient descent is dominated by a matrix-vector product, where a  $n \times n$  matrix is multiplied by a vector of length n. Computing this product requires n multiplications and n-1 additions per each of the n-rows in the matrix. This is a total of  $(n+n-1)n=2n^2-n$  operations per iteration. For  $k=\frac{\epsilon}{\log \rho||x^{(0)}-x^*||_2}$  iterations, the overall running time is

$$t \propto \frac{\epsilon(2n^2 - 2)}{\log \rho ||x^{(0)} - x^*||_2}$$

The risk function is defined as:

$$R(f(x) = i|x) = \sum_{j=1}^{c} L(f(x) = i, y = j)P(Y = j|x).$$

We can try to minimize R by selecting a policy that chooses class i if  $P(Y = i|x) \ge P(Y = j|x) \forall j$ . Then

$$R(f(x) = i|x) = L(f(x) = i, y = i)P(Y = i|x) + \sum_{j=1, j \neq i}^{c} L(f(x) = i, y = j \neq i)P(Y = j \neq i|x)$$

Without doubt, this becomes

$$R(f(x) = i|x) = 0 + \lambda_s(1 - P(Y = i|x))$$

$$R(f(x) = i|x) = \lambda_s(1 - P(Y = i|x))$$

Since 
$$P(Y = i|x) \ge P(Y = j|x) \ \forall j$$
, we can state  $(1 - P(Y = i|x)) \le (1 - P(Y = j|x)) \ \forall j$ .

Introducing doubt may allow us to minimize this risk function further.

Imposing the condition that we only choose doubt when  $P(Y=i|x) \leq 1 - \lambda_r/\lambda_s$ , we find that

$$\lambda_s(1 - P(Y = i|x)) \ge \lambda_r$$

$$R(f(x) = i|x) \ge \lambda_r$$
.

Our old risk function is no longer a minimum, and so choosing doubt will minimize the risk function. Otherwise, our old risk function is, in fact, minimized.

## b.)

If  $\lambda_r = 0$  then the second condition of part (1) of the policy is never satisfied (except when P(Y = i|x) = 1) since a probability can not be greater than 1. In this case, doubt will always be chosen This makes sense intuitively because there is no longer any penalty to choosing doubt.

If  $\lambda_r > \lambda_s$  then the second condition of part (1) of the policy is always satisfied, since a probability cannot be less than 0. This also makes sense intuitively, as it makes no sense to choose doubt if you will be more harshly penalized for it than a misclassification.

a.)

Using Gaussian Discriminant Analysis, we aim to maximize  $P(X = x | Y = i)\pi_i$  using Bayes decision rule. Given our Gaussiant probability distributions,  $P(x|\omega_i) \sim \mathcal{N}(\mu_i, \sigma^2)$ , it is equivalent to maximize  $Q_i(x) = \ln\left(\left(\sqrt{2\pi}\right)^d P(x)\pi_i\right)$  instead.

We can also express  $Q_i(x)$  as

$$Q_i(x) = \ln\left(\left(\sqrt{2\pi}\right)^d P(x)\pi_i\right) = -\frac{|x - \mu_i|^2}{2\sigma_i^2} - d\ln\sigma_i + \ln\pi_i.$$

The Bayes optimal decision boundary is given by

$$Q_1(x) - Q_2(x) = -\frac{|x - \mu_1|^2}{2\sigma_1^2} - d\ln\sigma_1 + \ln\pi_1 + \frac{|x - \mu_2|^2}{2\sigma_2^2} + d\ln\sigma_2 - \ln\pi_2 = 0.$$

Since the problem is one-dimensional, d = 1. Also,  $\sigma_1 = \sigma_2 = \sigma$ . The decision boundary equation simplifies to:

$$Q_1(x) - Q_2(x) = \frac{|x - \mu_2|^2 - |x - \mu_1|^2}{2\sigma^2} + \ln \pi_1 - \ln \pi_2 = 0.$$

We are given that  $\pi_1 = P(x|\omega_1) = \pi_2 = P(x|\omega_2) = \frac{1}{2}$ , so we find

$$Q_1(x) - Q_2(x) = \frac{|x - \mu_2|^2 - |x - \mu_1|^2}{2\sigma^2} = 0$$
$$|x - \mu_2|^2 - |x - \mu_1|^2 = 0$$
$$|x - \mu_2|^2 = |x - \mu_1|^2$$
$$|x - \mu_2| = |x - \mu_1|$$

For this to be true, either  $\mu_1 = \mu_2$ , or

$$x - \mu_2 = -x + \mu_1$$
$$2x = \mu_1 + \mu_2$$
$$x = \frac{\mu_1 + \mu_2}{2}$$

The Bayes decison rule which corresponds to this boundary is

$$r^*(x) = \begin{cases} 1 & Q_1(x) - Q_2(x) > 0 \\ 2 & \text{otherwise} \end{cases}$$

**b.**)

Using the given definition

$$P_{\ell} = P((\text{misclassified as }\omega_1)|\omega_2)P(\omega_2) + P((\text{misclassified as }\omega_2)|\omega_1)P(\omega_1)$$

We can say

$$P((\text{misclassified as }\omega_1)|\omega_2) = \int_{-\infty}^{\frac{\mu_1 + \mu_2}{2}} \frac{1}{(\sqrt{2\pi}\sigma)^d} e^{\frac{-|x-\mu_2|^2}{2\sigma^2}} dx$$

and

$$P((\text{misclassified as }\omega_2)|\omega_1) = \int_{\frac{\mu_1 + \mu_2}{2}}^{\infty} \frac{1}{(\sqrt{2\pi}\sigma)^d} e^{\frac{-|x - \mu_1|^2}{2\sigma^2}} dx$$

We can then use the fact that d = 1 to remove it from the equation.

If we use a change of variables such that  $z(x) = \frac{-x+\mu_2}{\sigma}$  on  $P((\text{misclassified as }\omega_1)|\omega_2)$ , we find

$$P((\text{misclassified as }\omega_1)|\omega_2) = \int_{-\infty}^{z\left(\frac{\mu_1+\mu_2}{2}\right)} \frac{1}{\sqrt{2\pi}\sigma} e^{\frac{-(-z)^2}{2}} (-\sigma) dz$$

$$P((\text{misclassified as }\omega_1)|\omega_2) = -\int_{\infty}^{\frac{\mu_2-\mu_1}{2\sigma}} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz$$

$$P((\text{misclassified as }\omega_1)|\omega_2) = \int_{\frac{\mu_2-\mu_1}{2\sigma}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz$$

We can use a similar change of variables on  $P((\text{misclassified as }\omega_2)|\omega_1)$ , where  $z(x)=\frac{x-\mu_1}{\sigma}$ . We find

$$P((\text{misclassified as }\omega_2)|\omega_1) = \int_{z\left(\frac{\mu_1 + \mu_2}{2}\right)}^{\infty} \frac{1}{\sqrt{2\pi}\sigma} e^{\frac{-z^2}{2}} \sigma dz$$

$$P((\text{misclassified as }\omega_2)|\omega_1) = \int_{\frac{\mu_2 - \mu_1}{2\sigma}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz$$

Then,

$$P_{\ell} = \frac{1}{2} \int_{\frac{\mu_2 - \mu_1}{2\sigma}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz + \frac{1}{2} \int_{\frac{\mu_2 - \mu_1}{2\sigma}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz$$

$$P_{\ell} = \int_{\frac{\mu_2 - \mu_1}{2\sigma}}^{\infty} \frac{1}{\sqrt{2\pi}} e^{\frac{-z^2}{2}} dz$$

$$P_{\ell} = \frac{1}{\sqrt{2\pi}} \int_{a}^{\infty} e^{\frac{-z^2}{2}} dz, \text{ where } a = \frac{\mu_2 - \mu_1}{2\sigma}$$

We can analyze each probability (P(X=1), P(X=2), P(X=3)), by treating each likelihood function as if it were only of two outcomes, P(X=i) and  $P(X \neq i)$ .

Then, we can use the binomial distribution to say

$$\mathcal{L}(p_i) = \binom{n}{k_i} p_i^{k_i} (1 - p_i)^{n - k_i}$$

$$\frac{d\mathcal{L}(p_i)}{dp_i} = \frac{n!}{k!(n - k)!} \left( k_i p_i^{k_i - 1} (1 - p_i)^{n - k_i} - p_i^{k_i} (n - k_i) (1 - p_i)^{n - k_i - 1} \right)$$

To maximize this likelihood function, we find  $\frac{d\mathscr{L}}{dp_i} = 0$ .

$$0 = \frac{n!}{k!(n-k)!} \left( k_i p_i^{k_i-1} (1-p_i)^{n-k_i} - p_i^{k_i} (n-k_i) (1-p_i)^{n-k_i-1} \right)$$

$$p_i^{k_i} (n-k_i) (1-p_i)^{n-k_i-1} = k_i p_i^{k_i-1} (1-p_i)^{n-k_i}$$

$$p_i (n-k_i) = k_i (1-p_i)$$

$$np_i - k_i p_i = k_i - k_i p_i$$

$$np_i = k_i$$

$$p_i = \frac{k_i}{n}$$

We find then, that  $p_1 = k_1/n$ ,  $p_2 = k_2/n$ , and  $p_3 = k_3/n$ .