

29.01.23

Definition: A block sequence $B = (b_1, \dots, b_\ell)$ is an ordered partition of the set of blocks.

Notation: $B_i = \bigcup_{j=1}^i b_j$
 $\star = (b_1, \dots, b_\ell)$

Definition: Let B be a block sequence. The update sequence $R_B = (r_1, \dots, r_{\ell+2})$ induced by B is defined

as follows: For every $i \in [1, \ell+2]$, $r_i = \{(v, P) : b(v, P) \in b_{i-1}, v \in V(P^u - P^o)\} \cup$

$\{(v, P) : b(v, P) \in b_{i-1}, v \in V(P^o \cap P^u)\} \cup \{(v, P) : b(v, P) \in b_{i-2}, v \in V(P^o - P^u)\}$, where

$b_{\ell+1} := \{\}$ and $b_{\ell+2} := \{\}$, $b_{\ell+1} := \{\}$, and $b_{\ell+2} := \{\}$.

* We now define, given a block sequence $B = (b_1, \dots, b_\ell)$, the update sequence $R_B = (r_1, \dots, r_{\ell+2})$

induced by B . The idea is update every block in b_i in rounds r_i, r_{i+1} and r_{i+2} according to

Lemma 1.

Lemma: Let $B = (b_1, \dots, b_\ell)$ be a block sequence and $R_B = (r_1, \dots, r_{\ell+2})$ be the induced update sequence.

Then, for every $P \in \mathcal{P}$, $i \in [1, \ell+2]$, $S \subseteq r_i$, P is transient for U_i^S .

Proof: Definition: Let $(v, P) \in V \times \mathcal{P}$. Then we call (v, P) ^{be an update}

- a type-1 update if $v \in V(P^u - P^o)$,

- 2 $V(P^o \cap P^u)$, and

- 3 $V(P^o - P^u)$.

By definition, r_i contains the type-1 updates for induced by blocks in b_i , the type-2...

Proof:

Let $P \in \mathcal{P}$.

Claim: Let $i \in [1, \ell+2]$ and $S \subseteq r_i$. Then:

(i) For every type-1 update $(v, P) \in S$, P is transient for U_i^S iff P is transient for $U_i^{S \setminus \{(v, P)\}}$.

(ii) \downarrow

Proof:

(i).

We show the following, which implies the claim:

(1) For every $S \subseteq r_1$, P is transient for $U_1^S = S$.

(2) For every $i \in [1]$ and $S \subseteq r_{i+1}$, P is transient for U_{i+1}^S .

(3) For every $S \subseteq r_{\ell+2}$, P is transient for $U_{\ell+2}^S$ iff P is transient for $U_{\ell+2}^{S \setminus \{(v, P)\}} = U_{\ell+1}^{S \setminus \{(v, P)\}}$.

(1) By definition, r_1 contains only type-1 updates. Hence, by claim (i), P is transient for U_1^S iff

P is transient for $U_1^{S \setminus \{(v, P)\}} = \{\}$, which in the case holds.

We might need a different definition, as we're ignoring capacity constraints here.

TODO I think this proof depends significantly on my definition of blocks.

(3). Let $S \subseteq r_{c+2}$. By definition, r_{c+2} contains only type-3 updates. Hence, by Claim ~ (iii), P is transient

for U_{c+2}^S iff P is transient for $U_{c+2}^{S^*} = U_{c+1}^{r_{c+1}}$.

(2). Let $i \in [1]$ and $S \subseteq r_{i+1}$. Suppose not. Then let $i \in [1]$ be minimum and $S \subseteq r_{i+1}$ be minimal such that

P is not transient for U_{i+1}^S . Since i is minimum, $S \neq \emptyset$. Moreover $U_{i+1}^S = U_{i+1}^{S^*} = U_i^{r_i}$. Then let (v, P') be any

element update in S . Since S is minimal, $P' = P$ and P is transient for $U_{i+1}^{S - \{(v, P)\}}$. Hence, by Claim ~, (v, P) is This requires that $v \in V(P)$.

a type-2 update.

By definition, all type-1 updates for $b(v, P)$ are contained in r_i, \dots

Note that Lemma ~ implies that, for every $P \in \mathcal{P}$, $i \in [1+2]$, $S \subseteq r_i$, \bar{I}_{P, U_i^S} exists. (s.t.) - Now

Definition: $\alpha_P((u, v), B) = \begin{cases} \text{active if } b(u, P) \notin B \text{ and } (u, v) \in E(P^0) \\ \text{active if } b(u, P) \in B \text{ and } (u, v) \in E(P^0) \\ \text{inactive else} \end{cases}$

Definition: A York sequence $B = (b_1, \dots, b_c)$ is feasible if for every $e \in E$ and every $i \in [1]$,

$c(e) \geq \sum_{P \in \mathcal{P}: \alpha_P(e, B_{i-1}) = \text{active or } \alpha_P(e, B_i) = \text{active}} d_P$.

Lemma: Let $e \in E$. Let $B \subseteq B$, $e \in E$, $P \in \mathcal{P}$. Then $e \in E(T_{P, U(B)})$ iff $\alpha_P(e, B) = \text{active}$.

Lemma: If a York sequence $B = (b_1, \dots, b_c)$ is feasible, then the induced update sequence $R_B = (r_1, \dots, r_{c+2})$

is feasible.

Proof.

TODO

(1)

We need a new name for the set of all York. Also, notation $U(B)$.